

# Grimoire's Standard Code Library<sup>\*</sup>

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<sup>\*</sup><https://github.com/kzoacn/Grimoire>

# 目录

<b>1</b>	<b>代数</b>	<b>5</b>
1.1	$O(n^2 \log n)$ 求线性递推数列第 $n$ 项	5
1.2	任意模数快速傅里叶变换	6
1.3	快速傅里叶变换	6
1.4	闪电数论变换与魔力 CRT	7
1.5	多项式求逆	8
1.6	多项式除法	8
1.7	多项式取指数取对数	9
1.8	快速沃尔什变换	10
1.9	单纯形	11
<b>2</b>	<b>数论</b>	<b>13</b>
2.1	大整数相乘取模	13
2.2	EX-GCD	13
2.3	Miller-rabin	13
2.4	Pollard-rho.cpp	14
2.5	非互质 CRT	14
2.6	非互质 CRT -zky	14
2.7	Pell 方程	15
2.8	Simpson	15
2.9	解一元三次方程	16
2.10	线段下整点	16
2.11	线性同余不等式	16
2.12	EX-BSGS -zzq	16
2.13	EX-BSGS -zky	17
2.14	分治乘法	18
2.15	组合数模 $p^k$	18
<b>3</b>	<b>图论</b>	<b>19</b>
3.1	图论基础	19
3.2	坚固无敌的点双 -zzq	19
3.3	坚固无敌的边双 -zzq	20
3.4	坚固无敌的点双 -jzh	21
3.5	坚固无敌的边双 -jzh	22
3.6	2-sat	23
3.7	闪电二分图匹配	24
3.8	一般图匹配	25

目录	3
3.9 一般图最大权匹配	26
3.10 有根树 hash	30
3.11 无向图最小割	31
3.12 必经点 Dominator-tree	31
3.13 K 短路	33
3.14 最大团搜索	36
3.15 极大团计数	37
3.16 欧拉回路	38
3.17 朱刘最小树形图	38
<b>4 数据结构</b>	<b>41</b>
4.1 Kd-tree	41
4.2 LCT	44
4.3 树状数组上二分第 k 大	46
4.4 Treap	46
4.5 FHQ-Treap	47
4.6 真-FHQTreap	49
4.7 带修改莫队上树	51
4.8 虚树	52
<b>5 字符串</b>	<b>53</b>
5.1 Manacher	53
5.2 指针版回文自动机	53
5.3 后缀数组	54
5.4 最小表示法	55
<b>6 计算几何</b>	<b>57</b>
6.1 点类	57
6.2 圆基础	59
6.3 点在多边形内	60
6.4 二维最小覆盖圆	60
6.5 圆并	61
6.6 经典阿波罗尼斯圆	64
6.7 半平面交	64
6.8 求凸包	65
6.9 凸包游戏	65
6.10 平面最近点	67
6.11 无敌面积并 (多圆多多边形), $n^3$	67
6.12 Farmland	69
6.13 三维基础	71
6.14 三维凸包	73
6.15 三角剖分与 V 图	74
6.16 三维最小覆盖球	76
6.17 空间四点外接球	78

<b>7 技巧</b>	<b>79</b>
7.1 无敌的读入优化 . . . . .	79
7.2 真正释放 STL 内存 . . . . .	79
7.3 梅森旋转算法 . . . . .	79
7.4 蔡勒公式 . . . . .	80
7.5 开栈 . . . . .	80
7.6 Size 为 k 的子集 . . . . .	80
7.7 长方体表面两点最短距离 . . . . .	80
7.8 经纬度求球面最短距离 . . . . .	81
7.9 32-bit/64-bit 随机素数 . . . . .	81
7.10 NTT 素数及其原根 . . . . .	81
7.11 Formulas . . . . .	81
7.11.1 Arithmetic Function . . . . .	81
7.11.2 Binomial Coefficients . . . . .	82
7.11.3 Fibonacci Numbers . . . . .	83
7.11.4 Stirling Cycle Numbers . . . . .	83
7.11.5 Stirling Subset Numbers . . . . .	83
7.11.6 Eulerian Numbers . . . . .	83
7.11.7 Harmonic Numbers . . . . .	83
7.11.8 Pentagonal Number Theorem . . . . .	84
7.11.9 Bell Numbers . . . . .	84
7.11.10 Bernoulli Numbers . . . . .	84
7.11.11 Tetrahedron Volume . . . . .	84
7.11.12 BEST Theorem . . . . .	84
7.11.13 重心 . . . . .	84
7.11.14 Others . . . . .	84
7.12 Java . . . . .	88

# Chapter 1

## 代数

### $O(n^2 \log n)$ 求线性递推数列第 $n$ 项

Given  $a_0, a_1, \dots, a_{m-1}$

$a_n = c_0 * a_{n-m} + \dots + c_{m-1} * a_0$

$a_0$  is the  $n$ th element,  $\dots$ ,  $a_{m-1}$  is the  $n + m - 1$ th element

```
1 void linear_recurrence(long long n, int m, int a[], int c[], int p) {
2     long long v[M] = {1 % p}, u[M << 1], msk = !!n;
3     for(long long i(n); i > 1; i >= 1) {
4         msk <= 1;
5     }
6     for(long long x(0); msk; msk >= 1, x <= 1) {
7         fill_n(u, m << 1, 0);
8         int b(!(n & msk));
9         x |= b;
10        if(x < m) {
11            u[x] = 1 % p;
12        }else {
13            for(int i(0); i < m; i++) {
14                for(int j(0), t(i + b); j < m; j++, t++) {
15                    u[t] = (u[t] + v[i] * v[j]) % p;
16                }
17            }
18            for(int i((m << 1) - 1); i >= m; i--) {
19                for(int j(0), t(i - m); j < m; j++, t++) {
20                    u[t] = (u[t] + c[j] * u[i]) % p;
21                }
22            }
23        }
24        copy(u, u + m, v);
25    }
26    //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
27    for(int i(m); i < 2 * m; i++) {
28        a[i] = 0;
29        for(int j(0); j < m; j++) {
30            a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31        }
32    }
33    for(int j(0); j < m; j++) {
34        b[j] = 0;
35        for(int i(0); i < m; i++) {
36            b[j] = (b[j] + v[i] * a[i + j]) % p;
37        }
38    }
39    for(int j(0); j < m; j++) {
40        a[j] = b[j];
41    }
```

42 }

## 任意模数快速傅里叶变换

```

1 // double 精度对  $10^9 + 7$  取模最多可以做到  $2^{20}$ 
2 const int MOD = 1000003;
3 const double PI = acos(-1);
4 typedef complex<double> Complex;
5 const int N = 65536, L = 15, MASK = (1 << L) - 1;
6 Complex w[N];
7 void FFTInit() {
8     for (int i = 0; i < N; ++i)
9         w[i] = Complex(cos(2 * i * PI / N), sin(2 * i * PI / N));
10 }
11 void FFT(Complex p[], int n) {
12     for (int i = 1, j = 0; i < n - 1; ++i) {
13         for (int s = n; j ^= s >>= 1, ~j & s;);
14         if (i < j) swap(p[i], p[j]);
15     }
16     for (int d = 0; (1 << d) < n; ++d) {
17         int m = 1 << d, m2 = m * 2, rm = n >> (d + 1);
18         for (int i = 0; i < n; i += m2) {
19             for (int j = 0; j < m; ++j) {
20                 Complex &p1 = p[i + j + m], &p2 = p[i + j];
21                 Complex t = w[rm * j] * p1;
22                 p1 = p2 - t, p2 = p2 + t;
23             }
24         }
25     }
26     Complex A[N], B[N], C[N], D[N];
27     void mul(int a[N], int b[N]) {
28         for (int i = 0; i < N; ++i) {
29             A[i] = Complex(a[i] >> L, a[i] & MASK);
30             B[i] = Complex(b[i] >> L, b[i] & MASK);
31         }
32         FFT(A, N), FFT(B, N);
33         for (int i = 0; i < N; ++i) {
34             int j = (N - i) % N;
35             Complex da = (A[i] - conj(A[j])) * Complex(0, -0.5),
36                   db = (A[i] + conj(A[j])) * Complex(0.5, 0),
37                   dc = (B[i] - conj(B[j])) * Complex(0, -0.5),
38                   dd = (B[i] + conj(B[j])) * Complex(0.5, 0);
39             C[j] = da * dd + da * dc * Complex(0, 1);
40             D[j] = db * dd + db * dc * Complex(0, 1);
41         }
42         FFT(C, N), FFT(D, N);
43         for (int i = 0; i < N; ++i) {
44             long long da = (long long)(C[i].imag() / N + 0.5) % MOD,
45                   db = (long long)(C[i].real() / N + 0.5) % MOD,
46                   dc = (long long)(D[i].imag() / N + 0.5) % MOD,
47                   dd = (long long)(D[i].real() / N + 0.5) % MOD;
48             a[i] = ((dd << (L * 2)) + ((db + dc) << L) + da) % MOD;
49         }
50     }
51 }

```

## 快速傅里叶变换

```

1 int prepare(int n) {
2     int len = 1;

```

```

3   for (; len <= 2 * n; len <<= 1);
4   for (int i = 0; i < len; i++) {
5       e[0][i] = Complex(cos(2 * pi * i / len), sin(2 * pi * i / len));
6       e[1][i] = Complex(cos(2 * pi * i / len), -sin(2 * pi * i / len));
7   }
8   return len;
9 }
10 void DFT(Complex *a, int n, int f) {
11     for (int i = 0, j = 0; i < n; i++) {
12         if (i > j) std::swap(a[i], a[j]);
13         for (int t = n >> 1; (j ^= t) < t; t >>= 1);
14     }
15     for (int i = 2; i <= n; i <<= 1)
16         for (int j = 0; j < n; j += i)
17             for (int k = 0; k < (i >> 1); k++) {
18                 Complex A = a[j + k];
19                 Complex B = e[f][n / i * k] * a[j + k + (i >> 1)];
20                 a[j + k] = A + B;
21                 a[j + k + (i >> 1)] = A - B;
22             }
23     if (f == 1) {
24         for (int i = 0; i < n; i++)
25             a[i].a /= n;
26     }
27 }

```

## 闪电数论变换与魔力 CRT

```

1 #define meminit(A, l, r) memset(A + (l), 0, sizeof(*A) * ((r) - (l)))
2 #define memcpy(B, A, l, r) memcpy(B, A + (l), sizeof(*A) * ((r) - (l)))
3 void DFT(int *a, int n, int f) { //f=1 逆 DFT
4     for (register int i = 0, j = 0; i < n; i++) {
5         if (i > j) std::swap(a[i], a[j]);
6         for (register int t = n >> 1; (j ^= t) < t; t >>= 1);
7     }
8     for (register int i = 2; i <= n; i <<= 1) {
9         static int exp[MAXN];
10        exp[0] = 1; exp[1] = fpm(PRT, (MOD - 1) / i, MOD);
11        if (f == 1) exp[1] = fpm(exp[1], MOD - 2, MOD);
12        for (register int k = 2; k < (i >> 1); k++) {
13            exp[k] = 1ll * exp[k - 1] * exp[1] % MOD;
14        }
15        for (register int j = 0; j < n; j += i) {
16            for (register int k = 0; k < (i >> 1); k++) {
17                register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
18                register long long B = 1ll * pB * exp[k];
19                pB = (pA - B) % MOD;
20                pA = (pA + B) % MOD;
21            }
22        }
23    }
24    if (f == 1) {
25        register int rev = fpm(n, MOD - 2, MOD);
26        for (register int i = 0; i < n; i++) {
27            a[i] = 1ll * a[i] * rev % MOD;
28            if (a[i] < 0) { a[i] += MOD; }
29        }
30    }
31 }

```

```

32 // 在不写高精度的情况下合并 FFT 所得结果对 MOD 取模后的答案
33 // 值得注意的是, 这个东西不能最后再合并, 而是应该每做一次多项式乘法就 CRT 一次
34 int CRT(int *a) {
35     static int x[3];
36     for (int i = 0; i < 3; i++) {
37         x[i] = a[i];
38         for (int j = 0; j < i; j++) {
39             int t = (x[i] - x[j] + FFT[i] -> MOD) % FFT[i] -> MOD;
40             if (t < 0) t += FFT[i] -> MOD;
41             x[i] = 1LL * t * inv[j][i] % FFT[i] -> MOD;
42         }
43     }
44     int sum = 1, ret = x[0] % MOD;
45     for (int i = 1; i < 3; i++) {
46         sum = 1LL * sum * FFT[i] -> MOD % MOD;
47         ret += 1LL * x[i] * sum % MOD;
48         if (ret >= MOD) ret -= MOD;
49     }
50     return ret;
51 }
52 for (int i = 0; i < 3; i++) // inv 数组的预处理过程, inverse(x, p) 表示求 x 在 p 下逆元
53     for (int j = 0; j < 3; j++)
54         inv[i][j] = inverse(FFT[i] -> MOD, FFT[j] -> MOD);

```

## 多项式求逆

Given polynomial  $a$  and  $n$ ,  $b$  is the polynomial such that  $a * b \equiv 1 \pmod{x^n}$

```

1 void getInv(int *a, int *b, int n) {
2     static int tmp[MAXN];
3     b[0] = fpm(a[0], MOD - 2, MOD);
4     for (int c = 2, M = 1; c < (n << 1); c <= 1) {
5         for (; M <= 3 * (c - 1); M <= 1);
6         meminit(b, c, M);
7         meminit(tmp, c, M);
8         memcpy(tmp, a, 0, c);
9         DFT(tmp, M, 0);
10        DFT(b, M, 0);
11        for (int i = 0; i < M; i++) {
12            b[i] = 1ll * b[i] * (2ll - 1ll * tmp[i] * b[i] % MOD + MOD) % MOD;
13        }
14        DFT(b, M, 1);
15        meminit(b, c, M);
16    }
17 }

```

## 多项式除法

$d$  is quotient and  $r$  is remainder

```

1 void divide(int n, int m, int *a, int *b, int *d, int *r) { // n、m 分别为多项式 A (被除数)
    ↪ 和 B (除数) 的指数 + 1
2     static int M, tA[MAXN], tB[MAXN], inv[MAXN], tD[MAXN];
3     for (; n > 0 && a[n - 1] == 0; n--);
4     for (; m > 0 && b[m - 1] == 0; m--);
5     for (int i = 0; i < n; i++) tA[i] = a[n - i - 1];
6     for (int i = 0; i < m; i++) tB[i] = b[m - i - 1];
7     for (M = 1; M <= n - m + 1; M <= 1);
8     if (m < M) meminit(tB, m, M);

```



```

9      getInv(tB, inv, M);
10     for (M = 1; M <= 2 * (n - m + 1); M <= 1);
11     meminit(inv, n - m + 1, M);
12     meminit(tA, n - m + 1, M);
13     DFT(inv, M, 0);
14     DFT(tA, M, 0);
15     for (int i = 0; i < M; i++) {
16         d[i] = 1ll * inv[i] * tA[i] % MOD;
17     }
18     DFT(d, M, 1);
19     std::reverse(d, d + n - m + 1);
20     for (M = 1; M <= n; M <= 1);
21     memcpy(tB, b, 0, m);
22     if (m < M) meminit(tB, m, M);
23     memcpy(tD, d, 0, n - m + 1);
24     meminit(tD, n - m + 1, M);
25     DFT(tD, M, 0);
26     DFT(tB, M, 0);
27     for (int i = 0; i < M; i++) {
28         r[i] = 1ll * tD[i] * tB[i] % MOD;
29     }
30     DFT(r, M, 1);
31     meminit(r, n, M);
32     for (int i = 0; i < n; i++) {
33         r[i] = (a[i] - r[i] + MOD) % MOD;
34     }
35 }

```

## 多项式取指数取对数

Given polynomial  $a$  and  $n$ ,  $b$  is the polynomial such that  $b \equiv e^a \pmod{x^n}$  or  $b \equiv \ln a \pmod{x^n}$

```

1 void getDiff(int *a, int *b, int n) { // 多项式取微分
2     for (int i = 0; i + 1 < n; i++) {
3         b[i] = 1ll * (i + 1) * a[i + 1] % MOD;
4     }
5     b[n - 1] = 0;
6 }
7 void getInt(int *a, int *b, int n) { // 多项式取积分, 积分常数为 0
8     static int inv[MAXN];
9     inv[1] = 1;
10    for (int i = 2; i < n; i++) {
11        inv[i] = 1ll * (MOD - MOD / i) * inv[MOD % i] % MOD;
12    }
13    b[0] = 0;
14    for (int i = 1; i < n; i++) {
15        b[i] = 1ll * a[i - 1] * inv[i] % MOD;
16    }
17 }
18 void getLn(int *a, int *b, int n) {
19     static int inv[MAXN], d[MAXN];
20     int M = 1;
21     for (; M <= 2 * (n - 1); M <= 1);
22     getInv(a, inv, n);
23     getDiff(a, d, n);
24     meminit(d, n, M);
25     meminit(inv, n, M);
26     DFT(d, M, 0); DFT(inv, M, 0);
27     for (int i = 0; i < M; i++) {

```

```

28     d[i] = 111 * d[i] * inv[i] % MOD;
29 }
30 DFT(d, M, 1);
31 getInt(d, b, n);
32 }
33 void getExp(int *a, int *b, int n) {
34     static int ln[MAXN], tmp[MAXN];
35     b[0] = 1;
36     for (int c = 2, M = 1; c < (n << 1); c <= 1) {
37         for (; M <= 2 * (c - 1); M <= 1);
38         int bound = std::min(c, n);
39         memcpy(tmp, a, 0, bound);
40         meminit(tmp, bound, M);
41         meminit(b, c, M);
42         getLn(b, ln, c);
43         meminit(ln, c, M);
44         DFT(b, M, 0);
45         DFT(tmp, M, 0);
46         DFT(ln, M, 0);
47         for (int i = 0; i < M; i++) {
48             b[i] = 111 * b[i] * (111 - ln[i] + tmp[i] + MOD) % MOD;
49         }
50         DFT(b, M, 1);
51         meminit(b, c, M);
52     }
53 }

```

## 快速沃尔什变换

```

1 void FWT(LL a[], int n, int ty){
2     for(int d=1; d<n; d<=1){
3         for(int m=(d<<1), i=0; i<n; i+=m){
4             if(ty==1){
5                 for(int j=0; j<d; j++){
6                     LL x=a[i+j], y=a[i+j+d];
7                     a[i+j]=x+y;
8                     a[i+j+d]=x-y;
9                     //xor:a[i+j]=x+y, a[i+j+d]=x-y;
10                    //and:a[i+j]=x+y;
11                    //or:a[i+j+d]=x+y;
12                }
13            }else{
14                for(int j=0; j<d; j++){
15                    LL x=a[i+j], y=a[i+j+d];
16                    a[i+j]=(x+y)/2;
17                    a[i+j+d]=(x-y)/2;
18                    //xor:a[i+j]=(x+y)/2, a[i+j+d]=(x-y)/2;
19                    //and:a[i+j]=x-y;
20                    //or:a[i+j+d]=y-x;
21                }
22            }
23        }
24    }
25 }
26 FWT(a, 1<<n, 1);
27 FWT(b, 1<<n, 1);
28 for(int i=0; i<(1<<n); i++)
29     c[i]=a[i]*b[i];

```

```
30 FWT(c,1<<n,-1);
```

## 单纯形

```
1 namespace LP{
2     const int maxn=233;
3     double a[maxn][maxn];
4     int Ans[maxn],pt[maxn];
5     int n,m;
6     void pivot(int l,int i){
7         double t;
8         swap(Ans[l+n],Ans[i]);
9         t=-a[l][i];
10        a[l][i]=-1;
11        for(int j=0;j<=n;j++)a[l][j]/=t;
12        for(int j=0;j<=m;j++){
13            if(a[j][i]&&j!=l){
14                t=a[j][i];
15                a[j][i]=0;
16                for(int k=0;k<=n;k++)a[j][k]+=t*a[l][k];
17            }
18        }
19    }
20    vector<double> solve(vector<vector<double> >A,vector<double>B,vector<double>C){
21        n=C.size();
22        m=B.size();
23        for(int i=0;i<C.size();i++)
24            a[0][i+1]=C[i];
25        for(int i=0;i<B.size();i++)
26            a[i+1][0]=B[i];
27
28        for(int i=0;i<m;i++)
29            for(int j=0;j<n;j++)
30                a[i+1][j+1]=-A[i][j];
31
32        for(int i=1;i<=n;i++)Ans[i]=i;
33
34        double t;
35        for(;;){
36            int l=0;t=-eps;
37            for(int j=1;j<=m;j++)if(a[j][0]<t)t=a[l=j][0];
38            if(!l)break;
39            int i=0;
40            for(int j=1;j<=n;j++)if(a[l][j]>eps){i=j;break;}
41            if(!i){
42                puts("Infeasible");
43                return vector<double>();
44            }
45            pivot(l,i);
46        }
47        for(;;){
48            int i=0;t=eps;
49            for(int j=1;j<=n;j++)if(a[0][j]>t)t=a[0][i=j];
50            if(!i)break;
51            int l=0;
52            t=1e30;
53            for(int j=1;j<=m;j++)if(a[j][i]<-eps){
54                double tmp;
55                tmp=-a[j][0]/a[j][i];
```

```
56         if(t>tmp)t=tmp,l=j;
57     }
58     if(!l){
59         puts("Unbounded");
60         return vector<double>();
61     }
62     pivot(l,i);
63 }
64 vector<double>x;
65 for(int i=n+1;i<=n+m;i++)pt[Ans[i]]=i-n;
66 for(int i=1;i<=n;i++)x.push_back(pt[i]?a[pt[i]][0]:0);
67 return x;
68 }
69 }
```

## Chapter 2

# 数论

### 大整数相乘取模

```
1 // x 与 y 须非负
2 long long mult(long long x, long long y, long long MODN) {
3     long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
4     return t < 0 ? t + MODN : t;
5 }
```

### EX-GCD

```
1 LL exgcd(LL a, LL b, LL &x, LL &y){
2     if(!b){
3         x=1; y=0; return a;
4     }else{
5         LL d=exgcd(b, a%b, x, y);
6         LL t=x; x=y; y=t-a/b*y;
7         return d;
8     }
9 }
```

### Miller-rabin

```
1 const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
2 bool check(long long n, int base) {
3     long long n2=n-1, res;
4     int s=0;
5     while(n2%2==0) n2>>=1, s++;
6     res=pw(base, n2, n);
7     if((res==1) || (res==n-1)) return 1;
8     while(s--){
9         res=mul(res, res, n);
10        if(res==n-1) return 1;
11    }
12    return 0; // n is not a strong pseudo prime
13 }
14 bool isprime(const long long &n) {
15     if(n==2)
16         return true;
17     if(n<2 || n%2==0)
18         return false;
19     for(int i=0; i<12&&BASE[i]<n; i++){
20         if(!check(n, BASE[i]))
21             return false;
22     }
23 }
```

```

23     return true;
24 }

```

## Pollard-rho.cpp

```

1 LL prho(LL n, LL c){
2     LL i=1, k=2, x=rand()%(n-1)+1, y=x;
3     while(1){
4         i++; x=(x*x%n+c)%n;
5         LL d=__gcd((y-x+n)%n, n);
6         if(d>1&&d<n) return d;
7         if(y==x) return n;
8         if(i==k) y=x, k<=1;
9     }
10 }
11 void factor(LL n, vector<LL>&fat){
12     if(n==1) return;
13     if(isprime(n)){
14         fat.push_back(n);
15         return;
16     } LL p=n;
17     while(p>=n) p=prho(p, rand()%(n-1)+1);
18     factor(p, fat);
19     factor(n/p, fat);
20 }

```

## 非互质 CRT

first is remainder, second is module

```

1 inline void fix(LL &x, LL y) {
2     x = (x % y + y) % y;
3 }
4 bool solve(int n, std::pair<LL, LL> a[],
5             std::pair<LL, LL> &ans) {
6     ans = std::make_pair(1, 1);
7     for (int i = 0; i < n; ++i) {
8         LL num, y;
9         euclid(ans.second, a[i].second, num, y);
10        LL divisor = std::__gcd(ans.second, a[i].second);
11        if ((a[i].first - ans.first) % divisor) {
12            return false;
13        }
14        num *= (a[i].first - ans.first) / divisor;
15        fix(num, a[i].second);
16        ans.first += ans.second * num;
17        ans.second *= a[i].second / divisor;
18        fix(ans.first, ans.second);
19    }
20    return true;
21 }

