

## 高频题的陨落篇

该篇主要讲解面试中的高频题目：

1. [Single Number I, II, III](#)
2. [Majority Number I, II, III](#)
3. [Best Time to Buy and Sale Stock I, II, III](#)
4. [Subarray I, II, III, IV, V, VI](#)
5. [2-Sum, 3-Sum, 3-Sum Closest, 4-Sum, k-Sum, k-Sum II](#)
6. [Quick Questions](#)
7. [Partition Array](#)

1.1 [Single Number I](#)，利用异或  $a \oplus a = 0$  的性质，扫描一遍即可。

```
public int singleNumber(int[] A) {  
    if(A == null || A.length == 0){  
        return -1;  
    }  
    int res = A[0];  
    for(int i = 1; i < A.length; i++){  
        res ^= A[i];  
    }  
    return res;  
}
```

1.2 [Single Number II](#)，提供2种解法。第一种，为  $nk + 1$  类型题目的通解；第二种，该方法比较巧妙也比较难以理解，在位运算篇有详细解释，有能力者可以参考。

```
public int singleNumber(int[] A) {  
    if(A == null || A.length == 0){  
        return -1;  
    }  
    int[] digits = new int[32];  
    for(int i = 0; i < 32; i++){  
        for(int j = 0; j < A.length; j++){  
            digits[i] += A[j] >> i & 1;  
        }  
    }  
}
```

```

    int res = 0;
    for(int i = 0; i < 32; i++){
        res += digits[i] % 3 << i;
    }
    return res;
}

public int singleNumber(int[] A) {
    int ones = 0;
    int twos = 0;
    for(int i = 0; i < A.length; i++){
        ones = (A[i] ^ ones) & (~twos);
        twos = (A[i] ^ twos) & (~ones);
    }
    return ones;
}

```

1.3 [Single Number III](#)，这道题是 $2n + 2$ 的问题。将所有数异或后得到的结果c是a异或b的结果，那么问题来了，我们该如何将a和b区别开至两组。由于c一定不为0的，所以c的某一位一定不为0，所以我们可以根据这一位将所有数分成2组，即可得到结果。

```

public List<Integer> singleNumberIII(int[] A) {
    List<Integer> res = new ArrayList<Integer>();
    if(A == null || A.length == 0){
        return res;
    }
    int c = 0;
    for(int num : A){
        c ^= num;
    }
    int digit = 0;
    for(int i = 0; i < 32; i++){
        if((c >> i & 1) == 1){
            digit = i;
        }
    }
    res.add(0);
    res.add(0);
    for(int num : A){
        if((num >> digit & 1) == 1){

```

```

        res.set(0, res.get(0) ^ num);
    } else {
        res.set(1, res.get(1) ^ num);
    }
}
return res;
}

```

2.1 [Majority Number I](#), 又到了一年一度的 Majority Number 选举, 每个数字都为自己投票, Majority Number 获得的票数最多。不同的数字会抵消 Majority Number 的投票。

```

public int majorityNumber(ArrayList<Integer> nums) {
    if(nums == null || nums.size() == 0){
        return -1;
    }
    int vote = nums.get(0);
    int count = 1;
    for(int i = 1; i < nums.size(); i++){
        if(nums.get(i) == vote){
            count++;
        } else if(count == 0){
            vote = nums.get(i);
            count++;
        } else {
            count--;
        }
    }
    return vote;
}

```

2.2 [Majority Number II](#), 使用2个候选变量储存结果即可。注意, 第一遍的count1和count2并不是2个候选数字在数组中出现的次数, 需要重新扫描一遍数组, 找出最终结果。

```

public int majorityNumber(ArrayList<Integer> nums) {
    if(nums == null || nums.size() == 0){
        return -1;
    }
    if(nums.size() < 3){
        return nums.get(0);
    }
}

