

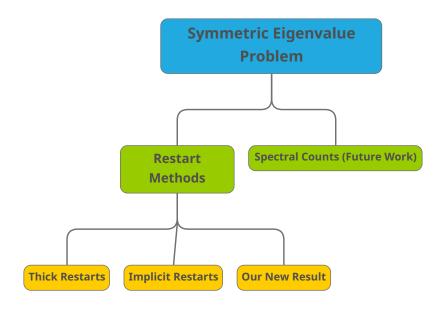


Symmetric Eigenvalue Problems

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Symmetric Eigenvalue Problem

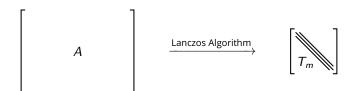
Given: Large, sparse, symmetric $A \in \mathbb{R}^{n \times n}$. Compute some approximate eigenpairs (λ, x) . That is, $Ax = \lambda x$.

If *n* is small, then *A* can be "tridiagonalized" to *T*. An $O(n^3)$ process.

$$A = \begin{bmatrix} \times \times \times \times \times \\ \times \times \times \times \times \\ \times \times \times \times \times \end{bmatrix} \longrightarrow T = Q^T A Q = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix} \longrightarrow \text{ Eigenpairs of } T$$

We assume n is too large for this to be done.

Instead, we'll "reduce" A to a **smaller** tridiagonal matrix T_m .



Lanczos Algorithm

Given symmetric $A \in \mathbb{R}^{n \times n}$, $m \ll n$, and unit vector $p_1 \in \mathbb{R}^n$, generate

$$AP_m = P_m T_m + f e_m^T.$$

$$[P, T, f] = \text{Lanczos}(A, p_1, m)$$
1: **for** $j = 1, 2, ..., m$ **do**
2: $f = Ap_j$
3: **if** $j > 1$ **then** $f = f - p_{j-1}\beta_{j-1}$
4: $\alpha_j = f^T p_j$
5: $f = f - p_j \alpha_j$ $f = f - P_j (P_j^T f)$
6: **if** $j < m$ **then** $\beta_j = ||f||$ $p_{j+1} = f/\beta_j$

$$\begin{bmatrix} & & & \\ & A & & \\ & & & \end{bmatrix} \begin{bmatrix} & & \\ & P_m & \\ & & \end{bmatrix} \begin{bmatrix} & & \\ & & \end{bmatrix} + \begin{bmatrix} & & \\ & & \\ & & \end{bmatrix}$$

Ritz Pairs of A

$$\left[\begin{array}{c} A \\ \end{array}\right] \left[\begin{array}{c} P_m \\ \end{array}\right] = \left[\begin{array}{c} P_m \\ \end{array}\right] \left[\begin{array}{c} T_m \\ \end{array}\right] + \left[\begin{array}{c} T_m \\ \end{array}\right]$$

If
$$(\theta_j, y_j)$$
 is an eigenpair of T_m $(T_m y_j = \theta_j y_j)$ then Ritz pairs of A are (θ_j, x_j) , where $x_j = P_m y_j$.

The residual is

$$||Ax_j - \theta_j x_j|| = ||AP_m y_j - \theta_j P_m y_j|| = ||fe_m^\mathsf{T} y_j|| = ||f|| \cdot |e_m^\mathsf{T} y_j|$$

What if residual is large? *Restart* Lanczos, "keeping" desired Ritz vectors.

Ritz Pairs of A

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What if residual is large? *Restart* Lanczos, "keeping" desired Ritz vectors.

Thick Restarts

- MATLAB post-2016
- Krylov-Schur [Ste02], [Wu&Sim00]

Implicit Restarts

- MATLAB pre-2016 (Octave, ARPACK)
- IRA / IRL [Sor92] / [Cal&Sor94]

TRL and IRL

Until convergence ...

- **1** Generate *m*-Lanczos factorization: $AP_m = P_m T_m + f e_m^T$.
- **2** Compute Ritz pairs: (θ_j, x_j) $x_j = P_m y_j$ and $T_m y_j = \theta_j y_j$.
- 3 Choose k < m "desired" Ritz pairs.
- "Restart" with these pairs; build a new m-Lanczos factorization.

