



# Numerical Methods for Elliptic Equations

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Fall 2001

## Solution Methods for Elliptic Equations

### Direct Methods

- Cramer's rule
- Gaussian elimination
- Tridiagonal matrix algorithm (Thomas algorithm)
- Block tridiagonal matrix algorithm

### Iterative Methods

- Jacobi, Gauss-Seidel iteration
- Successive overrelaxation (SOR)
- Multigrid method

A model equation for elliptic problems

$$\frac{\partial^2 f}{\partial x^2} + \frac{\partial^2 f}{\partial y^2} = S; \quad \nabla^2 f = S$$

Poisson equation (Laplace equation if  $S = 0$ ).

On the boundaries (BC)

$$f = f_0(x, y) \quad \text{Dirichlet}$$

$$\frac{\partial f}{\partial n} = g_0(x, y) \quad \text{Neumann}$$

## Examples of Elliptic Equations

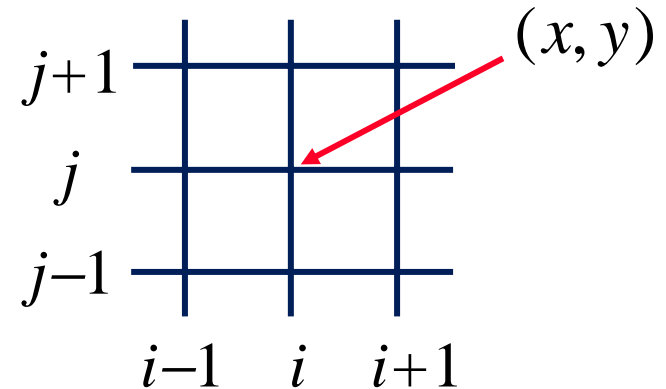
$$\nabla^2 T = -\frac{\dot{q}}{k} \quad \text{Steady conduction equation}$$

$$\nabla^2 \psi = -\omega \quad \text{2-D stream function equation}$$

$$\nabla_h^2 P_{i,j} = \frac{1}{\Delta t} \nabla_h \cdot \mathbf{u}_{i,j}^t \quad \begin{array}{l} \text{Projection method} \\ \text{(Step 2)} \end{array}$$

Solving the Poisson equation

$$\frac{\partial^2 f}{\partial x^2} + \frac{\partial^2 f}{\partial y^2} = S$$



Applying central differencing

$$\frac{f_{i+1,j} - 2f_{i,j} + f_{i-1,j}}{\Delta x^2} + \frac{f_{i,j+1} - 2f_{i,j} + f_{i,j-1}}{\Delta y^2} = S_{i,j}$$

for which the modified equation becomes

$$f_{xx} + f_{yy} = S - \frac{1}{12} [f_{xxxx} \Delta x^2 + f_{yyyy} \Delta y^2] + \dots$$

**If**  $\Delta x = \Delta y = h \Rightarrow f_{i+1,j} + f_{i-1,j} + f_{i,j+1} + f_{i,j-1} - 4f_{i,j} = h^2 S_{i,j}$

[illegible]

Sparse matrix: only 5 non-zero entries in each row

Ultimately, the difference form of the Poisson equation boils down to solving for

$$[A]\mathbf{f} = \mathbf{b}$$

Hence,

$$\mathbf{f} = [A]^{-1}\mathbf{b}$$

Direct method:

- Solving inverse matrix directly (Cramer's rule)
- Inverting matrix more cleverly (Gaussian elimination)
- Other (L-U decomposition, Thomas algorithm)

## Cramer's Rule

- Check elementary linear algebra book
- Extremely time consuming for large matrices

Operation count:  $(n+1)!$

For example,  $11 \times 11$  matrix:

$$(9 \times 9 + 1)! = 4.75 \times 10^{22}$$

Using 1 Gigafllops ( $10^9$  flops) computer

$$3.20 \times 10^{100} \text{ years}!!!!!!$$





## Gaussian Elimination

- Pivoting: rearranging equations to put the largest coefficient on the main diagonal.
- Eliminate the column below main diagonal.
- Repeat until the last equation is reached.
- Back-substitution

$$\begin{array}{rcl} a_{11}f_1 + a_{12}f_2 + \cdots = c_1 & & a_{11}f_1 + a_{12}f_2 + \cdots = c_1 \\ a_{21}f_1 + a_{22}f_2 + \cdots = c_2 & \Rightarrow & a'_{22}f_2 + \cdots = c_2 \\ \vdots & & \vdots \\ a_{n1}f_1 + a_{n2}f_2 + \cdots = c_n & & a'_{nn}f_n = c_n \end{array}$$

- Special case: tri-diagonal matrix - Thomas algorithm

## Discretized Poisson Equation

$$f_{i+1,j} + f_{i-1,j} + f_{i,j+1} + f_{i,j-1} - 4f_{i,j} = h^2 S_{i,j}$$

Rearranging for  $f_{i,j}$ 

$$f_{i,j} = \frac{1}{4} [f_{i+1,j} + f_{i-1,j} + f_{i,j+1} + f_{i,j-1} - h^2 S_{i,j}]$$

$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i+1,j}^n + f_{i-1,j}^n + f_{i,j+1}^n + f_{i,j-1}^n - h^2 S_{i,j}] \quad \text{Jacobi}$$

$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i+1,j}^n + f_{i-1,j}^{n+1} + f_{i,j+1}^n + f_{i,j-1}^{n+1} - h^2 S_{i,j}] \quad \text{Gauss-Seidel}$$

$$f_{i,j}^{n+1} = \frac{\beta}{4} [f_{i+1,j}^n + f_{i-1,j}^{n+1} + f_{i,j+1}^n + f_{i,j-1}^{n+1} - h^2 S_{i,j}] + (1-\beta)f_{i,j}^n \quad \text{SOR}$$

