

1. **(Sanity Check!)** Derive Chebyshev's inequality using Markov's inequality for random variable X .

Answer: We're interested in the probability $\Pr(|X - \mathbf{E}[X]| \geq k) = \Pr((X - \mathbf{E}[X])^2 \geq k^2)$. We simply apply Markov's inequality and we find that $\Pr((X - \mathbf{E}[X])^2 \geq k^2) \leq \text{Var}[X]/k^2$.

2. **(Balls and Bins Again)** For this problem we toss m balls into n bins.

- (a) What is the expected number of collisions?

Answer: Let X be the expected number of collisions and X_{ij} the indicator for the event that balls i and j collide. Then $X = \sum_{1 \leq i < j \leq m} X_{ij}$. Since $\mathbf{E}[X_{ij}] = 1/n$, we know that $\mathbf{E}[X] = \binom{m}{2} \times \frac{1}{n}$.

- (b) Now, let's define X to be the number of collisions. At what threshold of collisions c can we ensure that the probability of having more than c collisions is less than $1/2n$?

Answer: By Markov's inequality, $\Pr[X \geq c] \leq \mathbf{E}[X]/c$. Since we know $\mathbf{E}[X]$ from the previous part we simply plug it in: $\Pr[X \geq c] \leq m(m-1)/2nc$, and so we set $m(m-1)/2nc = 1/2n$ and find that $c = m(m-1)$.

3. **(Coupon Collector)** There are n baseball cards Jonny is trying to collect! Each day, Jonny buys a cereal box and dumps out all the cereal, searching for the baseball card. Each box contains each of the n cards with equal probability.

- (a) What's the expected number of days until Jonny gets one unique card? Two unique cards?

Answer: $1, 1 + \frac{n}{n-1}$.

- (b) Develop a formula for the expected number of days until Jonny gets k unique cards. On average, how many days will it take Jonny to collect all the baseball cards? Approximate your sum with an integral.

Answer: $\sum_{i=1}^k \frac{n}{n-i+1}$. On average, it will take $n \sum_{i=1}^n \frac{1}{i}$ days to collect all the baseball cards. This is approximately $n \ln n$ days.

- (c) Jonny and his sister, Jill, are working together to collect the cards. Each day, they both buy cereal boxes. On average, how many days will it take to collect k unique baseball cards? How many days until they collect them all? Be exact.

Answer: $\lceil 0.5 \sum_{i=1}^k \frac{n}{n-i+1} \rceil$. On average, it will take $\lceil 0.5n \sum_{i=1}^n \frac{1}{i} \rceil$ days to collect all the baseball cards. (Note: $\lceil x \rceil \neq \lfloor x \rfloor + 1$ when x is an integer.)

4. **(Coin flips)**

- (a) Suppose we flip a fair coin n times and we wish to understand the probability that we get at least $3n/4$ heads. Use Markov's inequality to come up with an upper bound for this probability.

Answer: Let X be a random variable for the number of heads. Let X_i be an indicator for the event that the i -th flip is heads. Since X_i is an indicator, $E[X_i] = \Pr(X_i = 1) = 1/2$. Since $X = \sum_{i=1}^n X_i$,

$$E[X] = E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i] = \frac{n}{2}$$

We want to bound the probability that $X \geq 3n/4$. Since X is nonnegative, we can use Markov's inequality to get

$$\Pr\left(X \geq \frac{3n}{4}\right) \leq \frac{E[X]}{\frac{3n}{4}} = \frac{\frac{n}{2}}{\frac{3n}{4}} = \frac{2}{3}$$

- (b) Use Markov's inequality to come up with a similar upper bound on the probability that the number of heads is at least n .

Answer: This time, we want to bound the probability that $X \geq n$. By Markov's inequality,

$$\Pr(X \geq n) \leq \frac{E[X]}{n} = \frac{\frac{n}{2}}{n} = \frac{1}{2}$$

- (c) Find the true probability that there are at least n heads in a sequence of n fair coin flips. Is the bound you derived in the previous part tight?

Answer: Since X can't be greater than n ,

$$\Pr(X \geq n) = \Pr(X = n) = \left(\frac{1}{2}\right)^n$$

So we can see that Markov's inequality gives a very loose bound; it bounds $\Pr(X \geq n)$ by a constant, whereas in reality this probability decreases exponentially as n increases.

5. (More coin flips)

- (a) Suppose we flip a biased coin 100 times and X is the number of heads we get. We know that $\text{Var}[X] = 16$. What are the possible values for the expected value of X ?

Answer: This is a binomial random variable, so we can use the properties we derived from last discussion. First of all, we know that the $\text{Var}[X] = np(1-p)$. This means that $16 = 100p(1-p)$. Secondly, we know that $E[X] = np$ which tells us that $p = E[X]/100$. We plug it into our first equation and find that

$$16 = 100(E[X]/100)(1 - E[X]/100)$$

$$0 = E[X]^2 - 100E[X] + 1600$$

$$0 = (E[X] - 20)(E[X] - 80)$$

So we know that $E[X] = 20$ or $E[X] = 80$.

- (b) Now suppose $E[X] = 20$. Use Chebyshev's inequality to derive an upper bound on $\Pr[X \geq 40]$.

Answer: According to Chebyshev's,

$$\Pr(|X - 20| \geq 20) \leq \text{Var}[X]/20^2$$

We simply plug in the variance and we find that

$$\Pr[X \geq 40] \leq \Pr(|X - 20| \geq 20) \leq 16/400 = 1/25$$

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