

Organization of Digital Computers Lab

EECS 112L

Lab 3

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1 - Objective

For this experiment we were tasked with translating MIPS assembly code into machine code, high level language code to assembly then machine code, and lastly converting assembly back into C code. After retrieving the machine code from our 4 problems, we load it into the instruction memory module for our MIPS processor from Lab 2. The following programs is what we have to translate.

3.1 Program 1

```
f = (g + h) - (i + j);
```

3.2 Program 2

```
if (i ≠ j) f = g + h; else f = g - h;
```

3.3 Program 3

```
while (save[i] ≠ k)  
    i += 1;
```

3.4 Program 4

```
sll $t0,$a1,2  
add $t0,$a0,$t0  
lw $t1,0($t0)  
lw $t2,4($t0)  
sw $t2,0($t0)  
sw $t1,4($t0)
```

2 - Procedure

Each program's translation will answer the following questions:

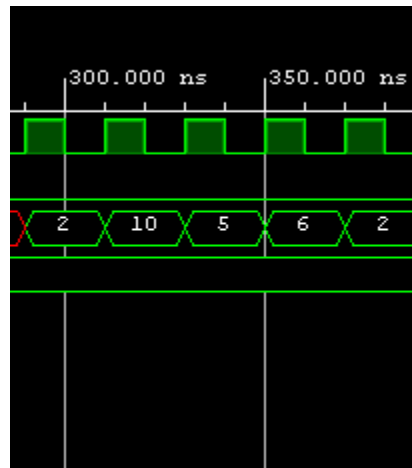
1. What MIPS assembly code might a C compiler produce?
2. From the assembly code, what machine code might a MIPS assembler produce?
3. What does this program do? (For program 4 only).

Each program is decoded from assembly code to machine code using the instruction set tables from lab 2 to get the opcode and function code

Table 4

| Name | format | Instruction[31:26] | Instruction[5:0] | ALUop | alu_control |
|------|--------|--------------------|------------------|-------|-------------|
| Add | R | 000000 | 100000 | 10 | 0010 |
| Addi | I | 001000 | - | 00 | 0010 |
| and | R | 000000 | 100100 | 10 | 0000 |
| andi | I | 001100 | - | 11 | 0000 |
| Beq | I | 000100 | - | 01 | 0110 |
| Lw | I | 100011 | - | 00 | 0010 |
| Nor | R | 000000 | 100111 | 10 | 1100 |
| Or | R | 000000 | 100101 | 10 | 0001 |
| Slt | R | 000000 | 101010 | 10 | 0111 |
| Sll | R | 110000 | 000000 | 10 | 1000 |
| Srl | R | 110000 | 000010 | 10 | 1001 |
| sra | R | 110000 | 000011 | 10 | 1010 |
| Sw | I | 101011 | - | 00 | 0010 |
| Sub | R | 000000 | 100010 | 10 | 0110 |
| xor | R | 000000 | 100110 | 10 | 0100 |
| Mult | R | 000000 | 011000 | 10 | 0101 |
| div | R | 000000 | 011010 | 10 | 1011 |
| jump | J | 000010 | - | - | - |

For all programs, we use the following data for each register and the arrays unless shown otherwise:



Here we have \$s0-\$s4 respectively

```
ram[0] = 32'd0 << 2;
ram[1] = 32'd500;
ram[2] = 32'd10;
ram[3] = 32'd243;
```

```
ram[9] = 32'd9 << 2;
ram[10] = 32'd289;
ram[11] = 32'd321;
```

end

ram[0:3] is the array save[]

ram[9:11] is the array v[]

Program 1

The following is the machine and assembly code for program 1

```
# Program 1
program_1_1: # this is the label for the assembly program
# ----- begin your MIPS assembly code -----
addi $s0, $zero, 2 #all addi instructions load constants for all
programs
addi $s1, $zero, 10
addi $s2, $zero, 5
addi $s3, $zero, 6
addi $s4, $zero, 2

add $t0, $s2, $s3 # g + h
add $t1, $s4, $s0 # i + j
```

```
sub $s1, $t0, $t1 # f = (g + h) - (i + j)
# ----- end -----
```

```
# Program 1.2
# ----- begin your MIPS Machine code -----
00100000000100000000000000000010
00100000000100010000000000000101
00100000000100100000000000000101
00100000000100110000000000000110
0010000000010100000000000000010
00000010010100110100000000100000
00000010100100000100100000100000
00000001000010011000100000100010
# ----- end -----
```

Program 2

The following is the machine code and assembly code for program 2:

```
# Program 2
program_2_1:
# ----- begin your MIPS assembly code -----
beq $s4, $s0, else #else { f = g - h }
add $s1, $s2, $s3 #if (i!=j) { f = g + h }
j exit
else:
    sub $s1, $s2, $s3
exit: # this will end up jumping to program 3
# ----- end -----
```

```
# Program 2.2
# ----- begin your MIPS Machine code -----
00010010000101000000000000000010
00000010010100111000100000100000
00001000000000000000000000010110
00000010010100111000100000100010
# ----- end -----
```

