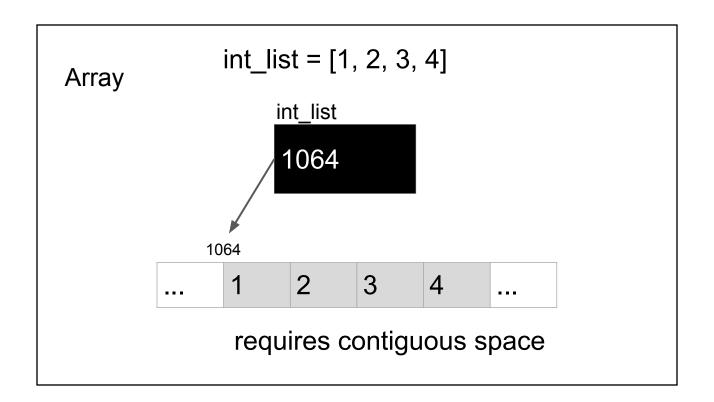
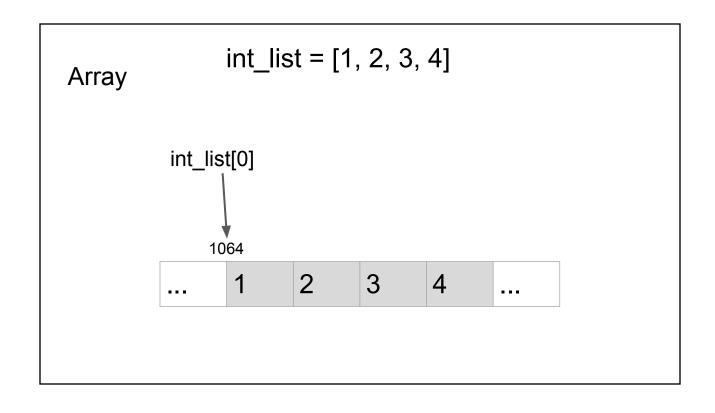
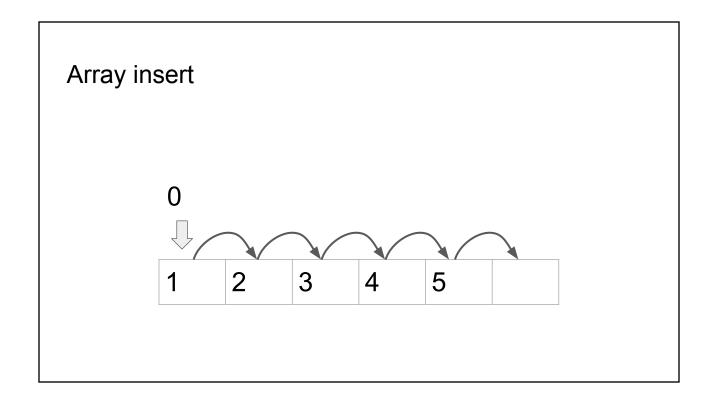
List

CPE202 Winter 2020

- Linear Data Structure
- collection of data



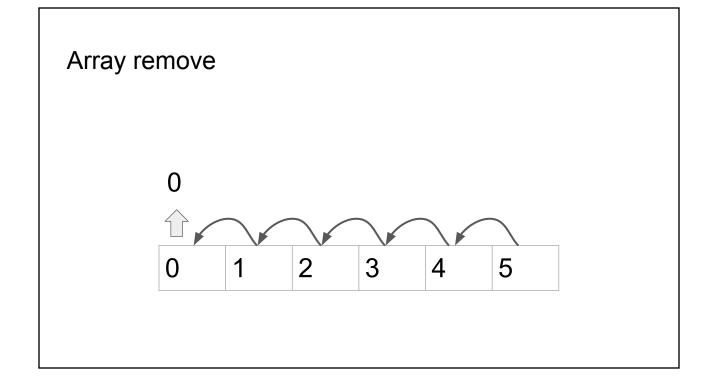




Array append (insert at the end)

6

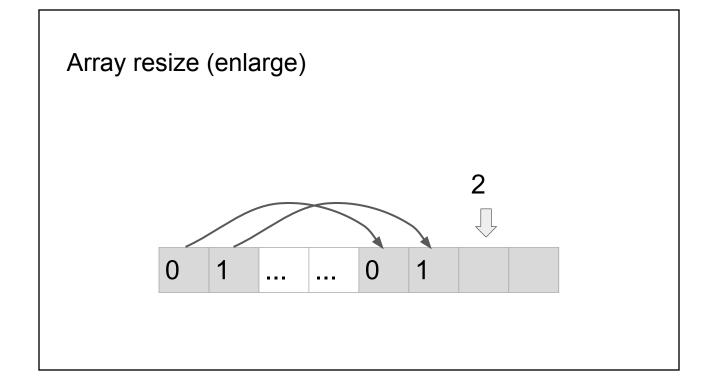
1
2
3
4
5



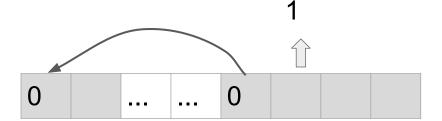
Array pop (remove from the end)