```

```

int vote1 = nums.get(0);
int vote2 = 0;
int count1 = 1;
int count2 = 0;
for(int i = 1; i < nums.size(); i++){
    if(nums.get(i) == vote1){
        count1++;
    } else if(nums.get(i) == vote2){
        count2++;
    } else if(count1 == 0){
        vote1 = nums.get(i);
        count1++;
    } else if(count2 == 0){
        vote2 = nums.get(i);
        count2++;
    } else {
        count1--;
        count2--;
    }
}
count1 = 0;
count2 = 0;
for(int num : nums){
    if(num == vote1){
        count1++;
    } else if(num == vote2){
        count2++;
    }
}
return count1 > count2 ? vote1 : vote2;
}

```

2.3 [Majority Number III](#)，我们使用一个HashMap来存储候选元素。这道题可以用来练习如何对HashMap进行遍历。同样注意，最后要重新遍历一遍找出最终结果。

```

public int majorityNumber(ArrayList<Integer> nums, int k) {
    if(nums == null || nums.size() == 0){
        return -1;
    }
    HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();

```

```

map.put(nums.get(0), 1);
for(int i = 1; i < nums.size(); i++){
    if(map.containsKey(nums.get(i))){
        int value = map.get(nums.get(i)) + 1;
        if(value == 0){
            map.remove(nums.get(i));
        } else {
            map.put(nums.get(i), value);
        }
    } else {
        if(map.size() < k - 1){
            map.put(nums.get(i), 1);
        } else {
            ArrayList<Integer> keyList = new ArrayList<Integer>();
            for(Map.Entry entry : map.entrySet()){
                int key = (int)entry.getKey();
                int value = (int)entry.getValue();
                value--;
                if(value == 0){
                    keyList.add(key);
                }
                map.put(key, value);
            }
            for(int key : keyList){
                map.remove(key);
            }
        }
    }
}

for(Map.Entry entry : map.entrySet()){
    entry.setValue(0);
}

int vote = 0;
int count = 0;
for(int i = 0; i < nums.size(); i++){
    if(map.containsKey(nums.get(i))){
        int value = map.get(nums.get(i)) + 1;
        if(value > count){
            vote = nums.get(i);
            count = value;
        }
    }
}

```

```

        }
        map.put(nums.get(i), value);
    }
}
return vote;
}

```

### 3.1 [Best Time to Buy and Sale Stock I](#)，局部最优和全局最优解法。

```

public int maxProfit(int[] prices) {
    if(prices == null || prices.length == 0){
        return 0;
    }
    int local = 0;
    int global = 0;
    for(int i = 0; i < prices.length - 1; i++){
        local = Math.max(local, 0) + prices[i + 1] - prices[i];
        global = Math.max(local, global);
    }
    return global;
}

```

### 3.2 [Best Time to Buy and Sale Stock II](#)，贪心算法。

```

public int maxProfit(int[] prices) {
    if(prices == null || prices.length == 0){
        return 0;
    }
    int res = 0;
    for(int i = 0; i < prices.length - 1; i++){
        res += Math.max(0, prices[i + 1] - prices[i]);
    }
    return res;
}

```

### 3.3 [Best Time to Buy and Sale Stock III](#)，提供2种解法。稳定版和通用版。

```

public int maxProfit(int[] prices) {
    if(prices == null || prices.length == 0){
        return 0;
    }
}

```

```

int local = 0;
int global = 0;
int[] left = new int[prices.length];
int[] right = new int[prices.length];
for(int i = 0; i < prices.length - 1; i++){
    local = Math.max(local, 0) + prices[i + 1] - prices[i];
    global = Math.max(local, global);
    left[i] = global;
}
local = 0;
global = 0;
for(int i = prices.length - 2; i >= 0; i--){
    local = Math.max(local, 0) + prices[i + 1] - prices[i];
    global = Math.max(local, global);
    right[i] = global;
}
int max = 0;
for(int i = 0; i < prices.length - 1; i++){
    max = Math.max(max, left[i] + right[i + 1]);
}
return max;
}

```

```

public int maxProfit(int[] prices) {
    if(prices == null || prices.length == 0){
        return 0;
    }
    int[] local = new int[3];
    int[] global = new int[3];
    for(int i = 0; i < prices.length - 1; i++){
        int diff = prices[i + 1] - prices[i];
        for(int j = 2; j >= 1; j--){
            local[j] = Math.max(global[j - 1] + Math.max(diff, 0), local[j] + diff);
            global[j] = Math.max(local[j], global[j]);
        }
    }
    return global[2];
}