```

## 非互质 CRT -zky

```

1 //merge Ax=B and ax=b to A'x=B'
2 LL china(int n, int *a, int *m){
3     LL M=1, d, x=0, y;

```

```

4     for(int i=0;i<n;i++){
5         M*=m[i];
6     for(int i=0;i<n;i++){
7         LL w=M/m[i];
8         d=exgcd(m[i],w,d,y);
9         y=(y%M+M)%M;
10        x=(x+y*w%M*a[i])%M;
11    }
12    while(x<0)x+=M;
13    return x;
14 }
15 void merge(LL &A,LL &B,LL a,LL b){
16     LL x,y;
17     sol(A,-a,b-B,x,y);
18     A=lcm(A,a);
19     B=(a*y+b)%A;
20     B=(B+A)%A;
21 }

```

## Pell 方程

```

1 //  $x_{k+1} = x_0 x_k + n y_0 y_k$ 
2 //  $y_{k+1} = x_0 y_k + y_0 x_k$ 
3 // n is not the index of which you want
4 pair<ll, ll> pell(ll n) {
5     static ll p[N], q[N], g[N], h[N], a[N];
6     p[1] = q[0] = h[1] = 1; p[0] = q[1] = g[1] = 0;
7     a[2] = (ll)(floor(sqrtl(n) + 1e-7L));
8     for(int i = 2; ; i++) {
9         g[i] = -g[i - 1] + a[i] * h[i - 1];
10        h[i] = (n - g[i] * g[i]) / h[i - 1];
11        a[i + 1] = (g[i] + a[2]) / h[i];
12        p[i] = a[i] * p[i - 1] + p[i - 2];
13        q[i] = a[i] * q[i - 1] + q[i - 2];
14        if(p[i] * p[i] - n * q[i] * q[i] == 1)
15            return {p[i], q[i]};
16    }
17 } //  $x^2 - n * y^2 = 1$  最小正整数根, n 为完全平方数时无解

```

## Simpson

```

1 // 三次函数, 两倍精度拟合
2 //  $error = \frac{(r-l)^5}{6480} |f^{(4)}|$ 
3 //  $\int_a^b f(x) dx \approx \frac{(b-a)}{8} [f(a) + 3f(\frac{2a+b}{3}) + 3f(\frac{a+2b}{3}) + f(b)]$ 
4 // 三次函数拟合  $error = \frac{1}{90} \frac{(r-l)^5}{2} |f^{(4)}|$ 
5 d simpson(d fl,d fr,d fmid,d l,d r) {
6     return (fl+fr+4.0*fmid)*(r-l)/6.0; }
7 d rsimpson(d slr,d fl,d fr,d fmid,d l,d r) {
8     d mid = (l+r)/2, fml = f((l+mid)/2), fmr = f((mid+r)/2);
9     d slm = simpson(fl,fmid,fml,l,mid);
10    d smr = simpson(fmid,fr,fmr,mid,r);
11    if(fabs(slr - smr - slm) / slr < eps) return slm + smr;
12    return rsimpson(slm,fl,fmid,fml,l,mid)+
13        rsimpson(smr,fmid,fr,fmr,mid,r);
14 }

```

## 解一元三次方程

听说极端情况精度不够

```

1 double a(p[3]), b(p[2]), c(p[1]), d(p[0]);
2 double k(b / a), m(c / a), n(d / a);
3 double p(-k * k / 3. + m);
4 double q(2. * k * k * k / 27 - k * m / 3. + n);
5 Complex omega[3] = {Complex(1, 0), Complex(-0.5, 0.5 * sqrt(3)), Complex(-0.5, -0.5 *
    ↪ sqrt(3))};
6 Complex r1, r2;
7 double delta(q * q / 4 + p * p * p / 27);
8 if (delta > 0) {
9     r1 = cubrt(-q / 2. + sqrt(delta));
10    r2 = cubrt(-q / 2. - sqrt(delta));
11 } else {
12     r1 = pow(-q / 2. + pow(Complex(delta), 0.5), 1. / 3);
13     r2 = pow(-q / 2. - pow(Complex(delta), 0.5), 1. / 3);
14 }
15 for(int _(0); _ < 3; _++) {
16     Complex x = -k / 3. + r1 * omega[_ * 1] + r2 * omega[_ * 2 % 3];
17 }

```

## 线段下整点

solve for  $\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor$ ,  $n, m, a, b > 0$

```

1 LL solve(LL n, LL a, LL b, LL m) {
2     if(b==0) return n*(a/m);
3     if(a>=m) return n*(a/m)+solve(n, a%m, b, m);
4     if(b>=m) return (n-1)*n/2*(b/m)+solve(n, a, b%m, m);
5     return solve((a+b*n)/m, (a+b*n)%m, m, b);
6 }

```

## 线性同余不等式

```

1 // Find the minimal non-negative solutions for  $l \leq d \cdot x \bmod m \leq r$ 
2 //  $0 \leq d, l, r < m; l \leq r, O(\log n)$ 
3 ll cal(ll m, ll d, ll l, ll r) {
4     if (l == 0) return 0;
5     if (d == 0) return MXL; // 无解
6     if (d * 2 > m) return cal(m, m - d, m - r, m - l);
7     if ((l - 1) / d < r / d) return (l - 1) / d + 1;
8     ll k = cal(d, (-m % d + d) % d, l % d, r % d);
9     return k == MXL ? MXL : (k * m + l - 1) / d + 1; // 无解 2
10 }

```

## EX-BSGS -zzq

```

1 /*
2  * EX_BSGS
3  *  $a^x = b \pmod p$ 
4  * p may not be a prime
5  */
6
7 ll qpow(ll a, ll x, ll Mod) {
8     ll res = 1;

```



```

9     for (; x; x >>= 1) {
10         if (x & 1) res = res * a % Mod;
11         a = a * a % Mod;
12     }
13     return res;
14 }
15
16 std::unordered_map<int, int> mp;
17
18 ll exbsgs(ll a, ll b, ll p) {
19     if (b == 1) return 0;
20     ll t, d = 1, k = 0;
21     while ((t = std::__gcd(a, p)) != 1) {
22         if (b % t) return -1;
23         ++k, b /= t, p /= t, d = d * (a / t) % p;
24         if (b == d) return k;
25     }
26     mp.clear();
27     ll m = std::ceil(std::sqrt(p));
28     ll a_m = qpow(a, m, p);
29     ll mul = b;
30     for (ll j = 1; j <= m; ++j) {
31         mul = mul * a % p;
32         mp[mul] = j;
33     }
34     for (ll i = 1; i <= m; ++i) {
35         d = d * a_m % p;
36         if (mp.count(d)) return i * m - mp[d] + k;
37     }
38     return -1;
39 }

```

## EX-BSGS -zky

```

1 LL BSGS(LL a, LL b, LL p) {
2     LL m = sqrt(p) + .5, v = inv(pw(a, m, p), p), e = 1;
3     map<LL, LL> hash; hash[1] = 0;
4     for (int i = 1; i <= m; ++i)
5         e = e * a % p, hash[e] = i;
6     for (int i = 0; i <= m; ++i) {
7         if (hash.count(b)) return i * m + hash[b];
8         b = b * v % p;
9     } return -1;
10 }
11
12 LL solve2(LL a, LL b, LL p) {
13     // a^x = b (mod p)
14     b %= p;
15     LL e = 1 % p;
16     for (int i = 0; i < 100; ++i) {
17         if (e == b) return i;
18         e = e * a % p;
19     }
20     int r = 0;
21     while (gcd(a, p) != 1) {
22         LL d = gcd(a, p);
23         if (b % d) return -1;
24         p /= d; b /= d; b = b * inv(a / d, p);
25         r++;

```

```

26     }LL res=BSGS(a,b,p);
27     if(res==-1)return -1;
28     return res+r;
29 }

```

## 分治乘法

```

1 (a+b)(c+d) = ac+(bc+ad)+bd = 2ac-(a-b)(c-d)+2bd
2
3 x = x^m m=(n+1)/2
4 (ax+b)(cx+d) = x^2ac + x(bc+ad) + bd = x^2ac + x(ac + bd - (a-b)(c-d)) + bd

```

## 组合数模 $p^k$

```

1 LL prod=1,P;
2 pair<LL,LL> comput(LL n,LL p,LL k){
3     if(n<=1)return make_pair(0,1);
4     LL ans=1,cnt=0;
5     ans=pow(prod,n/P,P);
6     cnt=n/p;
7     pair<LL,LL>res=comput(n/p,p,k);
8     cnt+=res.first;
9     ans=ans*res.second%P;
10    for(int i=n-n%P+1;i<=n;i++)if(i%p){
11
12        ans=ans*i%P;
13    }
14    return make_pair(cnt,ans);
15 }
16 pair<LL,LL> calc(LL n,LL p,LL k){
17     prod=1;P=pow(p,k,1e18);
18     for(int i=1;i<P;i++)if(i%p)prod=prod*i%P;
19     pair<LL,LL> res=comput(n,p,k);
20     // res.second=res.second*pow(p,res.first%k,P)%P;
21     // res.first-=res.first%k;
22     return res;
23 }
24 LL calc(LL n,LL m,LL p,LL k){
25     pair<LL,LL>A,B,C;
26     LL P=pow(p,k,1e18);
27     A=calc(n,p,k);
28     B=calc(m,p,k);
29     C=calc(n-m,p,k);
30     LL ans=1;
31     ans=pow(p,A.first-B.first-C.first,P);
32     ans=ans*A.second%P*inv(B.second,P)%P*inv(C.second,P)%P;
33     return ans;
34 }

```

# Chapter 3

## 图论

### 图论基础

```
1 struct Graph { // Remember to call .init()!
2     int e, nxt[M], v[M], adj[N], n;
3     bool base;
4     __inline void init(bool _base, int _n = 0) {
5         n = _n; base = _base;
6         e = 0; memset(adj + base, -1, sizeof(*adj) * n);
7     }
8     __inline int new_node() {
9         adj[n + base] = -1;
10        return n++ + base;
11    }
12    __inline void ins(int u0, int v0) { // directional
13        v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
14    }
15    __inline void bi_ins(int u0, int v0) { // bi-directional
16        ins(u0, v0); ins(v0, u0);
17    }
18 };
```

### 坚固无敌的点双 -zzq

```
1 typedef std::pair<int, int> pii;
2 #define mkpair std::make_pair
3 int n, m;
4 std::vector<int> G[MAXN];
5 int dfn[MAXN], low[MAXN], bcc_id[MAXN], bcc_cnt, stamp;
6 bool iscut[MAXN];
7 std::vector<int> bcc[MAXN]; // Unnecessary
8 pii stk[MAXN]; int stk_top;
9 // Use a handwritten structure to get higher efficiency
10 void Tarjan(int now, int fa) {
11     int child = 0;
12     dfn[now] = low[now] = ++stamp;
13     for (int to: G[now]) {
14         if (!dfn[to]) {
15             stk[++stk_top] = mkpair(now, to); ++child;
16             Tarjan(to, now);
17             low[now] = std::min(low[now], low[to]);
18             if (low[to] >= dfn[now]) {
19                 iscut[now] = 1;
20                 bcc[++bcc_cnt].clear();
21                 while (1) {
22                     pii tmp = stk[stk_top--];
```

```

23         if (bcc_id[tmp.first] != bcc_cnt) {
24             bcc[bcc_cnt].push_back(tmp.first);
25             bcc_id[tmp.first] = bcc_cnt;
26         }
27         if (bcc_id[tmp.second] != bcc_cnt) {
28             bcc[bcc_cnt].push_back(tmp.second);
29             bcc_id[tmp.second] = bcc_cnt;
30         }
31         if (tmp.first == now && tmp.second == to)
32             break;
33     }
34 }
35 }
36 else if (dfn[to] < dfn[now] && to != fa) {
37     stk[++stk_top] = mkpair(now, to);
38     low[now] = std::min(low[now], dfn[to]);
39 }
40 }
41 if (!fa && child == 1) iscut[now] = 0;
42 }
43
44 void PBCC() {
45     memset(dfn, 0, sizeof dfn);
46     memset(low, 0, sizeof low);
47     memset(iscut, 0, sizeof iscut);
48     memset(bcc_id, 0, sizeof bcc_id);
49     stamp = bcc_cnt = stk_top = 0;
50     for (int i = 1; i <= n; ++i)
51         if (!dfn[i]) Tarjan(i, 0);
52 }

```

## 坚固无敌的边双 -zzq

```

1  int n, m;
2  int head[MAXN], nxt[MAXM << 1], to[MAXM << 1], ed;
3  // Opposite edge exists, set head[] to -1.
4  int dfn[MAXN], low[MAXN], bcc_id[MAXN], bcc_cnt, stamp;
5  bool isbridge[MAXM << 1], vis[MAXN];
6  std::vector<int> bcc[MAXN];
7  void Tarjan(int now, int fa) {
8     dfn[now] = low[now] = ++stamp;
9     for (int i = head[now]; ~i; i = nxt[i]) {
10         if (!dfn[to[i]]) {
11             Tarjan(to[i], now);
12             low[now] = std::min(low[now], low[to[i]]);
13             if (low[to[i]] > dfn[now])
14                 isbridge[i] = isbridge[i ^ 1] = 1;
15         }
16         else if (dfn[to[i]] < dfn[now] && to[i] != fa)
17             low[now] = std::min(low[now], dfn[to[i]]);
18     }
19 }
20 void DFS(int now) {
21     vis[now] = 1;
22     bcc[bcc_id[now] = bcc_cnt].push_back(now);
23     for (int i = head[now]; ~i; i = nxt[i]) {
24         if (isbridge[i]) continue;
25         if (!vis[to[i]]) DFS(to[i]);
26     }

```

```

27 }
28 void EBCC() {
29     memset(dfn, 0, sizeof dfn);
30     memset(low, 0, sizeof low);
31     memset(isbridge, 0, sizeof isbridge);
32     memset(bcc_id, 0, sizeof bcc_id);
33     bcc_cnt = stamp = 0;
34     for (int i = 1; i <= n; ++i)
35         if (!dfn[i]) Tarjan(i, 0);
36     memset(vis, 0, sizeof vis);
37     for (int i = 1; i <= n; ++i)
38         if (!vis[i]) {
39             ++bcc_cnt; DFS(i);
40         }
41 }

```

## 坚固无敌的点双 -jzh

```

1  const bool BCC_VERTEX = 0, BCC_EDGE = 1;
2  struct BCC { // N = NO + MO. Remember to call init(&raw_graph).
3      Graph *g, forest; // g is raw graph ptr.
4      int dfn[N], DFN, low[N];
5      int stack[N], top;
6      int expand_to[M]; // Where edge i is expanded to in expanded graph.
7      // Vertex i expanded to i.
8      int compress_to[N]; // Where vertex i is compressed to.
9      bool cut[N], compress_cut[N], branch[M], vis[N], flag;
10     //std::vector<int> BCC_component[N]; // Cut vertex belongs to none.
11     __inline void init(Graph *raw_graph) {
12         g = raw_graph;
13     }
14     void DFS(int u, int pe) {
15         dfn[u] = low[u] = ++DFN; cut[u] = false;
16         if (!g->adj[u]) {
17             cut[u] = 1;
18             compress_to[u] = forest.new_node();
19             compress_cut[compress_to[u]] = 1;
20         }
21         for (int e = g->adj[u]; ~e; e = g->nxt[e]) {
22             int v = g->v[e];
23             if ((e ^ pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
24                 stack[top++] = e;
25                 low[u] = std::min(low[u], dfn[v]);
26             }
27             else if (!dfn[v]) {
28                 stack[top++] = e; branch[e] = 1;
29                 DFS(v, e);
30                 low[u] = std::min(low[v], low[u]);
31                 if (low[v] >= dfn[u]) {
32                     if ((pe == -1 && flag || pe != -1) && !cut[u]) {
33                         cut[u] = 1;
34                         compress_to[u] = forest.new_node();
35                         compress_cut[compress_to[u]] = 1;
36                     }
37                     int cc = forest.new_node();
38                     if (cut[u]) forest.bi_ins(compress_to[u], cc);
39                     compress_cut[cc] = 0;
40                     //BCC_component[cc].clear();
41                     do {

```

```

42         int cur_e = stack[--top];
43         compress_to[expand_to[cur_e]] = cc;
44         compress_to[expand_to[cur_e^1]] = cc;
45         if (branch[cur_e]) {
46             int v = g->v[cur_e];
47             if (cut[v]) {
48                 forest.bi_ins(cc, compress_to[v]);
49             } else {
50                 //BCC_component[cc].push_back(v);
51                 compress_to[v] = cc;
52             }
53         }
54     } while (stack[top] != e);
55     if (pe == -1 && !flag) {
56         compress_to[u] = cc;
57     }
58 }
59 }
60 }
61 }
62 inline bool dfs(int u, int pe) {
63     vis[u] = 1;
64     int d = 0;
65     for (int e = g->adj[u]; ~e; e = g->nxt[e]) {
66         int v = g->v[e];
67         if (!vis[v]) {
68             ++d; dfs(v, e);
69         }
70     }
71     return pe == -1 ? d > 1 : 0;
72 }
73 void solve() {
74     forest.init(g->base);
75     int n = g->n;
76     for (int i = 0; i < g->e; i++) {
77         expand_to[i] = g->new_node();
78     }
79     memset(vis + g->base, 0, sizeof(*vis) * n);
80     memset(branch, 0, sizeof(*branch) * g->e);
81     memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
82     for (int i = 0; i < n; i++)
83         if (!dfn[i + g->base]) {
84             top = 0;
85             flag = dfs(i + g->base, -1);
86             DFS(i + g->base, -1);
87         }
88 }
89 } bcc;

```

## 坚固无敌的边双 -jzh

```

1 struct BCC {
2     Graph *g, forest;
3     int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs_clock;
4     // tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
5     pair<int, int> ori[M]; // bridge in raw_graph(raw node)
6     bool is_bridge[M];
7     __inline void init(Graph *raw_graph) {
8         g = raw_graph;

```

```

9      memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
10     memset(vis + g -> base, 0, sizeof(*vis) * g -> n);
11 }
12 void tarjan(int u, int from) {
13     dfn[u] = low[u] = ++dfs_clock; vis[u] = 1; stack[++top] = u;
14     for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
15         if ((p ^ 1) == from) continue;
16         int v = g -> v[p];
17         if (vis[v]) {
18             if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19         } else {
20             tarjan(v, p);
21             low[u] = min(low[u], low[v]);
22             if (low[v] > dfn[u]) is_bridge[p / 2] = true;
23         }
24     }
25     if (dfn[u] != low[u]) return;
26     tot[forest.new_node()] = 0;
27     do {
28         belong[stack[top]] = forest.n;
29         vis[stack[top]] = 2;
30         tot[forest.n]++;
31         --top;
32     } while (stack[top + 1] != u);
33 }
34 void solve() {
35     forest.init(g -> base);
36     int n = g -> n;
37     for (int i = 0; i < n; ++i)
38         if (!vis[i + g -> base]) {
39             top = dfs_clock = 0;
40             tarjan(i + g -> base, -1);
41         }
42     for (int i = 0; i < g -> e / 2; ++i)
43         if (is_bridge[i]) {
44             int e = forest.e;
45             forest.bi_ins(belong[g -> v[i * 2]], belong[g -> v[i * 2 + 1]], g -> w[i *
46 ↪ 2]);
47             ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
48             ori[e + 1] = make_pair(g -> v[i * 2], g -> v[i * 2 + 1]);
49         }
50 } bcc;

```

## 2-sat

```

1 //清点清边要两倍
2 int stamp, comps, top;
3 int dfn[N], low[N], comp[N], stack[N];
4 void add(int x, int a, int y, int b) {
5     edge[x << 1 | a].push_back(y << 1 | b);
6 }
7 void tarjan(int x) {
8     dfn[x] = low[x] = ++stamp;
9     stack[top++] = x;
10    for (int i = 0; i < (int)edge[x].size(); ++i) {
11        int y = edge[x][i];
12        if (!dfn[y]) {
13            tarjan(y);

```

```

14         low[x] = std::min(low[x], low[y]);
15     } else if (!comp[y]) {
16         low[x] = std::min(low[x], dfn[y]);
17     }
18 }
19 if (low[x] == dfn[x]) {
20     comps++;
21     do {
22         int y = stack[--top];
23         comp[y] = comps;
24     } while (stack[top] != x);
25 }
26 }
27 bool solve() {
28     int counter = n + n + 1;
29     stamp = top = comps = 0;
30     std::fill(dfn, dfn + counter, 0);
31     std::fill(comp, comp + counter, 0);
32     for (int i = 0; i < counter; ++i) {
33         if (!dfn[i]) {
34             tarjan(i);
35         }
36     }
37     for (int i = 0; i < n; ++i) {
38         if (comp[i << 1] == comp[i << 1 | 1]) {
39             return false;
40         }
41         answer[i] = (comp[i << 1 | 1] < comp[i << 1]);
42     }
43     return true;
44 }

```

## 闪电二分图匹配

```

1 int matchx[N], matchy[N], level[N];
2 vector<int> edge[N];
3 bool dfs(int x) {
4     for (int i = 0; i < (int)edge[x].size(); ++i) {
5         int y = edge[x][i];
6         int w = matchy[y];
7         if (w == -1 || level[x] + 1 == level[w] && dfs(w)) {
8             matchx[x] = y; matchy[y] = x;
9             return true;
10        }
11    }
12    level[x] = -1;
13    return false;
14 }
15 int solve() {
16     memset(matchx, -1, sizeof(*matchx) * n);
17     memset(matchy, -1, sizeof(*matchy) * m);
18     for (int ans = 0; ; ) {
19         std::vector<int> q;
20         for (int i = 0; i < n; ++i) {
21             if (matchx[i] == -1) {
22                 level[i] = 0;
23                 q.push_back(i);
24             } else level[i] = -1;
25         }

```



```

26     for (int head = 0; head < (int)q.size(); ++head) {
27         int x = q[head];
28         for (int i = 0; i < (int)edge[x].size(); ++i) {
29             int y = edge[x][i];
30             int w = matchy[y];
31             if (w != -1 && level[w] < 0) {
32                 level[w] = level[x] + 1;
33                 q.push_back(w);
34             }
35         }
36     }
37     int delta = 0;
38     for (int i = 0; i < n; ++i)
39         if (matchx[i] == -1 && dfs(i)) ++delta;
40     if (delta == 0) return ans; else ans += delta;
41 }
42 }

```

## 一般图匹配

```

1 // 0-base, match[u] is linked to u
2 vector<int> lnk[MAXN];
3 int match[MAXN], Queue[MAXN], pred[MAXN], base[MAXN], head, tail, sta, fin, nbase;
4 bool inQ[MAXN], inB[MAXN];
5 inline void push(int u) {
6     Queue[tail++] = u; inQ[u] = 1;
7 }
8 inline int pop() {
9     return Queue[head++];
10 }
11 inline int FindCA(int u, int v) {
12     static bool inP[MAXN];
13     fill(inP, inP + n, false);
14     while (1) {
15         u = base[u]; inP[u] = 1;
16         if (u == sta) break;
17         u = pred[match[u]];
18     }
19     while (1) {
20         v = base[v];
21         if (inP[v]) break;
22         v = pred[match[v]];
23     }
24     return v;
25 }
26 inline void RT(int u) {
27     int v;
28     while (base[u] != nbase) {
29         v = match[u];
30         inB[base[u]] = inB[base[v]] = 1;
31         u = pred[v];
32         if (base[u] != nbase) pred[u] = v;
33     }
34 }
35 inline void BC(int u, int v) {
36     nbase = FindCA(u, v);
37     fill(inB, inB + n, 0);
38     RT(u); RT(v);
39     if (base[u] != nbase) pred[u] = v;

