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CSE 250A HW3
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1. (a) when t=1, P(X_2=j|X_1=i)=[A]_{ij}
 Suppose P(Xk+1=j|X,=i)=[Ak]; for some k >1,
 then P(X_{k+2}=j \mid X_1=i) = \sum_{i=1}^{m} P(X_{k+2}=j, X_{k+1}=l \mid X_1=i)
                           = \sum_{k=1}^{\infty} P(X_{k+1} = l \mid X_1 = i) P(X_{k+2} = j \mid X_{k+1} = l \mid X_1 = i) Product Rule_{j}
= \sum_{k=1}^{\infty} [A^k]_{i,k} P(X_{k+2} = j \mid X_{k+1} = l) \quad (by d. sep (1))
X_{1} = \sum_{k=1}^{\infty} [A^k]_{i,k} P(X_{k+2} = j \mid X_{k+1} = l) \quad (by d. sep (1))
                           = 2 [A] [[A]
                            = [A"] ii
By induction, we have P(X_{t+1}=j \mid X_1=i)=[A^t]_{ij}, \forall t \ge 1.
       int prob(int[][] A, int t, intj, inti) {
(b)
           if (t==1) return A[i][j];
           int res = 0;
           for (int k=0; k < A. length; k++)
              res += prob(A, t-1, k, i) * prob(A, 1, j, k);
          return res;
 Since [A^t]_{ij} = \sum_{k=1}^{\infty} [A^{t-1}]_{ik}[A]_{kj}, we use recursion to compute each term in summation,
 each row * col is O(m), prob(A, t-1, k,i) needs O(t) to reduce to base case,
  and sum takes m steps, so it's O(m2+).
(c) Similarly, we split the Ak to Ak/2 and Ak/2 in recursive steps:
       int prob2 (int[][]A, int t, intj, inti) {
          if (t==1) return Alij[j];
          int res = 0;
          许((+&1)==1) {
            for (int k=0; k<A.length; k++){
                res += prob2(A, f-1, k, i) * prob2(A.1.j. k);
         ] else {
             for (int k=0; k < A. length; k++) }
                res t = prob2(A, +12, K, i)*prob2(A, +12, j, k);
                                                                                     hase case (t=1), so it is O(m²)
         return res;
This is O(m3 logt) since recursion takes (og(+) but in summation we don't always have
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(d) Instead of going over the entire row, we only look at nonzero entries.
We add following to function in (b):
     Set (Integer) nonzero = new HashSet ()();
     for (int k=0; k<A.length; k++)
if (A[i][k]!=0) nonzero.add(k);
then change (*) to for (int k: nonzero)
This is O(smt) since the row *col becomes O(s).
where P(X_{t+1}=j) = \sum_{i=1}^{m} P(X_{t+1}=j, X_{i}=i)
                     = \sum_{i=1}^{m} P(X_{i+1} = j | X_1 = i) P(X_1 = i)
                    = \sum_{i=1}^{\infty} [A^{i}]_{ij} P(X_{i}=i)
 So P(X_{i=1}|X_{i+1=j}) = \frac{[A^{i}]_{ij} P(X_{i=1})}{\sum_{k=1}^{m} [A^{i}]_{kj} P(X_{i=k})}
2.(a) P(Y_1|X_1) = \sum_{x} P(Y_1, X_0 = x | X_1)
                   = \overline{2}P(Y_1|X_0=Y_1,X_1)P(X_0=X|X_1)
                  = \overline{2} P(Y_1) X_0 = Y_1 X_1) P(X_0 = X)
 (b) P(Y,)= = P(Y,, X,=X)
              =\frac{1}{2}P(Y, |X|=x)P(X_i=x)
 (C) P(Xn | Y ..., Yn-1) = P(Xn)
              By disep(3), Xn->Yil & is blocked by Yn, Vi.
 (d) P(Yn | Xn, Y1, ..., Yn-1) = = P(Yn, Xn-1=x | Xn, Y1, ..., Yn-1)
                              = = P(Yn 1 Xn, Y,, ..., Yn-1, Xn-1=X) P(Xn-1=X | Xn, Y,, ..., Yn-1)
                              = 2 P(Yn) Xn, Xn-1) P(Xn-1 = X [Y1, ..., Yn-1)
                              TYn - Y ..... Yn | Xn, Xn-1 by d. sep (1), Xn-1-)Xn | Y ..... Yn-1 by d. sep (3)]
 (e) P(Yn|Y,, ", Yn) = = P(Yn, Xn=X|Y,, ..., Yn-1)
                         = IP(Xn= Y | Y1, ..., Yn-1) P(Yn | Xn= X. Y1, ..., Yn-1)
                         = = P(Xn=X) P(In | Xn=X, Y, ..., In-1)
by (c) known from (d)
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(b) P(B|A,C,D,E,F) = P(B|A,C,D) \cap B \to E,F|A|CD  by d.sep(2)_{B \to E} blocked by C, B \to E blocked by A.

(c) P(B,E,F|A,(D)) = P(B|E,F,A,C,D) P(E|F,A,C,D) P(F|A,C,D) = P(B|A,C,D) P(E|C) P(F|A)

by (a) \in \to A.F.D(C \in F\to C.D|A)

\bullet \to A.F.D(C \in F\to C.D|A)

by (a) \in \to A.F.D(C \in F\to C.D|A)

\bullet \to A.F.D(C \in F\to C.D|A)

by (a) \in \to A.F.D(C \in F\to C.D)

\bullet \to A.F.D(C \in F\to C.D|A)

\bullet \to A.F.D(C \in F\to C.D(A)

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