Program 3

The following is the machine code and assembly code for program 3:

```
# Program 3
program_3_1:
# ----- begin your MIPS assembly code -----
addi $s1, $zero, 10 #loading k
addi $t0, $zero, 1 #initialize i to 0
addi $t1, $zero, 0 # temp to store calculated address
lw  $s2, 0(0x0) #base address for array
loop:
    sll $t1, $t0, 2 #temp reg $t1 = i*4
    add $t1, $t1, $s2 #add base address
    lw $t2, 0x0($t1) #temp reg $t2 = save[i]
    beq $t2, $s1, exit
    addi $t0, $t0, 1
    j loop
exit: # in this case we will exit to program 4
# ----- end -----
```

```
# Program 3.2
# ----- begin your MIPS Machine code -----
0010000000001000100000000000001010
00100000000010000000000000000001
00100000000010010000000000000000
10001100000100100000000000000000
11000001000000000100100010000000
00000010010010010100100000100000
10001101001010100000000000000100
00010001010100010000000000000010
00100001000010000000000000000001
000010000000000000000000000011010
# ----- end -----
```

Program 4

The following is the machine code and C code for program 4:

```
# Program 4.3
void program_4_3(int v[], int k){
```

```

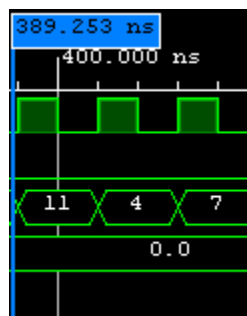
# ----- begin your C code -----
int temp;
temp = v[k];
v[k] = v[k+1];
v[k+1] = temp;
# ----- end -----
    return;
}

# Program 4
# ----- begin your MIPS Machine code -----
00100000000001010000000000000001 #k = 1
100011000000001000000000000100100 #a0 = mem[9]
11000000101000000100000010000000 #sll $t0, $a1, 2 (k+1)
00000000100010000100000000100000 #add base address
10001101000010010000000000000000
100011010000101000000000000000100
10101101000010100000000000000000
101011010000100100000000000000100
# ----- end -----

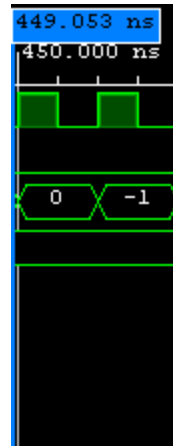
```

3 - Simulation Results

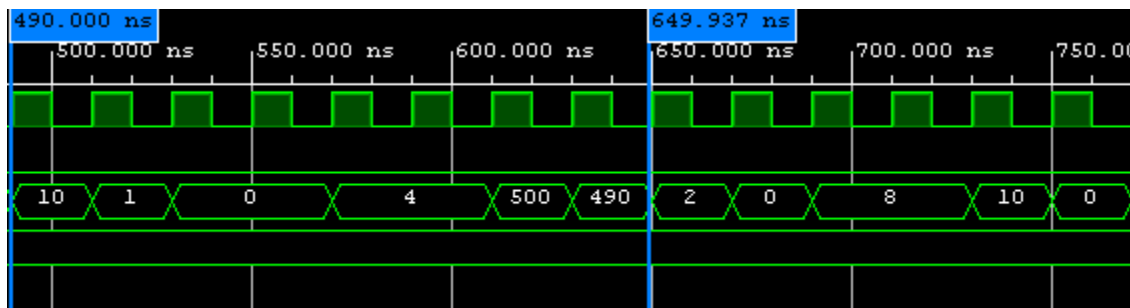
The following are screenshots of the results of the programs above



These are the results of program 1,
 11 corresponds to $g + h$,
 4 corresponds to $i + j$,
 7 corresponds to $f = (g+h) - (i + j)$;



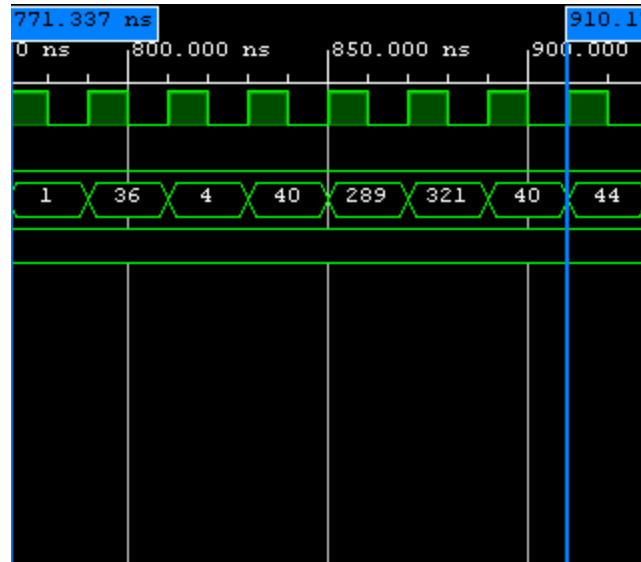
Above is the result of program 2 when branch is taken
 The first output is showing that branch was equal to zero
 and then computed the second output $s_2 - s_3$ or $5 - 6$



Program 3:

Here we are indexing through the array $save[3] = \{500, 10, 243\}$

First we are checking index $i = 1 * 4$ and checking if $save[i] = k$, if this is true then we exit the program. In this case we have $k = 10$, and we can see that at index 2 (assuming indexing on 1 for this program) is equal to 10 and the 0 indicates a jump we have taken to exit to the next program.



Program 4:

Here program 4 performs a simple swap of the content of two registers
 First we get the index of our first argument k and $k + 1$ and calculate their respective addresses in the array by multiplying by 4.

The array v here has 289 and 321 as elements. In the output waveform we see the data being loaded in then saved in swapped memory locations.