```

```

40     if (base[v] != nbase) pred[v] = u;
41     for (int i = 0; i < n; ++i)
42         if (inB[base[i]]) {
43             base[i] = nbase;
44             if (!inQ[i]) push(i);
45         }
46 }
47 bool FindAP(int u) {
48     bool found = false;
49     for (int i = 0; i < n; ++i) {
50         pred[i] = -1; base[i] = i; inQ[i] = 0;
51     }
52     sta = u; fin = -1; head = tail = 0; push(sta);
53     while (head < tail) {
54         int u = pop();
55         for (int i = (int)lnk[u].size() - 1; i >= 0; --i) {
56             int v = lnk[u][i];
57             if (base[u] != base[v] && match[u] != v) {
58                 if (v == sta || match[v] >= 0 && pred[match[v]] >= 0) BC(u, v);
59                 else if (pred[v] == -1) {
60                     pred[v] = u;
61                     if (match[v] >= 0) push(match[v]);
62                     else {
63                         fin = v;
64                         return true;
65                     }
66                 }
67             }
68         }
69     }
70     return found;
71 }
72 inline void AP() {
73     int u = fin, v, w;
74     while (u >= 0) {
75         v = pred[u]; w = match[v];
76         match[v] = u; match[u] = v;
77         u = w;
78     }
79 }
80 inline int FindMax() {
81     for (int i = 0; i < n; ++i) match[i] = -1;
82     for (int i = 0; i < n; ++i)
83         if (match[i] == -1 && FindAP(i)) AP();
84     int ans = 0;
85     for (int i = 0; i < n; ++i) ans += (match[i] != -1);
86     return ans;
87 }

```

## 一般图最大权匹配

```

1 //maximum weight blossom, change g[u][v].w to INF - g[u][v].w when minimum weight blossom
  ↳ is needed
2 //type of ans is long long
3 //replace all int to long long if weight of edge is long long
4
5 struct WeightGraph {
6     static const int INF = INT_MAX;
7     static const int MAXN = 400;

```

```

8      struct edge{
9          int u, v, w;
10         edge() {}
11         edge(int u, int v, int w): u(u), v(v), w(w) {}
12     };
13     int n, n_x;
14     edge g[MAXN * 2 + 1][MAXN * 2 + 1];
15     int lab[MAXN * 2 + 1];
16     int match[MAXN * 2 + 1], slack[MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 + 1];
17     int flower_from[MAXN * 2 + 1][MAXN+1], S[MAXN * 2 + 1], vis[MAXN * 2 + 1];
18     vector<int> flower[MAXN * 2 + 1];
19     queue<int> q;
20     inline int e_delta(const edge &e){ // does not work inside blossoms
21         return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
22     }
23     inline void update_slack(int u, int x){
24         if(!slack[x] || e_delta(g[u][x]) < e_delta(g[slack[x]][x]))
25             slack[x] = u;
26     }
27     inline void set_slack(int x){
28         slack[x] = 0;
29         for(int u = 1; u <= n; ++u)
30             if(g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
31                 update_slack(u, x);
32     }
33     void q_push(int x){
34         if(x <= n)q.push(x);
35         else for(size_t i = 0; i < flower[x].size(); i++)
36             q_push(flower[x][i]);
37     }
38     inline void set_st(int x, int b){
39         st[x]=b;
40         if(x > n) for(size_t i = 0; i < flower[x].size(); ++i)
41             set_st(flower[x][i], b);
42     }
43     inline int get_pr(int b, int xr){
44         int pr = find(flower[b].begin(), flower[b].end(), xr) - flower[b].begin();
45         if(pr % 2 == 1){
46             reverse(flower[b].begin() + 1, flower[b].end());
47             return (int)flower[b].size() - pr;
48         } else return pr;
49     }
50     inline void set_match(int u, int v){
51         match[u]=g[u][v].v;
52         if(u > n){
53             edge e=g[u][v];
54             int xr = flower_from[u][e.u], pr=get_pr(u, xr);
55             for(int i = 0; i < pr; ++i)
56                 set_match(flower[u][i], flower[u][i ^ 1]);
57             set_match(xr, v);
58             rotate(flower[u].begin(), flower[u].begin()+pr, flower[u].end());
59         }
60     }
61     inline void augment(int u, int v){
62         for(;;){
63             int xnv=st[match[u]];
64             set_match(u, v);
65             if(!xnv)return;
66             set_match(xnv, st[pa[xnv]]);
67             u=st[pa[xnv]], v=xnv;

```

```

68     }
69 }
70 inline int get_lca(int u, int v){
71     static int t=0;
72     for(++t; u || v; swap(u, v)){
73         if(u == 0)continue;
74         if(vis[u] == t)return u;
75         vis[u] = t;
76         u = st[match[u]];
77         if(u) u = st[pa[u]];
78     }
79     return 0;
80 }
81 inline void add_blossom(int u, int lca, int v){
82     int b = n + 1;
83     while(b <= n_x && st[b]) ++b;
84     if(b > n_x) ++n_x;
85     lab[b] = 0, S[b] = 0;
86     match[b] = match[lca];
87     flower[b].clear();
88     flower[b].push_back(lca);
89     for(int x = u, y; x != lca; x = st[pa[y]]) {
90         flower[b].push_back(x),
91         flower[b].push_back(y = st[match[x]]),
92         q_push(y);
93     }
94     reverse(flower[b].begin() + 1, flower[b].end());
95     for(int x = v, y; x != lca; x = st[pa[y]]) {
96         flower[b].push_back(x),
97         flower[b].push_back(y = st[match[x]]),
98         q_push(y);
99     }
100     set_st(b, b);
101     for(int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w = 0;
102     for(int x = 1; x <= n; ++x) flower_from[b][x] = 0;
103     for(size_t i = 0; i < flower[b].size(); ++i){
104         int xs = flower[b][i];
105         for(int x = 1; x <= n_x; ++x)
106             if(g[b][x].w == 0 || e_delta(g[xs][x]) < e_delta(g[b][x]))
107                 g[b][x] = g[xs][x], g[x][b] = g[x][xs];
108         for(int x = 1; x <= n; ++x)
109             if(flower_from[xs][x]) flower_from[b][x] = xs;
110     }
111     set_slack(b);
112 }
113 inline void expand_blossom(int b){ // S[b] == 1
114     for(size_t i = 0; i < flower[b].size(); ++i)
115         set_st(flower[b][i], flower[b][i]);
116     int xr = flower_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
117     for(int i = 0; i < pr; i += 2){
118         int xs = flower[b][i], xns = flower[b][i + 1];
119         pa[xs] = g[xns][xs].u;
120         S[xs] = 1, S[xns] = 0;
121         slack[xs] = 0, set_slack(xns);
122         q_push(xns);
123     }
124     S[xr] = 1, pa[xr] = pa[b];
125     for(size_t i = pr + 1; i < flower[b].size(); ++i){
126         int xs = flower[b][i];
127         S[xs] = -1, set_slack(xs);

```

```

128     }
129     st[b] = 0;
130 }
131 inline bool on_found_edge(const edge &e){
132     int u = st[e.u], v = st[e.v];
133     if(S[v] == -1){
134         pa[v] = e.u, S[v] = 1;
135         int nu = st[match[v]];
136         slack[v] = slack[nu] = 0;
137         S[nu] = 0, q_push(nu);
138     }else if(S[v] == 0){
139         int lca = get_lca(u, v);
140         if(!lca) return augment(u, v), augment(v, u), true;
141         else add_blossom(u, lca, v);
142     }
143     return false;
144 }
145 inline bool matching(){
146     memset(S + 1, -1, sizeof(int) * n_x);
147     memset(slack + 1, 0, sizeof(int) * n_x);
148     q = queue<int>();
149     for(int x = 1; x <= n_x; ++x)
150         if(st[x] == x && !match[x]) pa[x]=0, S[x]=0, q_push(x);
151     if(q.empty())return false;
152     for(;;){
153         while(q.size()){
154             int u = q.front();q.pop();
155             if(S[st[u]] == 1)continue;
156             for(int v = 1; v <= n; ++v)
157                 if(g[u][v].w > 0 && st[u] != st[v]){
158                     if(e_delta(g[u][v]) == 0){
159                         if(on_found_edge(g[u][v]))return true;
160                     }else update_slack(u, st[v]);
161                 }
162         }
163         int d = INF;
164         for(int b = n + 1; b <= n_x; ++b)
165             if(st[b] == b && S[b] == 1)d = min(d, lab[b]/2);
166         for(int x = 1; x <= n_x; ++x)
167             if(st[x] == x && slack[x]){
168                 if(S[x] == -1)d = min(d, e_delta(g[slack[x]][x]));
169                 else if(S[x] == 0)d = min(d, e_delta(g[slack[x]][x])/2);
170             }
171         for(int u = 1; u <= n; ++u){
172             if(S[st[u]] == 0){
173                 if(lab[u] <= d)return 0;
174                 lab[u] -= d;
175             }else if(S[st[u]] == 1)lab[u] += d;
176         }
177         for(int b = n+1; b <= n_x; ++b)
178             if(st[b] == b){
179                 if(S[st[b]] == 0) lab[b] += d * 2;
180                 else if(S[st[b]] == 1) lab[b] -= d * 2;
181             }
182         q=queue<int>();
183         for(int x = 1; x <= n_x; ++x)
184             if(st[x] == x && slack[x] && st[slack[x]] != x && e_delta(g[slack[x]][x]) ==
↪ 0)
185                 if(on_found_edge(g[slack[x]][x]))return true;
186         for(int b = n + 1; b <= n_x; ++b)

```

```

187         if(st[b] == b && S[b] == 1 && lab[b] == 0) expand_blossom(b);
188     }
189     return false;
190 }
191 inline pair<long long, int> solve(){
192     memset(match + 1, 0, sizeof(int) * n);
193     n_x = n;
194     int n_matches = 0;
195     long long tot_weight = 0;
196     for(int u = 0; u <= n; ++u) st[u] = u, flower[u].clear();
197     int w_max = 0;
198     for(int u = 1; u <= n; ++u)
199         for(int v = 1; v <= n; ++v){
200             flower_from[u][v] = (u == v ? u : 0);
201             w_max = max(w_max, g[u][v].w);
202         }
203     for(int u = 1; u <= n; ++u) lab[u] = w_max;
204     while(matching()) ++n_matches;
205     for(int u = 1; u <= n; ++u)
206         if(match[u] && match[u] < u)
207             tot_weight += g[u][match[u]].w;
208     return make_pair(tot_weight, n_matches);
209 }
210 inline void init(){
211     for(int u = 1; u <= n; ++u)
212         for(int v = 1; v <= n; ++v)
213             g[u][v] = edge(u, v, 0);
214 }
215 };

```

## 有根树 hash

```

1  const unsigned long long MAGIC = 4423;
2
3  unsigned long long magic[N];
4  std::pair<unsigned long long, int> hash[N];
5
6  void solve(int root) {
7      magic[0] = 1;
8      for (int i = 1; i <= n; ++i) {
9          magic[i] = magic[i - 1] * MAGIC;
10     }
11     std::vector<int> queue;
12     queue.push_back(root);
13     for (int head = 0; head < (int)queue.size(); ++head) {
14         int x = queue[head];
15         for (int i = 0; i < (int)son[x].size(); ++i) {
16             int y = son[x][i];
17             queue.push_back(y);
18         }
19     }
20     for (int index = n - 1; index >= 0; --index) {
21         int x = queue[index];
22         hash[x] = std::make_pair(0, 0);
23
24         std::vector<std::pair<unsigned long long, int> > value;
25         for (int i = 0; i < (int)son[x].size(); ++i) {
26             int y = son[x][i];
27             value.push_back(hash[y]);

```

```

28     }
29     std::sort(value.begin(), value.end());
30
31     hash[x].first = hash[x].first * magic[1] + 37;
32     hash[x].second++;
33     for (int i = 0; i < (int)value.size(); ++i) {
34         hash[x].first = hash[x].first * magic[value[i].second] + value[i].first;
35         hash[x].second += value[i].second;
36     }
37     hash[x].first = hash[x].first * magic[1] + 41;
38     hash[x].second++;
39 }
40 }

```

## 无向图最小割

```

1  /*
2  * Stoer Wagner 全局最小割  $O(V^3)$ 
3  * lbase, 点数 n, 邻接矩阵 edge[MAXN][MAXN]
4  * 返回值为全局最小割
5  */
6
7  int StoerWagner() {
8      static int v[MAXN], wage[MAXN];
9      static bool vis[MAXN];
10     for (int i = 1; i <= n; ++i) v[i] = i;
11     int res = INF;
12     for (int nn = n; nn > 1; --nn) {
13         memset(vis, 0, sizeof(bool) * (nn + 1));
14         memset(wage, 0, sizeof(int) * (nn + 1));
15         int pre, last = 1; // vis[1] = 1;
16         for (int i = 1; i < nn; ++i) {
17             pre = last; last = 0;
18             for (int j = 2; j <= nn; ++j) if (!vis[j]) {
19                 wage[j] += edge[v[pre]][v[j]];
20                 if (!last || wage[j] > wage[last]) last = j;
21             }
22             vis[last] = 1;
23         }
24         res = std::min(res, wage[last]);
25         for (int i = 1; i <= nn; ++i) {
26             edge[v[i]][v[pre]] += edge[v[last]][v[i]];
27             edge[v[pre]][v[i]] += edge[v[last]][v[i]];
28         }
29         v[last] = v[nn];
30     }
31     return res;
32 }

```

## 必经点 Dominator-tree

```

1  //solve(s, n, raw_g): s is the root and base accords to base of raw_g
2  //idom[x] will be x if x does not have a dominator, and will be -1 if x is not reachable from
   ↪ s.
3  struct dominator_tree {
4      int base, dfn[N], sdom[N], idom[N], id[N], f[N], fa[N], smin[N], stamp;
5      Graph *g;
6      void predfs(int u) {

```

```

7      id[dfn[u] = stamp++] = u;
8      for (int i = g -> adj[u]; ~i; i = g -> nxt[i]) {
9          int v = g -> v[i];
10         if (dfn[v] < 0) f[v] = u, predfs(v);
11     }
12 }
13 int getfa(int u) {
14     if (fa[u] == u) return u;
15     int ret = getfa(fa[u]);
16     if (dfn[sdom[smin[fa[u]]]] < dfn[sdom[smin[u]]])
17         smin[u] = smin[fa[u]];
18     return fa[u] = ret;
19 }
20 void solve (int s, int n, Graph *raw_graph) {
21     g = raw_graph;
22     base = g -> base;
23     memset(dfn + base, -1, sizeof(*dfn) * n);
24     memset(idom + base, -1, sizeof(*idom) * n);
25     static Graph pred, tmp;
26     pred.init(base, n);
27     for (int i = 0; i < n; ++i) {
28         for (int p = g -> adj[i + base]; ~p; p = g -> nxt[p])
29             pred.ins(g -> v[p], i + base);
30     }
31     stamp = 0; tmp.init(base, n); predfs(s);
32     for (int i = 0; i < stamp; ++i) {
33         fa[id[i]] = smin[id[i]] = id[i];
34     }
35     for (int o = stamp - 1; o >= 0; --o) {
36         int x = id[o];
37         if (o) {
38             sdom[x] = f[x];
39             for (int i = pred.adj[x]; ~i; i = pred.nxt[i]) {
40                 int p = pred.v[i];
41                 if (dfn[p] < 0) continue;
42                 if (dfn[p] > dfn[x]) {
43                     getfa(p);
44                     p = sdom[smin[p]];
45                 }
46                 if (dfn[sdom[x]] > dfn[p]) sdom[x] = p;
47             }
48             tmp.ins(sdom[x], x);
49         }
50         while (~tmp.adj[x]) {
51             int y = tmp.v[tmp.adj[x]];
52             tmp.adj[x] = tmp.nxt[tmp.adj[x]];
53             getfa(y);
54             if (x != sdom[smin[y]]) idom[y] = smin[y];
55             else idom[y] = x;
56         }
57         for (int i = g -> adj[x]; ~i; i = g -> nxt[i])
58             if (f[g -> v[i]] == x) fa[g -> v[i]] = x;
59     }
60     idom[s] = s;
61     for (int i = 1; i < stamp; ++i) {
62         int x = id[i];
63         if (idom[x] != sdom[x]) idom[x] = idom[idom[x]];
64     }
65 }

```



```
66 };
```

## K 短路

```
1 //需保证 GivenEdge 里面边的顺序和 Edge 中一样
2 //两个优先队列要考虑大根还是小根
3 //heap 总是小根堆
4 //dij 不能求正权最长路
5 //INF or -INF
6
7 typedef long long LL;
8 MAXN, MAXK, MAXN, INF //int or LL, it depends
9 const int MAXNODE = MAXN + MAXM * 2; // m + nlgm ???
10 bool used[MAXN];
11 int n, m, cnt, S, T, Kth, N; // m is number of all edges
12 int rt[MAXN], seq[MAXN], adj[MAXN], from[MAXN], dep[MAXN];
13 LL dist[MAXN], w[MAXM], ans[MAXK];
14 struct GivenEdge { //edge given from origin input
15     int u, v, w;
16     GivenEdge() {};
17     GivenEdge(int _u, int _v, int _w): u(_u), v(_v), w(_w) {};
18 } edge[MAXM];
19 struct Edge {
20     int v, nxt, w;
21     Edge() {};
22     Edge(int _v, int _nxt, int _w): v(_v), nxt(_nxt), w(_w) {};
23 } e[MAXM];
24 inline void addedge(int u, int v, int w) {
25     e[++cnt] = Edge(v, adj[u], w); adj[u] = cnt;
26 }
27 inline void dij(int S) { //dij in original graph, spfa if needed
28     for (int i = 1; i <= N; ++i) {
29         dist[i] = INF; dep[i] = INF; used[i] = false; from[i] = 0;
30     }
31     static priority_queue<pair<LL, int>, vector<pair<LL, int>, greater<pair<LL, int>>> >
    hp;
32     while (!hp.empty()) hp.pop();
33     hp.push(make_pair(dist[S] = 0, S));
34     dep[S] = 1;
35     while (!hp.empty()) {
36         pair<LL, int> now = hp.top(); hp.pop();
37         int u = now.second;
38         if (used[u]) continue;
39         else used[u] = true;
40         for (int p = adj[u]; p; p = e[p].nxt) {
41             int v = e[p].v;
42             if (dist[u] + e[p].w < dist[v]) { //different when max or min
43                 dist[v] = dist[u] + e[p].w;
44                 dep[v] = dep[u] + 1;
45                 from[v] = p;
46                 hp.push(make_pair(dist[v], v));
47             }
48         }
49     }
50     for (int i = 1; i <= m; ++i) w[i] = 0;
51     for (int i = 1; i <= N; ++i)
52         if (from[i]) w[from[i]] = -1;
53     for (int i = 1; i <= m; ++i) {
54         if (~w[i] && dist[edge[i].u] < INF && dist[edge[i].v] < INF) {
```

```

55         w[i] = -dist[edge[i].u] + (dist[edge[i].v] + edge[i].w);    //different when max
    ↪ or min
56     } else {
57         w[i] = -1;
58     }
59 }
60 }
61 inline bool cmp_dep(int p, int q) {
62     return dep[p] < dep[q];
63 }
64 struct Heap {
65     LL key;
66     int id, lc, rc, dist;
67     Heap() {};
68     Heap(LL k, int i, int l, int r, int d): key(k), id(i), lc(l), rc(r), dist(d) {};
69     inline void clear() {
70         key = 0;
71         id = lc = rc = dist = 0;
72     }
73 } hp[MAXNODE];
74
75 inline int merge_simple(int u, int v) {
76     if (!u) return v;
77     if (!v) return u;
78     if (hp[u].key > hp[v].key) {
79         swap(u, v);
80     }
81     hp[u].rc = merge_simple(hp[u].rc, v);
82     if (hp[hp[u].lc].dist < hp[hp[u].rc].dist) {
83         swap(hp[u].lc, hp[u].rc);
84     }
85     hp[u].dist = hp[hp[u].rc].dist + 1;
86     return u;
87 }
88
89 inline int merge_full(int u, int v) {
90     if (!u) return v;
91     if (!v) return u;
92     if (hp[u].key > hp[v].key) {
93         swap(u, v);
94     }
95     int nownode = ++cnt;
96     hp[nownode] = hp[u];
97     hp[nownode].rc = merge_full(hp[nownode].rc, v);
98     if (hp[hp[nownode].lc].dist < hp[hp[nownode].rc].dist) {
99         swap(hp[nownode].lc, hp[nownode].rc);
100     }
101     hp[nownode].dist = hp[hp[nownode].rc].dist + 1;
102     return nownode;
103 }
104 priority_queue<pair<LL, int>, vector<pair<LL, int> >, greater<pair<LL, int> > > Q;
105 int main() {
106     scanf("%d%d%d", &n, &m, &Kth);
107     for (int i = 1; i <= m; ++i) {
108         int u, v, w;
109         scanf("%d%d%d", &u, &v, &w);
110         edge[i] = {u, v, w};
111     }
112     N = ; S = ; T = ;
113     memset(adj, 0, sizeof(*adj) * (N + 1));

```

```

114 cnt = 0;
115 for (int i = 1; i <= m; ++i) {
116     addedge(edge[i].v, edge[i].u, edge[i].w); // important!!! reverse the edge
117 }
118 dij(T);
119 if (dist[S] == INF) { //must judge before building heaps; -INF if max kth
120     ...
121     return 0;
122 }
123 for (int i = 1; i <= N; ++i) {
124     seq[i] = i;
125 }
126 sort(seq + 1, seq + N + 1, cmp_dep);
127
128 cnt = 0;
129 memset(adj, 0, sizeof(*adj) * (N + 1));
130 memset(rt, 0, sizeof(*rt) * (N + 1));
131 for (int i = 1; i <= m; ++i) {
132     addedge(edge[i].u, edge[i].v, edge[i].w);
133 }
134 rt[T] = cnt = 0; // now cnt is total nodes in heaps
135 hp[0].dist = -1;
136 for (int i = 1; i <= N; ++i) {
137     int u = seq[i], v = edge[from[u]].v;
138     rt[u] = 0;
139     for (int p = adj[u]; p; p = e[p].nxt) {
140         if (~w[p]) {
141             hp[++cnt] = Heap(w[p], p, 0, 0, 0);
142             rt[u] = merge_simple(rt[u], cnt);
143         }
144     }
145     if (i == 1) continue;
146     rt[u] = merge_full(rt[u], rt[v]);
147 }
148 while (!Q.empty()) Q.pop();
149 Q.push(make_pair(dist[S], 0));
150 edge[0].v = S;
151 for (int kth = 1; kth <= Kth; ++kth) {
152     if (Q.empty()) {
153         ans[kth] = -1;
154         continue;
155     }
156     pair<LL, int> now = Q.top(); Q.pop();
157     ans[kth] = now.first;
158     int p = now.second;
159     if (hp[p].lc) {
160         Q.push(make_pair(+hp[hp[p].lc].key + now.first - hp[p].key,
161 ↪ hp[p].lc)); //different when max or min
162     }
163     if (hp[p].rc) {
164         Q.push(make_pair(+hp[hp[p].rc].key + now.first - hp[p].key,
165 ↪ hp[p].rc)); //different when max or min
166     }
167     if (rt[edge[hp[p].id].v]) {
168         Q.push(make_pair(hp[rt[edge[hp[p].id].v]].key + now.first,
169 ↪ rt[edge[hp[p].id].v])); //different when max or min
170     }
171 }
172 ...
173 for (int i = 1; i <= cnt; ++i) {