```

#### 4.1 Subarray I – Maximum Subarray，局部最优和全局最优以及Divide & Conquer解法。

##### // Solution 1 - Dynamic Programming

```
public int maxSubArray(int[] A) {  
    if(A == null || A.length == 0){  
        return 0;  
    }  
    int local = A[0];  
    int global = A[0];  
    for(int i = 1; i < A.length; i++){  
        local = Math.max(0, local) + A[i];  
        global = Math.max(local, global);  
    }  
    return global;  
}
```

##### // Solution 2 - Divide and Conquer

```
public int maxSubArray(int[] A) {  
    if(A == null || A.length == 0){  
        return 0;  
    }  
    return helper(A, 0, A.length - 1);  
}  
  
private int helper(int[] A, int left, int right){  
    if(left == right){  
        return A[left];  
    }  
    int mid = left + (right - left) / 2;  
    int leftSub = helper(A, left, mid);  
    int rightSub = helper(A, mid + 1, right);  
    int leftMax = A[mid];  
    int temp = 0;  
    for(int i = mid; i >= left; i--){  
        temp += A[i];  
        leftMax = Math.max(temp, leftMax);  
    }  
    int rightMax = A[mid + 1];  
    temp = 0;  
    for(int j = mid + 1; j <= right; j++){  
        temp += A[j];  
        rightMax = Math.max(temp, rightMax);  
    }  
    return Math.max(leftMax, rightMax);  
}
```



```

        rightMax = Math.max(temp, rightMax);
    }
    return Math.max(Math.max(leftSub, rightSub), leftMax + rightMax);
}

```

4.2 [Subarray II – Maximum Product Subarray](#)，变形的局部最优和全局最优解法。

```

public int maxProduct(int[] A) {
    if(A == null || A.length == 0){
        return 0;
    }
    int localMax = A[0];
    int localMin = A[0];
    int global = A[0];
    for(int i = 1; i < A.length; i++){
        int tempLocalMax = localMax;
        localMax = Math.max(A[i], Math.max(A[i] * localMin, A[i] * localMax));
        localMin = Math.min(A[i], Math.min(A[i] * localMin, A[i] * tempLocalMax));
        global = Math.max(global, localMax);
    }
    return global;
}

```

4.3 [Subarray III – Minimum Subarray](#)，将nums数组元素取相反数，再按[Maximum Subarray](#)的方法求最大，最后返回-global即可。

```

public int minSubArray(ArrayList<Integer> nums) {
    if(nums == null || nums.size() == 0){
        return 0;
    }
    int local = -nums.get(0);
    int global = -nums.get(0);
    for(int i = 1; i < nums.size(); i++){
        local = Math.max(local, 0) - nums.get(i);
        global = Math.max(local, global);
    }
    return -global;
}

```

4.4 Subarray IV – Maximum Subarray II, 同[Best Time to Buy and Sale Stock III](#)稳定版的原理一样, 两次扫描得到从左至右和从右至左的最优解, 再最终扫描一次得到结果。

```
public int maxTwoSubArrays(ArrayList<Integer> nums) {
    if(nums == null || nums.size() < 2){
        return 0;
    }
    int[] left = new int[nums.size()];
    left[0] = nums.get(0);
    int local = nums.get(0);
    int global = nums.get(0);
    for(int i = 1; i < nums.size(); i++){
        local = Math.max(local, 0) + nums.get(i);
        global = Math.max(local, global);
        left[i] = global;
    }
    int[] right = new int[nums.size()];
    right[nums.size() - 1] = nums.get(nums.size() - 1);
    local = nums.get(nums.size() - 1);
    global = nums.get(nums.size() - 1);
    for(int i = nums.size() - 2; i >= 0; i--){
        local = Math.max(local, 0) + nums.get(i);
        global = Math.max(local, global);
        right[i] = global;
    }
    int max = Integer.MIN_VALUE;
    for(int i = 0; i < nums.size() - 1; i++){
        max = Math.max(max, left[i] + right[i + 1]);
    }
    return max;
}
```