## A One Dimensional Example

An equation in the form

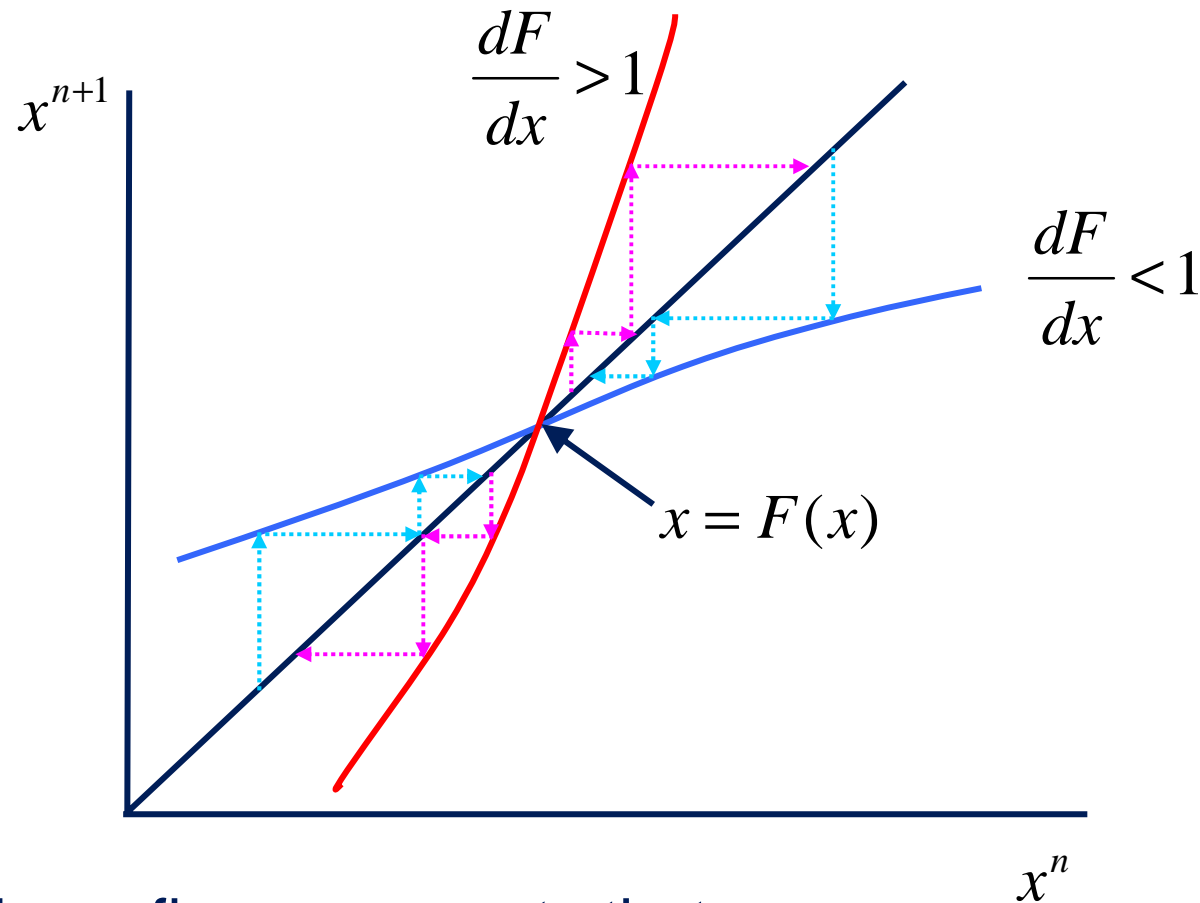
$$x = F(x)$$

can be solved by iterative procedure:

$$x^{n+1} = F(x^n)$$

for which convergence is achieved if

$$x^{n+1} \approx x^n \quad \text{or} \quad \left| \frac{x^{n+1}}{x^n} - 1 \right| < \varepsilon$$



The above figure suggests that

$$\left| \frac{dF}{dx} \right| \leq 1 \quad \text{for convergence}$$

For multi-dimensional problems

$$\mathbf{f}^{n+1} = [M] \mathbf{f}^n$$

It can be shown that the space of  $\mathbf{f}$  can be “spanned” by eigenvectors

$$[M] \mathbf{q}_i = \lambda_i \mathbf{q}_i, \quad i = 1, \dots, M \times N$$

such that

$$\mathbf{f} = y_1 \mathbf{q}_1 + y_2 \mathbf{q}_2 + \dots + y_{M \times N} \mathbf{q}_{M \times N}$$

Hence,

$$\begin{aligned} [M] \mathbf{f} &= [M] (y_1 \mathbf{q}_1 + y_2 \mathbf{q}_2 + \dots + y_{M \times N} \mathbf{q}_{M \times N}) \\ &= y_1 [M] \mathbf{q}_1 + y_2 [M] \mathbf{q}_2 + \dots \\ &= y_1 \lambda_1 \mathbf{q}_1 + y_2 \lambda_2 \mathbf{q}_2 + \dots \end{aligned}$$

Therefore, the problem

$$\mathbf{f}^{n+1} = [\mathbf{M}] \mathbf{f}^n$$

is equivalent to

$$y_1^{n+1} \mathbf{q}_1 + y_2^{n+1} \mathbf{q}_2 + \cdots = y_1^n \lambda_1 \mathbf{q}_1 + y_2^n \lambda_2 \mathbf{q}_2 + \cdots$$

or

$$y_1^{n+1} = \lambda_1 y_1^n$$

$$y_2^{n+1} = \lambda_2 y_2^n$$

$$\vdots$$

yielding a convergence condition

$$\lambda_{\max} \leq 1$$

General iterative procedure

$$[A]\mathbf{f} = \mathbf{b}$$

Let  $[A] = [A_1] - [A_2]$

$$[A_1]\mathbf{f} = [A_2]\mathbf{f} + \mathbf{b}$$

An iterative scheme is constructed as

$$[A_1]\mathbf{f}^{n+1} = [A_2]\mathbf{f}^n + \mathbf{b}$$

For example,

$$[A_1] = [D] = -\frac{1}{4}[I], \quad [A_2] = [B] = [A_1] - [D] \quad \text{Jacobi}$$

$$[A_1] = [D] - [L], \quad [A_2] = [U] \quad \text{Gauss-Seidel}$$

$$[A_1]\mathbf{f}^{n+1} = [A_2]\mathbf{f}^n + \mathbf{b}$$

Requirements:

1.  $[A_1]$  should be invertible.
2. Iteration should converge, i.e.

$$\lim_{n \rightarrow \infty} \mathbf{f}^n = \mathbf{f}$$

Define error at n-th iteration

$$\boldsymbol{\varepsilon}^n = \mathbf{f} - \mathbf{f}^n$$

$$[A_1]\boldsymbol{\varepsilon}^{n+1} = [A_2]\boldsymbol{\varepsilon}^n \quad \boldsymbol{\varepsilon}^{n+1} = [A_1]^{-1}[A_2]\boldsymbol{\varepsilon}^n$$

$$\boldsymbol{\varepsilon}^n = ([A_1]^{-1}[A_2])^n \boldsymbol{\varepsilon}^0$$



Condition for convergence

$$\lim_{n \rightarrow \infty} \epsilon^n = 0$$

which requires  $\lim_{n \rightarrow \infty} ([A_1]^{-1} [A_2])^n = 0$

This is achieved if the modulus of the eigenvalues are less than unity.