```

```

171     hp[i].clear();
172 }
173 }

```

## 最大团搜索

```

1 // Super Fast Maximum Clique
2 // To Build Graph: Maxclique(Edges, Number of Nodes)
3 // To Get Answer: mcqdyn(AnswerNodes Index Array, AnswerLength)
4 typedef bool BB[N];
5 struct Maxclique {
6     const BB* e; int pk, level; const float Tlimit;
7     struct Vertex{ int i, d; Vertex(int i):i(i),d(0){} };
8     typedef vector<Vertex> Vertices; typedef vector<int> ColorClass;
9     Vertices V; vector<ColorClass> C; ColorClass QMAX, Q;
10    static bool desc_degree(const Vertex &vi, const Vertex &vj){
11        return vi.d > vj.d;
12    }
13    void init_colors(Vertices &v){
14        const int max_degree = v[0].d;
15        for(int i = 0; i < (int)v.size(); i++) v[i].d = min(i, max_degree) + 1;
16    }
17    void set_degrees(Vertices &v){
18        for(int i = 0, j; i < (int)v.size(); i++)
19            for(v[i].d = j = 0; j < (int)v.size(); j++)
20                v[i].d += e[v[i].i][v[j].i];
21    }
22    struct StepCount{ int i1, i2; StepCount():i1(0),i2(0){} };
23    vector<StepCount> S;
24    bool cut1(const int pi, const ColorClass &A){
25        for(int i = 0; i < (int)A.size(); i++) if (e[pi][A[i]]) return true;
26        return false;
27    }
28    void cut2(const Vertices &A, Vertices &B){
29        for(int i = 0; i < (int)A.size() - 1; i++)
30            if(e[A.back().i][A[i].i])
31                B.push_back(A[i].i);
32    }
33    void color_sort(Vertices &R){
34        int j = 0, maxno = 1, min_k = max((int)QMAX.size() - (int)Q.size() + 1, 1);
35        C[1].clear(), C[2].clear();
36        for(int i = 0; i < (int)R.size(); i++) {
37            int pi = R[i].i, k = 1;
38            while(cut1(pi, C[k])) k++;
39            if(k > maxno) maxno = k, C[maxno + 1].clear();
40            C[k].push_back(pi);
41            if(k < min_k) R[j++].i = pi;
42        }
43        if(j > 0) R[j - 1].d = 0;
44        for(int k = min_k; k <= maxno; k++)
45            for(int i = 0; i < (int)C[k].size(); i++)
46                R[j].i = C[k][i], R[j++].d = k;
47    }
48    void expand_dyn(Vertices &R){// diff -> diff with no dyn
49        S[level].i1 = S[level].i1 + S[level - 1].i1 - S[level].i2;//diff
50        S[level].i2 = S[level - 1].i1;//diff
51        while((int)R.size()) {
52            if((int)Q.size() + R.back().d > (int)QMAX.size()){
53                Q.push_back(R.back().i); Vertices Rp; cut2(R, Rp);

```

```

54         if((int)Rp.size()){
55             if((float)S[level].i1 / ++pk < Tlimit) degree_sort(Rp); //diff
56             color_sort(Rp);
57             S[level].i1++, level++; //diff
58             expand_dyn(Rp);
59             level--; //diff
60         }
61         else if((int)Q.size() > (int)QMAX.size()) QMAX = Q;
62         Q.pop_back();
63     }
64     else return;
65     R.pop_back();
66 }
67 }
68 void mcqdyn(int* maxclique, int &sz){
69     set_degrees(V); sort(V.begin(), V.end(), desc_degree); init_colors(V);
70     for(int i = 0; i < (int)V.size() + 1; i++) S[i].i1 = S[i].i2 = 0;
71     expand_dyn(V); sz = (int)QMAX.size();
72     for(int i = 0; i < (int)QMAX.size(); i++) maxclique[i] = QMAX[i];
73 }
74 void degree_sort Vertices &R){
75     set_degrees(R); sort(R.begin(), R.end(), desc_degree);
76 }
77 Maxclique(const BB* conn, const int sz, const float tt = 0.025) \
78 : pk(0), level(1), Tlimit(tt){
79     for(int i = 0; i < sz; i++) V.push_back(Vertex(i));
80     e = conn, C.resize(sz + 1), S.resize(sz + 1);
81 }
82 };

```

## 极大团计数

```

1 Bool g[][]为图的邻接矩阵，图点的标号由1至n。
2 void dfs(int size){
3     int i, j, k, t, cnt, best = 0;
4     bool bb;
5     if (ne[size]==ce[size]){
6         if (ce[size]==0) ++ans;
7         return;
8     }
9     for (t=0, i=1; i<=ne[size]; ++i) {
10         for (cnt=0, j=ne[size]+1; j<=ce[size]; ++j)
11             if (!g[list[size][i]][list[size][j]]) ++cnt;
12         if (t==0 || cnt<best) t=i, best=cnt;
13     }
14     if (t && best<=0) return;
15     for (k=ne[size]+1; k<=ce[size]; ++k) {
16         if (t>0){
17             for (i=k; i<=ce[size]; ++i) if (!g[list[size][t]][list[size][i]]) break;
18             swap(list[size][k], list[size][i]);
19         }
20         i=list[size][k];
21         ne[size+1]=ce[size+1]=0;
22         for (j=1; j<k; ++j) if (g[i][list[size][j]])
↪ list[size+1][++ne[size+1]]=list[size][j];
23         for (ce[size+1]=ne[size+1], j=k+1; j<=ce[size]; ++j)
24             if (g[i][list[size][j]]) list[size+1][++ce[size+1]]=list[size][j];
25         dfs(size+1);
26         ++ne[size];

```

```

27     --best;
28     for (j=k+1, cnt=0; j<=ce[size]; ++j) if (!g[i][list[size][j]]) ++cnt;
29     if (t==0 || cnt<best) t=k, best=cnt;
30     if (t && best<=0) break;
31 }
32 }
33 void work(){
34     int i;
35     ne[0]=0; ce[0]=0;
36     for (i=1; i<=n; ++i) list[0][++ce[0]]=i;
37     ans=0;
38     dfs(0);
39 }

```

## 欧拉回路

```

1 //从一个奇度点 dfs, sqn 即为回路/路径
2 //first 存点, second 存边的编号, 正反边编号一致
3 //清空 cur、used 数组
4 void getCycle(int u) {
5     for(int &i=cur[u]; i < (int)adj[u].size(); ++ i) {
6         int id = adj[u][i].second;
7         if (used[id]) continue;
8         used[id] = true;
9         getCycle(adj[u][i].first);
10    }
11    sqn.push_back(u);
12 }

```

## 朱刘最小树形图

```

1 struct D_MT {
2     struct Edge {
3         int u, v, w;
4         inline Edge() {}
5         inline Edge(int _u, int _v, int _w):u(_u), v(_v), w(_w) {}
6     };
7     int nn, mm, n, m, vis[maxn], pre[maxn], id[maxn], in[maxn];
8     Edge edges[maxn], bac[maxn];
9     void init(int _n) {
10         n = _n; m = 0;
11     }
12     void AddEdge(int u, int v, int w) {
13         edges[m++] = Edge(u, v, w);
14     }
15     int work(int root) {
16         int ret = 0;
17         while(true) {
18             for (int i = 0; i < n; i++) in[i] = inf + 1;
19             for (int i = 0; i < m; i++) {
20                 int u = edges[i].u, v = edges[i].v;
21                 if(edges[i].w < in[v] && u != v){
22                     in[v] = edges[i].w;
23                     pre[v] = u;
24                 }
25             }
26             for (int i = 0; i < n; i++) {
27                 if(i == root) continue;

```

```
28         if(in[i] == inf + 1) return inf;
29     }
30     int cnt = 0;
31     for (int i = 0; i < n; i++) id[i] = vis[i] = -1;
32     in[root] = 0;
33     for (int i = 0; i < n; i++) {
34         ret += in[i];
35         int v = i;
36         while (vis[v] != i && id[v] == -1 && v != root ){
37             vis[v] = i; v = pre[v];
38         }
39         if (v != root && id[v] == -1) {
40             for (int u = pre[v]; u != v; u = pre[u]) id[u] = cnt;
41             id[v] = cnt++;
42         }
43     }
44     if (!cnt) break;
45     for (int i = 0; i < n; i++)
46         if (id[i] == -1) id[i] = cnt++;
47     for (int i = 0; i < m; i++){
48         int u = edges[i].u, v = edges[i].v;
49         edges[i].v = id[v]; edges[i].u = id[u];
50         if(id[u] != id[v]) edges[i].w -= in[v];
51     }
52     n = cnt; root = id[root];
53 }
54 return ret;
55 }
56 } MT;
```





# Chapter 4

## 数据结构

### Kd-tree

```
1 int n;
2 LL norm(const LL &x) {
3     // For manhattan distance
4     //return std::abs(x);
5     // For euclid distance
6     return x * x;
7 }
8
9 struct P{
10     int a[2],val;
11     int id;
12     int& operator[](int s){return a[s];}
13     const int& operator[](int s)const{return a[s];}
14
15     LL dis(const P &b)const{
16         LL ans=0;
17         for (int i = 0; i < 2; ++i) {
18             ans += norm(a[i] - b[i]);
19         }
20         return ans;
21     }
22 }p[maxn];
23
24 bool operator==(const P &a,const P &b){
25     for(int i=0;i<DIM;i++)
26         if(a[i]!=b[i])
27             return false;
28     return true;
29 }
30 bool byVal(P a,P b){
31     return a.val!=b.val ? a.val<b.val : a.id<b.id;
32 }
33
34 struct Rec{
35     int mn[DIM],mx[DIM];
36     Rec(){}
37     Rec(const P &p){
38         for(int i=0;i<DIM;i++){
39             mn[i]=mx[i]=p[i];
40         }
41     }
42     void add(const P &p){
43         for(int i=0;i<DIM;i++){
44             mn[i]=min(p[i],mn[i]);
```

```

45         mx[i]=max(p[i],mx[i]);
46     }
47 }
48
49 LL dis(const P &p) {
50     LL ans = 0;
51     for (int i = 0; i < 2; ++i) {
52         // For minimum distance
53         ans += norm(min(max(p[i], mn[i]), mx[i]) - p[i]);
54         // For maximum distance
55         //ans += std::max(norm(max[i] - p[i]), norm(min[i] - p[i]));
56     }
57     return ans;
58 }
59 };
60 inline Rec operator+(const Rec &ls,const Rec &rs){
61     static Rec rec;
62     for(int i=0;i<DIM;i++){
63         rec.mn[i]=min(ls.mn[i],rs.mn[i]);
64         rec.mx[i]=max(ls.mx[i],rs.mx[i]);
65     }
66     return rec;
67 }
68 struct node{
69     Rec rec;
70     P sep;
71     int sum,siz;
72     node *c[2];
73     node *rz(){
74         sum=sep.val;
75         rec=Rec(sep);
76         siz=1;
77         if(c[0]){
78             sum+=c[0]->sum;
79             rec=rec+c[0]->rec;
80             siz+=c[0]->siz;
81         }
82         if(c[1]){
83             sum+=c[1]->sum;
84             rec=rec+c[1]->rec;
85             siz+=c[1]->siz;
86         }
87         return this;
88     }
89     node(){sum=0;siz=1;c[0]=c[1]=0;}
90 }*root,*re,pool[maxn],*cur=pool;
91 node *sta[maxn];
92 P tmp[maxn];
93 int D,si;
94 void init(){
95     si=0;
96     cur=pool;
97     root=0;
98 }
99 bool cmp(const P &A,const P &B){
100
101     if(!(A[D]==B[D]))
102         return A[D]<B[D];
103     return A.id<B.id;
104 }

```

```

105 int top;
106 node *newnode(){
107     if(si)return sta[si--];
108     return cur++;
109 }
110 node* build(P *p,int l,int r,int d){
111     int mid=(l+r)>>1;D=d;
112     nth_element(p+l,p+mid,p+r+1,cmp);
113     node *t=newnode();
114     t->sep=p[mid];
115     if(l<=mid-1)
116         t->c[0]=build(p,l,mid-1,d^1);
117     if(mid+1<=r)
118         t->c[1]=build(p,mid+1,r,d^1);
119     return t->rz();
120 }
121 void dfs(node *&t){
122     if(t->c[0])dfs(t->c[0]);
123     tmp[++top]=t->sep;
124     if(t->c[1])dfs(t->c[1]);
125     sta[++si]=t;*t=node();
126     //delete t;
127 }
128 node* rebuild(node *&t){
129     if(!t)return 0;
130     top=0;dfs(t);
131     return build(tmp,1,top,0);
132 }
133 #define siz(x) (x?x->siz:0)
134 void Add(node *&t,const P &p,int d=0){//调用前 re=0; 调用后 rebuild(re);
135     D=d;
136     if(!t){
137         t=newnode();
138         t->sep=p;t->rz();
139         return;
140     }
141     if(t->sep==p){
142         t->sep.val+=p.val;
143         t->rz();
144         return;
145     }
146     if(p[D]<t->sep[D])
147         Add(t->c[0],p,d^1);
148     else
149         Add(t->c[1],p,d^1);
150
151     t->rz();
152
153     if(max(siz(t->c[0]),siz(t->c[1]))>0.7*t->siz)
154         re=t;
155 }
156 int ans;
157
158 bool Out(const Rec &a,const Rec &b){
159     for(int i=0;i<DIM;i++){
160         int l=max(a.mn[i],b.mn[i]);
161         int r=min(a.mx[i],b.mx[i]);
162         if(l>r)
163             return true;
164     }

```

```

165     return false;
166 }
167 bool In(const Rec &a,const Rec &b){
168     for(int i=0;i<DIM;i++){
169         if(a.mn[i]<b.mn[i])
170             return false;
171         if(a.mx[i]>b.mx[i])
172             return false;
173     }
174     return true;
175 }
176
177 bool In(const P &a,const Rec &b){
178     for(int i=0;i<DIM;i++){
179         if(!(b.mn[i]<=a[i]&&a[i]<=b.mx[i]))
180             return false;
181     }
182     return true;
183 }
184
185 void Q(node *t,const Rec &R){
186     if(Out(t->rec,R))return ;
187     if(In(t->rec,R)){
188         ans+=t->sum;
189         return;
190     }
191     if(In(t->sep,R))
192         ans+=t->sep.val;
193     if(t->c[0])
194         Q(t->c[0],R);
195     if(t->c[1])
196         Q(t->c[1],R);
197 }
198
199 priority_queue<pair<long long, int> > kNN;
200 void query(node *t, const P &p, int k, int d = 0) { //用钱清空 kNN
201     D=d;
202     if (!t || ((int)kNN.size() == k && t->rec.dis(p) > kNN.top().first)) {
203         return;
204     }
205     kNN.push(make_pair(t->sep.dis(p), t->sep.id));
206     if ((int)kNN.size() > k) {
207         kNN.pop();
208     }
209     if (cmp(p, t->sep)) {
210         query(t->c[0], p, k, d ^ 1);
211         query(t->c[1], p, k, d ^ 1);
212     } else {
213         query(t->c[1], p, k, d ^ 1);
214         query(t->c[0], p, k, d ^ 1);
215     }
216 }

```

## LCT

```

1 struct LCT{
2     struct node{
3         bool rev;
4         int mx,val;

```

```

5     node *f,*c[2];
6     bool d(){return this==f->c[1];}
7     bool rt(){return !f||(f->c[0]!=this&&f->c[1]!=this);}
8     void sets(node *x,int d){pd();if(x)x->f=this;c[d]=x;rz();}
9     void makerv(){rev^=1;swap(c[0],c[1]);}
10    void pd(){
11        if(rev){
12            if(c[0])c[0]->makerv();
13            if(c[1])c[1]->makerv();
14            rev=0;
15        }
16    }
17    void rz(){
18        mx=val;
19        if(c[0])mx=max(mx,c[0]->mx);
20        if(c[1])mx=max(mx,c[1]->mx);
21    }
22    }nd[1e4+1];
23    void rot(node *x){
24        node *y=x->f;if(!y->rt())y->f->pd();
25        y->pd();x->pd();bool d=x->d();
26        y->sets(x->c[!d],d);
27        if(y->rt())x->f=y->f;
28        else y->f->sets(x,y->d());
29        x->sets(y,!d);
30    }
31    void splay(node *x){
32        while(!x->rt())
33            if(x->f->rt())rot(x);
34            else if(x->d()==x->f->d())rot(x->f),rot(x);
35            else rot(x),rot(x);
36    }
37    node* access(node *x){
38        node *y=0;
39        for(;x;x=x->f){
40            splay(x);
41            x->sets(y,1);y=x;
42        }return y;
43    }
44    void makert(node *x){
45        access(x)->makerv();
46        splay(x);
47    }
48    void link(node *x,node *y){
49        makert(x);
50        x->f=y;
51        access(x);
52    }
53    void cut(node *x,node *y){
54        makert(x);access(y);splay(y);
55        y->c[0]=x->f=0;
56        y->rz();
57    }
58    void link(int x,int y){link(nd+x,nd+y);}
59    void cut(int x,int y){cut(nd+x,nd+y);}
60 }T;

```

## 树状数组上二分第 k 大

```

1 int find(int k){
2     int cnt=0,ans=0;
3     for(int i=22;i>=0;i--){
4         ans+=(1<<i);
5         if(ans>n || cnt+d[ans]>=k)ans-=(1<<i);
6         else cnt+=d[ans];
7     }
8     return ans+1;
9 }

```

## Treap

```

1 #include<bits/stdc++.h>
2 using namespace std;
3 const int maxn=1e5+5;
4 #define sz(x) (x?x->siz:0)
5 struct Treap{
6     struct node{
7         int key,val;
8         int siz,s;
9         node *c[2];
10        node(int v=0){
11            val=v;
12            key=rand();
13            siz=1,s=1;
14            c[0]=c[1]=0;
15        }
16        void rz(){siz=s;if(c[0])siz+=c[0]->siz;if(c[1])siz+=c[1]->siz;}
17    }pool[maxn],*cur,*root;
18    Treap(){cur=pool;}
19    node* newnode(int val){return *cur=node(val),cur++;}
20    void rot(node *&t,int d){
21        if(!t->c[d])t=t->c[!d];
22        else{
23            node *p=t->c[d];t->c[d]=p->c[!d];
24            p->c[!d]=t;t->rz();p->rz();t=p;
25        }
26    }
27    void insert(node *&t,int x){
28        if(!t){t=newnode(x);return;}
29        if(t->val==x){t->s++;t->siz++;return;}
30        insert(t->c[x>t->val],x);
31        if(t->key<t->c[x>t->val]->key)
32            rot(t,x>t->val);
33        else t->rz();
34    }
35    void del(node *&t,int x){
36        if(!t)return;
37        if(t->val==x){
38            if(t->s>1){t->s--;t->siz--;return;}
39            if(!t->c[0]||!t->c[1]){
40                if(!t->c[0])t=t->c[1];
41                else t=t->c[0];
42                return;
43            }
44            int d=t->c[0]->key<t->c[1]->key;
45            rot(t,d);

```

```

46         del(t,x);
47         return;
48     }
49     del(t->c[x>t->val],x);
50     t->rz();
51 }
52 int pre(node *t,int x){
53     if(!t)return INT_MIN;
54     int ans=pre(t->c[x>t->val],x);
55     if(t->val<x)ans=max(ans,t->val);
56     return ans;
57 }
58 int nxt(node *t,int x){
59     if(!t)return INT_MAX;
60     int ans=nxt(t->c[x>=t->val],x);
61     if(t->val>x)ans=min(ans,t->val);
62     return ans;
63 }
64 int rank(node *t,int x){
65     if(!t)return 0;
66     if(t->val==x)return sz(t->c[0]);
67     if(t->val<x)return sz(t->c[0])+t->s+rank(t->c[1],x);
68     if(t->val>x)return rank(t->c[0],x);
69 }
70 int kth(node *t,int x){
71     if(sz(t->c[0])>=x)return kth(t->c[0],x);
72     if(sz(t->c[0])+t->s>=x)return t->val;
73     return kth(t->c[1],x-t->s-sz(t->c[0]));
74 }
75 void deb(node *t){
76     if(!t)return;
77     deb(t->c[0]);
78     printf("%d ",t->val);
79     deb(t->c[1]);
80 }
81 void insert(int x){insert(root,x);}
82 void del(int x){del(root,x);}
83 int pre(int x){return pre(root,x);}
84 int nxt(int x){return nxt(root,x);}
85 int rank(int x){return rank(root,x);}
86 int kth(int x){return kth(root,x);}
87 void deb(){deb(root);puts("");}
88 }T;

```

## FHQ-Treap

```

1 #include<bits/stdc++.h>
2 using namespace std;
3 typedef long long LL;
4 const int maxn=1e5+5;
5 int in(){
6     int r=0,f=1;char c=getchar();
7     while(!isdigit(c))f=c=='-'?-1:f,c=getchar();
8     while(isdigit(c))r=r*10+c-'0',c=getchar();
9     return r*f;
10 }
11 int n,m;
12 #define sz(x) (x?x->siz:0)
13 struct node{

```

```

14     int siz,key;
15     LL val,sum;
16     LL mu,a,d;
17     node *c[2],*f;
18     void split(int ned,node *&p,node *&q);
19     node* rz(){
20         sum=val;siz=1;
21         if(c[0])sum+=c[0]->sum,siz+=c[0]->siz;
22         if(c[1])sum+=c[1]->sum,siz+=c[1]->siz;
23         return this;
24     }
25     void make(LL _mu,LL _a,LL _d){
26         sum=sum*_mu+_a*siz+_d*siz*(siz-1)/2;
27         val=val*_mu+_a+_d*sz(c[0]);
28         mu*=_mu;a=a*_mu+_a;d=d*_mu+_d;
29     }
30     void pd(){
31         if(mu==1&&a==0&&d==0)return;
32         if(c[0])c[0]->make(mu,a,d);
33         if(c[1])c[1]->make(mu,a+d+d*sz(c[0]),d);
34         mu=1;a=d=0;
35     }
36     node(){mu=1;}
37 }nd[maxn*2],*root;
38 node *merge(node *p,node *q){
39     if(!p||!q)return p?p->rz():(q?q->rz():0);
40     p->pd();q->pd();
41     if(p->key<q->key){
42         p->c[1]=merge(p->c[1],q);
43         return p->rz();
44     }else{
45         q->c[0]=merge(p,q->c[0]);
46         return q->rz();
47     }
48 }
49 void node::split(int ned,node *&p,node *&q){
50     if(!ned){p=0;q=this;return;}
51     if(ned==siz){p=this;q=0;return;}
52     pd();
53     if(sz(c[0])>=ned){
54         c[0]->split(ned,p,q);c[0]=0;rz();
55         q=merge(q,this);
56     }else{
57         c[1]->split(ned-sz(c[0])-1,p,q);c[1]=0;rz();
58         p=merge(this,p);
59     }
60 }
61 int tot;
62 void C(int l,int r,int v){
63     node *p,*q,*x,*y;
64     root->split(l-1,p,q);
65     q->split(r-l+1,x,y);
66     x->make(0,v,0);x->pd();
67     root=merge(p,merge(x,y));
68 }
69 void A(int l,int r,int d){
70     node *p,*q,*x,*y;
71     root->split(l-1,p,q);
72     q->split(r-l+1,x,y);
73     x->make(1,d,d);x->pd();

```



```

74     root=merge(p,merge(x,y));
75 }
76 void I(int ps,int v){
77     node *p,*q;
78     root->split(ps-1,p,q);
79     node *x=nd(++tot);
80     x->key=rand();x->val=v;x->rz();
81     root=merge(merge(p,x),q);
82 }
83 LL Q(int l,int r){
84     node *p,*q,*x,*y;
85     root->split(l-1,p,q);
86     q->split(r-l+1,x,y);
87     LL ans=x->sum;
88     root=merge(p,merge(x,y));
89     return ans;
90 }
91 int main(){
92     // freopen("bzoj3188.in","r",stdin);
93     n=in();m=in();
94     for(int i=1;i<=n;i++){
95         nd[i].val=in();
96         nd[i].key=rand();
97         nd[i].rz();
98         root=merge(root,nd+i);
99     }tot=n;
100     while(m--){
101         int ty=in();
102         int l,r;
103         if(ty==1){
104             l=in();r=in();
105             C(l,r,in());
106         }else if(ty==2){
107             l=in();r=in();
108             A(l,r,in());
109         }else if(ty==3){
110             int ps=in();
111             I(ps,in());
112         }else if(ty==4){
113             l=in();r=in();
114             printf("%lld\n",Q(l,r));
115         }
116     }
117     return 0;
118 }

```

## 真-FHQTreap

```

1  const int mo=1e9+7;
2  int rnd(){
3      static int x=1;
4      return x=(x*23333+233);
5  }
6  int rnd(int n){
7      int x=rnd();
8      if(x<0)x=-x;
9      return x%n+1;
10 }
11 struct node{