4.5 Subarray V – Maximum Subarray Difference, 与上一题思路一样, 只是这次要同时保存左右两边的最大和最小值, 最终结果可能为左边最大-右边最小或者右边最大-左边最小。

```
public int maxDiffSubArrays(ArrayList<Integer> nums) {
    if(nums == null || nums.size() < 2){
        return 0;
    }
    int len = nums.size();
```

```

int[] leftMax = new int[len];
leftMax[0] = nums.get(0);
int[] leftMin = new int[len];
leftMin[0] = nums.get(0);
int localMax = nums.get(0);
int globalMax = nums.get(0);
int localMin = nums.get(0);
int globalMin = nums.get(0);
for(int i = 1; i < len; i++){
    localMax = Math.max(localMax, 0) + nums.get(i);
    globalMax = Math.max(localMax, globalMax);
    leftMax[i] = globalMax;
    localMin = Math.min(localMin, 0) + nums.get(i);
    globalMin = Math.min(localMin, globalMin);
    leftMin[i] = globalMin;
}
int[] rightMax = new int[len];
rightMax[len - 1] = nums.get(len - 1);
int[] rightMin = new int[len];
rightMin[len - 1] = nums.get(len - 1);
localMax = nums.get(len - 1);
globalMax = nums.get(len - 1);
localMin = nums.get(len - 1);
globalMin = nums.get(len - 1);
for(int i = len - 2; i >= 0; i--){
    localMax = Math.max(localMax, 0) + nums.get(i);
    globalMax = Math.max(localMax, globalMax);
    rightMax[i] = globalMax;
    localMin = Math.min(localMin, 0) + nums.get(i);
    globalMin = Math.min(localMin, globalMin);
    rightMin[i] = globalMin;
}
int max = 0;
for(int i = 0; i < len - 1; i++){
    max = Math.max(Math.max(Math.abs(leftMax[i] - rightMin[i + 1]),
Math.abs(rightMax[i + 1] - leftMin[i])), max);
}
return max;
}

```

4.6 Subarray VI – Maximum Subarray III, 该题就是[Best Time to Buy and Sale Stock III](#)中使用的通用解法。由于本人能力有限，理解的不是很透彻。仅在此提供一个解法，希望有朝一人能被有缘人发现，为其所用。

$d[i][j]$  means the maximum sum we can get by selecting  $j$  subarrays from the first  $i$  elements.

$d[i][j] = \max\{d[p][j-1] + \text{maxSubArray}(p+1, i)\}$

we iterate  $p$  from  $i-1$  to  $j-1$ , so we can record the max subarray we get at current  $p$ , this value can be used to calculate the max subarray from  $p-1$  to  $i$  when  $p$  becomes  $p-1$ .

```
public int maxSubArray(ArrayList<Integer> nums, int k) {
    if (nums.size() < k){
        return 0;
    }
    int len = nums.size();
    int[][] d = new int[len + 1][k + 1];
    for (int i = 0; i <= len; i++){
        d[i][0] = 0;
    }
    for (int j = 1; j <= k; j++){
        for (int i = j; i <= len; i++){
            d[i][j] = Integer.MIN_VALUE;
            int endMax = 0;
            int max = Integer.MIN_VALUE;
            for (int p = i - 1; p >= j - 1; p--){
                endMax = Math.max(nums.get(p), endMax + nums.get(p));
                max = Math.max(endMax, max);
                if (d[i][j] < d[p][j - 1] + max)
                    d[i][j] = d[p][j - 1] + max;
            }
        }
    }
    return d[len][k];
}
```

5.1 2-Sum, 巧妙利用HashMap和补集原理来寻找要求的2个数字。

```
public int[] twoSum(int[] numbers, int target) {
    if(numbers == null || numbers.length < 2){
        return null;
    }
    int[] res = new int[2];
    HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();
    for(int i = 0; i < numbers.length; i++){
        if(map.containsKey(target - numbers[i])){
            res[0] = map.get(target - numbers[i]) + 1;
            res[1] = i + 1;
            return res;
        }
        map.put(numbers[i], i);
    }
    return res;
}
```

