#### CS711008Z Algorithm Design and Analysis

Lecture 7. Kruskal's algorithm for Minimum Spanning Tree <sup>1</sup>

Dongbo Bu

Institute of Computing Technology Chinese Academy of Sciences, Beijing, China

 $<sup>^1</sup> The$  slides were made based on Chapter 5 of Algorithms, and Data Structure by Ellis Horowitz.

#### Kruskal's MST algorithm

#### Kruskal's algorithm [1956]

 Basic idea: during the execution, F is always an acyclic forest, and the safe edge added to F is always a least-weight edge connecting two distinct components.

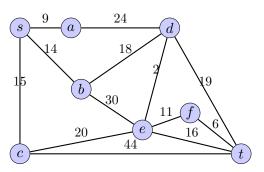


Figure 1: Joseph Kruskal

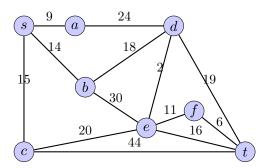
# Kruskal's algorithm [1956]

```
MST-Kruskal(G, W)
 1: F = \{\};
 2: for all vertex v \in V do
 3: MakeSet(v);
 4: end for
 5: sort the edges of E into nondecreasing order by weight W;
 6: for each edge (u,v) \in E in the order do
    if FINDSet(u) \neq FINDSet(v) then
   F = F \cup \{(u, v)\};
    Union (u, v);
10: end if
11: end for
```

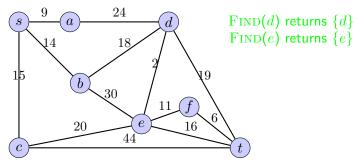
Here, Union-Find structure is used to detect whether a set of edges form a cycle.



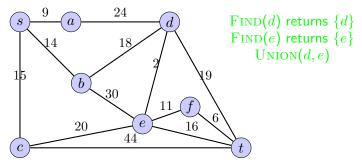
Step 1



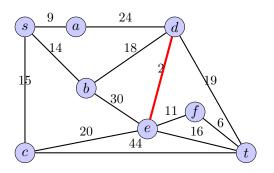
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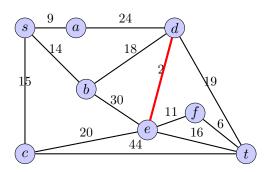
Step 1



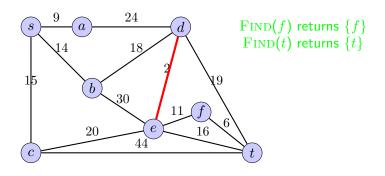
Step 1



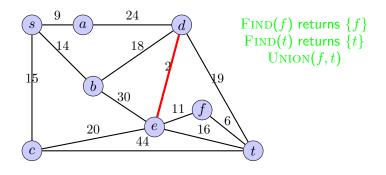
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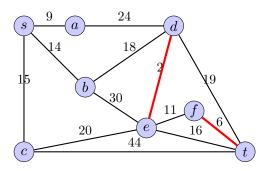
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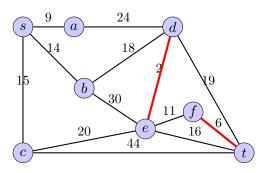
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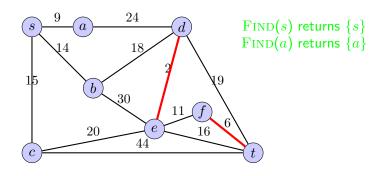
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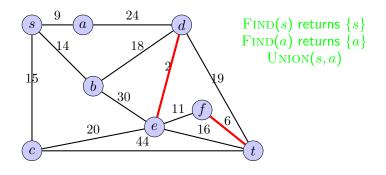
Step 3



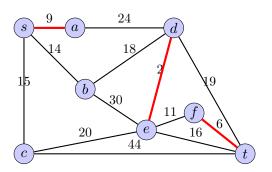
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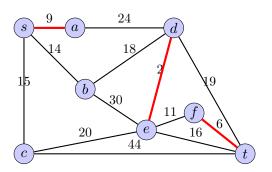
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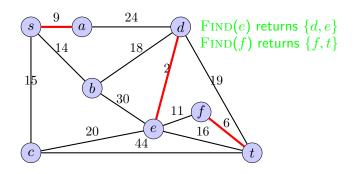
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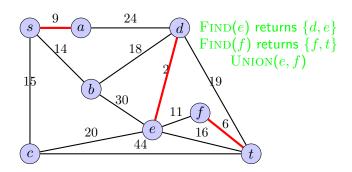
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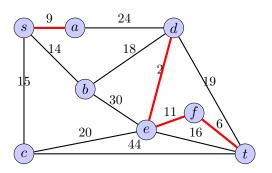
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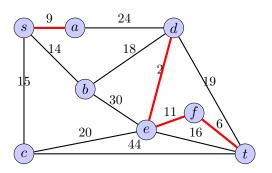
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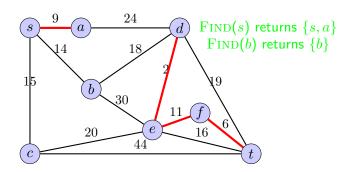
Step 4