5

0 1 2 3 4 5

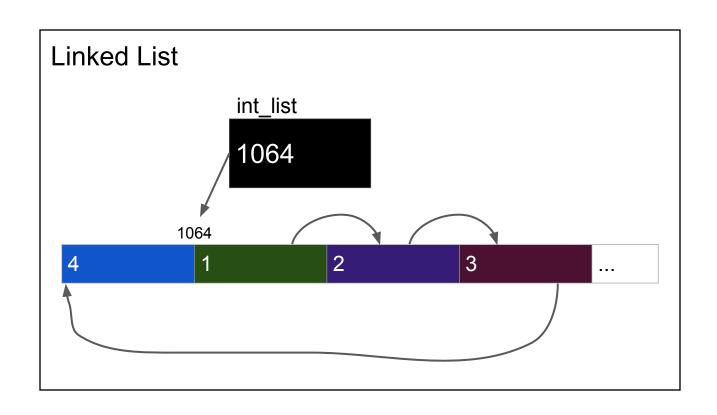


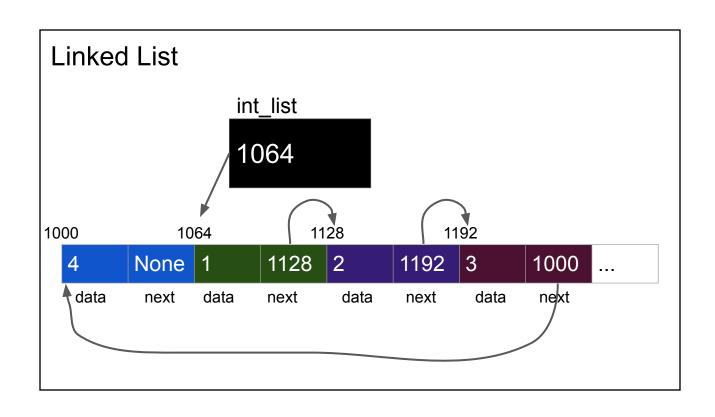
Array resize (shrink)

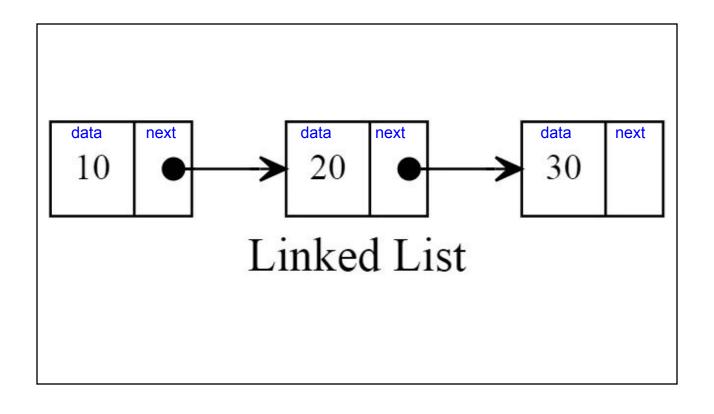


Time Complexity of Linked List v.s. Array

Operation	Array	
Accessing i th element	O(1)	
Appending an element	O(1)	
Prepending an element	O(N)	
Search	O(N) if not sorted	
Deletion	O(N) Usually involves shifting items to the left	
Resizing	O(N) on average in case of dynamic array	

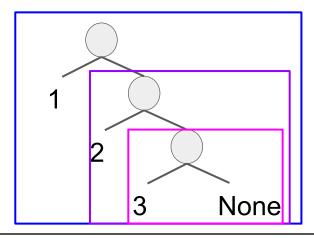






class Node: "An Int List is one of - None, or - Node(data, next) : An Int List Attributes: data (int) : int value in this example but can be any data type next (Node) : an object of Node class (an Int list) "" def __init__(self, data, nxt=None): self.data = data # a number self.next = nxt # a reference to the next node

int_list = Node(1, Node(2, Node(3, None)))