5.2 3-Sum, 遍历数组num一遍, 使用排序+夹逼的方式找出2Sum等于-num[i]的结果即可。

```
public ArrayList<ArrayList<Integer>> threeSum(int[] num) {
    ArrayList<ArrayList<Integer>> res = new ArrayList<ArrayList<Integer>>();
    if(num == null || num.length < 3){
        return res;
    }
    Arrays.sort(num);
    for(int i = num.length - 1; i >= 2; i--){
        if(i != num.length - 1 && num[i] == num[i + 1]){
            continue;
        }
        ArrayList<ArrayList<Integer>> item = findTwoSum(num, i - 1, -num[i]);
        for(int j = 0; j < item.size(); j++){
            item.get(j).add(num[i]);
        }
        res.addAll(item);
    }
    return res;
}
```

```

private ArrayList<ArrayList<Integer>> findTwoSum(int[] num, int end, int target){
    int left = 0;
    int right = end;
    ArrayList<ArrayList<Integer>> res = new ArrayList<ArrayList<Integer>>();
    while(left < right){
        if(num[left] + num[right] == target){
            ArrayList<Integer> item = new ArrayList<Integer>();
            item.add(num[left]);
            item.add(num[right]);
            res.add(item);
            left++;
            right--;
            while(left < right && num[left] == num[left - 1]){
                left++;
            }
            while(left < right && num[right] == num[right + 1]){
                right--;
            }
        } else if(num[left] + num[right] < target){
            left++;
        } else {
            right--;
        }
    }
    return res;
}

```

5.3 [3-Sum Closest](#), 与3-Sum基本一致，只是找closest的时候细节需要稍作变动即可。

```

public int threeSumClosest(int[] num, int target) {
    if(num == null || num.length < 3){
        return -1;
    }
    Arrays.sort(num);
    int closest = Integer.MAX_VALUE;
    for(int i = 0; i < num.length - 2; i++){
        int value = twoSumClosest(num, target - num[i], i + 1);
        if(Math.abs(value) < Math.abs(closest)){
            closest = value;
        }
    }
}

```

```

    }
    return target + closest;
}

private int twoSumClosest(int[] num, int target, int start){
    int left = start;
    int right = num.length - 1;
    int closest = Integer.MAX_VALUE;
    while(left < right){
        if(num[left] + num[right] == target){
            return 0;
        }
        int value = num[left] + num[right] - target;
        if(Math.abs(value) < Math.abs(closest)){
            closest = value;
        }
        if(value < 0){
            left++;
        } else {
            right--;
        }
    }
    return closest;
}

```

5.4 4-Sum，这道题提供2种解法。第一种，即是3-Sum再套一层循环，容易理解且代码容易写，但时间复杂度为 $O(n^3)$ ；第二种，二分思想很好，合理利用了TreeMap数据结构，但代码较为复杂冗长，时间复杂度可以降至 $O(n^2 \log n)$ 。

```

public ArrayList<ArrayList<Integer>> fourSum(int[] num, int target) {
    ArrayList<ArrayList<Integer>> res = new ArrayList<ArrayList<Integer>>();
    if(num == null || num.length < 4){
        return res;
    }
    Arrays.sort(num);
    for(int i = num.length - 1; i >= 3; i--){
        if(i == num.length - 1 || num[i] != num[i + 1]){
            ArrayList<ArrayList<Integer>> curRes = threeSum(num, 0, i - 1,
target - num[i]);

```

```

        for(int j = 0; j < curRes.size(); j++){
            curRes.get(j).add(num[i]);
        }
        res.addAll(curRes);
    }
}
return res;
}

public ArrayList<ArrayList<Integer>> threeSum(int[] num, int l, int r, int target) {
    ArrayList<ArrayList<Integer>> res = new ArrayList<ArrayList<Integer>>();
    if(num == null || num.length < 3){
        return res;
    }
    for(int i = r; i >= 2; i--){
        if(i == r || num[i] != num[i + 1]){
            ArrayList<ArrayList<Integer>> curRes = twoSum(num, 0, i - 1, target
- num[i]);

            for(int j = 0; j < curRes.size(); j++){
                curRes.get(j).add(num[i]);
            }
            res.addAll(curRes);
        }
    }
    return res;
}

public ArrayList<ArrayList<Integer>> twoSum(int[] num, int l, int r, int target) {
    ArrayList<ArrayList<Integer>> res = new ArrayList<ArrayList<Integer>>();
    if(num == null || num.length < 2){
        return res;
    }
    while(l < r){
        if(num[l] + num[r] == target){
            ArrayList<Integer> item = new ArrayList<Integer>();
            item.add(num[l]);
            item.add(num[r]);
            res.add(item);
            l++;
            r--;
        }
    }
}