Therefore, the convergence condition becomes:

$$\rho = |\lambda_i|_{\max} \leq 1$$

$\rho$  Spectral radius of convergence

$\lambda_i$  Eigenvalues of matrix  $[A_1]^{-1} [A_2]$

Jacobi Iteration Method for Poisson Equation  
(2nd order central difference with uniform mesh)

$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i+1,j}^n + f_{i-1,j}^n + f_{i,j+1}^n + f_{i,j-1}^n - h^2 S_{i,j}]$$

$$[A_1]^{-1} = [I]$$

and using a discrete analog of separation of variables,  
it can be shown that the eigenvalues of  $[A_2]$  are

$$\lambda_{mn} = \frac{1}{2} \left[ \cos \frac{m\pi}{M} + \cos \frac{n\pi}{N} \right], \quad m = 1, \dots, M-1, \quad n = 1, \dots, N-1$$

Therefore,  $|\lambda_{mn}| \leq 1$  and the Jacobi method **converges**.

For a large matrix

$$\lambda_{mn,\max} = \frac{1}{2} \left[ \cos \frac{\pi}{M} + \cos \frac{\pi}{N} \right]$$
$$\cong 1 - \frac{1}{4} \left[ \frac{\pi^2}{M^2} + \frac{\pi^2}{N^2} \right] + \dots$$

largest eigenvalues for  
 $m = 1, n = 1$

Thus, for a large matrix,  $|\lambda_{mn,\max}|$  is only slightly less than unity.

$\Rightarrow$  Very slow convergence

## Gauss-Seidel Iteration Method for Poisson Equation (2nd order central difference with uniform mesh)

$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i+1,j}^n + f_{i-1,j}^{n+1} + f_{i,j+1}^n + f_{i,j-1}^{n+1} - h^2 S_{i,j}]$$

$$[A_1] = [D] - [L], \quad [A_2] = [U]$$

$$\underbrace{\begin{bmatrix} 4 & 0 & 0 & 0 & \cdots & 0 \\ -1 & 4 & 0 & 0 & \cdots & 0 \\ 0 & -1 & 4 & 0 & \cdots & 0 \\ 0 & 0 & -1 & 4 & \cdots & 0 \\ \vdots & \vdots & \vdots & & \cdots & \vdots \\ 0 & \cdots & \cdots & 0 & -1 & 4 \end{bmatrix}}_{A_1} \underbrace{\begin{bmatrix} f_{1,1}^{n+1} \\ f_{1,2}^{n+1} \\ \vdots \\ \vdots \\ f_{I,J}^{n+1} \end{bmatrix}} = \underbrace{\begin{bmatrix} 0 & 1 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & & \vdots \\ \vdots & \vdots & & 0 & & \vdots \\ 0 & 0 & 0 & 0 & \cdots & 0 \end{bmatrix}}_{A_2} \underbrace{\begin{bmatrix} f_{1,1}^n \\ f_{1,2}^n \\ \vdots \\ \vdots \\ f_{M,N}^n \end{bmatrix}} - h^2 \underbrace{\begin{bmatrix} S_{1,1} \\ S_{1,2} \\ \vdots \\ \vdots \\ S_{M,N} \end{bmatrix}}$$

It can be shown that the eigenvalues of matrix  $[A_1]^{-1}[A_2]$   
Are simply square of the eigenvalues of Jacobi method

$$\lambda_{mn,\max} = \frac{1}{4} \left[ \cos \frac{\pi}{M} + \cos \frac{\pi}{N} \right]^2$$

Thus, Gauss-Seidel method is twice as fast as the Jacobi method.

## Successive Overrelaxation

Consider the Gauss-Seidel method

$$([D] - [L])\mathbf{f}^{n+1} = [U]\mathbf{f}^n + \mathbf{b}$$

If Gauss-Seidel is an attempt to change the solution as

$$\mathbf{f}^{n+1} = \mathbf{f}^n + \mathbf{d}$$

Can we accelerate the change by a parameter?

$$\mathbf{f}^{n+1} = \mathbf{f}^n + \beta \mathbf{d}, \quad \beta > 1$$

Hence, SOR first uses Gauss-Seidel to compute intermediate solution,  $\tilde{\mathbf{f}}$

$$([D] - [L])\tilde{\mathbf{f}} = [U]\mathbf{f}^n + \mathbf{b} \quad \text{or} \quad [D]\tilde{\mathbf{f}} = [L]\tilde{\mathbf{f}} + [U]\mathbf{f}^n + \mathbf{b}$$

Then accelerate the next iteration solution

$$\mathbf{f}^{n+1} = \mathbf{f}^n + \beta(\tilde{\mathbf{f}} - \mathbf{f}^n) = \beta\tilde{\mathbf{f}} + (1 - \beta)\mathbf{f}^n$$

Combining the two steps

$$\mathbf{f}^{n+1} = \frac{\beta}{c} ([L]\tilde{\mathbf{f}} + [U]\mathbf{f}^n + \mathbf{b}) + (1 - \beta)\mathbf{f}^n$$

$$\text{Note that} \quad [D] = c[I] \quad [L]\tilde{\mathbf{f}} = [L]\mathbf{f}^{n+1}$$

Example: Poisson Equation (2<sup>nd</sup> order CS, uniform mesh)

$$f_{i+1,j} + f_{i-1,j} + f_{i,j+1} + f_{i,j-1} - 4f_{i,j} = h^2 S_{i,j}$$

G-S

$$[D]\tilde{\mathbf{f}} = [L]\tilde{\mathbf{f}} + [U]\mathbf{f}^n + \mathbf{b}$$

$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i-1,j}^{n+1} + f_{i,j-1}^{n+1} + f_{i+1,j}^n + f_{i,j+1}^n - h^2 S_{i,j}]$$

Combining

$$\mathbf{f}^{n+1} = \frac{\beta}{c} ([L]\tilde{\mathbf{f}} + [U]\mathbf{f}^n + \mathbf{b}) + (1 - \beta)\mathbf{f}^n$$

$$f_{i,j}^{n+1} = \frac{\beta}{4} [f_{i-1,j}^{n+1} + f_{i,j-1}^{n+1} + f_{i+1,j}^n + f_{i,j+1}^n - h^2 S_{i,j}] + (1 - \beta)f_{i,j}^n$$



## Convergence of SOR

From  $[D]\tilde{\mathbf{f}} = [L]\tilde{\mathbf{f}} + [U]\mathbf{f}^n + \mathbf{b}$

$$\mathbf{f}^{n+1} = \mathbf{f}^n + \beta(\tilde{\mathbf{f}} - \mathbf{f}^n)$$

Eliminating  $\tilde{\mathbf{f}}$  and solving for  $\mathbf{f}^{n+1}$

$$\mathbf{f}^{n+1} = \overbrace{\left([I] - \beta[D]^{-1}[L]\right)^{-1} \left\{ (1-\beta)[I] + \beta[D]^{-1}[U] \right\}}^{[M]_{SOR}} \mathbf{f}^n + \left([I] - \beta[D]^{-1}[L]\right)^{-1} \beta[D]^{-1} \mathbf{b}$$