Thick Restarting Lanczos

$$Ax_{j} - \theta_{j}x_{j} = fe_{m}^{T}y_{j} \quad \Rightarrow \quad Ax_{j} = \underbrace{\beta_{m}p_{m+1}}_{f}e_{m}^{T}y_{j} + \theta_{j}x_{j} = \underbrace{\overline{\beta}_{j}}_{\beta_{m}e_{m}^{T}y_{j}}p_{m+1} + \theta_{j}x_{j}$$

Let
$$\overline{P}_k = [x_1 \dots x_k]$$
 and $\overline{P}_{k+1} = [x_1 \dots x_k p_{m+1}]$.

One can show that

$$\overline{P}_{k+1}^{T}A\overline{P}_{k+1} := \overline{T}_{k+1} = \begin{bmatrix} \theta_1 & & & \overline{\beta}_1 \\ & \theta_2 & & \overline{\beta}_2 \\ & & \ddots & \vdots \\ \overline{\beta}_1 & \overline{\beta}_2 & \dots & \overline{\beta}_k & \overline{\alpha}_{k+1} \end{bmatrix}$$

where $\alpha_{k+1} = p_{m+1}^T A p_{m+1}$.

Thick Restarting Lanczos

Now the "restart" makes sense:

Start Lanczos with vector p_{m+1} , and do m - (k+1) steps.

The first k+1 chunk is $A\overline{P}_{k+1} = \overline{P}_{k+1} \overline{T}_{k+1} + \beta_{k+1} p_{k+2} e_{k+1}^T$.

Thick Restarting Lanczos

Now the "restart" makes sense:

Start Lanczos with vector p_{m+1} , and do m - (k+1) steps.

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Question: Can the "arrowhead" \overline{T}_{k+1} be "tridiagonalized" efficiently?

Is there a benefit?

Arrow → Tridiagonal Conversion

$$\left[\begin{array}{c} A \end{array} \right] \left[\begin{array}{c} \overline{P}_m \end{array} \right] \ = \ \left[\begin{array}{c} \overline{P}_m \end{array} \right] \left[\begin{array}{c} \overline{T}_m \end{array} \right] + \left[\begin{array}{c} \overline{T}_m \end{array} \right]$$

We showed [Bag,Mon,Per25+] that we can tridiagonalize \overline{T}_{k+1} via

$$\widetilde{T}_{k+1} = \overline{Q}_{k+1}^T \overline{T}_{k+1} \overline{Q}_{k+1}$$
 $\overline{Q}_{k+1} = \begin{bmatrix} 0 & 0 \\ \overline{Q}_k & \vdots \\ 0 & \dots & 0 & 1 \end{bmatrix}$

$$\left[\begin{array}{c} A \end{array}\right] \left[\begin{array}{c} \overline{P}_m \overline{Q}_m \end{array}\right] = \left[\begin{array}{c} \overline{P}_m \overline{Q}_m \end{array}\right] \left[\begin{array}{c} \overline{T}_m \end{array}\right] + \left[\begin{array}{c} \overline{f} \\ \overline{f} \end{array}\right]$$

Equating the first k columns: $A\overline{P}_k\overline{Q}_k = \overline{P}_k\overline{Q}_k\widetilde{T}_k + \widetilde{t}_{k+1,k}p_{m+1}e_k^T$

Another Restart Method — IRL

TRL Summary: Restart with *k* Ritz vectors.

Another Variation: Can we choose a new (better) initial vector \hat{p}_1 ?

Goal:
$$\widehat{p}_1 \approx \underbrace{c_1 x_1 + \dots + c_k x_k}_{\text{large } c_i} + \underbrace{c_{k+1} x_{k+1} + \dots + c_n x_n}_{\text{small } c_i}$$

IRL applies a polynomial filter to update the starting vector.

$$\widehat{p}_1 = q(A)p_1 = c_1q(\lambda_1)x_1 + \dots + c_kq(\lambda_k)x_k + c_{k+1}q(\lambda_{k+1})x_{k+1} + \dots + c_nq(\lambda_n)x_n$$

Which polynomial q?

Applying Shifts

Say (μ, x) is an approximate eigenpair of A. Then

$$x^T(A-\mu I)p_1 \approx 0.$$

So $(A - \mu I)p_1$ has (approx.) no components in direction of x.

Choose μ to be a "shift" from the undesired part of the spectrum.

Result: A vector $(A - \mu I)p_1$ enriched in direction of desired eigenvectors.

However, notice that

$$(A - \mu I)P_m = P_m(T_m - \mu I) + fe_m^T.$$

So, we shift T_m instead of A.

Applying Shifts

Given the undesired eigenvalues $\theta_{k+1}, \ldots, \theta_m$ of T_m , we consider

$$q(T_m) = (T_m - \theta_{k+1}I) \dots (T_m - \theta_mI)$$

(which is not formed explicitly!).