```

```

12     int siz,key;
13     int val;
14     LL sum;
15     node *c[2];
16     node* rz(){
17         sum=val;siz=1;
18         if(c[0])sum+=c[0]->sum,siz+=c[0]->siz;
19         if(c[1])sum+=c[1]->sum,siz+=c[1]->siz;
20         return this;
21     }
22     node(){
23     }
24     node(int v){
25         siz=1;key=rnd();
26         val=v;sum=v;
27         c[0]=c[1]=0;
28     }
29 }pool[maxn*8],*root,*cur=pool,*old_root,*stop;
30 node *newnode(int v=0){
31     *cur=node(v);
32     return cur++;
33 }
34 node *old_merge(node *p,node *q){
35     if(!p&&!q)return 0;
36     node *u=0;
37     if(!p||!q)return u=p?p->rz():(q?q->rz():0);
38     if(rnd(sz(p)+sz(q))<sz(p)){
39         u=p;
40         u->c[1]=old_merge(u->c[1],q);
41     }else{
42         u=q;
43         u->c[0]=old_merge(p,u->c[0]);
44     }
45     return u->rz();
46 }
47 node *merge(node *p,node *q){
48     if(!p&&!q)return 0;
49     node *u=newnode();
50     if(!p||!q)return u=p?p->rz():(q?q->rz():0);
51     if(rnd(sz(p)+sz(q))<sz(p)){
52         *u=*p;
53         u->c[1]=merge(u->c[1],q);
54     }else{
55         *u=*q;
56         u->c[0]=merge(p,u->c[0]);
57     }
58     return u->rz();
59 }
60 node *split(node *u,int l,int r){
61     if(l>r||!u)return 0;
62     node *x=0;
63     if(l==1&&r==sz(u)){
64         x=newnode();
65         *x=*u;
66         return x->rz();
67     }
68     int lsz=sz(u->c[0]);
69     if(r<=lsz)
70         return split(u->c[0],l,r);
71     if(l>lsz+1)

```

```

72     return split(u->c[1],l-lsz-1,r-lsz-1);
73     x=newnode();
74     *x=*u;
75     x->c[0]=split(u->c[0],l,lsz);
76     x->c[1]=split(u->c[1],1,r-lsz-1);
77     return x->rz();
78 }

```

## 带修改莫队上树

```

1 bool operator<(qes a,qes b){
2     if(dfn[a.x]/B!=dfn[b.x]/B) return dfn[a.x]/B<dfn[b.x]/B;
3     if(dfn[a.y]/B!=dfn[b.y]/B) return dfn[a.y]/B<dfn[b.y]/B;
4     if(a.tm/B!=b.tm/B) return a.tm/B<b.tm/B;
5     return a.tm<b.tm;
6 }
7 void vxor(int x){
8     if(vis[x]) ans-=(LL)W[cnt[col[x]]]*V[col[x]],cnt[col[x]]--;
9     else cnt[col[x]]++,ans+=(LL)W[cnt[col[x]]]*V[col[x]];
10    vis[x]^=1;
11 }
12 void change(int x,int y){
13     if(vis[x]){
14         vxor(x);col[x]=y;vxor(x);
15     }else col[x]=y;
16 }
17 void TimeMachine(int tar){//XD
18     for(int i=now+1;i<=tar;i++) change(C[i].x,C[i].y);
19     for(int i=now;i>tar;i--) change(C[i].x,C[i].pre);
20     now=tar;
21 }
22 void vxor(int x,int y){
23     while(x!=y) if(dep[x]>dep[y]) vxor(x),x=fa[x];
24     else vxor(y),y=fa[y];
25 }
26 for(int i=1;i<=q;i++){
27     int ty=getint(),x=getint(),y=getint();
28     if(ty&&dfn[x]>dfn[y]) swap(x,y);
29     if(ty==0) C[++Csize]=(oper){x,y,pre[x],i},pre[x]=y;
30     else Q[Qsize+1]=(qes){x,y,Qsize+1,Csize},Qsize++;
31 }sort(Q+1,Q+1+Qsize);
32 int u=Q[1].x,v=Q[1].y;
33 TimeMachine(Q[1].tm);
34 vxor(Q[1].x,Q[1].y);
35 int LCA=lca(Q[1].x,Q[1].y);
36 vxor(LCA);anss[Q[1].id]=ans;vxor(LCA);
37 for(int i=2;i<=Qsize;i++){
38     TimeMachine(Q[i].tm);
39     vxor(Q[i-1].x,Q[i].x);
40     vxor(Q[i-1].y,Q[i].y);
41     int LCA=lca(Q[i].x,Q[i].y);
42     vxor(LCA);
43     anss[Q[i].id]=ans;
44     vxor(LCA);
45 }

```

## 虚树

```
1 int a[maxn*2], sta[maxn*2];
2 int top=0, k;
3 void build(){
4     top=0;
5     sort(a, a+k, bydfn);
6     k=unique(a, a+k)-a;
7     sta[top++]=1; _n=k;
8     for(int i=0; i<k; i++){
9         int LCA=lca(a[i], sta[top-1]);
10        while(dep[LCA]<dep[sta[top-1]]){
11            if(dep[LCA]>=dep[sta[top-2]]){
12                add_edge(LCA, sta[--top]);
13                if(sta[top-1]!=LCA) sta[top++]=LCA;
14                break;
15            }add_edge(sta[top-2], sta[top-1]); top--;
16        }if(sta[top-1]!=a[i]) sta[top++]=a[i];
17    }
18    while(top>1)
19        add_edge(sta[top-2], sta[top-1]), top--;
20    for(int i=0; i<k; i++) inr[a[i]]=1;
21 }
```

## Chapter 5

# 字符串

### Manacher

```
1 //prime is the origin string(0-base)
2 //-10,-1,-20 are added to s
3 //length of s is exactly 2 * l + 3
4 inline void manacher(char prime[]) {
5     int l = strlen(prime), n = 0;
6     s[n++] = -10;
7     s[n++] = -1;
8     for (int i = 0; i < l; ++i) {
9         s[n++] = prime[i];
10        s[n++] = -1;
11    }
12    s[n++] = -20; f[0] = 1;
13    int mx = 0, id = 0;
14    for (int i = 1; i + 1 < n; ++i) {
15        f[i] = i > mx ? 1 : min(f[id * 2 - i], mx - i + 1);
16        while (s[i + f[i]] == s[i - f[i]]) ++f[i];
17        if (i + f[i] - 1 > mx) {
18            mx = i + f[i] - 1;
19            id = i;
20        }
21    }
22 }
```

### 指针版回文自动机

```
1 /*
2  * Palindrome Automaton - pointer version
3  * PAMPAMPAM? PAMPAMPAM!
4  */
5
6 namespace PAM {
7     struct Node *pool_pointer;
8     struct Node {
9         Node *fail, *to[26];
10        int cnt, len;
11
12        Node() {}
13        Node(int len): len(len) {
14            memset(to, 0, sizeof(to));
15            fail = 0;
16            cnt = 0;
17        }
18 }
```

```

19     void *operator new (size_t) {
20         return pool_pointer++;
21     }
22 } pool[100005], *root[2], *last;
23 int pam_len, str[100005];
24
25 void init() {
26     pool_pointer = pool;
27     root[0] = new Node(0);
28     root[1] = new Node(-1);
29     root[0]->fail = root[1]->fail = root[1];
30     str[pam_len = 0] = -1; // different from all characters
31     last = root[0];
32 }
33
34 void extend(char ch) {
35     static Node *p, *np, *q;
36
37     int x = str[++pam_len] = ch - 'a';
38
39     p = last;
40     while (str[pam_len - p->len - 1] != x)
41         p = p->fail;
42     if (!p->to[x]) {
43         np = new Node(p->len + 2), q = p->fail;
44         while (str[pam_len - q->len - 1] != x) q = q->fail;
45         np->fail = q->to[x] ? q->to[x] : root[0];
46         p->to[x] = np;
47     }
48     last = p->to[x];
49     ++last->cnt;
50 }
51 }

```

## 后缀数组

```

1 const int maxl=1e5+1e4+5;
2 const int maxn=maxl*2;
3 int a[maxn],x[maxn],y[maxn],c[maxn],sa[maxn],rank[maxn],height[maxn];
4 void calc_sa(int n){
5     int m=alphabet,k=1;
6     memset(c,0,sizeof(*c)*(m+1));
7     for(int i=1;i<=n;i++)c[x[i]=a[i]]++;
8     for(int i=1;i<=m;i++)c[i]+=c[i-1];
9     for(int i=1;i<=n;i++)sa[c[x[i]]--]=i;
10    for(;k<=n;k<=1){
11        int tot=k;
12        for(int i=n-k+1;i<=n;i++)y[i-n+k]=i;
13        for(int i=1;i<=n;i++)
14            if(sa[i]>k)y[++tot]=sa[i]-k;
15        memset(c,0,sizeof(*c)*(m+1));
16        for(int i=1;i<=n;i++)c[x[i]]++;
17        for(int i=1;i<=m;i++)c[i]+=c[i-1];
18        for(int i=n;i>=1;i--)sa[c[x[y[i]]]--]=y[i];
19        for(int i=1;i<=n;i++)y[i]=x[i];
20        tot=1;x[sa[1]]=1;
21        for(int i=2;i<=n;i++){
22            if(max(sa[i],sa[i-1])+k>n||y[sa[i]]!=y[sa[i-1]]||y[sa[i]+k]!=y[sa[i-1]+k])
23                ++tot;

```

```

24         x[sa[i]]=tot;
25     }
26     if(tot==n)break;else m=tot;
27 }
28 }
29 void calc_height(int n){
30     for(int i=1;i<=n;i++)rank[sa[i]]=i;
31     for(int i=1;i<=n;i++){
32         height[rank[i]]=max(0,height[rank[i-1]]-1);
33         if(rank[i]==1)continue;
34         int j=sa[rank[i]-1];
35         while(max(i,j)+height[rank[i]]<=n&&a[i+height[rank[i]]]==a[j+height[rank[i]]])
36             ++height[rank[i]];
37     }
38 }

```

## 最小表示法

```

1 int solve(char *text, int length) { // 0-base , 多解答案为起点最小
2     int i = 0, j = 1, delta = 0;
3     while (i < length && j < length && delta < length) {
4         char tokeni = text[(i + delta) % length];
5         char tokenj = text[(j + delta) % length];
6         if (tokeni == tokenj) {
7             delta++;
8         } else {
9             if (tokeni > tokenj) {
10                 i += delta + 1;
11             } else {
12                 j += delta + 1;
13             }
14             if (i == j) {
15                 j++;
16             }
17             delta = 0;
18         }
19     }
20     return std::min(i, j);
21 }

```





## Chapter 6

# 计算几何

### 点类

```
1 int sgn(double x){return (x>eps)-(x<-eps);}
2 int sgn(double a,double b){return sgn(a-b);}
3 double sqr(double x){return x*x;}
4 struct P{
5     double x,y;
6     P(){}
7     P(double x,double y):x(x),y(y){}
8     double len2(){
9         return sqr(x)+sqr(y);
10    }
11    double len(){
12        return sqrt(len2());
13    }
14    void print(){
15        printf("%.3f,%.3f\n",x,y);
16    }
17    P turn90(){return P(-y,x);}
18    P norm(){return P(x/len(),y/len());}
19 };
20 bool operator==(P a,P b){
21     return !sgn(a.x-b.x) and !sgn(a.y-b.y);
22 }
23 P operator+(P a,P b){
24     return P(a.x+b.x,a.y+b.y);
25 }
26 P operator-(P a,P b){
27     return P(a.x-b.x,a.y-b.y);
28 }
29 P operator*(P a,double b){
30     return P(a.x*b,a.y*b);
31 }
32 P operator/(P a,double b){
33     return P(a.x/b,a.y/b);
34 }
35 double operator^(P a,P b){
36     return a.x*b.x + a.y*b.y;
37 }
38 double operator*(P a,P b){
39     return a.x*b.y - a.y*b.x;
40 }
41 double det(P a,P b,P c){
42     return (b-a)*(c-a);
43 }
44 double dis(P a,P b){
```

```

45     return (b-a).len();
46 }
47 double Area(vector<P>poly){
48     double ans=0;
49     for(int i=1;i<poly.size();i++)
50         ans+=(poly[i]-poly[0])*(poly[(i+1)%poly.size()]-poly[0]);
51     return fabs(ans)/2;
52 }
53 struct L{
54     P a,b;
55     L(){}
56     L(P a,P b):a(a),b(b){}
57     P v(){return b-a;}
58 };
59 bool onLine(P p,L l){
60     return sgn((l.a-p)*(l.b-p))==0;
61 }
62 bool onSeg(P p,L s){
63     return onLine(p,s) and sgn((s.b-s.a)^(p-s.a))>=0 and sgn((s.a-s.b)^(p-s.b))>=0;
64 }
65 bool parallel(L l1,L l2){
66     return sgn(l1.v()*l2.v())==0;
67 }
68 P intersect(L l1,L l2){
69     double s1=det(l1.a,l1.b,l2.a);
70     double s2=det(l1.a,l1.b,l2.b);
71     return (l2.a*s2-l2.b*s1)/(s2-s1);
72 }
73 P project(P p,L l){
74     return l.a+l.v()*((p-l.a)^l.v())/l.v().len2();
75 }
76 double dis(P p,L l){
77     return fabs((p-l.a)*l.v())/l.v().len();
78 }
79 int dir(P p,L l){
80     int t=sgn((p-l.b)*(l.b-l.a));
81     if(t<0)return -1;
82     if(t>0)return 1;
83     return 0;
84 }
85 bool segIntersect(L l1,L l2){//strictly
86     if(dir(l2.a,l1)*dir(l2.b,l1)<0&&dir(l1.a,l2)*dir(l1.b,l2)<0)
87         return true;
88     return false;
89 }
90 bool in_tri(P pt,P *p){//change p
91     if((p[1]-p[0])*(p[2]-p[0])<0)
92         reverse(p,p+3);
93     for(int i=0;i<3;i++){
94         if(dir(pt,L(p[i],p[(i+1)%3]))==1)
95             return false;
96     }
97     return true;
98 }
99
100 vector<P> convexCut(const vector<P>&ps, L l) { // 用半平面 l 的逆时针方向去切凸多边形
101     vector<P> qs;
102     int n = ps.size();
103     for (int i = 0; i < n; ++i) {
104         Point p1 = ps[i], p2 = ps[(i + 1) % n];

```

```

105         int d1 = sgn(l.b * (p1 - l.a)), d2 = sign(l.b * (p2 - l.a));
106         if (d1 >= 0) qs.push_back(p1);
107         if (d1 * d2 < 0) qs.push_back(intersect(L(p1, p2 - p1), l));
108     }
109     return qs;
110 }

```

## 圆基础

```

1 struct C{
2     P o;
3     double r;
4     C(){}
5     C(P _o, double _r):o(_o),r(_r){}
6 };
7 // 求圆与直线的交点
8 //turn90() P(-y,x)
9 double fix(double x){return x>=0?x:0;}
10 bool intersect(C a, L l, P &p1, P &p2) {
11     double x = ((l.a - a.o)^(l.b - l.a)),
12     y = (l.b - l.a).len2(),
13     d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
14     if (sgn(d) < 0) return false;
15     d = max(d, 0.0);
16     P p = l.a - ((l.b - l.a) * (x / y)), delta = (l.b - l.a) * (sqrt(d) / y);
17     p1 = p + delta, p2 = p - delta;
18     return true;
19 }
20 // 求圆与圆的交点，注意调用前要先判定重圆
21 bool intersect(C a, C b, P &p1, P &p2) {
22     double s1 = (a.o - b.o).len();
23     if (sgn(s1 - a.r - b.r) > 0 || sgn(s1 - fabs(a.r - b.r)) < 0) return false;
24     double s2 = (a.r * a.r - b.r * b.r) / s1;
25     double aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
26     P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
27     P delta = (b.o - a.o).norm().turn90() * sqrt(fix(a.r * a.r - aa * aa));
28     p1 = o + delta, p2 = o - delta;
29     return true;
30 }
31 // 求点到圆的切点，按关于点的顺时针方向返回两个点
32 bool tang(const C &c, const P &p0, P &p1, P &p2) {
33     double x = (p0 - c.o).len2(), d = x - c.r * c.r;
34     if (d < eps) return false; // 点在圆上认为没有切点
35     P p = (p0 - c.o) * (c.r * c.r / x);
36     P delta = ((p0 - c.o) * (-c.r * sqrt(d) / x)).turn90();
37     p1 = c.o + p + delta;
38     p2 = c.o + p - delta;
39     return true;
40 }
41 // 求圆到圆的外公切线，按关于 c1.o 的顺时针方向返回两条线
42 vector<L> extan(const C &c1, const C &c2) {
43     vector<L> ret;
44     if (sgn(c1.r - c2.r) == 0) {
45         P dir = c2.o - c1.o;
46         dir = (dir * (c1.r / dir.len())).turn90();
47         ret.push_back(L(c1.o + dir, c2.o + dir));
48         ret.push_back(L(c1.o - dir, c2.o - dir));
49     } else {
50         P p = (c1.o * -c2.r + c2.o * c1.r) / (c1.r - c2.r);

```

```

51     P p1, p2, q1, q2;
52     if (tang(c1, p, p1, p2) && tang(c2, p, q1, q2)) {
53 //         if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
54         ret.push_back(L(p1, q1));
55         ret.push_back(L(p2, q2));
56     }
57 }
58 return ret;
59 }
60 // 求圆到圆的内共切线, 按关于 c1.o 的顺时针方向返回两条线
61 vector<L> intan(const C &c1, const C &c2) {
62     vector<L> ret;
63     P p = (c1.o * c2.r + c2.o * c1.r) / (c1.r + c2.r);
64     P p1, p2, q1, q2;
65     if (tang(c1, p, p1, p2) && tang(c2, p, q1, q2)) { // 两圆相切认为没有切线
66         ret.push_back(L(p1, q1));
67         ret.push_back(L(p2, q2));
68     }
69     return ret;
70 }

```

## 点在多边形内

```

1 bool inPoly(P p,vector<P>poly){
2     int cnt=0;
3     for(int i=0;i<poly.size();i++){
4         P a=poly[i],b=poly[(i+1)%poly.size()];
5         if(onSeg(p,L(a,b)))
6             return false;
7         int x=sgn(det(a,p,b));
8         int y=sgn(a.y-p.y);
9         int z=sgn(b.y-p.y);
10        cnt+=(x>0&&y<=0&&z>0);
11        cnt-=(x<0&&z<=0&&y>0);
12    }
13    return cnt;
14 }

```

## 二维最小覆盖圆

```

1 struct line{
2     point p,v;
3 };
4 point Rev(point v){return point(-v.y,v.x);}
5 point operator*(line A,line B){
6     point u=B.p-A.p;
7     double t=(B.v*u)/(B.v*A.v);
8     return A.p+A.v*t;
9 }
10 point get(point a,point b){
11     return (a+b)/2;
12 }
13 point get(point a,point b,point c){
14     if(a==b)return get(a,c);
15     if(a==c)return get(a,b);
16     if(b==c)return get(a,b);
17     line AB0=(line){(a+b)/2,Rev(a-b)};
18     line BC0=(line){(c+b)/2,Rev(b-c)};

```

```

19     return ABO*BCO;
20 }
21 int main(){
22     scanf("%d",&n);
23     for(int i=1;i<=n;i++)scanf("%lf%lf",&p[i].x,&p[i].y);
24     random_shuffle(p+1,p+1+n);
25     O=p[1];r=0;
26     for(int i=2;i<=n;i++){
27         if(dis(p[i],O)<r+1e-6)continue;
28         O=get(p[1],p[i]);r=dis(O,p[i]);
29         for(int j=1;j<i;j++){
30             if(dis(p[j],O)<r+1e-6)continue;
31             O=get(p[i],p[j]);r=dis(O,p[i]);
32             for(int k=1;k<j;k++){
33                 if(dis(p[k],O)<r+1e-6)continue;
34                 O=get(p[i],p[j],p[k]);r=dis(O,p[i]);
35             }
36         }
37     }printf("%.21f %.21f %.21f\n",O.x,O.y,r);
38     return 0;
39 }s

```

## 圆并

```

1 double ans[2001];
2 struct Point {
3     double x, y;
4     Point(){}
5     Point(const double & x, const double & y) : x(x), y(y) {}
6     void scan() {scanf("%lf%lf", &x, &y);}
7     double sqrlen() {return sqr(x) + sqr(y);}
8     double len() {return sqrt(sqrlen());}
9     Point rev() {return Point(y, -x);}
10    void print() {printf("%f %f\n", x, y);}
11    Point zoom(const double & d) {double lambda = d / len(); return Point(lambda * x, lambda
    ↪ * y);}
12 } dvd, a[2001];
13 Point centre[2001];
14 double atan2(const Point & x) {
15     return atan2(x.y, x.x);
16 }
17 Point operator - (const Point & a, const Point & b) {
18     return Point(a.x - b.x, a.y - b.y);
19 }
20 Point operator + (const Point & a, const Point & b) {
21     return Point(a.x + b.x, a.y + b.y);
22 }
23 double operator * (const Point & a, const Point & b) {
24     return a.x * b.y - a.y * b.x;
25 }
26 Point operator * (const double & a, const Point & b) {
27     return Point(a * b.x, a * b.y);
28 }
29 double operator % (const Point & a, const Point & b) {
30     return a.x * b.x + a.y * b.y;
31 }
32 struct circle {
33     double r; Point o;
34     circle() {}

```

```

35     void scan() {
36         o.scan();
37         scanf("%lf", &r);
38     }
39 } cir[2001];
40 struct arc {
41     double theta;
42     int delta;
43     Point p;
44     arc() {}
45     arc(const double & theta, const Point & p, int d) : theta(theta), p(p), delta(d) {}
46 } vec[4444];
47 int nV;
48 inline bool operator < (const arc & a, const arc & b) {
49     return a.theta + eps < b.theta;
50 }
51 int cnt;
52 inline void psh(const double t1, const Point p1, const double t2, const Point p2) {
53     if(t2 + eps < t1)
54         cnt++;
55     vec[nV++] = arc(t1, p1, 1);
56     vec[nV++] = arc(t2, p2, -1);
57 }
58 inline double cub(const double & x) {
59     return x * x * x;
60 }
61 inline void combine(int d, const double & area, const Point & o) {
62     if(sign(area) == 0) return;
63     centre[d] = 1 / (ans[d] + area) * (ans[d] * centre[d] + area * o);
64     ans[d] += area;
65 }
66 bool equal(const double & x, const double & y) {
67     return x + eps > y and y + eps > x;
68 }
69 bool equal(const Point & a, const Point & b) {
70     return equal(a.x, b.x) and equal(a.y, b.y);
71 }
72 bool equal(const circle & a, const circle & b) {
73     return equal(a.o, b.o) and equal(a.r, b.r);
74 }
75 bool f[2001];
76 int main() {
77     //freopen("hdu4895.in", "r", stdin);
78     int n, m, index;
79     while(EOF != scanf("%d%d%d", &m, &n, &index)) {
80         index--;
81         for(int i(0); i < m; i++) {
82             a[i].scan();
83         }
84         for(int i(0); i < n; i++) {
85             cir[i].scan(); //n 个圆
86         }
87         for(int i(0); i < n; i++) { //这一段在去重圆 能加速 删掉不会错
88             f[i] = true;
89             for(int j(0); j < n; j++) if(i != j) {
90                 if(equal(cir[i], cir[j]) and i < j or !equal(cir[i], cir[j]) and cir[i].r <
91                 ↪ cir[j].r + eps and (cir[i].o - cir[j].o).sqrln() < sqr(cir[i].r - cir[j].r) + eps) {
92                     f[i] = false;
93                     break;
94                 }
95             }
96         }
97     }
98 }