Convergence depends on the eigenvalues of  $[M]_{SOR}$

For the discretized Poisson operator, it can be shown that

$$\mu^{1/2} = \frac{1}{2} \left[ \lambda \beta + \sqrt{\lambda^2 \beta^2 - 4(\beta - 1)} \right]$$

where  $\lambda$  is an eigenvalue of the Jacobi matrix

$$[M]_j = [D]^{-1} ([L] + [U])$$

Note that  $\mu = \lambda^2$  if  $\beta = 1$  (Gauss-Seidel)

Minimum  $\mu$  occurs at

$$\beta_{\text{opt}} = \frac{2}{1 + \sqrt{1 - \lambda_{\text{max}}^2}}$$

Typically

$$\beta_{\text{opt}} \approx 1.7 \sim 1.9$$

$$|\lambda_{\text{max}}| \approx 1$$

For problems with irregular geometry and non-uniform mesh,  $\beta_{\text{opt}}$  must be found by trial and error.

### Typical Comparison Chart

	$\lambda_{\text{max}} (\mu_{\text{max}})$	Iterations
Jacobi	0.9945	1250
Gauss-Seidel	0.9890	625
SOR	0.7906	29

Ferziger, J. H., *Numerical Method for Engineering Application* (1981)

In large computer application (vector or parallel platform),  
SOR faces difficulties in using constantly updated values.

Remedy: Two separate grid system (red & black)

do i = 1, nx, 2

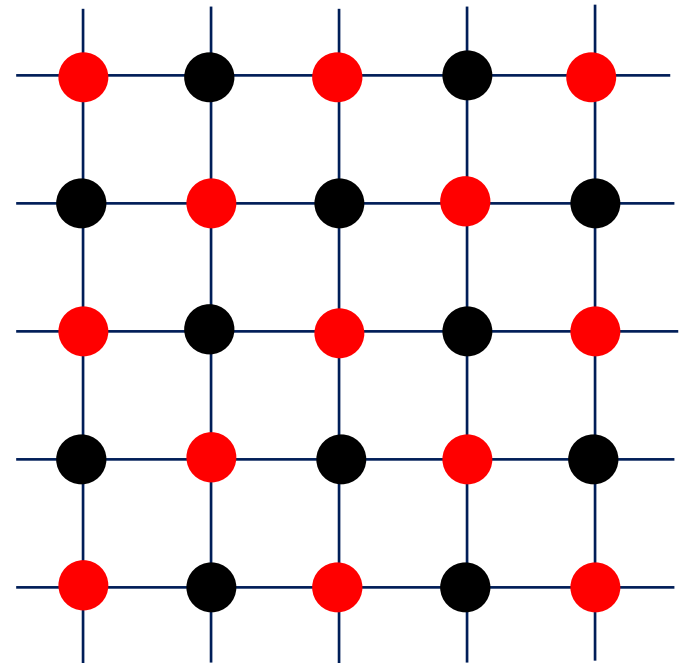
$$f_{i,j} = \frac{\beta}{4} [f_{i-1,j} + f_{i,j-1} + f_{i+1,j} + f_{i,j+1} - h^2 S_{i,j}] + (1 - \beta) f_{i,j}$$

enddo

do i = 2, nx, 2

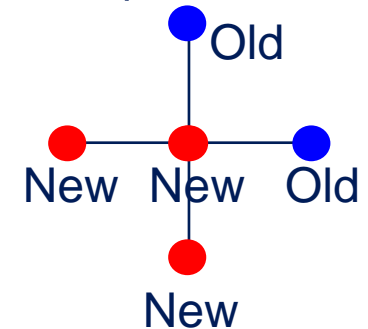
$$f_{i,j} = \frac{\beta}{4} [f_{i-1,j} + f_{i,j-1} + f_{i+1,j} + f_{i,j+1} - h^2 S_{i,j}] + (1 - \beta) f_{i,j}$$

enddo



## Line Relaxation Method (Line Gauss-Seidel Method)

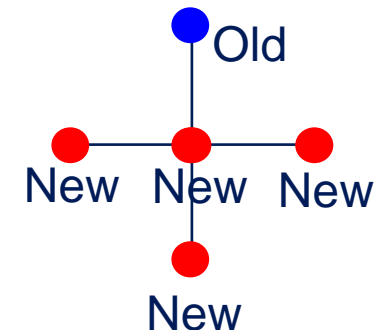
$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i-1,j}^{n+1} + f_{i,j-1}^{n+1} + f_{i+1,j}^n + f_{i,j+1}^n - h^2 S_{i,j}]$$



## Adding one more coupling (E)

$$f_{i,j}^{n+1} = \frac{1}{4} [f_{i-1,j}^{n+1} + f_{i,j-1}^{n+1} + f_{i+1,j}^{n+1} + f_{i,j+1}^n - h^2 S_{i,j}]$$

$$-\frac{1}{4} f_{i-1,j}^{n+1} + f_{i,j}^{n+1} - \frac{1}{4} f_{i+1,j}^{n+1} = \frac{1}{4} [f_{i,j-1}^{n+1} + f_{i,j+1}^n - h^2 S_{i,j}]$$



⇒ Thomas algorithm

SLOR = Line Relaxation + Overrelaxation

Apply line relaxation for intermediate solution

$$-\frac{1}{4} \tilde{f}_{i-1,j} + \tilde{f}_{i,j} - \frac{1}{4} \tilde{f}_{i+1,j} = \frac{1}{4} [f_{i,j-1}^{n+1} + f_{i,j+1}^n - h^2 S_{i,j}]$$

and then overrelax

$$f_{i,j}^{n+1} = \beta \tilde{f}_{i,j} + (1 - \beta) f_{i,j}^n$$

which is no more complicated than line relaxation.

## Notes on SLOR

- Exact eigenvalues are unknown.
- To ensure convergence,  $\beta \leq 2$
- Converges approximately twice as fast as Gauss-Seidel.
- May be faster than pointwise SOR, but each iteration takes longer with Thomas algorithm.
- Improved convergence is due to the direct effect of the boundary condition in each row.