Use the implicit *QR* algorithm with shifts $\theta_{k+1}, \dots, \theta_m$ to obtain

$$T_m^+ := Q_m^+ T_m Q_m^+, \qquad q(T_m) = Q_m^+ R_m^+.$$

Choose a shift μ .

Create a rotation matrix
$$\widetilde{G}_1 = \begin{bmatrix} c & -s \\ s & c \end{bmatrix}$$
 such that $\widetilde{G}_1^T \begin{bmatrix} t_{11} - \mu \\ t_{21} \end{bmatrix} = \begin{bmatrix} \star \\ 0 \end{bmatrix}$.

Choose a shift μ .

Create a rotation matrix
$$\widetilde{G}_1 = \begin{bmatrix} c & -s \\ s & c \end{bmatrix}$$
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Action of G2

Choose a shift μ .

Create a rotation matrix
$$\widetilde{G}_1 = \left[\begin{array}{cc} c & -s \\ s & c \end{array} \right]$$
 such that $\widetilde{G}_1^T \left[\begin{array}{cc} t_{11} - \mu \\ t_{21} \end{array} \right] = \left[\begin{array}{cc} \star \\ 0 \end{array} \right]$.

$$\begin{bmatrix}
* & \times \\
\times & \times & \times \\
& \times & \times
\end{bmatrix}$$
Action of G_2

$$\begin{bmatrix}
* & \times & & & \\
\times & \times & \times & \bullet & \\
& \times & \times & \times & \bullet \\
& & * & * & *
\end{bmatrix}$$

$$\begin{bmatrix}
* & * & & & \\
* & * & \times & & \\
& & \times & \times & \times & \bullet \\
& & \times & \times & \times & \bullet \\
& & \times & \times & \times & \times
\end{bmatrix}$$
Action of G_2
Action of G_3
Action of G_4

So after one shift μ , we obtain

$$T_m^+ := Q_m^+ T_m Q_m^+, \qquad Q_m^+ = G_1 \dots G_{m-1}.$$

If μ is an (undesired) eigenvalue of T_m (an "exact shift"):

Then
$$t_{m,m}^+ = \mu$$
 and $t_{m,m-1}^+ = t_{m-1,m}^+ = 0$.

So after one shift μ , we obtain

$$T_m^+ := Q_m^+ T_m Q_m^+, \qquad Q_m^+ = G_1 \dots G_{m-1}.$$

If μ is an (undesired) eigenvalue of T_m (an "exact shift"):

Then
$$t_{m,m}^+ = \mu$$
 and $t_{m,m-1}^+ = t_{m-1,m}^+ = 0$.

Repeat for p := m - k exact shifts:

Mechanics of IRL

Now, restart Lanczos with f^+ and build to size m.

Revisiting the initial vector:

From
$$AP_m = P_m T_m + f e_m^T$$
 one can show $q(A)P_m e_1 = P_m \underbrace{q(T_m)}_{Q_m^+ R_m^+} e_1$.

Then
$$q(A)p_1 = P_m Q_m^+ r_{11}^+ e_1 \implies q(A)p_1 = r_{11}^+ \widehat{p}_1.$$

Therefore, $\widehat{p}_1 \propto q(A)p_1$ is a "polynomial filter" update to p_1 .

Which has stronger components in desired Ritz vectors.

Equivalence

TRL and IRL (exact shifts) are "mathematically equivalent".

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By considering the k-Lanczos factorizations ...

TRL
$$A\overline{P}_{k}\overline{Q}_{k} = \overline{P}_{k}\overline{Q}_{k}\widetilde{T}_{k} + \widetilde{t}_{k+1,k}p_{m+1}e_{k}^{T}$$
IRL
$$AP_{k}^{+} = P_{k}^{+}T_{k}^{+} + q_{m,k}^{+}p_{m+1}e_{k}^{T}$$

$$\begin{bmatrix} A \end{bmatrix} \begin{bmatrix} \overline{P}_{k}\overline{Q}_{k} \end{bmatrix} = \begin{bmatrix} \overline{P}_{k}\overline{Q}_{k} \end{bmatrix} \begin{bmatrix} \overline{T}_{k} \end{bmatrix} + \begin{bmatrix} \overline{T}_{k} \end{bmatrix} \begin{bmatrix} \overline{T}_{k} \end{bmatrix} + \begin{bmatrix} \overline{T}_{k} \end{bmatrix}$$

$$\begin{bmatrix} A \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} = \begin{bmatrix} P_{k}^{+} \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} + \begin{bmatrix} P_{k}^{+} \end{bmatrix}$$

Equivalence

TRL and IRL (exact shifts) are "mathematically equivalent".

