[3 points] Fill in the blanks in the following to fulfill the comments.
 import java.util.regex.Pattern;

```
import java.util.regex.Matcher;
class Args {
   /** A pattern that matches a time containing minutes and hours on a 24-hour
    * clock. Minutes must be two-digit decimal numerals 00 through
    * 59, and hours must consist of one or two digits (0 through
    * 23). There must be two captured groups: one for hours and one
    * for minutes. Time 0:00 (or 00:00) denotes midnight. Times such as
    * 24:00, 25:01, 12:60, or 002:22 are not valid.
   static final Pattern PATN =
      Pattern.compile("______
                                                       _____");
   /** Prints the number of minutes after midnight denoted by TEXT, a
    * time in the format described by PATN.
   public static void printMinutes(String text) {
      Matcher matcher = PATN.matcher(text);
      matcher.matches();
      int hours = _____;
      int minutes = ______;
      System.out.println(_______);
   }
}
```

[3 points]

In the following, assume in each case that y is a non-negative Java int variable.

a. What can you say about y if the following expression is true?

y == (((y & 0xaaaaaaaa) >> 1) | (y & 0xaaaaaaaa))
(Choose one)

O y == 0.
O y is odd.
O y is divisible by 5.
O y is divisible by 3.
O y is divisible by 10.
O None of the above.
O All of the above.

- b. Does the answer to (a) change if we change >> to >>>? ○ Yes. ○ No.
- c. For what ranges of y is the following true?

$$y * 144 == ((y << 4) ^ (y << 7))$$

(Pick the largest range that applies.)

- 0 <= y < 16.</pre>
 0 <= y <= 16.</p>
- 0 <= y <= 8.
- 0 <= y <= 4.
- All non-negative integers less than 2²⁴.

3. [4 points] Fill in the skeleton below to fulfill its comment. You may use any of the classes in package java.util.*. Do not introduce any arrays. Your solution must take $\Theta(N)$ time, where N is the length of the input list. You may make no assumptions about the dynamic type of the input list, and therefore **you cannot assume that** .get runs in constant time, and will need an iterator. You need not use all lines of the skeleton.

```
import java.util.*;
class tri {
   /** Non-destructively return the list of lists of values
       [[a[0]], [a[1], a[2]], [a[3], a[4], a[5]], ...],
       where a[i] denotes the value at index i in A. Each list gets
       longer by one, except that the last should only be as long
       as needed to hold the remaining values of A.
      For example, given the input [ 1, 2, 3, 4, 5, 6, 7, 8 ], the
       result would be [[1], [2, 3], [4, 5, 6], [7, 8]]. */
    static List < List < Integer >> triangle (List < Integer > A) {
       _____ result = new
       for (int k = 1; _____; k += 1) {
       7
       return result;
   }
}
```

4.	[8 points] Fill in the appropriate bubbles for the following.
a.	Which of the following can be caused by providing a heuristic for A* search that underestimates the distance to the target node? Mark all that apply.
	 □ Cause the search to fail to reach the target. □ Cause the search to find a path to the target that is not the shortest. □ Cause the search to consider more nodes than necessary. □ Force the search into an infinite loop. □ None of the above.
b.	Which of the following can be caused by providing a heuristic for A^* search that overestimates the distance to the target node? Mark all that apply.
	 □ Cause the search to fail to reach the target. □ Cause the search to find a path to the target that is not the shortest. □ Cause the search to consider more nodes than necessary. □ Force the search into an infinite loop. □ None of the above.
C.	Consider a game-tree search that starts from a position where the maximizing player is to move and searches down $2N$ moves (N for each player, so that it is again the maximizing player's turn in the bottom positions). Suppose that during the search, the alpha-beta criteria causes the maximizing player to prune a certain branch. If we repeat the search, going two moves deeper (one for each side, for a total $2N + 2$), must that same branch be pruned? Choose one answer.
	 ○ Yes; extending the search will yield results consistent with the first search. ○ Yes, as long as the deeper search does not discover a win at level 2N + 2. ○ No; a position at level 2N (where the search stopped before) may have increased its value. ○ No; a position at level 2N may have decreased its value. ○ No, unless the deeper search discovers a win at level 2N + 2.
d.	Consider a complete binary max-heap with $2^n-1>16$ values, all distinct, stored in the usual way in an array, \mathbb{H} in positions $\mathbb{H}[1\\ 2^n-1]$. Where can the third-smallest value be? Choose the best answer.
	○ Anywhere in H except H[1]. ○ Anywhere in H[2^{n-1} $2^n - 1$]. ○ Anywhere in H[2^{n-2} $2^n - 1$]. ○ Anywhere in H[2^{n-3} $2^n - 1$]. ○ Anywhere in H[2^{n-1} $2^{n-1} + 2$]. ○ Anywhere in H[2^{n-2} $2^{n-1} - 1$]. ○ Anywhere in H[2^{n-3} $2^{n-2} - 1$]. ○ Anywhere in H[2^{n-3} $2^{n-2} - 1$].