ADI for elliptic equation is analogous to ADI in parabolic equation

$$\frac{\partial f}{\partial t} = \alpha \left[ \frac{\partial^2 f}{\partial x^2} + \frac{\partial^2 f}{\partial y^2} \right] - S$$

In discrete form

$$f_{i,j}^{n+1} - f_{i,j}^n = \alpha \Delta t [\delta_{xx} f + \delta_{yy} f - S]$$

$$\delta_{xx} ( ) = \frac{1}{\Delta x^2} [( )_{i-1,j} - 2( )_{i,j} + ( )_{i+1,j}]$$

$$\delta_{yy} ( ) = \frac{1}{\Delta y^2} [( )_{i,j-1} - 2( )_{i,j} + ( )_{i,j+1}]$$

and take it to the limit to obtain the steady solution.  $\left( \frac{\partial f}{\partial t} = 0 \right)$



ADI for  $f_{i,j}^{n+1} - f_{i,j}^n = \alpha \Delta t [\delta_{xx} f + \delta_{yy} f] - S$

is written as

$$f^{n+1/2} - f^n = \frac{\alpha \Delta t}{2h^2} \left[ (f_{i+1,j}^{n+1/2} - 2f_{i,j}^{n+1/2} + f_{i-1,j}^{n+1/2}) + (f_{i,j+1}^n - 2f_{i,j}^n + f_{i,j-1}^n) - S_{i,j} \right]$$

$$f^{n+1} - f^{n+1/2} = \frac{\alpha \Delta t}{2h^2} \left[ (f_{i+1,j}^{n+1/2} - 2f_{i,j}^{n+1/2} + f_{i-1,j}^{n+1/2}) + (f_{i,j+1}^{n+1} - 2f_{i,j}^{n+1} + f_{i,j-1}^{n+1}) - S_{i,j} \right]$$

or

$$\begin{aligned} (1 - \rho_n \delta_{xx}) f^{n+1/2} &= (1 + \rho_n \delta_{yy}) f^n - \frac{\alpha \Delta t}{2h^2} S_{i,j} \\ (1 - \rho_n \delta_{yy}) f^{n+1} &= (1 + \rho_n \delta_{xx}) f^{n+1/2} - \frac{\alpha \Delta t}{2h^2} S_{i,j} \end{aligned} \quad \left( \rho_n = \frac{\alpha \Delta t_n}{2} \right)$$

## Convergence of ADI

- Iteration parameter  $\rho_n = \frac{\alpha \Delta t_n}{2}$  usually varies with iteration
- For example, Wachspress

$$\frac{\rho_k}{h^2} = \frac{\alpha \Delta t_k}{2h^2} = b \left( \frac{a}{b} \right)^{(k-1)/(n-1)}, \quad k = 1, \dots, n$$

$a$  lower bound eigenvalue

$b$  upper bound eigenvalue

- Comparison with SOR is difficult
- ADI can be efficient if appropriate parameters are found.

Dirichlet conditions are easily implemented.

For Neumann condition, the simplest approach is

$$\frac{\partial f}{\partial n} = 0 \Rightarrow f_{i,0} - f_{i,1} = 0 \quad (1\text{st order})$$

Update interior points  $f_{i,1}, f_{i,2}, f_{i,3}, \dots$  and then set  $f_{i,0} = f_{i,1}$

$\Rightarrow$  This generally does not converge.

Instead, incorporate BC directly into the equations

$$f_{i,1} = \frac{1}{4} [f_{i-1,1} + f_{i+1,1} + f_{i,2} + f_{i,0} - h^2 S_{i,j}]$$

$$f_{i,1} = \frac{1}{3} [f_{i-1,1} + f_{i+1,1} + f_{i,2} - h^2 S_{i,j}]$$

## Basic Idea

Suppose we want to solve a 1-D elliptic equation

$$\frac{\partial^2 f}{\partial x^2} - g = 0$$

An iterative process is analogous to solving an unsteady problem (recall ADI)

$$\frac{\partial f}{\partial t} = \alpha \left( \frac{\partial^2 f}{\partial x^2} - g \right)$$

And take it to the limit  $t \rightarrow \infty$

Analytic Solution Behavior  $\frac{\partial f}{\partial t} = \alpha \left( \frac{\partial^2 f}{\partial x^2} - g \right)$

The equation can be solved by Fourier series

$$f = \sum_k a_k(t) e^{ikx}$$

and assume that  $g$  can be expanded as

$$g = \sum_k b_k k^2 e^{ikx}$$

Substituting into the equation,

$$\sum_k \frac{da_k}{dt} e^{ikx} = -\alpha \sum_k (a_k(t) - b_k) k^2 e^{ikx}$$

Solving for each  $k$

$$\frac{da_k}{dt} = -\alpha k^2 (a_k(t) - b_k)$$

$$a_k(t) - b_k = (a_k(0) - b_k) e^{-\alpha k^2 t} \quad \left( \varepsilon_k = \varepsilon_0 e^{-\alpha k^2 t} \right)$$

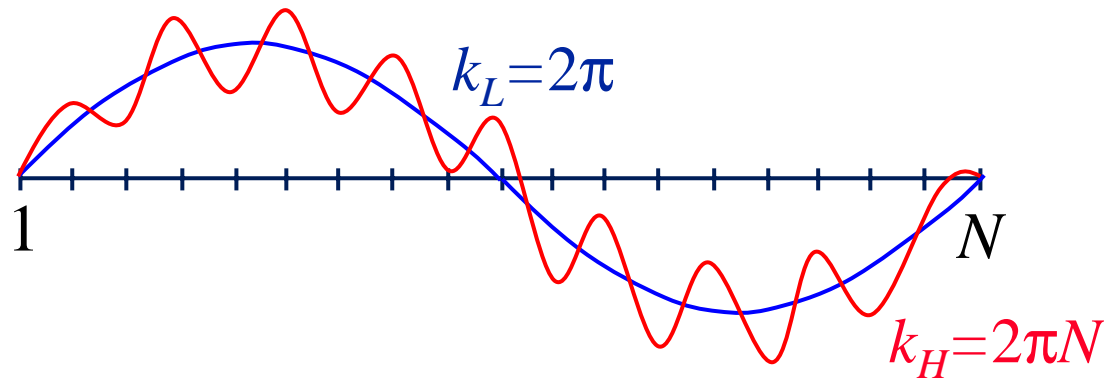
Therefore,  $a_k(t \rightarrow \infty) = b_k$

$$\text{Rate of convergence} \sim \alpha k^2 \quad \left( \tau_c = \frac{1}{\alpha k^2} \right)$$

High wave number modes damps out faster.

For iterative method to solve elliptic equations, there are high wave number and low wave number components of errors that need to be damped out.