By considering the k-Lanczos factorizations ...

TRL
$$A\overline{P}_{k}\overline{Q}_{k} = \overline{P}_{k}\overline{Q}_{k}\widetilde{T}_{k} + \widetilde{t}_{k+1,k}p_{m+1}e_{k}^{T}$$
IRL
$$AP_{k}^{+} = P_{k}^{+}T_{k}^{+} + q_{m,k}^{+}p_{m+1}e_{k}^{T}$$

$$\begin{bmatrix} A \end{bmatrix} \begin{bmatrix} \overline{P}_{k}\overline{Q}_{k} \end{bmatrix} = \begin{bmatrix} \overline{P}_{k}\overline{Q}_{k} \end{bmatrix} \begin{bmatrix} \overline{T}_{k} \end{bmatrix} + \begin{bmatrix} \overline{T}_{k} \end{bmatrix} \begin{bmatrix} \overline{T}_{k} \end{bmatrix} + \begin{bmatrix} \overline{T}_{k} \end{bmatrix}$$

$$\begin{bmatrix} A \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} = \begin{bmatrix} P_{k}^{+} \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} \begin{bmatrix} P_{k}^{+} \end{bmatrix} + \begin{bmatrix} P_{k}^{+} \end{bmatrix}$$

Our New Result:

$$\overline{P}_k \overline{Q}_k = \pm P_k^+, \qquad \widetilde{T}_k = \pm T_k^+, \qquad \overline{f} = \pm f^+.$$

$$\widetilde{T}_{\nu} = \pm T_{\nu}^{+}$$

$$\overline{f} = \pm f^+$$
.

In Summary ...

Our new method [Bag,Mon,Per25+] (arrowhead to tridiagonal) ...

- Maintains tridiagonal structure.
- Not sensitive to numerical round-off.
- **Equivalent** to a new method [Mas&Dor18].
- Provides new, **stronger** insight: TRL = IRL (not just equivalent!)
- Adapts to the bidiagonal case & singular value problem.
 A seminar presentation (and another paper!) to come ...

New Topic — Spectral Counts

Let $A \in \mathbb{R}^{n \times n}$ be a (large!) symmetric matrix and $[a, b] \subset \mathbb{R}$.

Question: How many eigenvalues of A are in [a, b]? c.f. [Nap16]

Our (Future) Interest: Accelerating algorithms with this count.

Potential Insight: Form the projection matrix P

$$P := \sum_{\substack{i \\ \lambda_i \in [a,b]}} u_i u_i^T$$

which satisfies $P^2 = P$, so eigenvalues of P are 0 and 1.

Key Result: trace(P) is the number of eigenvalues of A in [a, b].

Approximate trace(P)

Let h(t) be the indicator function

$$h(t) = \begin{cases} 1 & t \in [a, b] \\ 0 & \text{otherwise} \end{cases}$$

and consider its representation in the Chebyshev basis

$$h(t) \approx \sum_{j=0}^{\rho} \gamma_j T_j(t).$$

Linearly map eigenvalues of A and [a, b] into [-1, 1]. Then

$$P \approx h(A) \approx \sum_{j=0}^{p} \gamma_j T_j(A)$$

has eigenvalues of 1 if $\lambda_i \in [a, b]$ and eigenvalues of 0 otherwise.

Therefore: $trace(P) = \# eigenvalues of A in [a, b] \approx trace(h(A))$.

Key Theorem

Hutchinson's Theorem

Suppose each component of $v \in \mathbb{R}^n$ is drawn independently from a distribution with mean 0 and variance 1; that is, $\mathbb{E}[v] = 0$ and $\mathbb{E}[vv^T] = I_n$.

Then for any symmetric matrix A and matrix function f defined on A,

$$\mathbb{E}\big[v^T f(A)v\big] = \operatorname{trace} f(A).$$

In particular, we'll use $v \sim \text{unif}\{-1,1\}^n$, called Rademacher vectors (for the least variance estimator).