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e.	A student creates a hashing function $h(x)$ for integers in an externally chained table with 1000000 buckets that returns (int) (1000000 * y), where y is the fifth value generated by the .nextDouble method on a Random object that has been seeded with the value x. Is this likely to work reasonably well? Choose the best answer.
	\bigcirc Yes, because all the values of h will be distinct for distinct inputs. \bigcirc Yes, with relatively high probability, distinct values of x will hash to different values.
	\bigcirc No, because a hashing function must not be non-deterministic. \bigcirc No, because for some values of x , the random-number generator will have a short period.
	O No, because of patterns in the numbers generated by the generator.
f.	If we apply LZW encoding to encode the string "aaaa", where 'a' is repeated M times, how many codewords will be in the output?
	\bigcirc About $M/2$. \bigcirc About \sqrt{M} . \bigcirc About M . \bigcirc About $\lg M$.
g.	Consider implementing a deque (double-ended queue) with a linked list versus using the best representation you can think of in an array. Which of the following are true for an arbitrary, sufficiently long mix of N operations (insertions in front, insertions at the end, deletions from the front, deletions from the end) when comparing times from the linked-list representation to the times from the array representation? Choose all that apply.
	☐ The times required for executing the entire sequence will be within a constant factor of each other.
	\square The times spent performing the first $M < N$ operations will be within a constant
	factor of each other. The times spent on any given insertion will be within a constant factor of each
	other.
	☐ The times spent on any given deletion will be within a constant factor of each other.

☐ The time spent on a particular deletion may be much larger for the array repre-

sentation because of having to move the remaining elements.

5. [6 points] In the following questions, notations such as $A \subseteq B$ mean that every function in the set of functions A is also in the set of functions B. Likewise, A = B means that A and B are the same set of functions, and $A \cup B$ is the union of the sets of functions A and B. Fill in the appropriate bubbles.

```
a. True or false: \Theta(f(N)) \subseteq O(f(N)).
   O True O False
b. True or false: \Omega(f(N)) \cup O(f(N)) includes all functions of N.
   ○ True ○ False
c. True or false: 1/2^N \in \Theta(1).
   ○ True ○ False
d. What is the worst-case running time of the call funcD(N, true) as a function of N?
   Assume p takes constant time.
   public void funcD(int n, boolean which) {
        if (n < 1) {
              doSomething();
        } else if (which && p(n)) {
              funcD(n / 2, true);
              funcD(n, false);
        } else {
              funcD(n - 1, false);
        }
   }
   \bigcirc \ \Theta(1) \ \bigcirc \ \Theta(\lg N) \ \bigcirc \ \Theta(N) \ \bigcirc \ \Theta(N \lg N) \ \bigcirc \ \Theta(N^2) \ \bigcirc \ \Theta(N^3) \ \bigcirc \ \Theta(2^N)
```

e. What is the worst-case running time of the call funcE(N) as a function of N? Assume that the function h takes constant time.