Consider a domain  $L=1$  discretized by  $N$  points



For explicit time step, stability condition yields

$$\frac{\alpha \Delta t}{h^2} = \frac{1}{2}; \quad \Delta t = \frac{h^2}{2\alpha}$$

If the error at various wave number decays at

$$\varepsilon = \varepsilon_0 e^{-\alpha k^2 t} \quad \text{where} \quad t = n\Delta t = \frac{nh^2}{2\alpha}$$

$$\Rightarrow \varepsilon / \varepsilon_0 = \exp(-k^2 nh^2 / 2)$$

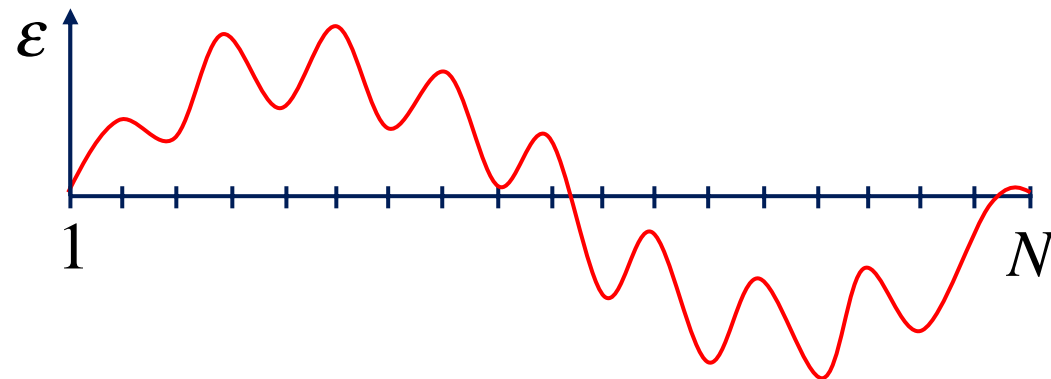
$$\frac{\varepsilon_H}{\varepsilon_0} = \exp(-n\pi^2 N^2 h^2) \quad \text{for } k_H = 2\pi N$$

$$\frac{\varepsilon_L}{\varepsilon_0} = \exp(-n\pi^2 h^2) \quad \text{for } k_L = 2\pi$$

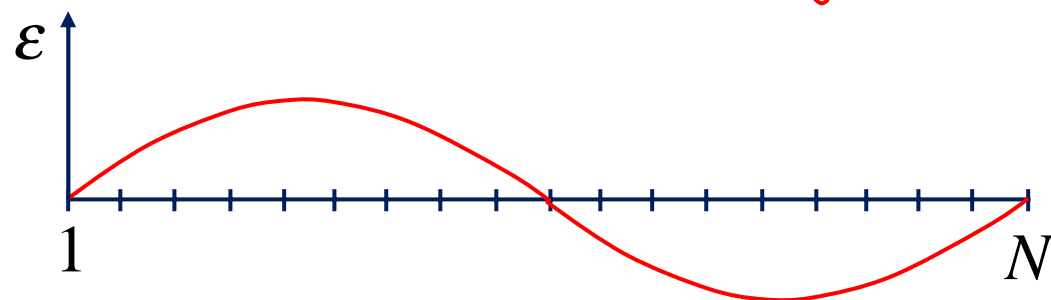
Short-wave errors  
decay much faster



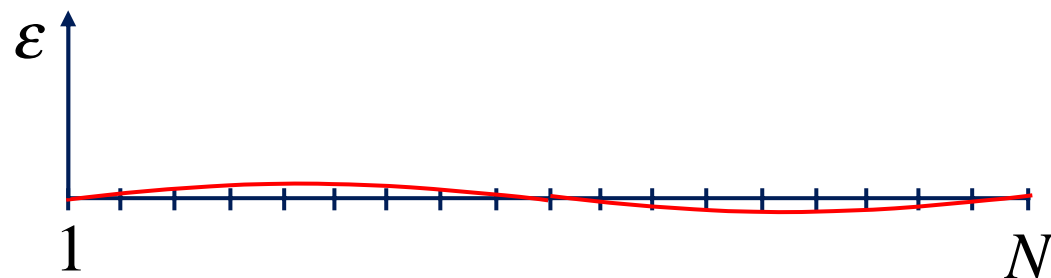
$t = 0$



$t = 5\Delta t$



$t = 500\Delta t$



## Multigrid Method – the Idea

- A low wave number component on a fine grid becomes a high wave number component on a coarse grid

$$\varepsilon_L / \varepsilon_0 = \exp(-n\pi^2 h^2) = \exp(-n\pi^2 L^2 / N^2)$$

- Use a coarse grid system to converge low wave number component of the solution rapidly
- Map it onto the fine grid system to converge high wave number component.

Suppose we are solving a Laplace equation:

$$Lf_{i,j} = \frac{f_{i+1,j} - 2f_{i,j} + f_{i-1,j}}{\Delta x^2} + \frac{f_{i,j+1} - 2f_{i,j} + f_{i,j-1}}{\Delta y^2} = R_{i,j}$$

and iterate until residual becomes smaller than tolerance

$$R_{i,j} \rightarrow 0 \quad \text{or} \quad R_{i,j} < \varepsilon$$

Define

$$f_{i,j} = f_{i,j}^n + \Delta f_{i,j}$$

$f_{i,j}$       Converged solution

$f_{i,j}^n$       Solution after n-th iteration

$\Delta f_{i,j}$       Correction

Since

$$Lf_{i,j} = Lf_{i,j}^n + L\Delta f_{i,j} = 0$$

and

$$Lf_{i,j}^n = R_{i,j}$$

we have

$$L\Delta f_{i,j} = -R_{i,j} \quad \text{Coarse grid correction}$$

Steps:

Fine grid solution

⇒ Interpolation onto coarse grid (restriction)

⇒ Coarse grid correction

⇒ Interpolated onto fine grid (prolongation)

