# Expressions and Statements

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### 1 Introduction

The previous handout considered a language which consisted only of expressions, no statements. In this handout, we use fundamentally the same language, but use statements instead of expressions for let and assign.

### 1.1 Language Used

We'll first define the abstract syntax for our language, shown below:

```
x \in \mathit{Variable} \qquad i \in \mathit{Integer} \tau \in \mathit{Type} ::= \mathsf{int} \mid \mathsf{bool} e \in \mathit{Exp} ::= x \mid i \mid \mathsf{true} \mid \mathsf{false} \mid e_1 \&\& e_2 \mid e_1 + e_2 \mid e_1 < e_2 s \in \mathit{Stmt} ::= \mathsf{let} \ x : \tau = e \mid x = e p \in \mathit{Program} ::= s \mid s \ p
```

Compared to the last handout, there are two notable changes:

- let and assignment have been made into statements.
- A program is now explicitly defined as either a statement, or a statement followed by another program. Phrased more simply, a program consists of one or more statements.

# 2 Type System

Now that we have the syntax defined, we can define the language's type system. Formerly, this entailed the definition of a single set of rules which operated over the language's expressions. However, since we now have both statements and expressions, we need to define *two* sets of rules: one for expressions and one for statements. Both sets of rules will need type environments (expressions have variable access, and all statements currently involve variables), so we'll define that first:

$$\Gamma \in \mathit{TypeEnv} = \mathit{Variable} \to \mathit{Type}$$

Note this is the same type environment definition from last time; this is still a mapping of variables to types. From here, we define the rules for expressions and statements in the subsections below.

#### 2.1 Rules for Expressions

#### 2.2 Rules for Statements