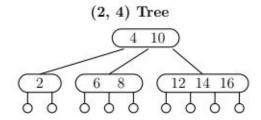
```
public void funcE(int n) { for (int i = n; i > 0; i /= 2) { for (int j = i; j > 0; j -= 1) { h(i, j, n); } } } }  \Theta(1) \bigcirc \Theta(\lg N) \bigcirc \Theta(N) \bigcirc \Theta(N \lg N) \bigcirc \Theta(N^2) \bigcirc \Theta(N^3) \bigcirc \Theta(2^N)
```

f. What is the worst-case running time of the call funcF(N, M) as a function of N? Assume calls to h take constant time and that M > 1 is a constant.

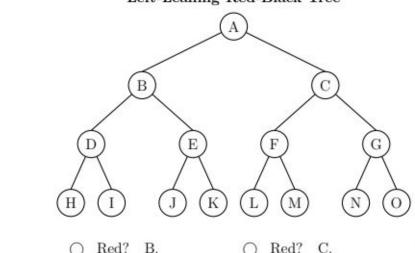
6. [1 point] What commercial product resulted from an accident with a batch of photosensitive glass?

7. [5 points]

a. Fill in the unique left-leaning red-black tree corresponding to the (2, 4) tree below. For each lettered node, indicate the key it contains and indicate whether it is red. Not all lettered nodes are actually needed. Leaving one blank indicates that it is actually an empty (null) node.



Left-Leaning Red-Black Tree



A	() Red?	В.	O Red? C	() Red
D	_ ○ Red?	E	O Red? F.	○ Red
G	_ Red?	Н.	O Red? I.	O Red
J.	2000 Company (1990)	K.	O Red? L.	O Red
M.	O Red?	N	O Red? O.	O Red

More parts on the following page.

- b. Consider a new kind a balanced tree: a red-green-black tree. This is a binary search tree whose nodes can be colored either black, red, or green, according to the following constraints, of which (a), (b), and (c) are the same as for red-black trees:
 - a. The root is black, and empty nodes are considered black as well.
 - b. The number of black nodes in any path from the root to the leaves is the same.
 - c. A black node may have only black and red children.
 - d. A red node may have only black and green children.
 - e. A green node may have only black children.

Why or why not?

 Except for null nodes and possibly the root, all black nodes must have two red children.

Like red-black trees, these trees correspond to (i.e., can represent) a kind of B-tree.

	Which kind?
	\bigcirc (2, 5) trees \bigcirc (3, 6) trees \bigcirc (3, 7) trees \bigcirc (4,8) trees.
c.	Assuming a search tree meets the constraints in part (b), will it be sufficiently balanced that searching for a value will require $O(\lg N)$ time for a tree with N keys?
	○ True ○ False

8. [6 points] Below you will find some intermediate steps in performing various sorting algorithms on the same input list. The steps do not necessarily represent consecutive steps in the algorithm (that is, many steps are missing), but they are in the correct sequence. For each of them, select the algorithm (by filling in the appropriate bubble) from among the following choices: insertion sort, straight selection sort, mergesort, quicksort, heapsort, LSD radix and MSD radix sort. For quicksort, the pivot is the median value of the first, last, and middle element (specifically, the index of the middle element is the average of the indices of the left and right elements, rounded down). Algorithms may appear twice.

In all these cases, the final step of the algorithm will be the sorted sequence:

43, 62, 191, 193, 265, 414, 482, 523, 592, 615, 674, 759, 834, 894, 920 but it might not be shown.