⇒ Converge on fine grid

Two-Level Multigrid Method -  $(n_x+1, n_y+1)$ ;  $(n_x/2+1, n_y/2+1)$ 

1. Do  $n$  iterations on the fine grid ( $n = 3 \sim 4$ ) using G-S

$$L f_{i,j}^n = R_{i,j}$$

2. If  $R_{i,j} < \varepsilon$  then stop.

3. Interpolate the residual  $R_{i,j}$  onto the coarse grid  $I_1^2 R_{i,j}$   
(restriction (injection if  $n_x/2$ ))

4. Iterations on the coarse grid

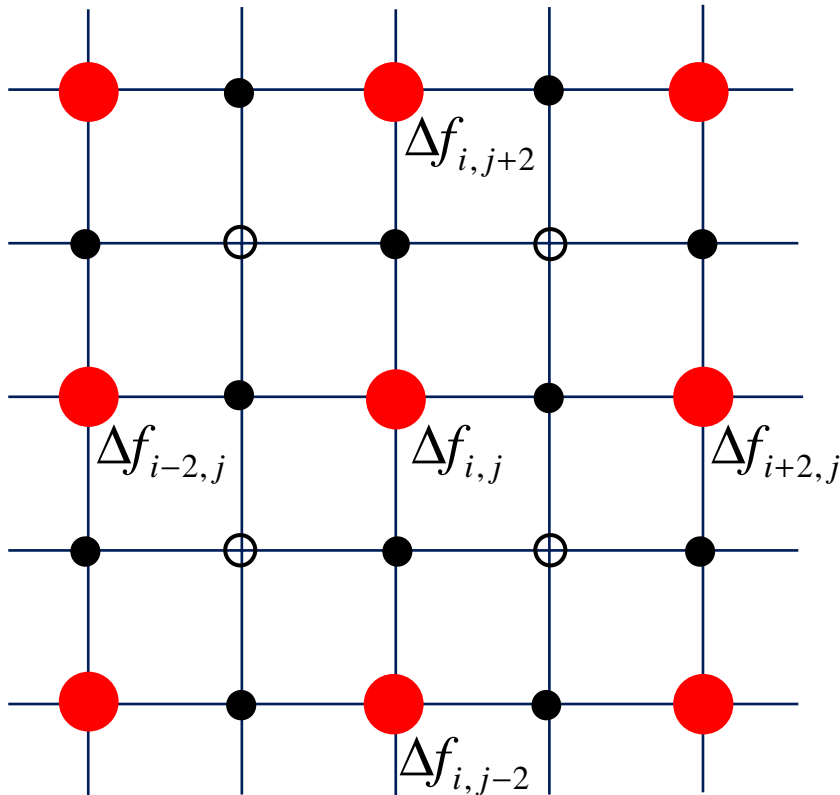
$$L \Delta(f_2)_{i,j} = -I_1^2 R_{i,j}$$

5. Interpolate the correction  $\Delta(f_2)_{i,j}$  onto the fine grid

$$\text{(prolongation)} \quad f_{i,j} = f_{i,j}^n + I_2^1 \Delta(f_2)_{i,j} \quad I_2^1 \Delta(f_2)_{i,j}$$

6. Go to 1.

## Details on Prolongation



Coarse Grid

Fine grid

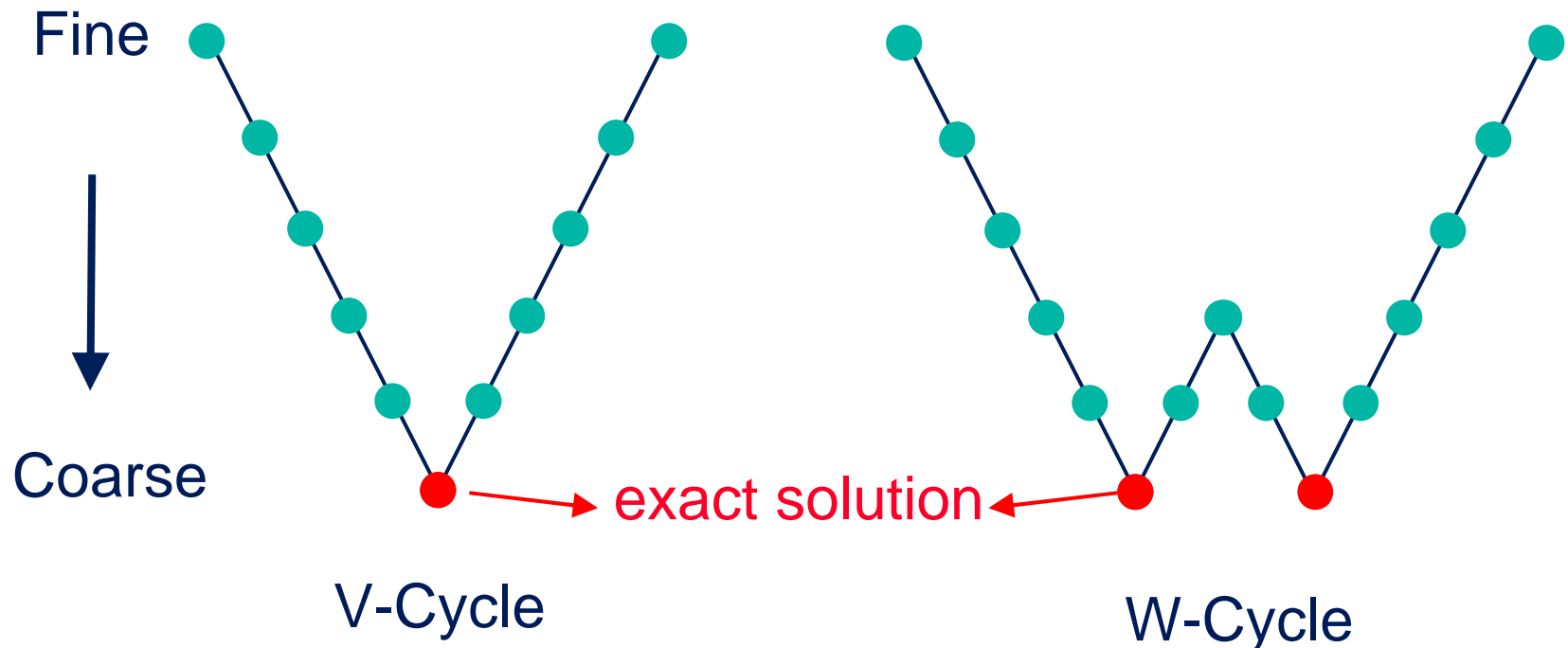
$$\Delta f_{i+1,j} = \frac{1}{2} (\Delta f_{i,j} + \Delta f_{i+2,j})$$

$$\Delta f_{i,j+1} = \frac{1}{2} (\Delta f_{i,j} + \Delta f_{i,j+2})$$

$$\begin{aligned} \Delta f_{i+1,j+1} &= \frac{1}{4} (\Delta f_{i+1,j} + \Delta f_{i,j+1} + \Delta f_{i+1,j+2} + \Delta f_{i+2,j+1}) \\ &= \frac{1}{4} (\Delta f_{i,j} + \Delta f_{i+2,j} + \Delta f_{i,j+2} + \Delta f_{i+2,j+2}) \end{aligned}$$

The process can be extended to multi-level grids  
 $(n_x+1, n_y+1)$ ,  $(n_x/2+1, n_y/2+1)$ ,  $(n_x/4+1, n_y/4+1)$ , ...,  $(3,3)$