```

```

94     }
95 }
96 int n1(0);
97 for(int i(0); i < n; i++)
98     if(f[i])
99         cir[n1++] = cir[i];
100 n = n1;//去重圆结束
101 fill(ans, ans + n + 1, 0);//ans[i] 表示被圆覆盖至少 i 次的面积
102 fill(centre, centre + n + 1, Point(0, 0));//centre[i] 表示上面 ans[i] 部分的重心
103 for(int i(0); i < m; i++)
104     combine(0, a[i] * a[(i + 1) % m] * 0.5, 1. / 3 * (a[i] + a[(i + 1) % m]));
105 for(int i(0); i < n; i++) {
106     dvd = cir[i].o - Point(cir[i].r, 0);
107     nV = 0;
108     vec[nV++] = arc(-pi, dvd, 1);
109     cnt = 0;
110     for(int j(0); j < n; j++) if(j != i) {
111         double d = (cir[j].o - cir[i].o).sqrln();
112         if(d < sqr(cir[j].r - cir[i].r) + eps) {
113             if(cir[i].r + i * eps < cir[j].r + j * eps)
114                 psh(-pi, dvd, pi, dvd);
115             }else if(d + eps < sqr(cir[j].r + cir[i].r)) {
116                 double lambda = 0.5 * (1 + (sqr(cir[i].r) - sqr(cir[j].r)) / d);
117                 Point cp(cir[i].o + lambda * (cir[j].o - cir[i].o));
118                 Point nor((cir[j].o - cir[i].o).rev().zoom(sqrt(sqr(cir[i].r) - (cp -
119 ↪ cir[i].o).sqrln())));
119                 Point frm(cp + nor);
120                 Point to(cp - nor);
121                 psh(atan2(frm - cir[i].o), frm, atan2(to - cir[i].o), to);
122             }
123         }
124     sort(vec + 1, vec + nV);
125     vec[nV++] = arc(pi, dvd, -1);
126     for(int j = 0; j + 1 < nV; j++) {
127         cnt += vec[j].delta;
128         //if(cnt == 1) {//如果只算 ans[1] 和 centre[1], 可以加这个 if 加速.
129             double theta(vec[j + 1].theta - vec[j].theta);
130             double area(sqr(cir[i].r) * theta * 0.5);
131             combine(cnt, area, cir[i].o + 1. / area / 3 * cub(cir[i].r) *
132 ↪ Point(sin(vec[j + 1].theta) - sin(vec[j].theta), cos(vec[j].theta) - cos(vec[j +
133 ↪ 1].theta)));
134             combine(cnt, -sqr(cir[i].r) * sin(theta) * 0.5, 1. / 3 * (cir[i].o +
135 ↪ vec[j].p + vec[j + 1].p));
136             combine(cnt, vec[j].p * vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j +
137 ↪ 1].p));
138         }
139     }
140     }
141 }
142 }
143 }
144 }
145 fclose(stdin);
146 return 0;

```

147 }

---

## 经典阿波罗尼斯圆

```

1 硬币问题：易知两两相切的圆半径为 r1, r2, r3, 求与他们都相切的圆的半径 r4
2 分母取负号，答案再取绝对值，为外切圆半径
3 分母取正号为内切圆半径
4 //  $r_4^{\pm} = \frac{r_1 r_2 r_3}{r_1 r_2 + r_1 r_3 + r_2 r_3 \pm 2\sqrt{r_1 r_2 r_3 (r_1 + r_2 + r_3)}}$ 

```

## 半平面交

```

1 struct P{
2     int quad() const { return sgn(y) == 1 || (sgn(y) == 0 && sgn(x) >= 0);}
3 };
4 struct L{
5     bool onLeft(const P &p) const { return sgn((b - a)*(p - a)) > 0; }
6     L push() const{ // push out eps
7         const double eps = 1e-10;
8         P delta = (b - a).turn90().norm() * eps;
9         return L(a - delta, b - delta);
10    }
11 };
12 bool sameDir(const L &l0, const L &l1) {
13     return parallel(l0, l1) && sgn((l0.b - l0.a)^(l1.b - l1.a)) == 1;
14 }
15 bool operator < (const P &a, const P &b) {
16     if (a.quad() != b.quad())
17         return a.quad() < b.quad();
18     else
19         return sgn((a*b)) > 0;
20 }
21 bool operator < (const L &l0, const L &l1) {
22     if (sameDir(l0, l1))
23         return l1.onLeft(l0.a);
24     else
25         return (l0.b - l0.a) < (l1.b - l1.a);
26 }
27 bool check(const L &u, const L &v, const L &w) {
28     return w.onLeft(intersect(u, v));
29 }
30 vector<P> intersection(vector<L> &l) {
31     sort(l.begin(), l.end());
32     deque<L> q;
33     for (int i = 0; i < (int)l.size(); ++i) {
34         if (i && sameDir(l[i], l[i - 1])) {
35             continue;
36         }
37         while (q.size() > 1
38             && !check(q[q.size() - 2], q[q.size() - 1], l[i]))
39             q.pop_back();
40         while (q.size() > 1
41             && !check(q[1], q[0], l[i]))
42             q.pop_front();
43         q.push_back(l[i]);
44     }
45     while (q.size() > 2
46         && !check(q[q.size() - 2], q[q.size() - 1], q[0]))

```



```

47         q.pop_back();
48     while (q.size() > 2
49         && !check(q[1], q[0], q[q.size() - 1]))
50         q.pop_front();
51     vector<P> ret;
52     for (int i = 0; i < (int)q.size(); ++i)
53         ret.push_back(intersect(q[i], q[(i + 1) % q.size()]));
54     return ret;
55 }

```

## 求凸包

```

1 vector<P> convex(vector<P>p){
2     sort(p.begin(),p.end());
3     vector<P>ans,S;
4     for(int i=0;i<p.size();i++){
5         while(S.size())>=2
6             && sgn(det(S[S.size()-2],S.back(),p[i]))<=0)
7             S.pop_back();
8         S.push_back(p[i]);
9     }//dw
10    ans=S;
11    S.clear();
12    for(int i=(int)p.size()-1;i>=0;i--){
13        while(S.size())>=2
14            && sgn(det(S[S.size()-2],S.back(),p[i]))<=0)
15            S.pop_back();
16        S.push_back(p[i]);
17    }//up
18    for(int i=1;i+1<S.size();i++)
19        ans.push_back(S[i]);
20    return ans;
21 }

```

## 凸包游戏

```

1  /*
2  给定凸包， $\log n$  内完成各种询问，具体操作有：
3  1. 判定一个点是否在凸包内
4  2. 询问凸包外的点到凸包的两个切点
5  3. 询问一个向量关于凸包的切点
6  4. 询问一条直线和凸包的交点
7  INF 为坐标范围，需要定义点类大于号
8  改成实数只需修改 sign 函数，以及把 long long 改为 double 即可
9  构造函数时传入凸包要求无重点，面积非空，以及 pair(x,y) 的最小点放在第一个
10 */
11 const int INF = 1000000000;
12 struct Convex
13 {
14     int n;
15     vector<Point> a, upper, lower;
16     Convex(vector<Point> _a) : a(_a) {
17         n = a.size();
18         int ptr = 0;
19         for(int i = 1; i < n; ++ i) if (a[ptr] < a[i]) ptr = i;
20         for(int i = 0; i <= ptr; ++ i) lower.push_back(a[i]);
21         for(int i = ptr; i < n; ++ i) upper.push_back(a[i]);
22         upper.push_back(a[0]);

```

```

23 }
24 int sign(long long x) { return x < 0 ? -1 : x > 0; }
25 pair<long long, int> get_tangent(vector<Point> &convex, Point vec) {
26     int l = 0, r = (int)convex.size() - 2;
27     for( ; l + 1 < r; ) {
28         int mid = (l + r) / 2;
29         if (sign((convex[mid + 1] - convex[mid]).det(vec)) > 0) r = mid;
30         else l = mid;
31     }
32     return max(make_pair(vec.det(convex[r]), r),
33               , make_pair(vec.det(convex[0]), 0));
34 }
35 void update_tangent(const Point &p, int id, int &i0, int &i1) {
36     if ((a[i0] - p).det(a[id] - p) > 0) i0 = id;
37     if ((a[i1] - p).det(a[id] - p) < 0) i1 = id;
38 }
39 void binary_search(int l, int r, Point p, int &i0, int &i1) {
40     if (l == r) return;
41     update_tangent(p, l % n, i0, i1);
42     int sl = sign((a[l % n] - p).det(a[(l + 1) % n] - p));
43     for( ; l + 1 < r; ) {
44         int mid = (l + r) / 2;
45         int smid = sign((a[mid % n] - p).det(a[(mid + 1) % n] - p));
46         if (smid == sl) l = mid;
47         else r = mid;
48     }
49     update_tangent(p, r % n, i0, i1);
50 }
51 int binary_search(Point u, Point v, int l, int r) {
52     int sl = sign((v - u).det(a[l % n] - u));
53     for( ; l + 1 < r; ) {
54         int mid = (l + r) / 2;
55         int smid = sign((v - u).det(a[mid % n] - u));
56         if (smid == sl) l = mid;
57         else r = mid;
58     }
59     return l % n;
60 }
61 // 判定点是否在凸包内, 在边界返回 true
62 bool contain(Point p) {
63     if (p.x < lower[0].x || p.x > lower.back().x) return false;
64     int id = lower_bound(lower.begin(), lower.end(),
65                         , Point(p.x, -INF)) - lower.begin();
66     if (lower[id].x == p.x) {
67         if (lower[id].y > p.y) return false;
68     } else if ((lower[id - 1] - p).det(lower[id] - p) < 0) return false;
69     id = lower_bound(upper.begin(), upper.end(), Point(p.x, INF),
70                     , greater<Point>()) - upper.begin();
71     if (upper[id].x == p.x) {
72         if (upper[id].y < p.y) return false;
73     } else if ((upper[id - 1] - p).det(upper[id] - p) < 0) return false;
74     return true;
75 }
76 // 求点 p 关于凸包的两个切点, 如果在凸包外则有序返回编号
77 // 共线的多个切点返回任意一个, 否则返回 false
78 bool get_tangent(Point p, int &i0, int &i1) {
79     if (contain(p)) return false;
80     i0 = i1 = 0;
81     int id = lower_bound(lower.begin(), lower.end(), p) - lower.begin();
82     binary_search(0, id, p, i0, i1);

```

```

83     binary_search(id, (int)lower.size(), p, i0, i1);
84     id = lower_bound(upper.begin(), upper.end(), p
85         , greater<Point>()) - upper.begin();
86     binary_search((int)lower.size() - 1, (int)lower.size() - 1 + id, p, i0, i1);
87     binary_search((int)lower.size() - 1 + id
88         , (int)lower.size() - 1 + (int)upper.size(), p, i0, i1);
89     return true;
90 }
91 // 求凸包上和向量 vec 叉积最大的点, 返回编号, 共线的多个切点返回任意一个
92 int get_tangent(Point vec) {
93     pair<long long, int> ret = get_tangent(upper, vec);
94     ret.second = (ret.second + (int)lower.size() - 1) % n;
95     ret = max(ret, get_tangent(lower, vec));
96     return ret.second;
97 }
98 // 求凸包和直线 u,v 的交点, 如果无严格相交返回 false.
99 //如果有则是和 (i,next(i)) 的交点, 两个点无序, 交在点上不确定返回前后两条线段其中之一
100 bool get_intersection(Point u, Point v, int &i0, int &i1) {
101     int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
102     if (sign((v - u).det(a[p0] - u)) * sign((v - u).det(a[p1] - u)) < 0) {
103         if (p0 > p1) swap(p0, p1);
104         i0 = binary_search(u, v, p0, p1);
105         i1 = binary_search(u, v, p1, p0 + n);
106         return true;
107     } else {
108         return false;
109     }
110 }
111 };

```

## 平面最近点

```

1 bool byY(P a,P b){return a.y<b.y;}
2 LL solve(P *p,int l,int r){
3     LL d=1LL<<62;
4     if(l==r)
5         return d;
6     if(l+1==r)
7         return dis2(p[l],p[r]);
8     int mid=(l+r)>>1;
9     d=min(solve(l,mid),d);
10    d=min(solve(mid+1,r),d);
11    vector<P>tmp;
12    for(int i=l;i<=r;i++)
13        if(sqr(p[mid].x-p[i].x)<=d)
14            tmp.push_back(p[i]);
15    sort(tmp.begin(),tmp.end(),byY);
16    for(int i=0;i<tmp.size();i++)
17        for(int j=i+1;j<tmp.size()&&j-i<10;j++)
18            d=min(d,dis2(tmp[i],tmp[j]));
19    return d;
20 }

```

## 无敌面积并 (多圆多多边形), $n^3$

```

1 /*
2  n^3 计算多边形圆的面积并
3  注意先去重圆

```

```

4
5  */
6
7 double form(double x){
8     while(x>=2*pi)x-=2*pi;
9     while(x<0)x+=2*pi;
10    return x;
11 }
12 double calcCir(C cir){
13     vector<double>ang;
14     ang.push_back(0);
15     ang.push_back(pi);
16     double ans=0;
17     for(int i=1;i<=n;i++){
18         if(cir==c[i])continue;
19         P p1,p2;
20         if(intersect(cir,c[i],p1,p2)){
21             ang.push_back(form(cir.ang(p1)));
22             ang.push_back(form(cir.ang(p2)));
23         }
24     }
25
26     for(int i=1;i<=m;i++){
27         vector<P>tmp;
28         tmp=intersect(poly[i],cir);
29         for(int j=0;j<tmp.size();j++){
30             ang.push_back(form(cir.ang(tmp[j])));
31         }
32     }
33     sort(ang.begin(),ang.end());
34     for(int i=0;i<ang.size();i++){
35         double t1=ang[i],t2=(i+1==ang.size()?ang[0]+2*pi:ang[i+1]);
36         P p=cir.at((t1+t2)/2);
37         int ok=1;
38         for(int j=1;j<=n;j++){
39             if(cir==c[j])continue;
40             if(inC(p,c[j],true)){
41                 ok=0;
42                 break;
43             }
44         }
45         for(int j=1;j<=m&&ok;j++){
46             if(inPoly(p,poly[j],true)){
47                 ok=0;
48                 break;
49             }
50         }
51         if(ok){
52             double r=cir.r,x0=cir.o.x,y0=cir.o.y;
53             ans+=(r*r*(t2-t1)+r*x0*(sin(t2)-sin(t1))-r*y0*(cos(t2)-cos(t1)))/2;
54         }
55     }
56 }
57 return ans;
58 }
59 P st;
60 bool bySt(P a,P b){
61     return dis(a,st)<dis(b,st);
62 }
63 double calcSeg(L l){

```

```

64     double ans=0;
65     vector<P>pt;
66     pt.push_back(l.a);
67     pt.push_back(l.b);
68     for(int i=1;i<=n;i++){
69         P p1,p2;
70         if(intersect(c[i],l,p1,p2)){
71             if(onSeg(p1,l))
72                 pt.push_back(p1);
73             if(onSeg(p2,l))
74                 pt.push_back(p2);
75         }
76     }
77     st=l.a;
78     sort(pt.begin(),pt.end(),bySt);
79     for(int i=0;i+1<pt.size();i++){
80         P p1=pt[i],p2=pt[i+1];
81         P p=(p1+p2)/2;
82         int ok=1;
83         for(int j=1;j<=n;j++){
84             if(sgn(dis(p,c[j].o),c[j].r)<0){
85                 ok=0;
86                 break;
87             }
88         }
89         if(ok){
90             double x1=p1.x,y1=p1.y,x2=p2.x,y2=p2.y;
91             double res=(x1*y2-x2*y1)/2;
92             ans+=res;
93         }
94     }
95     return ans;
96 }

```

## Farmland

```

1  const int N = 11111, M = 111111 * 4;
2
3  struct eglist {
4      int other[M], succ[M], last[M], sum;
5      void clear() {
6          memset(last, -1, sizeof(last));
7          sum = 0;
8      }
9      void addEdge(int a, int b) {
10         other[sum] = b, succ[sum] = last[a], last[a] = sum++;
11         other[sum] = a, succ[sum] = last[b], last[b] = sum++;
12     }
13 }e;
14
15 int n, m;
16 struct point {
17     int x, y;
18     point(int x, int y) : x(x), y(y) {}
19     point() {}
20     friend point operator -(point a, point b) {
21         return point(a.x - b.x, a.y - b.y);
22     }
23     double arg() {

```

```

24         return atan2(y, x);
25     }
26 }points[N];
27
28 vector<pair<int, double> > vecs;
29 vector<int> ee[M];
30 vector<pair<double, pair<int, int> > > edges;
31 double length[M];
32 int tot, father[M], next[M], visit[M];
33
34 int find(int x) {
35     return father[x] == x ? x : father[x] = find(father[x]);
36 }
37
38 long long det(point a, point b) {
39     return 1LL * a.x * b.y - 1LL * b.x * a.y;
40 }
41
42 double dist(point a, point b) {
43     return sqrt(1.0 * (a.x - b.x) * (a.x - b.x) + 1.0 * (a.y - b.y) * (a.y - b.y));
44 }
45
46 int main() {
47     scanf("%d %d", &n, &m);
48     e.clear();
49     for(int i = 1; i <= n; i++) {
50         scanf("%d %d", &points[i].x, &points[i].y);
51     }
52     for(int i = 1; i <= m; i++) {
53         int a, b;
54         scanf("%d %d", &a, &b);
55         e.addEdge(a, b);
56     }
57     for(int x = 1; x <= n; x++) {
58         vector<pair<double, int> > pairs;
59         for(int i = e.last[x]; ~i; i = e.succ[i]) {
60             int y = e.other[i];
61             pairs.push_back(make_pair((points[y] - points[x]).arg(), i));
62         }
63         sort(pairs.begin(), pairs.end());
64         for(int i = 0; i < (int)pairs.size(); i++) {
65             next[pairs[(i + 1) % (int)pairs.size()].second ^ 1] = pairs[i].second;
66         }
67     }
68     memset(visit, 0, sizeof(visit));
69     tot = 0;
70     for(int start = 0; start < e.sum; start++) {
71         if (visit[start])
72             continue;
73         long long total = 0;
74         int now = start;
75         vecs.clear();
76         while(!visit[now]) {
77             visit[now] = 1;
78             total += det(points[e.other[now ^ 1]], points[e.other[now]]);
79             vecs.push_back(make_pair(now / 2, dist(points[e.other[now ^ 1]],
↪ points[e.other[now]])));
80             now = next[now];
81         }
82         if (now == start && total > 0) {

```

```

83         ++tot;
84         for(int i = 0; i < (int)vecs.size(); i++) {
85             ee[vecs[i].first].push_back(tot);
86         }
87     }
88 }
89
90 for(int i = 0; i < e.sum / 2; i++) {
91     int a = 0, b = 0;
92     if (ee[i].size() == 0)
93         continue;
94     else if (ee[i].size() == 1) {
95         a = ee[i][0];
96     } else if (ee[i].size() == 2) {
97         a = ee[i][0], b = ee[i][1];
98     }
99     edges.push_back(make_pair(dist(points[e.other[i * 2]], points[e.other[i * 2 + 1]]),
100 ↪ make_pair(a, b)));
101 }
102 sort(edges.begin(), edges.end());
103 for(int i = 0; i <= tot; i++)
104     father[i] = i;
105 double ans = 0;
106 for(int i = 0; i < (int)edges.size(); i++) {
107     int a = edges[i].second.first, b = edges[i].second.second;
108     double v = edges[i].first;
109     if (find(a) != find(b)) {
110         ans += v;
111         father[father[a]] = father[b];
112     }
113 }
114 printf("%.5f\n", ans);
115 }

```

## 三维基础

```

1 struct P {
2     double x, y, z;
3     P(){}
4     P(double _x, double _y, double _z):x(_x),y(_y),z(_z){}
5     double len2(){
6         return (x*x+y*y+z*z);
7     }
8     double len(){
9         return sqrt(x*x+y*y+z*z);
10    }
11 };
12 bool operator==(P a,P b){
13     return sgn(a.x-b.x)==0 && sgn(a.y-b.y)==0 && sgn(a.z-b.z)==0 ;
14 }
15 bool operator<(P a,P b){
16     return sgn(a.x-b.x) ? a.x<b.x :(sgn(a.y-b.y)?a.y<b.y :a.z<b.z);
17 }
18 P operator+(P a,P b){
19     return P(a.x+b.x,a.y+b.y,a.z+b.z);
20 }
21 P operator-(P a,P b){
22     return P(a.x-b.x,a.y-b.y,a.z-b.z);
23 }

```

```

24 P operator*(P a,double b){
25     return P(a.x*b,a.y*b,a.z*b);
26 }
27 P operator/(P a,double b){
28     return P(a.x/b,a.y/b,a.z/b);
29 }
30 P operator*(const P &a, const P &b) {
31     return P(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
32 }
33 double operator^(const P &a, const P &b) {
34     return a.x*b.x+a.y*b.y+a.z*b.z;
35 }
36
37 double dis(P a,P b){return (b-a).len();}
38 double dis2(P a,P b){return (b-a).len2();}
39
40 // 平面法向量 : 平面上两个向量叉积
41 // 点共平面 : 平面上一点与之的向量点积法向量为 0
42 // 点在线段 ( 直线 ) 上 : 共线且两边点积非正
43 // 点在三角形内 ( 不包含边界, 需再判断是与某条边共线 )
44 bool in_tri(const P &a, const P &b, const P &c, const P &p) {
45     return sgn(((a - b)*(a - c)).len() - ((p - a)*(p - b)).len() - ((p - b)*(p - c)).len() -
46         ↪ ((p - c)*(p - a)).len()) == 0;
47 }
48 // 共平面的两点是否在这平面上一条直线的同侧
49 bool sameSide(const P &a, const P &b, const P &p0, const P &p1) {
50     return sgn(((a - b)*(p0 - b)) ^ ((a - b)*(p1 - b))) > 0;
51 }
52 // 两点在平面同侧 : 点积法向量符号相同
53 // 两直线平行 / 垂直 : 同二维
54 // 平面平行 / 垂直 : 判断法向量
55 // 线面垂直 : 法向量和直线平行
56 // 判断空间线段是否相交 : 四点共面两线段不平行相互在异侧
57 // 线段和三角形是否相交 : 线段在三角形平面不同侧 三角形任意两点在线段和第三点组成的平面的不同侧
58 // 求空间直线交点
59 P intersect(const P &a0, const P &b0, const P &a1, const P &b1) {
60
61     double t = ((a0.x - a1.x) * (a1.y - b1.y) - (a0.y - a1.y) * (a1.x - b1.x)) / ((a0.x - b0.x)
62         ↪ * (a1.y - b1.y) - (a0.y - b0.y) * (a1.x - b1.x));
63
64     //double t = ((a0.x - a1.x) * (a1.y - b1.y) - (a0.y - a1.y) * (a1.x - b1.x)) / ((a0.x -
65     ↪ b0.x) * (a1.y - b1.y) - (a0.y - b0.y) * (a1.x - b1.x));
66     return a0 + (b0 - a0) * t;
67 }
68 // 求平面和直线的交点
69 P intersect(const P &a, const P &b, const P &c, const P &l0, const P &l1) {
70
71     P p = (b-a)*(c-a); // 平面法向量
72     double t = (p^(a-l0)) / (p^(l1-l0));
73     return l0 + (l1 - l0) * t;
74 //     P p = pVec(a, b, c); // 平面法向量
75 //     double t = (p.x * (a.x - l0.x) + p.y * (a.y - l0.y) + p.z * (a.z - l0.z)) / (p.x *
76     ↪ (l1.x - l0.x) + p.y * (l1.y - l0.y) + p.z * (l1.z - l0.z));
77 //     return l0 + (l1 - l0) * t;
78 }
79 // 求平面交线 : 取不平行的一条直线的一个交点, 以及法向量叉积得到直线方向
80 // 点到直线距离 : 叉积得到三角形的面积除以底边
81 // 点到平面距离 : 点积法向量
82 // 直线间距离 : 平行时随便取一点求距离, 否则叉积方向向量得到方向点积计算长度

```



80 // 直线夹角 : 点积 平面夹角 : 法向量点积

## 三维凸包

```

1  int mark[1005][1005], n, cnt;;
2  double mix(const P &a, const P &b, const P &c) {
3      return a^(b*c);
4  }
5  double area(int a, int b, int c) {
6      return ((info[b] - info[a])*(info[c] - info[a])).len();
7  }
8  double volume(int a, int b, int c, int d) {
9      return mix(info[b] - info[a], info[c] - info[a], info[d] - info[a]);
10 }
11 struct Face {
12     int a, b, c; Face() {}
13     Face(int a, int b, int c): a(a), b(b), c(c) {}
14     int &operator [](int k) {
15         if (k == 0) return a; if (k == 1) return b; return c;
16     }
17 };
18 vector <Face> face;
19 inline void insert(int a, int b, int c) {
20     face.push_back(Face(a, b, c));
21 }
22 void add(int v) {
23     vector <Face> tmp; int a, b, c; cnt++;
24     for (int i = 0; i < SIZE(face); i++) {
25         a = face[i][0]; b = face[i][1]; c = face[i][2];
26         if (sgn(volume(v, a, b, c)) < 0)
27             mark[a][b] = mark[b][a] = mark[b][c] = mark[c][b] = mark[c][a] = mark[a][c] = cnt;
28         else tmp.push_back(face[i]);
29     } face = tmp;
30     for (int i = 0; i < SIZE(tmp); i++) {
31         a = face[i][0]; b = face[i][1]; c = face[i][2];
32         if (mark[a][b] == cnt) insert(b, a, v);
33         if (mark[b][c] == cnt) insert(c, b, v);
34         if (mark[c][a] == cnt) insert(a, c, v);
35     }
36 }
37 int Find() {
38     for (int i = 2; i < n; i++) {
39         P ndir = (info[0] - info[i])*(info[1] - info[i]);
40         if (ndir == P()) continue; swap(info[i], info[2]);
41         for (int j = i + 1; j < n; j++) if (sgn(volume(0, 1, 2, j)) != 0) {
42             swap(info[j], info[3]); insert(0, 1, 2); insert(0, 2, 1); return 1;
43         }
44     }
45     return 0;
46 }
47
48 // 求重心
49 double calcDist(const P &p, int a, int b, int c) {
50     return fabs(mix(info[a] - p, info[b] - p, info[c] - p) / area(a, b, c));
51 }
52 //compute the minimal distance of center of any faces
53 P findCenter() { //compute center of mass
54     double totalWeight = 0;
55     P center(.0, .0, .0);