Printing an object on the screen

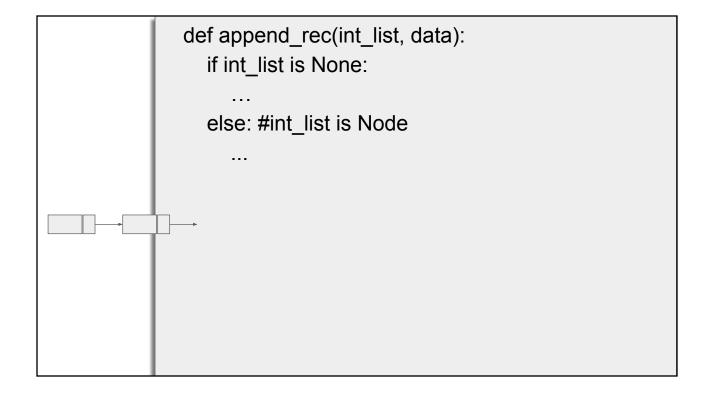
```
def __repr__(self):
return "Node(%d, %s)" % (self.data, self.nxt)
```

checking equality between objects

```
def __eq__(self, other):
    return isinstance(other, Node)\
    and self.data == other.data\
    and self.next == other.next
```

- Slower Access
 - only sequential access is possible





```
def remove_rec(int_list, target)->Node:
    if int_list is None:
    ...
    elif int_list.data == target:
    ...
    else: #int_list.data != target
    ...

return None
```

```
def get_rec(int_list, pos)->int:
    if int_list is None:
        raise IndexError()
    if ...:
        return int_list.data
    else:
    ...
```

```
def find_rec(int_list, target)->int:
    return find_helper(int_list, target, ?)

def find_helper(int_list, target, pos):
    if int_list is None:
    ...
    elif ...:
    ...
    else:
    ...
```

```
def truncate(int_list, val)->Node:

"""truncates an IntList at the val.

Args:

int_list (Node): an IntList

val (int): value to be found

Returns:

Node: an IntList (Node containing val) or None

"""

if int_list == None:

...

if val == int_list.data:

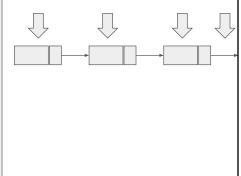
...

...
```

Iterator

```
def next(int_list)->(Node, int):
    """advances to the next element
    Args:
        int_list (Node): an IntList
    Returns:
        Node: the next element or None
        int: the current element's value
    Raises:
        StopIteration: when the end has been reached
    """
    if int_list is None:
        raise StopIteration
```

return int list.next, int list.data



```
def sum_list(int_list):
    """Sums all the numbers in the list and returns the sum.

Args:
    int_list (Node): None or An IntList
    Returns:
    int: the sum
    """

if node == None:
    return ?
return ?
```

```
def sum_accum(int_list, acc=0):
    """Sums all the numbers in the list and returns the sum.

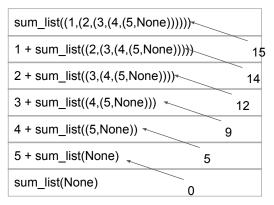
Args:
    int_list (Node): None or IntList
    acc (int): accumulator
Returns:
    int: the sum
    """

if int_list is None:
    return ?
return ?
```

Tail Recursion

- Tail Recursion
 - No computation after recursion.
 - Can be converted to iteration because the caller does not need to wait for the callee to return.
 - Some language compilers optimize tail recursion by converting it to iteration.
 - Python does not optimize it.

Tail Recursion

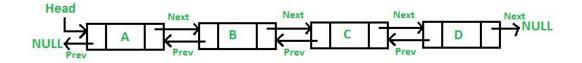


Non-Tail Recursive version

sum_accum((1,(2,(3,(4,(5,None))))), 0)		
sum_accum((2,(3,(4,(5,None)))), 1)		
sum_accum((3,(4,(5,None))), 3)		
sum_accum((4,(5,None)), 6)		
sum_accum((5,None), 10)		
sum_accum(None, 15)		
15		

Tail Recursive version

Doubly Linked List



class Node:

- "An Int List is one of
 - None, or
 - Node(data, next, prev) : An Int List

Attributes:

```
data (int): int value in this example but can be any data type
```

next (Node): an object of Node class (an Int list) prev (Node): an object of Node class (an Int list)

""

```
def __init__(self, data, nxt=None, prev=None):
```

self.data = data # a number

self.next = nxt # a reference to the next node

self.prev = prev # a reference to the previous node

Time Complexity of Linked List v.s. Array

Operation	Linked List	Array
Accessing i th element	O(N)	O(1)
Appending an element	O(1) may require 2 variables	O(1)
Prepending an element	O(1) may require 2 variables	O(N)
Search	O(N) on average	O(N) if not sorted
Deletion	O(N) on average	O(N) Usually involves shifting items to the left
Resizing	-	O(N) on average in case of dynamic array