Therefore, $\operatorname{trace}(P) \approx \operatorname{trace}(h(A)) = \mathbb{E}[v^T h(A) v].$

Computing trace $(P) \approx \mathbb{E}[v^T h(A)v]$

From before ...

eigenvalues of A in
$$[a, b]$$
 = trace (P)
 \approx trace $(h(A))$
 $= \mathbb{E}[v^T h(A) v]$
 $= \mathbb{E}\left[\sum_{j=0}^p \gamma_j v^T T_j(A) v\right]$
 $= n \cdot \mathbb{E}\left[\sum_{j=0}^p \gamma_j v^T T_j(A) v\right]$ (normalize v)
 $\approx \frac{n}{n_v} \sum_{k=1}^n \sum_{j=0}^p \gamma_j v_k^T T_j(A) v_k$

Future Work

 $T_p(A)v$ is a (modest size) polynomial of A times v ...

Then $w = T_p(A)v$ has components in the eigenvectors of A from [a, b].

Future work: Optimizing existing algorithms by using w.

- Solving A(x + w) = b + Aw instead of Ax = b.
- Symmetric eigenvalue problems.
- Singular value thresholding (since $\sigma(A) = \sqrt{\lambda(A^T A)}$).

• ...

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Thank you!

Questions?

Lanczos Recurrence

Three term recurrence: Arnoldi gives $T_{ij} = p_i^T A p_j$, for $1 \le i \le j$.

If A is symmetric, then $T_{ij} = (Ap_i)^T p_j$.

For i < j-1, then $Ap_i \in \text{span}\{p_1, \dots, p_{j-1}\}$ and p_j is perpendicular to this span, so it must be that $(Ap_i)^T p_j = 0$.

Therefore, *T* has zero 2nd off-diagonals.

Ritz Extraction

Let \mathcal{K}_m be an m-dimensional subspace of \mathbb{R}^n .

We say (θ, x) , where $x \in \mathcal{K}_m$, is a *Ritz pair* of A if $Ax - \theta x \perp \mathcal{K}_m$.

Let $P_m \in \mathbb{R}^{n \times m}$ be an ONB for \mathcal{K}_m .

Then $\exists y \in \mathbb{R}^m$ such that $x = P_m y$.

Let $T_m = P_m^T A P_m \in \mathbb{R}^{m \times m}$.

Then (θ, x) is a Ritz pair iff (θ, y) is an eigenpair of T_m .

TRL Step

Let $\Theta_k = \operatorname{diag}(\theta_1, \dots, \theta_k)$ and $Y_k = Y_{:,1:k}$. Then $A\overline{P}_k = \overline{P}_k \Theta_k + \beta_m p_{m+1} e_m^T Y_k$, or equivalently $A\overline{P}_k = \overline{P}_{k+1} \overline{T}_{k+1,k}$. Orthogonalize Ap_{m+1} against \overline{P}_{k+1} :

$$r = (I - \overline{P}_{k+1} \overline{P}_{k+1}^{T}) A p_{m+1} = A p_{m+1} - [\overline{P}_{k} \quad p_{m+1}] \begin{bmatrix} \overline{P}_{k}^{T} \\ p_{m+1}^{T} \end{bmatrix} A p_{m+1}$$

$$= A p_{m+1} - \overline{P}_{k} \overline{P}_{k}^{T} A p_{m+1} - p_{m+1} p_{m+1}^{T} A p_{m+1}$$

$$= A p_{m+1} - \overline{P}_{k} (A \overline{P}_{k})^{T} p_{m+1} - p_{m+1} p_{m+1}^{T} A p_{m+1}$$

$$= A p_{m+1} - \overline{P}_{k} (\Theta_{k} \overline{P}_{k}^{T} p_{m+1}) - \overline{P}_{k} \beta_{m} Y_{k}^{T} e_{m} p_{m+1}^{T} p_{m+1} - p_{m+1} p_{m+1}^{T} A p_{m+1}$$

$$= A p_{m+1} - p_{m+1} p_{m+1}^{T} A p_{m+1} - \overline{P}_{k} \beta_{m} Y_{k}^{T} e_{m}.$$

Let
$$p_{k+2} = r/\|r\| = r/\beta_{k+1}$$
, and $\alpha_{k+1} = p_{m+1}^T A p_{m+1}$. Then

$$Ap_{m+1} = \beta_{k+1}p_{k+2} + \alpha_{k+1}p_{m+1} + \overline{P}_k\overline{\beta}$$

where $\overline{\beta}$ is a vector of $(\overline{\beta}_1, \dots, \overline{\beta}_k)$ and $\overline{\beta}_i = \beta_m y_i^T e_m$.