```

```

56     P first = info[face[0][0]];
57     for (int i = 0; i < SIZE(face); ++i) {
58         P p = (info[face[i][0]]+info[face[i][1]]+info[face[i][2]]+first)*.25;
59         double weight = mix(info[face[i][0]] - first, info[face[i][1]] - first,
↪ info[face[i][2]] - first);
60         totalWeight += weight; center = center + p * weight;
61     }
62     center = center / totalWeight;
63     return center;
64 }
65 double minDis(P p) {
66     double res = 1e100; //compute distance
67     for (int i = 0; i < SIZE(face); ++i)
68         res = min(res, calcDist(p, face[i][0], face[i][1], face[i][2]));
69     return res;
70 }
71
72 void findConvex(P *info,int n) {
73     sort(info, info + n); n = unique(info, info + n) - info;
74     face.clear(); random_shuffle(info, info + n);
75     if(!Find())return abort();
76     memset(mark, 0, sizeof(mark)); cnt = 0;
77     for (int i = 3; i < n; i++) add(i);
78 }
79 // 三维绕轴旋转, 大拇指指向 axis 向量方向, 四指弯曲方向转 w 弧度
80 P rotate(const P& s, const P& axis, double w) {
81     double x = axis.x, y = axis.y, z = axis.z;
82     double s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
83     cosw = cos(w), sinw = sin(w);
84     double a[4][4];
85     memset(a, 0, sizeof a);
86     a[3][3] = 1;
87     a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
88     a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
89     a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
90     a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
91     a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
92     a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
93     a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
94     a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
95     a[2][2] = ((x * x + y * y) * cosw + z * z) / s1;
96     double ans[4] = {0, 0, 0, 0}, c[4] = {s.x, s.y, s.z, 1};
97     for (int i = 0; i < 4; ++ i)
98         for (int j = 0; j < 4; ++ j)
99             ans[i] += a[j][i] * c[j];
100     return P(ans[0], ans[1], ans[2]);
101 }

```

## 三角剖分与 V 图

```

1  /*
2  Delaunay Triangulation 随机增量算法 :
3  节点数至少为点数的 6 倍, 空间消耗较大注意计算内存使用
4  建图的过程在 build 中, 注意初始化内存池和初始三角形的坐标范围 (Triangulation::LOTS)
5  Triangulation::find 返回包含某点的三角形
6  Triangulation::add_point 将某点加入三角剖分
7  某个 Triangle 在三角剖分中当且仅当它的 has_children 为 0
8  如果要找到三角形 u 的邻域, 则枚举它的所有 u.edge[i].tri, 该条边的两个点为 u.p[(i+1)%3],
↪ u.p[(i+2)%3]

```

```

9 通过三角剖分构造 v 图：连接相邻三角形外接圆圆心即可
10 复杂度好像是  $O(n \log n)$ 
11 */
12 const int N = 100000 + 5, MAX_TRIS = N * 6;
13 const double eps = 1e-6, PI = acos(-1.0);
14 struct P {
15     double x,y; P():x(0),y(0){}
16     P(double x, double y):x(x),y(y){}
17     bool operator==(P const& that) const {return x==that.x&&y==that.y;}
18 };
19 inline double sqr(double x) { return x*x; }
20 double dist_sqr(P const& a, P const& b){return sqr(a.x-b.x)+sqr(a.y-b.y);}
21 bool in_circumcircle(P const& p1, P const& p2, P const& p3, P const& p4) { //p4 in C(p1,p2,p3)
22     double u11 = p1.x - p4.x, u21 = p2.x - p4.x, u31 = p3.x - p4.x;
23     double u12 = p1.y - p4.y, u22 = p2.y - p4.y, u32 = p3.y - p4.y;
24     double u13 = sqr(p1.x) - sqr(p4.x) + sqr(p1.y) - sqr(p4.y);
25     double u23 = sqr(p2.x) - sqr(p4.x) + sqr(p2.y) - sqr(p4.y);
26     double u33 = sqr(p3.x) - sqr(p4.x) + sqr(p3.y) - sqr(p4.y);
27     double det = -u13*u22*u31 + u12*u23*u31 + u13*u21*u32 - u11*u23*u32 - u12*u21*u33 +
    ↪ u11*u22*u33;
28     return det > eps;
29 }
30 double side(P const& a, P const& b, P const& p) { return (b.x-a.x)*(p.y-a.y) -
    ↪ (b.y-a.y)*(p.x-a.x); }
31 typedef int SideRef; struct Triangle; typedef Triangle* TriangleRef;
32 struct Edge {
33     TriangleRef tri; SideRef side; Edge() : tri(0), side(0) {}
34     Edge(TriangleRef tri, SideRef side) : tri(tri), side(side) {}
35 };
36 struct Triangle {
37     P p[3]; Edge edge[3]; TriangleRef children[3]; Triangle() {}
38     Triangle(P const& p0, P const& p1, P const& p2) {
39         p[0] = p0; p[1] = p1; p[2] = p2;
40         children[0] = children[1] = children[2] = 0;
41     }
42     bool has_children() const { return children[0] != 0; }
43     int num_children() const {
44         return children[0] == 0 ? 0
45             : children[1] == 0 ? 1
46             : children[2] == 0 ? 2 : 3;
47     }
48     bool contains(P const& q) const {
49         double a=side(p[0],p[1],q), b=side(p[1],p[2],q), c=side(p[2],p[0],q);
50         return a >= -eps && b >= -eps && c >= -eps;
51     }
52 } triange_pool[MAX_TRIS], *tot_triangles;
53 void set_edge(Edge a, Edge b) {
54     if (a.tri) a.tri->edge[a.side] = b;
55     if (b.tri) b.tri->edge[b.side] = a;
56 }
57 class Triangulation {
58 public:
59     Triangulation() {
60         const double LOTS = 1e6; //初始为极大三角形
61         the_root = new(tot_triangles++)
    ↪ Triangle(P(-LOTS,-LOTS),P(+LOTS,-LOTS),P(0,+LOTS));
62     }
63     TriangleRef find(P p) const { return find(the_root,p); }
64     void add_point(P const& p) { add_point(find(the_root,p),p); }
65 private:

```

```

66     TriangleRef the_root;
67     static TriangleRef find(TriangleRef root, P const& p) {
68         for( ; ; ) {
69             if (!root->has_children()) return root;
70             else for (int i = 0; i < 3 && root->children[i] ; ++i)
71                 if (root->children[i]->contains(p))
72                     {root = root->children[i]; break;}
73         }
74     }
75     void add_point(TriangleRef root, P const& p) {
76         TriangleRef tab,tbc,tca;
77         tab = new(tot_triangles++) Triangle(root->p[0], root->p[1], p);
78         tbc = new(tot_triangles++) Triangle(root->p[1], root->p[2], p);
79         tca = new(tot_triangles++) Triangle(root->p[2], root->p[0], p);
80         set_edge(Edge(tab,0),Edge(tbc,1)); set_edge(Edge(tbc,0),Edge(tca,1));
81         set_edge(Edge(tca,0),Edge(tab,1)); set_edge(Edge(tab,2),root->edge[2]);
82         set_edge(Edge(tbc,2),root->edge[0]); set_edge(Edge(tca,2),root->edge[1]);
83         root->children[0]=tab; root->children[1]=tbc; root->children[2]=tca;
84         flip(tab,2); flip(tbc,2); flip(tca,2);
85     }
86     void flip(TriangleRef tri, SideRef pi) {
87         TriangleRef trj = tri->edge[pi].tri; int pj = tri->edge[pi].side;
88         if(!trj || !in_circumcircle(tri->p[0],tri->p[1],tri->p[2],trj->p[pj])) return;
89         TriangleRef trk = new(tot_triangles++) Triangle(tri->p[(pi+1)%3], trj->p[pj],
90         ↪ tri->p[pi]);
91         TriangleRef trl = new(tot_triangles++) Triangle(trj->p[(pj+1)%3], tri->p[pi],
92         ↪ trj->p[pj]);
93         set_edge(Edge(trk,0), Edge(trl,0));
94         set_edge(Edge(trk,1), tri->edge[(pi+2)%3]); set_edge(Edge(trk,2),
95         ↪ trj->edge[(pj+1)%3]);
96         set_edge(Edge(trl,1), trj->edge[(pj+2)%3]); set_edge(Edge(trl,2),
97         ↪ tri->edge[(pi+1)%3]);
98         tri->children[0]=trk; tri->children[1]=trl; tri->children[2]=0;
99         trj->children[0]=trk; trj->children[1]=trl; trj->children[2]=0;
100        flip(trk,1); flip(trk,2); flip(trl,1); flip(trl,2);
101    }
102};
103int n; P ps[N];
104void build(){
105    tot_triangles = triange_pool; cin >> n;
106    for(int i = 0; i < n; ++ i) scanf("%lf%lf",&ps[i].x,&ps[i].y);
107    random_shuffle(ps, ps + n); Triangulation tri;
108    for(int i = 0; i < n; ++ i) tri.add_point(ps[i]);
109}

```

## 三维最小覆盖球

```

1 bool equal(const double & x, const double & y) {
2     return x + eps > y and y + eps > x;
3 }
4 double operator % (const Point & a, const Point & b) {
5     return a.x * b.x + a.y * b.y + a.z * b.z;
6 }
7 Point operator * (const Point & a, const Point & b) {
8     return Point(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
9 }
10 struct Circle {
11     double r; Point o;
12 };

```

```

13 struct Plane {
14     Point nor;
15     double m;
16     Plane(const Point & nor, const Point & a) : nor(nor){
17         m = nor % a;
18     }
19 };
20 Point intersect(const Plane & a, const Plane & b, const Plane & c) {
21     Point c1(a.nor.x, b.nor.x, c.nor.x), c2(a.nor.y, b.nor.y, c.nor.y), c3(a.nor.z,
    ↪ b.nor.z, c.nor.z), c4(a.m, b.m, c.m);
22     return 1 / ((c1 * c2) % c3) * Point((c4 * c2) % c3, (c1 * c4) % c3, (c1 * c2) % c4);
23 }
24 bool in(const Point & a, const Circle & b) {
25     return sign((a - b.o).len() - b.r) <= 0;
26 }
27 bool operator < (const Point & a, const Point & b) {
28     if(!equal(a.x, b.x)) {
29         return a.x < b.x;
30     }
31     if(!equal(a.y, b.y)) {
32         return a.y < b.y;
33     }
34     if(!equal(a.z, b.z)) {
35         return a.z < b.z;
36     }
37     return false;
38 }
39 bool operator == (const Point & a, const Point & b) {
40     return equal(a.x, b.x) and equal(a.y, b.y) and equal(a.z, b.z);
41 }
42 vector<Point> vec;
43 Circle calc() {
44     if(vec.empty()) {
45         return Circle(Point(0, 0, 0), 0);
46     } else if(1 == (int)vec.size()) {
47         return Circle(vec[0], 0);
48     } else if(2 == (int)vec.size()) {
49         return Circle(0.5 * (vec[0] + vec[1]), 0.5 * (vec[0] - vec[1]).len());
50     } else if(3 == (int)vec.size()) {
51         double r((vec[0] - vec[1]).len() * (vec[1] - vec[2]).len() * (vec[2] - vec[0]).len()
    ↪ / 2 / fabs(((vec[0] - vec[2]) * (vec[1] - vec[2])).len()));
52         return Circle(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
53                                 Plane(vec[2] - vec[1], 0.5 * (vec[2] + vec[1])),
54                                 Plane((vec[1] - vec[0]) * (vec[2] - vec[0]), vec[0])), r);
55     } else {
56         Point o(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
57                           Plane(vec[2] - vec[0], 0.5 * (vec[2] + vec[0])),
58                           Plane(vec[3] - vec[0], 0.5 * (vec[3] + vec[0]))));
59         return Circle(o, (o - vec[0]).len());
60     }
61 }
62 Circle miniBall(int n) {
63     Circle res(calc());
64     for(int i(0); i < n; i++) {
65         if(!in(a[i], res)) {
66             vec.push_back(a[i]);
67             res = miniBall(i);
68             vec.pop_back();
69             if(i) {
70                 Point tmp(a[i]);

```

```

71         memmove(a + 1, a, sizeof(Point) * i);
72         a[0] = tmp;
73     }
74 }
75 }
76 return res;
77 }
78 int main() {
79     int n;
80     sort(a, a + n);
81     n = unique(a, a + n) - a;
82     vec.clear();
83     printf("%.10f\n", miniBall(n).r);
84 }

```

## 空间四点外接球

```

1 // 注意, 无法处理小于四点的退化情况
2 pair<P, double> ball(P outer[4]) {
3     P res; double radius;
4     P q[3]; double m[3][3], sol[3], L[3], det;
5     int i, j; res.x = res.y = res.z = radius = 0;
6     for (i=0; i<3; ++i) q[i]=outer[i+1]-outer[0], sol[i]=(q[i]^q[i]);
7     for (i=0; i<3; ++i) for (j=0; j<3; ++j) m[i][j]=(q[i]^q[j])*2;
8     det= m[0][0]*m[1][1]*m[2][2]
9     + m[0][1]*m[1][2]*m[2][0]
10    + m[0][2]*m[2][1]*m[1][0]
11    - m[0][2]*m[1][1]*m[2][0]
12    - m[0][1]*m[1][0]*m[2][2]
13    - m[0][0]*m[1][2]*m[2][1];
14    if ( fabs(det)<1e-10) return;
15    for (j=0; j<3; ++j) {
16        for (i=0; i<3; ++i) m[i][j]=sol[i];
17        L[j]=( m[0][0]*m[1][1]*m[2][2]
18        + m[0][1]*m[1][2]*m[2][0]
19        + m[0][2]*m[2][1]*m[1][0]
20        - m[0][2]*m[1][1]*m[2][0]
21        - m[0][1]*m[1][0]*m[2][2]
22        - m[0][0]*m[1][2]*m[2][1]
23        ) / det;
24        for (i=0; i<3; ++i) m[i][j]=(q[i]^q[j])*2;
25    }
26    res=outer[0];
27    for (i=0; i<3; ++i ) res = res + q[i] * L[i];
28    radius=dis(res, outer[0]);
29    return make_pair(res, radius);
30 }

```

# Chapter 7

## 技巧

### 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
2 // 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
4 namespace Reader {
5     const int L = (1 << 15) + 5;
6     char buffer[L], *S, *T;
7     __inline bool getchar(char &ch) {
8         if (S == T) {
9             T = (S = buffer) + fread(buffer, 1, L, stdin);
10            if (S == T) {
11                ch = EOF;
12                return false;
13            }
14        }
15        ch = *S++;
16        return true;
17    }
18    __inline bool getint(int &x) {
19        char ch; bool neg = 0;
20        for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
21        if (ch == EOF) return false;
22        x = ch - '0';
23        for (; getchar(ch), ch >= '0' && ch <= '9'; )
24            x = x * 10 + ch - '0';
25        if (neg) x = -x;
26        return true;
27    }
28 }
```

### 真正释放 STL 内存

```
1 template <typename T>
2 __inline void clear(T& container) {
3     container.clear(); // 或者删除了一堆元素
4     T(container).swap(container);
5 }
```

### 梅森旋转算法

```
1 #include <random>
2
3 int main() {
```

```

4      std::mt19937 g(seed); // std::mt19937_64
5      std::cout << g() << std::endl;
6  }

```

## 蔡勒公式

```

1  int solve(int year, int month, int day) {
2      int answer;
3      if (month == 1 || month == 2) {
4          month += 12;
5          year--;
6      }
7      if ((year < 1752) || (year == 1752 && month < 9) ||
8          (year == 1752 && month == 9 && day < 3)) {
9          answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4 + 5) % 7;
10     } else {
11         answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4
12                 - year / 100 + year / 400) % 7;
13     }
14     return answer;
15 }

```

## 开栈

```

1  register char *_sp __asm__("rsp");
2  int main() {
3      const int size = 400 << 20; // 400MB
4      static char *sys, *mine(new char[size] + size - 4096);
5      sys = _sp; _sp = mine; _main(); _sp = sys;
6  }

```

## Size 为 k 的子集

```

1  void solve(int n, int k) {
2      for (int comb = (1 << k) - 1; comb < (1 << n); ) {
3          // ...
4          int x = comb & -comb, y = comb + x;
5          comb = (((comb & ~y) / x) >> 1) | y;
6      }
7  }

```

## 长方体表面两点最短距离

```

1  int r;
2  void turn(int i, int j, int x, int y, int z, int x0, int y0, int L, int W, int H) {
3      if (z==0) { int R = x*x+y*y; if (R<r) r=R;
4      } else {
5          if(i>=0 && i< 2) turn(i+1, j, x0+L+z, y, x0+L-x, x0+L, y0, H, W, L);
6          if(j>=0 && j< 2) turn(i, j+1, x, y0+W+z, y0+W-y, x0, y0+W, L, H, W);
7          if(i<=0 && i>-2) turn(i-1, j, x0-z, y, x-x0, x0-H, y0, H, W, L);
8          if(j<=0 && j>-2) turn(i, j-1, x, y0-z, y-y0, x0, y0-H, L, H, W);
9      }
10 }
11 int main(){
12     int L, H, W, x1, y1, z1, x2, y2, z2;
13     cin >> L >> W >> H >> x1 >> y1 >> z1 >> x2 >> y2 >> z2;

```



```
14     if (z1!=0 && z1!=H) if (y1==0 || y1==W)
15         swap(y1,z1), std::swap(y2,z2), std::swap(W,H);
16     else swap(x1,z1), std::swap(x2,z2), std::swap(L,H);
17     if (z1==H) z1=0, z2=H-z2;
18     r=0x3fffffff;
19     turn(0,0,x2-x1,y2-y1,z2,-x1,-y1,L,W,H);
20     cout<<r<<endl;
21 }
```

经纬度求球面最短距离