```



```

        while(l < r && num[l] == num[l - 1]){
            l++;
        }
        while(l < r && num[r] == num[r + 1]){
            r--;
        }
    } else if(num[l] + num[r] < target){
        l++;
    } else {
        r--;
    }
}
return res;
}

```

```
import java.util.TreeMap;
```

```

public class Solution {
    class Pair {
        int a;
        int ai;
        int b;
        int bi;

        public Pair(int a, int ai, int b, int bi){
            this.a = a;
            this.ai = ai;
            this.b = b;
            this.bi = bi;
        }

        boolean same(Pair p){
            return p != null && p.a == a && p.b == b;
        }
    }

    public List<List<Integer>> fourSum(int[] num, int target) {
        List<List<Integer>> res = new ArrayList<>();
        if(num.length < 4){
            return res;
        }
    }
}

```

```

}
Arrays.sort(num);
TreeMap<Integer, List<Pair>> map = new TreeMap<>();
for(int i = 0; i < num.length; i++){
    for(int j = i + 1; j < num.length; j++){
        Pair pair = new Pair(num[i], i, num[j], j);
        int sum = num[i] + num[j];
        List<Pair> list;
        if(map.containsKey(sum)){
            list = map.get(sum);
        } else {
            list = new ArrayList<>();
            map.put(sum, list);
        }
        list.add(pair);
    }
}
Integer first = map.firstKey();
Integer last = map.lastKey();
while(first != null && last != null && first <= last){
    if(first + last > target){
        last = map.lowerKey(last);
    } else if(first + last < target){
        first = map.higherKey(first);
    } else {
        Pair lastA = null;
        for(Pair a : map.get(first)){
            if(a.same(lastA)){
                continue;
            }
            lastA = a;
        }
        Pair lastB = null;
        for(Pair b: map.get(last)){
            if(a.bi < b.ai){
                if(b.same(lastB)){
                    continue;
                }
                lastB = b;
            }
            res.add(Arrays.asList(new Integer[]{a.a, a.b, b.a, b.b}));
        }
    }
}

```

```

        }
    }
    last = map.lowerKey(last);
    first = map.higherKey(first);
}
}
return res;
}
}

```

5.5 [k-Sum](#),  $res[i][j][v]$  means the way of selecting  $i$  elements from the first  $j$  elements so that their sum equals to  $v$ . Then we have:

$$res[i][j][v] = res[i - 1][j - 1][v - A[j - 1]] + d[i][j - 1][v]$$

It means two operations, select the  $j$ th element and not select the  $j$ th element.

```

public int kSum(int A[], int k, int target) {
    if(A == null || A.length < k){
        return 0;
    }
    int[][][] res = new int[k + 1][A.length + 1][target + 1];
    for(int i = 1; i <= A.length; i++){
        if(A[i - 1] <= target){
            for(int j = i; j <= A.length; j++){
                res[1][j][A[i - 1]] = 1;
            }
        }
    }
    for(int i = 2; i <= k; i++){
        for(int j = i; j <= A.length; j++){
            for(int v = 1; v <= target; v++){
                res[i][j][v] = 0;
                if(i < j){
                    res[i][j][v] += res[i][j - 1][v];
                }
                if(A[j - 1] <= v){
                    res[i][j][v] += res[i - 1][j - 1][v - A[j - 1]];
                }
            }
        }
    }
}