TRL Step

$$\begin{array}{rcl}
A\overline{P}_{k+1} & = & \left[\overline{P}_{k+1}\overline{T}_{k+1,k} \mid Ap_{m+1}\right] \\
& = & \left[\overline{P}_{k+1}\overline{T}_{k+1,k} \mid \beta_{k+1}p_{k+2} + \alpha_{k+1}p_{m+1} + \overline{P}_{k}\overline{\beta}\right] \\
& = & \left[\overline{P}_{k+1}\overline{T}_{k+1,k} \mid \overline{P}_{k+1}\left[\begin{array}{c}\overline{\beta} \\ \alpha_{k+1}\end{array}\right]\right] + \beta_{k+1}p_{k+2}e_{k+1}^{T} \\
& = & \overline{P}_{k+1}\overline{T}_{k+1} + \beta_{k+1}p_{k+2}e_{k+1}^{T}
\end{array}$$

which also shows that $\overline{P}_{k+1}^T A \overline{P}_{k+1} = \overline{T}_{k+1}$ by orthogonality.

TRL Truncation

From $A\overline{P}_{k+1} = \overline{P}_{k+1}\overline{T}_{k+1} + \beta_{k+1}p_{k+2}e_{k+1}^T$, then

$$A\overline{P}_{k+1}\overline{Q}_{k+1} = \overline{P}_{k+1}\overline{Q}_{k+1}\widetilde{T}_{k+1} + \beta_{k+1}p_{k+2}e_{k+1}^T\overline{Q}_{k+1}.$$

Note that $e_{k+1}^T \overline{Q}_{k+1} = e_{k+1}^T$ and the first k columns of $p_{k+2} e_{k+1}^T$ are zeros. Therefore,

$$A \begin{bmatrix} \overline{P}_{k} & p_{m+1} \end{bmatrix} \begin{bmatrix} \overline{Q}_{k} & 0 \\ 0 & 1 \end{bmatrix} = \underbrace{\begin{bmatrix} \overline{P}_{k} & p_{m+1} \end{bmatrix} \begin{bmatrix} \overline{Q}_{k} & 0 \\ 0 & 1 \end{bmatrix}}_{= \begin{bmatrix} \overline{P}_{k} \overline{Q}_{k} & p_{m+1} \\ 0 & 0 \end{bmatrix}} \begin{bmatrix} \widetilde{T}_{k} & \widetilde{t}_{k+1,k} e_{k} \\ \widetilde{t}_{k+1,k} e_{k}^{T} & \widetilde{t}_{k,k} \end{bmatrix} + \beta_{k+1} p_{k+2} e_{k+1}^{T}$$

which gives us

$$A\overline{P}_{k}\overline{Q}_{k} = \overline{P}_{k}\overline{Q}_{k}\widetilde{T}_{k} + \widetilde{t}_{k+1,k}p_{m+1}e_{k}^{T}.$$

IRL Truncation

From $AP_m = P_m T_m + f e_m^T$, then $AP_m Q_m^+ = P_m Q_m^+ T_m^+ + f e_m^T Q_m^+$.

$$AP_{m} \left[\begin{array}{c|c} Q_{k}^{+} & Q^{+} \end{array} \right] = P_{m} \left[\begin{array}{c|c} Q_{k}^{+} & Q^{+} \end{array} \right] \left[\begin{array}{c|c} T_{k}^{+} & t_{k,k+1}^{+} e_{k} e_{1}^{T} \\ t_{k+1,k}^{+} e_{1} e_{k}^{T} & \Theta \end{array} \right] + f e_{m}^{T} \left[\begin{array}{c|c} Q_{k}^{+} & Q^{+} \end{array} \right]$$

Then
$$AP_mQ_k^+ = P_mQ_k^+T_k^+ + P_mQ^+t_{k+1,k}^+e_1e_k^T + f\underbrace{e_m^TQ_k^+}_{=q_{m,k}^+e_k^T}.$$

Mathematically, $t_{k+1,k}^+ = 0$. But we need to include this term, so

$$AP_{m}Q_{k}^{+} = P_{m}Q_{k}^{+}T_{k}^{+} + \left(P_{m}Q_{m}^{+}e_{k+1}t_{k+1,k}^{+} + q_{m,k}^{+}f\right)e_{k}^{T}$$

since $Q^+e_1 := Q_m^+e_{k+1}$.

Example

 200×200 symmetric matrix

$$\mu^*=$$
 15 (exact)

(Left) Varying p, $n_v = 80$ fixed

(Right) Varying n_v , p=20 fixed