Input List:

```
674, 193, 523, 482, 920, 834, 191, 592, 615, 43, 265, 414, 759, 894,
               O Select
                         Merge
                                   O Quick
                                             ( Heap
                                                      O LSD

 a. () Insert.

                                                               \bigcirc MSD
    674, 193, 523, 482, 920, 834, 191, 592, 615, 43, 265, 414, 759, 894,
     43, 62, 193, 191, 265, 482, 414, 523, 592, 674, 615, 759, 834, 894, 920
         62, 191, 193, 265, 482, 414, 523, 592, 674, 615, 759, 834, 894, 920
     43, 62, 191, 193, 265, 414, 482, 523, 592, 615, 674, 759, 834, 894, 920
  b. O Insert.

    Select

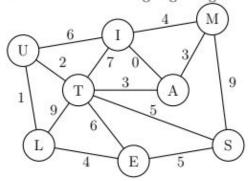
                         Merge
                                   O Quick
                                             ( Heap
                                                      O LSD
                                                               O MSD
    674, 193, 523, 482, 920, 834, 191, 592, 615, 43, 265, 414, 759, 894,
     62, 193, 523, 482, 191, 43, 265, 414, 592, 834, 920, 674, 759, 894, 615
     62, 193, 523, 482, 191, 43, 265, 414, 592, 615, 674, 834, 759, 894, 920
     62, 193, 523, 482, 191, 43, 265, 414, 592, 615, 674, 759, 834, 894, 920
     62, 193, 523, 482, 191, 43, 265, 414, 592, 615, 674, 759, 834, 894, 920
     62, 193, 191, 43, 265, 414, 523, 482, 592, 615, 674, 759, 834, 894, 920
  c. O Insert.
               O Select
                        Merge
                                   O Quick
                                            O Heap
                                                      O LSD
                                                              O MSD
    674, 193, 523, 482, 920, 834, 191, 592, 615, 43, 265, 414, 759, 894,
     62, 674, 894, 615, 482, 759, 834, 193, 592, 43, 265, 414, 523, 191, 920
     62, 674, 834, 615, 482, 759, 191, 193, 592, 43, 265, 414, 523, 894, 920
     62, 674, 759, 615, 482, 523, 191, 193, 592, 43, 265, 414, 834, 894, 920
     62, 674, 523, 615, 482, 414, 191, 193, 592, 43, 265, 759, 834, 894, 920
    265, 615, 523, 592, 482, 414, 191, 193, 62, 43, 674, 759, 834, 894, 920
```

Continued on next page.

d.	d. \bigcirc Insert.		\bigcirc S	elect	\bigcirc N	1erge	\bigcirc (Quick	\bigcirc I	Ieap	\bigcirc L	SD	O MS	SD	
	674,	193,	523,	482,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	674,	523,	482,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	482,	523,	674,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	482,	523,	674,	834,	920,	191,	592,	615,	43,	265,	414,	759,	894,	62
	191,	193,	482,	523,	674,	834,	920,	592,	615,	43,	265,	414,	759,	894,	62
	43,	191,	193,	265,	414,	482,	523,	592,	615,	674,	834,	920,	759,	894,	62
e.	() In	sert.	\bigcirc S	elect	\bigcirc V	1erge	\bigcirc (Quick	\bigcirc I	Ieap	\bigcirc L	SD	O MS	D	
	674,	193,	523,	482,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	920,	191,	482,	592,	62,	193,	523,	43,	674,	834,	414,	894,	615,	265,	759
	414,	615,	920,	523,	834,	43,	759,	62,	265,	674,	482,	191,	592,	193,	894
	43,	62,	191,	193,	265,	414,	482,	523,	592,	615,	674,	759,	834,	894,	920
f.	() In	sert.	\bigcirc S	elect	\bigcirc N	lerge [\bigcirc (Quick	\bigcirc I	Heap	O L	SD	O MS	SD	
	674,	193,	523,	482,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	674,	523,	482,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	674,	482,	523,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	482,	523,	674,	920,	834,	191,	592,	615,	43,	265,	414,	759,	894,	62
	193,	482,	523,	674,	191,	592,	834,	920,	615,	43,	265,	414,	759,	894,	62
	191.	193.	482.	523.	592.	674,	834.	920.	615.	43.	265.	414.	759,	894.	62

9. [5 points]

Consider the following edge-weighted undirected graph:



In the following questions, break ties in the algorithms in favor of vertices that come first in alphabetical order, or edges that are listed first.

- a. Run Prim's algorithm starting from node L. Indicate the vertices that are in the fringe in order of decreasing priority just after node T is removed from the fringe and processed. Break ties in priority by having the node that is first alphabetically have highest priority.

 - O A, E, S, I

 - A, E, S
 - O A, I, S, M
- b. What are the edges in the shortest-path tree from L? Mark all that apply.