```

```

    }
}
return res[k][A.length][target];
}

```

5.6 k-Sum II, 回溯黄金模板, 类似[Combination Sum](#), 只是加了要求k个数的限定条件。

```

public ArrayList<ArrayList<Integer>> kSumII(int A[], int k, int target) {
    ArrayList<ArrayList<Integer>> res = new ArrayList<ArrayList<Integer>>();
    if(A == null || A.length < k){
        return res;
    }
    Arrays.sort(A);
    ArrayList<Integer> item = new ArrayList<Integer>();
    helper(A, k, 0, target, 0, item, res);
    return res;
}

private void helper(int[] A, int k, int num, int target, int start, ArrayList<Integer>
item, ArrayList<ArrayList<Integer>> res){
    if(target < 0){
        return;
    }
    if(num == k && target == 0){
        res.add(new ArrayList<Integer>(item));
        return;
    }
    for(int i = start; i < A.length; i++){
        if(i > start && A[i] == A[i - 1]){
            continue;
        }
        item.add(A[i]);
        helper(A, k, num + 1, target - A[i], i + 1, item, res);
        item.remove(item.size() - 1);
    }
}

```

6.1 `Pow(x, n)`，提供2种解法。第一种，位运算；第二种，二分法。

**// Solution 1 - Bit Operation**

```
public double pow(double x, int n) {
    if(n == 0){
        return 1.0;
    }
    double res = 1.0;
    if(n < 0){
        if(x >= 1.0 / Double.MAX_VALUE || x <= 1.0 / Double.MIN_VALUE){
            x = 1.0 / x;
        } else {
            return Double.MAX_VALUE;
        }
        if(n == Integer.MIN_VALUE){
            res *= x;
            n++;
        }
    }
    n = Math.abs(n);
    boolean isNeg = false;
    if(x < 0 && n % 2 == 1){
        isNeg = true;
    }
    x = Math.abs(x);
    while(n > 0){
        if((n & 1) == 1){
            if(res > Double.MAX_VALUE / x){
                return Double.MAX_VALUE;
            }
            res *= x;
        }
        x *= x;
        n >>= 1;
    }
    return isNeg ? -res : res;
}
```

### // Solution 2 - Binary Store

```
public double pow(double x, int n) {
    if(n == 0){
        return 1.0;
    }
    double half = pow(x, n / 2);
    if((n & 1) == 0){
        return half * half;
    } else if(n > 0){
        return half * half * x;
    } else {
        return half * half / x;
    }
}
```

6.2 Sqrt, 提供3种做法。第一种, 二分查找; 第二种, 公式法; 第三种, 位运算。

### // Solution 1 - Binary Search

```
public int sqrt(int x) {
    if(x < 0){
        return -1;
    }
    if(x == 0 || x == 1){
        return x;
    }
    int left = 1;
    int right = x / 2 + 1;
    while(left <= right){
        int mid = left + (right - left) / 2;
        if(mid <= x / mid && x / (mid + 1) < (mid + 1)){
            return mid;
        }
        if(mid > x / mid){
            right = mid - 1;
        } else {
            left = mid + 1;
        }
    }
    return -1;
}
```

// **Solution 2 - Formula**

```
public int sqrt(int x) {  
    if(x < 0){  
        return -1;  
    }  
    if(x == 0 || x == 1){  
        return x;  
    }  
    double lastY = 0;  
    double y = 1;  
    while(y != lastY){  
        lastY = y;  
        y = (y + x / y) / 2;  
    }  
    return (int)y;  
}
```

// **Solution 3 - Bit Manipulation**

```
public int sqrt(int x) {  
    int res = 0;  
    int bit = 1 << 16;  
    while(bit > 0){  
        res |= bit;  
        if(res == x / res){  
            break;  
        } else if(res > x / res){  
            res ^= bit;  
        }  
        bit >>= 1;  
    }  
    return res;  
}
```

6.3 **Factorial Trailing Zero**，由于尾数中的0只可能由是 $2 * 5$ 形成，而2的数量肯定是充足的，所以只要计算5的数量即可得到最终尾数中0的数量。

```
public int trailingZeroes(int n) {  
    int res = 0;  
    while(n > 0){  
        n /= 5;  
        res += n;  
    }  
    return res;  
}
```

```

        res += n;
    }
    return res;
}