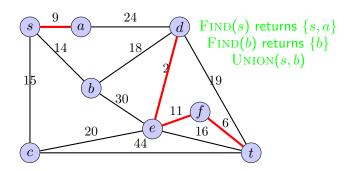
Step 5



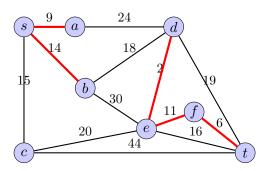
Step 5 Edge weight: 2,6,9,11,14,15,16,18,19,20,24,30,44 Disjoint sets:  $\{a,s\},\{b\},\{c\},\{d,e,f,t\}$ 



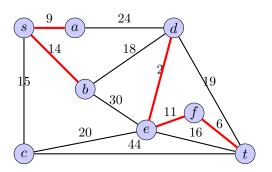
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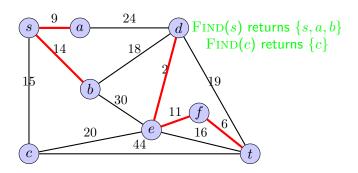
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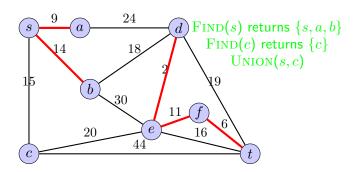
Step 6



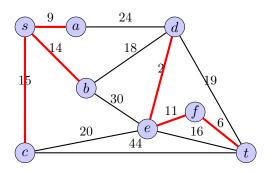
Step 6



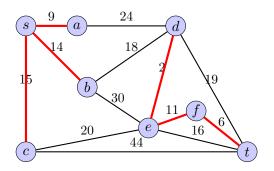
Step 6



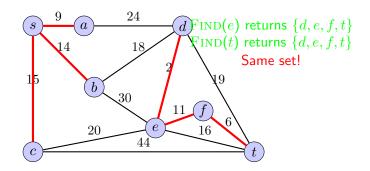
 $\begin{array}{c} {\bf Step~6}\\ {\bf Edge~weight:}~~2,6,9,11,14,15,16,18,19,20,24,30,44\\ {\bf Disjoint~sets:}~~\{a,s,b,c\},\{d,e,f,t\} \end{array}$ 



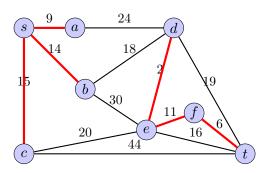
Step 7 
Edge weight: 2,6,9,11,14,15,16,18,19,20,24,30,44 
Disjoint sets:  $\{a,s,b,c\},\{d,e,f,t\}$ 



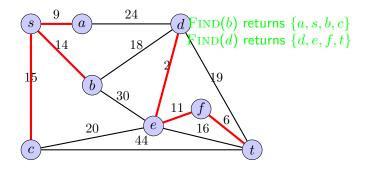
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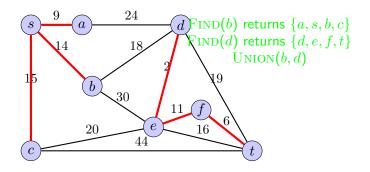
Step 8



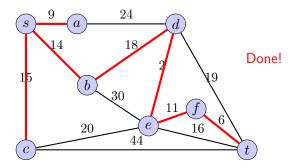
 $\begin{array}{c} {\bf Step~8}\\ {\bf Edge~weight:}~~2,6,9,11,14,15,16,18,19,20,24,30,44\\ {\bf Disjoint~sets:}~~\{a,s,b,c\},\{d,e,f,t\} \end{array}$ 



 $\begin{array}{c} {\bf Step~8}\\ {\bf Edge~weight:}~~2,6,9,11,14,15,16,18,19,20,24,30,44\\ {\bf Disjoint~sets:}~~\{a,s,b,c\},\{d,e,f,t\} \end{array}$ 



Step 8 Edge weight: 2,6,9,11,14,15,16,18,19,20,24,30,44 Disjoint sets:  $\{a,s,b,c,d,e,f,t\}$ 



# Time complexity of KRUSKAL'S MST algorithm

Operation	Array	Tree	Link-by-size	Link-by-size +
				path compression
MakeSet	1	1	1	1
FIND	1	n	$\log n$	$lg^*n$
Union	n	1	$\log n$	$lg^*n$
Kruskal's MST	$O(n^2)$	O(mn)	$O(m \log n)$	$O(mlg^*n)$

Kruskal's MST algorithm: n MakeSet, n-1 Union, and m Find.