 - $\square \ (I, \ M) \ \square \ (I, \ U) \ \square \ (I, \ T) \ \square \ (L, \ T) \ \square \ (L, \ U) \ \square \ (M, \ S)$
 - \square (S, T) \square (T, U)
- c. Dijkstra's algorithm assumes that all edge weights are non-negative. To illustrate why this is so, suppose we run Dijkstra's algorithm from node L, and terminate the algorithm as soon as it removes node I from the fringe. Which edge's weight could we change to -6 to produce an *incorrect* shortest path to I? Assume that nodes are never updated after being removed from the fringe. Choose one.
 - $\bigcirc \ \, (A,\ I) \ \, \bigcirc \ \, (A,\ M) \ \, \bigcirc \ \, (A,\ T) \ \, \bigcirc \ \, (E,\ L) \ \, \bigcirc \ \, (I,\ U) \ \, \bigcirc \ \, (I,\ T)$
 - \bigcirc (L, T) \bigcirc (M, S) \bigcirc (S, T)

d. Suppose that when running Kruskal's algorithm on the graph above, we fail to check whether an edge joins a subtree to itself. What is the first edge to be added to the resulting tree erroneously? Choose one.

 $\bigcirc \ \, (A,\ I) \ \, \bigcirc \ \, (A,\ M) \ \, \bigcirc \ \, (A,\ T) \ \, \bigcirc \ \, (E,\ L) \ \, \bigcirc \ \, (E,\ S) \ \, \bigcirc \ \, (E,\ T)$

 $\bigcirc \ (I,\ M) \ \bigcirc \ (I,\ U) \ \bigcirc \ (I,\ T) \ \bigcirc \ (L,\ T) \ \bigcirc \ (L,\ U) \ \bigcirc \ (M,\ S)$

○ (S, T) ○ (T, U)

10. [6 points] You are given the following interface to a general directed graph. The vertices of these graphs are simply positive integers in some contiguous range starting at 1.

```
public interface DiGraph {
    /** Returns the number of vertices in this DiGraph. The vertices
    * are represented by integers 1..size(). */
    int size();

    /** Returns a list of the successor vertices of V in this
        * graph. */
    List<Integer> successors(int v);

    /** Returns a list of the predecessor vertices of V in this
        * graph. */
    List<Integer> predecessors(int v);

    /** Returns a view of this graph in which all edges are
        * reversed. Changes to this graph will be reflected in the
        * result and vice-versa. */
    public DiGraph reverse();
}
```

- a. Implement the class Traverser. The public traverse method of Traverser performs a depth-first traversal of a DiGraph by performing a depth-first traversal starting from each vertex in turn, in order, and traversing only vertices that have not been processed in the previous traversals. As it does so, it calls the Traverser's previsit method on the vertices it traverses in preorder, and calls the postvisit method on the vertices it traverses in postorder. You need not use all the spaces provided.
- b. Assuming that the implementation of Traverser is correct, implement the method Algs.topoSort, using the Traverser class. You need not use all the spaces provided.

publ		tract class Traverser {
		tion to take on visiting vertex V in G in preorder. */
		ted abstract void previsit(DiGraph G, int v);
		tion to take on visiting vertex V in G in postorder. */
	protect	ted abstract void postvisit(DiGraph G, int v);
		rform a depth-first traversal of G, calling .previsit and ostvisit on all traversed nodes, as described in part (a). */
		void traverse(DiGraph G) {
		olean[] marks = new boolean[G.size() + 1];
	for	r () {
	}	
	}	
		rform a depth-first traversal of G starting from vertex V
		d traversing only vertices w such that marked[w] is false.
		t marked[w] to true for all visited vertices, and call
		revisit on all traversed nodes in preorder and .postvisit on all aversed nodes in postorder, as described in part (a). */
		e void traverse(DiGraph G, int v, boolean[] marks) {
	_	() {
		\/
		for () {
		}
	}	
}	}	Page 17 / 19 — 🤁 🕂

<pre>import java.util.*;</pre>
<pre>public class Algs {</pre>
<pre>/** Return a list of the vertices of G, sorted topologically. Assumes * that G is acyclic. */ public static List<integer> topoSort(DiGraph G) { TopoTraverser topo = new TopoTraverser();</integer></pre>
return}
<pre>private static class TopoTraverser extends Traverser {</pre>
}
}