### Various Strategies



Example: 4-Level V-Cycle

**Fine  $\Rightarrow$  Coarse**

1. Iterate on Grid 1 for n times  $Lf_{i,j}^n = 0 \Rightarrow R_{i,j}^1$

2. Restriction by injection to Grid 2  $I_1^2 R_{i,j}^1$

3. Iterate on Grid 2  $L\Delta(f_2)_{i,j} = -I_1^2 R_{i,j}^1$

$$R_{i,j}^2 = I_1^2 R_{i,j}^1 + L\Delta(f_2)_{i,j}^n \quad (\text{Save } \Delta(f_2)_{i,j}, I_1^2 R_{i,j}^1)$$

4. Restriction by injection to Grid 3  $I_2^3 R_{i,j}^2$

5. Iterate on Grid 3  $L\Delta(f_3)_{i,j} = -I_2^3 R_{i,j}^2$

$$R_{i,j}^3 = I_2^3 R_{i,j}^2 + L\Delta(f_3)_{i,j}^n \quad (\text{Save } \Delta(f_3)_{i,j}, I_2^3 R_{i,j}^2)$$

6. Restriction by injection to Grid 4  $I_3^4 R_{i,j}^3$

7. Iterate on Grid 4  $L\Delta(f_4)_{i,j} = -I_3^4 R_{i,j}^3$



**Coarse  $\Rightarrow$  Fine**

1. Prolongate from Grid 4 to Grid 3  $I_4^3 \Delta(f_4)_{i,j}^n$

$$\Delta(f_3)_{i,j}^n = \Delta(f_3)_{i,j}^n + I_4^3 \Delta(f_4)_{i,j}^n$$

2. Iterate on Grid 3  $L\Delta(f_3)_{i,j} = -I_2^3 R_{i,j}^2$  (saved)

3. Prolongate from Grid 3 to Grid 2  $I_3^2 \Delta(f_3)_{i,j}^n$

$$\Delta(f_2)_{i,j}^n = \Delta(f_2)_{i,j}^n + I_3^2 \Delta(f_3)_{i,j}^n$$

4. Iterate on Grid 2  $L\Delta(f_2)_{i,j} = -I_1^2 R_{i,j}^1$  (saved)

5. Prolongate from Grid 2 to Grid 1  $I_2^1 \Delta(f_2)_{i,j}^n$

$$f_{i,j}^n = f_{i,j}^n + I_2^1 \Delta(f_2)_{i,j}^n$$

6. Iterate on Grid 1  $Lf_{i,j} = 0$

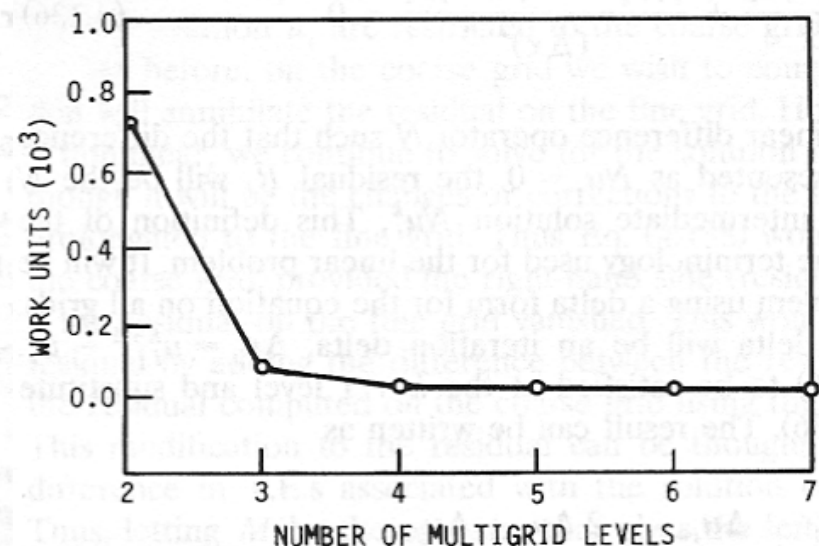
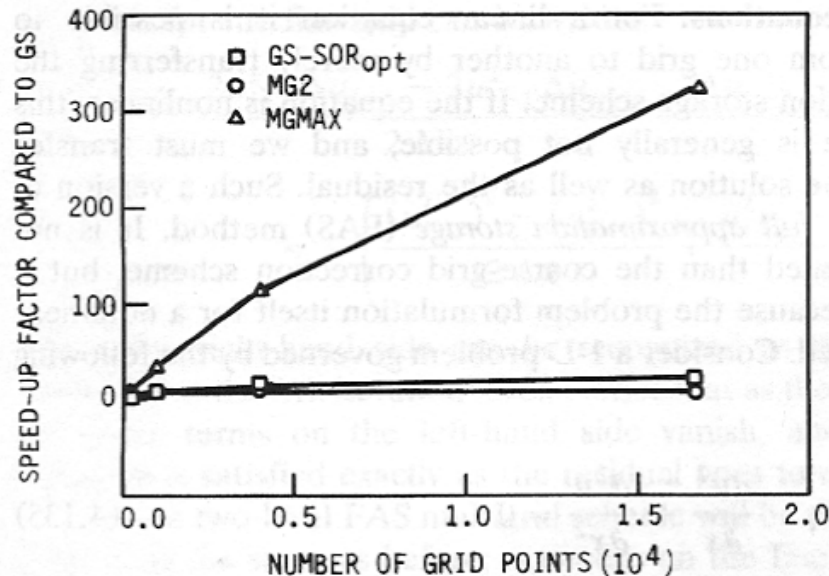
7. If  $R_{i,j} < \varepsilon$  then stop. Else, repeat the entire cycle.

Test Problem (Tannehill, p. 170)

- Laplace equation in a square domain
- Dirichlet conditions on four boundaries
- 5 Levels of resolution  
9×9, 17×17, 33×33, 65×65, 129×129
- 4 Methods
  1. Conventional Gauss-Seidel (GS)
  2. Gauss-Seidel-SOR with optimal  $\omega$  (GS<sub>opt</sub>)
  3. Multigrid with 2-level grids (MG2)
  4. Multigrid with maximum levels of grid (MGMAX)

## Multigrid Method - 16

Grid size	GS	$GS_{\omega_{opt}}$	MG2	MGMAX
9×9	62	19	19	17
17×17	215	40	40	19
33×33	715	75	95	20
65×65	2282	137	258	20
129×129	6826	282	732	21



The multigrid method is the fastest technique currently available for general elliptic equations, and are said to compare favorably with the fastest direct solvers for special problems.

MUDPACK

<http://www.scd.ucar.edu/css/software/mudpack/>

FISHPACK

<http://www.scd.ucar.edu/css/software/fishpack/>