```

6.4 [Check Power of 2 - O\(1\)](#)，知之为知之，不知为不知。

```

public boolean checkPowerOf2(int n) {
    if(n == 0 || n == Integer.MIN_VALUE){
        return false;
    }
    return (n & (n - 1)) == 0;
}

```

7.1 [Partition Array](#)，这道题是Quick Sort的一个Subroutine，应该好好练习。

```

public int partitionArray(ArrayList<Integer> nums, int k) {
    if(nums == null || nums.size() == 0){
        return 0;
    }
    int left = 0;
    int right = nums.size() - 1;
    while(true){
        while(nums.get(left) < k){
            left++;
            if(left == nums.size()){
                break;
            }
        }
        while(nums.get(right) >= k){
            right--;
            if(right == 0){
                break;
            }
        }
        if(left >= right){
            break;
        }
        int temp = nums.get(left);
        nums.set(left, nums.get(right));
        nums.set(right, temp);
    }
}

```



```

    }
    if(left == right && right == 0){
        return 0;
    }
    if(left == right && left == nums.size() - 1){
        return nums.size();
    }
    return right + 1;
}

```

7.2 [Sort Letters by Case](#)，思路与[Partition Array](#)一样。

```

public void sortLetters(char[] chars) {
    if(chars == null || chars.length <= 1){
        return;
    }
    int left = 0;
    int right = chars.length - 1;
    while(true){
        while('a' <= chars[left] && chars[left] <= 'z'){
            left++;
            if(left == chars.length - 1){
                break;
            }
        }
        while('A' <= chars[right] && chars[right] <= 'Z'){
            right--;
            if(right == 0){
                break;
            }
        }
        if(left >= right){
            break;
        }
        char temp = chars[left];
        chars[left] = chars[right];
        chars[right] = temp;
    }
}

```

7.3 [Sort Colors](#), 提供2种解法。第一种, Counting Sort思想; 第二种, Two Pointers思想。

**// Solution 1 - Counting Sort**

```
public void sortColors(int[] A) {
    if(A == null || A.length == 0){
        return;
    }
    int[] num = new int[3];
    int[] res = new int[A.length];
    for(int i = 0; i < A.length; i++){
        num[A[i]]++;
    }
    for(int i = 1; i < num.length; i++){
        num[i] += num[i - 1];
    }
    for(int i = A.length - 1; i >= 0; i--){
        res[--num[A[i]]] = A[i];
    }
    for(int i = 0; i < A.length; i++){
        A[i] = res[i];
    }
}
```

**// Solution 2 - Double Pointers**

```
public void sortColors(int[] A) {
    if(A == null || A.length == 0){
        return;
    }
    int idx0 = 0;
    int idx1 = 0;
    for(int i = 0; i < A.length; i++){
        if(A[i] == 0){
            A[i] = 2;
            A[idx1++] = 1;
            A[idx0++] = 0;
        } else if(A[i] == 1){
            A[i] = 2;
            A[idx1++] = 1;
        }
    }
}
```

7.4 [Interleaving Positive and Negative Numbers](#), 一定要完美的岔开正数和负数。如果正数比负数多, 那么开头第一个数一定要为正数, 反之为负数。

```
public int[] rerange(int[] A) {
    if(A == null || A.length <= 1){
        return A;
    }
    int posNum = 0;
    for(int num : A){
        posNum = num > 0 ? posNum + 1 : posNum;
    }
    int posIdx = posNum > (A.length - posNum) ? 0 : 1;
    int negIdx = posNum > (A.length - posNum) ? 1 : 0;
    while(true){
        while(A[posIdx] > 0){
            posIdx += 2;
            if(posIdx > A.length - 1){
                break;
            }
        }
        while(A[negIdx] < 0){
            negIdx += 2;
            if(negIdx > A.length - 1){
                break;
            }
        }
        if(posIdx > A.length - 1 || negIdx > A.length - 1){
            break;
        }
        int temp = A[posIdx];
        A[posIdx] = A[negIdx];
        A[negIdx] = temp;
    }
    return A;
}
```