```
1 double sphereDis(double lon1, double lat1, double lon2, double lat2, double R) {
2     return R * acos(cos(lat1) * cos(lat2) * cos(lon1 - lon2) + sin(lat1) * sin(lat2));
3 }
```

32-bit/64-bit 随机素数

32-bit	64-bit
73550053	1249292846855685773
148898719	1701750434419805569
189560747	3605499878424114901
459874703	5648316673387803781
1202316001	6125342570814357977
1431183547	6215155308775851301
1438011109	6294606778040623451
1538762023	6347330550446020547
1557944263	7429632924303725207
1981315913	8524720079480389849

NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3

Formulas

Arithmetic Function

$$\sigma_k(n) = \sum_{d|n} d^k = \prod_{i=1}^{\omega(n)} \frac{p_i^{(a_i+1)k} - 1}{p_i^k - 1}$$
$$J_k(n) = n^k \prod_{p|n} (1 - \frac{1}{p^k})$$

$J_k(n)$  is the number of  $k$ -tuples of positive integers all less than or equal to  $n$  that form a coprime  $(k + 1)$ -tuple together with  $n$ .

$$\sum_{\delta|n} J_k(\delta) = n^k$$
$$\sum_{\delta|n} \delta^s J_r(\delta) J_s(\frac{n}{\delta}) = J_{r+s}(n)$$

$$\begin{aligned}
\sum_{\delta|n} \varphi(\delta) d\left(\frac{n}{\delta}\right) &= \sigma(n), \quad \sum_{\delta|n} |\mu(\delta)| = 2^{\omega(n)} \\
\sum_{\delta|n} 2^{\omega(\delta)} &= d(n^2), \quad \sum_{\delta|n} d(\delta^2) = d^2(n) \\
\sum_{\delta|n} d\left(\frac{n}{\delta}\right) 2^{\omega(\delta)} &= d^2(n), \quad \sum_{\delta|n} \frac{\mu(\delta)}{\delta} = \frac{\varphi(n)}{n} \\
\sum_{\delta|n} \frac{\mu(\delta)}{\varphi(\delta)} &= d(n), \quad \sum_{\delta|n} \frac{\mu^2(\delta)}{\varphi(\delta)} = \frac{n}{\varphi(n)}
\end{aligned}$$

$$\begin{aligned}
&n|\varphi(a^n - 1) \\
&\sum_{\substack{1 \leq k \leq n \\ \gcd(k, n)=1}} f(\gcd(k-1, n)) = \varphi(n) \sum_{d|n} \frac{(\mu * f)(d)}{\varphi(d)} \\
&\varphi(\text{lcm}(m, n)) \varphi(\gcd(m, n)) = \varphi(m) \varphi(n) \\
&\sum_{\delta|n} d^3(\delta) = \left(\sum_{\delta|n} d(\delta)\right)^2 \\
&d(uv) = \sum_{\delta | \gcd(u, v)} \mu(\delta) d\left(\frac{u}{\delta}\right) d\left(\frac{v}{\delta}\right) \\
&\sigma_k(u) \sigma_k(v) = \sum_{\delta | \gcd(u, v)} \delta^k \sigma_k\left(\frac{uv}{\delta^2}\right) \\
&\mu(n) = \sum_{k=1}^n [\gcd(k, n) = 1] \cos 2\pi \frac{k}{n} \\
&\varphi(n) = \sum_{k=1}^n [\gcd(k, n) = 1] = \sum_{k=1}^n \gcd(k, n) \cos 2\pi \frac{k}{n} \\
&\begin{cases} S(n) = \sum_{k=1}^n (f * g)(k) \\ \sum_{k=1}^n S(\lfloor \frac{n}{k} \rfloor) = \sum_{i=1}^n f(i) \sum_{j=1}^{\lfloor \frac{n}{i} \rfloor} (g * 1)(j) \end{cases} \\
&\begin{cases} S(n) = \sum_{k=1}^n (f \cdot g)(k), g \text{ completely multiplicative} \\ \sum_{k=1}^n S(\lfloor \frac{n}{k} \rfloor) g(k) = \sum_{k=1}^n (f * 1)(k) g(k) \end{cases}
\end{aligned}$$

## Binomial Coefficients

$$\begin{aligned}
\binom{n}{k} &= (-1)^k \binom{k-n-1}{k} \\
\sum_{k \leq n} \binom{r+k}{k} &= \binom{r+n+1}{n} \\
\sum_{k=0}^n \binom{k}{m} &= \binom{n+1}{m+1} \\
\sqrt{1+z} &= 1 + \sum_{k=1}^{\infty} \frac{(-1)^{k-1}}{k \times 2^{2k-1}} \binom{2k-2}{k-1} z^k \\
\sum_{k=0}^r \binom{r-k}{m} \binom{s+k}{n} &= \binom{r+s+1}{m+n+1} \\
C_{n,m} &= \binom{n+m}{m} - \binom{n+m}{m-1}, n \geq m \\
\binom{n}{k} &\equiv [n \& k = k] \pmod{2}
\end{aligned}$$

## Fibonacci Numbers

$$F(z) = \frac{z}{1 - z - z^2}$$

$$f_n = \frac{\phi^n - \hat{\phi}^n}{\sqrt{5}}, \phi = \frac{1 + \sqrt{5}}{2}, \hat{\phi} = \frac{1 - \sqrt{5}}{2}$$

$$\sum_{k=1}^n f_k = f_{n+2} - 1$$

$$\sum_{k=1}^n f_k^2 = f_n f_{n+1}$$

$$\sum_{k=0}^n f_k f_{n-k} = \frac{1}{5}(n-1)f_n + \frac{2}{5}n f_{n-1}$$

$$f_n^2 + (-1)^n = f_{n+1} f_{n-1}$$

$$f_{n+k} = f_n f_{k+1} + f_{n-1} f_k$$

$$f_{2n+1} = f_n^2 + f_{n+1}^2$$

$$(-1)^k f_{n-k} = f_n f_{k-1} - f_{n-1} f_k$$

$$\text{Modulo } f_n, f_{mn+r} \equiv \begin{cases} f_r, & m \bmod 4 = 0; \\ (-1)^{r+1} f_{n-r}, & m \bmod 4 = 1; \\ (-1)^n f_r, & m \bmod 4 = 2; \\ (-1)^{r+1+n} f_{n-r}, & m \bmod 4 = 3. \end{cases}$$

## Stirling Cycle Numbers

$$\begin{bmatrix} n+1 \\ k \end{bmatrix} = n \begin{bmatrix} n \\ k \end{bmatrix} + \begin{bmatrix} n \\ k-1 \end{bmatrix}, \quad \begin{bmatrix} n+1 \\ 2 \end{bmatrix} = n! H_n$$

$$x^n = \sum_k \begin{bmatrix} n \\ k \end{bmatrix} (-1)^{n-k} x^k, \quad x^{\bar{n}} = \sum_k \begin{bmatrix} n \\ k \end{bmatrix} x^k$$

## Stirling Subset Numbers

$$\left\{ \begin{matrix} n+1 \\ k \end{matrix} \right\} = k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n \\ k-1 \end{matrix} \right\}$$

$$x^n = \sum_k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} x^k = \sum_k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} (-1)^{n-k} x^{\bar{k}}$$

$$m! \left\{ \begin{matrix} n \\ m \end{matrix} \right\} = \sum_k \binom{m}{k} k^n (-1)^{m-k}$$

## Eulerian Numbers

$$\left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle = (k+1) \left\langle \begin{matrix} n-1 \\ k \end{matrix} \right\rangle + (n-k) \left\langle \begin{matrix} n-1 \\ k-1 \end{matrix} \right\rangle$$

$$x^n = \sum_k \left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle \binom{x+k}{n}$$

$$\left\langle \begin{matrix} n \\ m \end{matrix} \right\rangle = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k$$

## Harmonic Numbers

$$\sum_{k=1}^n H_k = (n+1)H_n - n$$

$$\sum_{k=1}^n kH_k = \frac{n(n+1)}{2}H_n - \frac{n(n-1)}{4}$$

$$\sum_{k=1}^n \binom{k}{m} H_k = \binom{n+1}{m+1} \left( H_{n+1} - \frac{1}{m+1} \right)$$

## Pentagonal Number Theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{n=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2}$$

$$p(n) = p(n-1) + p(n-2) - p(n-5) - p(n-7) + \cdots$$

$$f(n, k) = p(n) - p(n-k) - p(n-2k) + p(n-5k) + p(n-7k) - \cdots$$

## Bell Numbers

$$B_{n+1} = \sum_{k=0}^n \binom{n}{k} B_k$$

$$B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$$

## Bernoulli Numbers

$$B_n = 1 - \sum_{k=0}^{n-1} \binom{n}{k} \frac{B_k}{n-k+1}$$

$$G(x) = \sum_{k=0}^{\infty} \frac{B_k}{k!} x^k = \frac{1}{\sum_{k=0}^{\infty} \frac{x^k}{(k+1)!}}$$

$$S_m(n) = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m-k+1}$$

## Tetrahedron Volume

$$V = \frac{\sqrt{4u^2v^2w^2 - \sum_{cyc} u^2(v^2 + w^2 - U^2)^2 + \prod_{cyc} (v^2 + w^2 - U^2)}}{12}$$

## BEST Thoerem

Counting the number of different Eulerian circuits in directed graphs.

$$\text{ec}(G) = t_w(G) \prod_{v \in V} (\deg(v) - 1)!$$

When calculating  $t_w(G)$  for directed multigraphs, the entry  $q_{i,j}$  for distinct  $i$  and  $j$  equals  $-m$ , where  $m$  is the number of edges from  $i$  to  $j$ , and the entry  $q_{i,i}$  equals the indegree of  $i$  minus the number of loops at  $i$ . It is a property of Eulerian graphs that  $\text{tv}(G) = \text{tw}(G)$  for every two vertices  $v$  and  $w$  in a connected Eulerian graph  $G$ .

## 重心

半径为  $r$ , 圆心角为  $\theta$  的扇形重心与圆心的距离为  $\frac{4r \sin(\theta/2)}{3\theta}$   
 半径为  $r$ , 圆心角为  $\theta$  的圆弧重心与圆心的距离为  $\frac{4r \sin^3(\theta/2)}{3(\theta - \sin(\theta))}$

## Others

$$S_j = \sum_{k=1}^n x_k^j$$

$$h_m = \sum_{1 \leq j_1 < \cdots < j_m \leq n} x_{j_1} \cdots x_{j_m}$$

$$H_m = \sum_{1 \leq j_1 \leq \dots \leq j_m \leq n} x_{j_1} \cdots x_{j_m}$$

$$h_n = \frac{1}{n} \sum_{k=1}^n (-1)^{k+1} S_k h_{n-k}$$

$$H_n = \frac{1}{n} \sum_{k=1}^n S_k H_{n-k}$$

$$\sum_{k=0}^n k c^k = \frac{n c^{n+2} - (n+1) c^{n+1} + c}{(c-1)^2}$$

$$n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \left(1 + \frac{1}{12n} + \frac{1}{288n^2} + O\left(\frac{1}{n^3}\right)\right)$$

$$\begin{aligned} & \max\{x_a - x_b, y_a - y_b, z_a - z_b\} - \min\{x_a - x_b, y_a - y_b, z_a - z_b\} \\ &= \frac{1}{2} \sum_{cyc} |(x_a - y_a) - (x_b - y_b)| \end{aligned}$$

$$(a+b)(b+c)(c+a) = \frac{(a+b+c)^3 - a^3 - b^3 - c^3}{3}$$

### Integrals of Rational Functions

$$\int \frac{1}{1+x^2} dx = \tan^{-1} x \quad (1)$$

$$\int \frac{1}{a^2+x^2} dx = \frac{1}{a} \tan^{-1} \frac{x}{a} \quad (2)$$

$$\int \frac{x}{a^2+x^2} dx = \frac{1}{2} \ln|a^2+x^2| \quad (3)$$

$$\int \frac{x^2}{a^2+x^2} dx = x - a \tan^{-1} \frac{x}{a} \quad (4)$$

$$\int \frac{x^3}{a^2+x^2} dx = \frac{1}{2}x^2 - \frac{1}{2}a^2 \ln|a^2+x^2| \quad (5)$$

$$\int \frac{1}{ax^2+bx+c} dx = \frac{2}{\sqrt{4ac-b^2}} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}} \quad (6)$$

$$\int \frac{1}{(x+a)(x+b)} dx = \frac{1}{b-a} \ln \frac{a+x}{b+x}, \quad a \neq b \quad (7)$$

$$\int \frac{x}{(x+a)^2} dx = \frac{a}{a+x} + \ln|a+x| \quad (8)$$

$$\int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln|ax^2+bx+c| - \frac{b}{a\sqrt{4ac-b^2}} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}} \quad (9)$$

### Integrals with Roots

$$\int \frac{x}{\sqrt{x \pm a}} dx = \frac{2}{3} (x \mp 2a) \sqrt{x \pm a} \quad (10)$$

$$\int \sqrt{\frac{x}{a-x}} dx = -\sqrt{x(a-x)} - a \tan^{-1} \frac{\sqrt{x(a-x)}}{x-a} \quad (11)$$

$$\int \sqrt{\frac{x}{a+x}} dx = \sqrt{x(a+x)} - a \ln[\sqrt{x} + \sqrt{x+a}] \quad (12)$$

$$\int x\sqrt{ax+bx} dx = \frac{2}{15a^2} (-2b^2+abx+3a^2x^2)\sqrt{ax+bx} \quad (13)$$

$$\int \sqrt{x(ax+b)} dx = \frac{1}{4a^{3/2}} \left[ (2ax+b)\sqrt{x(ax+b)} - b^2 \ln \left| a\sqrt{x} + \sqrt{a(ax+b)} \right| \right] \quad (14)$$

$$\int \sqrt{x^3(ax+b)} dx = \left[ \frac{b}{12a} - \frac{b^2}{8a^2x} + \frac{x}{3} \right] \sqrt{x^3(ax+b)} + \frac{b^3}{\dots} \ln|a\sqrt{x} + \sqrt{a(ax+b)}| \quad (15)$$

$$\int \sqrt{x^2 \pm a^2} dx = \frac{1}{2}x\sqrt{x^2 \pm a^2} \pm \frac{1}{2}a^2 \ln|x + \sqrt{x^2 \pm a^2}| \quad (16)$$

$$\int \sqrt{a^2 - x^2} dx = \frac{1}{2}x\sqrt{a^2 - x^2} + \frac{1}{2}a^2 \tan^{-1} \frac{x}{\sqrt{a^2 - x^2}} \quad (17)$$

$$\int x\sqrt{x^2 \pm a^2} dx = \frac{1}{3} (x^2 \pm a^2)^{3/2} \quad (18)$$

$$\int \frac{1}{\sqrt{x^2 \pm a^2}} dx = \ln|x + \sqrt{x^2 \pm a^2}| \quad (19)$$

$$\int \frac{1}{\sqrt{a^2 - x^2}} dx = \sin^{-1} \frac{x}{a} \quad (20)$$

$$\int \frac{x}{\sqrt{x^2 \pm a^2}} dx = \sqrt{x^2 \pm a^2} \quad (21)$$

$$\int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2} \quad (22)$$

$$\int \frac{x^2}{\sqrt{x^2 \pm a^2}} dx = \frac{1}{2}x\sqrt{x^2 \pm a^2} \mp \frac{1}{2}a^2 \ln|x + \sqrt{x^2 \pm a^2}| \quad (23)$$

$$\int \sqrt{ax^2+bx+c} dx = \frac{b+2ax}{4a} \sqrt{ax^2+bx+c} + \frac{4ac-b^2}{8a^{3/2}} \ln \left| 2ax+b+2\sqrt{a(ax^2+bx+c)} \right| \quad (24)$$

$$\int x\sqrt{ax^2+bx+c} = \frac{1}{48a^{5/2}} \left( 2\sqrt{a}\sqrt{ax^2+bx+c} \times (-3b^2+2abx+8a(c+ax^2)) + 3(b^3-4abc) \ln \left| b+2ax+2\sqrt{a}\sqrt{ax^2+bx+c} \right| \right) \quad (25)$$

$$\int \frac{1}{\sqrt{ax^2+bx+c}} dx = \frac{1}{\sqrt{a}} \ln \left| 2ax+b+2\sqrt{a(ax^2+bx+c)} \right| \quad (26)$$

$$\int \frac{x}{\sqrt{ax^2+bx+c}} dx = \frac{1}{a} \sqrt{ax^2+bx+c} - \frac{b}{2a^{3/2}} \ln \left| 2ax+b+2\sqrt{a(ax^2+bx+c)} \right| \quad (27)$$

$$\int \frac{dx}{(a^2+x^2)^{3/2}} = \frac{x}{a^2\sqrt{a^2+x^2}} \quad (28)$$

### Integrals with Logarithms

$$\int \frac{\ln ax}{x} dx = \frac{1}{2} (\ln ax)^2 \quad (29)$$

$$\int \ln(ax+b) dx = \left( x + \frac{b}{a} \right) \ln(ax+b) - x, a \neq 0 \quad (30)$$

$$\int \ln(x^2+a^2) dx = x \ln(x^2+a^2) + 2a \tan^{-1} \frac{x}{a} - 2x \quad (31)$$

$$\int \ln(x^2-a^2) dx = x \ln(x^2-a^2) + a \ln \frac{x+a}{x-a} - 2x \quad (32)$$

$$\int \ln(ax^2+bx+c) dx = \frac{1}{a} \sqrt{4ac-b^2} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}} - 2x + \left( \frac{b}{2a} + x \right) \ln(ax^2+bx+c) \quad (33)$$

$$\int x \ln(ax+b) dx = \frac{bx}{2a} - \frac{1}{4}x^2 + \frac{1}{2} \left( x^2 - \frac{b^2}{a^2} \right) \ln(ax+b) \quad (34)$$

$$\int x \ln(a^2-b^2x^2) dx = -\frac{1}{2}x^2 + \frac{1}{2} \left( x^2 - \frac{a^2}{b^2} \right) \ln(a^2-b^2x^2) \quad (35)$$

### Integrals with Exponentials

$$\int x^n e^{ax} dx = \frac{x^n e^{ax}}{a} - \frac{n}{a} \int x^{n-1} e^{ax} dx \quad (36)$$

$$\int x e^{-ax^2} dx = -\frac{1}{2a} e^{-ax^2} \quad (37)$$

### Integrals with Trigonometric Functions

$$\int \sin^3 ax dx = -\frac{3 \cos ax}{4a} + \frac{\cos 3ax}{12a} \quad (38)$$

$$\int \cos^2 ax dx = \frac{x}{2} + \frac{\sin 2ax}{4a} \quad (39)$$

$$\int \cos^3 ax dx = \frac{3 \sin ax}{4a} + \frac{\sin 3ax}{12a} \quad (40)$$

$$\int \cos ax \sin bxdx = \frac{\cos[(a-b)x]}{2(a-b)} - \frac{\cos[(a+b)x]}{2(a+b)}, a \neq b \quad (41)$$

$$\int \sin^2 ax \cos bxdx = -\frac{\sin[(2a-b)x]}{4(2a-b)} + \frac{\sin bx}{2b} - \frac{\sin[(2a+b)x]}{4(2a+b)} \quad (42)$$

$$\int \sin^2 x \cos x dx = \frac{1}{3} \sin^3 x \quad (43)$$

$$\int \cos^2 ax \sin bxdx = \frac{\cos[(2a-b)x]}{4(2a-b)} - \frac{\cos bx}{2b} - \frac{\cos[(2a+b)x]}{4(2a+b)} \quad (44)$$

$$\int \cos^2 ax \sin axdx = -\frac{1}{3a} \cos^3 ax \quad (45)$$

$$\int \sin^2 ax \cos^2 bxdx = \frac{x}{4} - \frac{\sin 2ax}{8a} - \frac{\sin[2(a-b)x]}{16(a-b)} + \frac{\sin 2bx}{8b} - \frac{\sin[2(a+b)x]}{16(a+b)} \quad (46)$$

$$\int \sin^2 ax \cos^2 axdx = \frac{x}{8} - \frac{\sin 4ax}{32a} \quad (47)$$

$$\int \tan axdx = -\frac{1}{a} \ln \cos ax \quad (48)$$

$$\int \tan^2 axdx = -x + \frac{1}{a} \tan ax \quad (49)$$

$$\int \tan^3 axdx = \frac{1}{a} \ln \cos ax + \frac{1}{2a} \sec^2 ax \quad (50)$$

$$\int \sec xdx = \ln |\sec x + \tan x| = 2 \tanh^{-1} \left( \tan \frac{x}{2} \right) \quad (51)$$

$$\int \sec^2 axdx = \frac{1}{a} \tan ax \quad (52)$$

$$\int \sec^3 x dx = \frac{1}{2} \sec x \tan x + \frac{1}{2} \ln |\sec x + \tan x| \quad (53)$$

$$\int \sec x \tan xdx = \sec x \quad (54)$$

$$\int \sec^2 x \tan xdx = \frac{1}{2} \sec^2 x \quad (55)$$

$$\int \sec^n x \tan xdx = \frac{1}{n} \sec^n x, n \neq 0 \quad (56)$$

$$\int \csc xdx = \ln \left| \tan \frac{x}{2} \right| = \ln |\csc x - \cot x| + C \quad (57)$$

$$\int \csc^2 axdx = -\frac{1}{a} \cot ax \quad (58)$$

$$\int \csc^3 xdx = -\frac{1}{2} \cot x \csc x + \frac{1}{2} \ln |\csc x - \cot x| \quad (59)$$

$$\int \csc^n x \cot xdx = -\frac{1}{n} \csc^n x, n \neq 0 \quad (60)$$

$$\int \sec x \csc xdx = \ln |\tan x| \quad (61)$$

#### Products of Trigonometric Functions and Monomials

$$\int x \cos xdx = \cos x + x \sin x \quad (62)$$

$$\int x \cos axdx = \frac{1}{a^2} \cos ax + \frac{x}{a} \sin ax \quad (63)$$

$$\int x^2 \cos xdx = 2x \cos x + (x^2 - 2) \sin x \quad (64)$$

$$\int x^2 \cos axdx = \frac{2x \cos ax}{a^2} + \frac{a^2 x^2 - 2}{a^3} \sin ax \quad (65)$$

$$\int x \sin xdx = -x \cos x + \sin x \quad (66)$$

$$\int x \sin axdx = -\frac{x \cos ax}{a} + \frac{\sin ax}{a^2} \quad (67)$$

$$\int x^2 \sin xdx = (2 - x^2) \cos x + 2x \sin x \quad (68)$$

$$\int x^2 \sin axdx = \frac{2 - a^2 x^2}{a^3} \cos ax + \frac{2x \sin ax}{a^2} \quad (69)$$

#### Products of Trigonometric Functions and Exponentials

$$\int e^x \sin xdx = \frac{1}{2} e^x (\sin x - \cos x) \quad (70)$$

$$\int e^{bx} \sin axdx = \frac{1}{a^2 + b^2} e^{bx} (b \sin ax - a \cos ax) \quad (71)$$

$$\int e^x \cos xdx = \frac{1}{2} e^x (\sin x + \cos x) \quad (72)$$

$$\int e^{bx} \cos axdx = \frac{1}{a^2 + b^2} e^{bx} (a \sin ax + b \cos ax) \quad (73)$$

$$\int xe^x \sin xdx = \frac{1}{2} e^x (\cos x - x \cos x + x \sin x) \quad (74)$$

$$\int xe^x \cos xdx = \frac{1}{2} e^x (x \cos x - \sin x + x \sin x) \quad (75)$$

## Java

```
1 import java.io.*;
2 import java.util.*;
3 import java.math.*;
4 public class Main {
5     public static void main(String[] args) {
6         InputStream inputStream = System.in;
7         OutputStream outputStream = System.out;
8         InputReader in = new InputReader(inputStream);
9         PrintWriter out = new PrintWriter(outputStream);
10    }
11 }
12 public static class edge implements Comparable<edge>{
13     public int u,v,w;
14     public int compareTo(edge e){
15         return w-e.w;
16     }
17 }
18 public static class cmp implements Comparator<edge>{
19     public int compare(edge a,edge b){
20         if(a.w<b.w)return 1;
21         if(a.w>b.w)return -1;
22         return 0;
23     }
24 }
25 class InputReader {
26     public BufferedReader reader;
27     public StringTokenizer tokenizer;
28
29     public InputReader(InputStream stream) {
30         reader = new BufferedReader(new InputStreamReader(stream), 32768);
31         tokenizer = null;
32     }
33
34     public String next() {
35         while (tokenizer == null || !tokenizer.hasMoreTokens()) {
36             try {
37                 tokenizer = new StringTokenizer(reader.readLine());
38             } catch (IOException e) {
39                 throw new RuntimeException(e);
40             }
41         }
42         return tokenizer.nextToken();
43     }
44
45     public int nextInt() {
46         return Integer.parseInt(next());
47     }
48
49     public long nextLong() {
50         return Long.parseLong(next());
51     }
52 }
```



Other methods may have slightly different rounding semantics. For example, the result of the `pow` method using the [specified algorithm](#) can occasionally differ from the rounded mathematical result by more than one unit in the last place, one *ulp*.

Two types of operations are provided for manipulating the scale of a `BigDecimal`: scaling/rounding operations and decimal point motion operations. Scaling/rounding operations (`setScale` and `round`) return a `BigDecimal` whose value is approximately (or exactly) equal to that of the operand, but whose scale or precision is the specified value; that is, they increase or decrease the precision of the stored number with minimal effect on its value. Decimal point motion operations (`movePointLeft` and `movePointRight`) return a `BigDecimal` created from the operand by moving the decimal point a specified distance in the specified direction.

For the sake of brevity and clarity, pseudo-code is used throughout the descriptions of `BigDecimal` methods. The pseudo-code expression `(i + j)` is shorthand for "a `BigDecimal` whose value is that of the `BigDecimal` `i` added to that of the `BigDecimal` `j`." The pseudo-code expression `(i == j)` is shorthand for "true if and only if the `BigDecimal` `i` represents the same value as the `BigDecimal` `j`." Other pseudo-code expressions are interpreted similarly. Square brackets are used to represent the particular `BigInteger` and scale pair defining a `BigDecimal` value; for example `[19, 2]` is the `BigDecimal` numerically equal to 0.19 having a scale of 2.

Note: care should be exercised if `BigDecimal` objects are used as keys in a [SortedMap](#) or elements in a [SortedSet](#) since `BigDecimal`'s *natural ordering* is *inconsistent with equals*. See [Comparable](#), [SortedMap](#) or [SortedSet](#) for more information.

All methods and constructors for this class throw `NullPointerException` when passed a null object reference for any input parameter.

#### See Also:

[BigInteger](#), [MathContext](#), [RoundingMode](#), [SortedMap](#), [SortedSet](#), [Serialized Form](#)

## Field Summary

### Fields

Modifier and Type	Field and Description
static <a href="#">BigDecimal</a>	<b>ONE</b> The value 1, with a scale of 0.
static int	<b>ROUND_CEILING</b> Rounding mode to round towards positive infinity.
static int	<b>ROUND_DOWN</b> Rounding mode to round towards zero.
static int	<b>ROUND_FLOOR</b> Rounding mode to round towards negative infinity.
static int	<b>ROUND_HALF_DOWN</b> Rounding mode to round towards "nearest neighbor" unless both neighbors are equidistant, in which case round down.
static int	<b>ROUND_HALF_EVEN</b>

Rounding mode to round towards the "nearest neighbor" unless both neighbors are equidistant, in which case, round towards the even neighbor.

static int

**ROUND\_HALF\_UP**

Rounding mode to round towards "nearest neighbor" unless both neighbors are equidistant, in which case round up.

static int

**ROUND\_UNNECESSARY**

Rounding mode to assert that the requested operation has an exact result, hence no rounding is necessary.

static int

**ROUND\_UP**

Rounding mode to round away from zero.

static **BigDecimal** **TEN**

The value 10, with a scale of 0.

static **BigDecimal** **ZERO**

The value 0, with a scale of 0.

## Constructor Summary

### Constructors

#### Constructor and Description

**BigDecimal**(**BigInteger** val)

Translates a **BigInteger** into a **BigDecimal**.

**BigDecimal**(**BigInteger** unscaledVal, int scale)

Translates a **BigInteger** unscaled value and an int scale into a **BigDecimal**.

**BigDecimal**(**BigInteger** unscaledVal, int scale, **MathContext** mc)

Translates a **BigInteger** unscaled value and an int scale into a **BigDecimal**, with rounding according to the context settings.

**BigDecimal**(**BigInteger** val, **MathContext** mc)

Translates a **BigInteger** into a **BigDecimal** rounding according to the context settings.

**BigDecimal**(char[] in)

Translates a character array representation of a **BigDecimal** into a **BigDecimal**, accepting the same sequence of characters as the **BigDecimal(String)** constructor.

**BigDecimal**(char[] in, int offset, int len)

Translates a character array representation of a **BigDecimal** into a **BigDecimal**, accepting the same sequence of characters as the **BigDecimal(String)** constructor, while allowing a sub-array to be specified.

**BigDecimal**(char[] in, int offset, int len, **MathContext** mc)

Translates a character array representation of a **BigDecimal** into a **BigDecimal**, accepting the same sequence of characters as the **BigDecimal(String)** constructor, while allowing a sub-array to be specified and with rounding according to the context settings.

**BigDecimal(char[] in, MathContext mc)**

Translates a character array representation of a `BigDecimal` into a `BigDecimal`, accepting the same sequence of characters as the **BigDecimal(String)** constructor and with rounding according to the context settings.

**BigDecimal(double val)**

Translates a double into a `BigDecimal` which is the exact decimal representation of the double's binary floating-point value.

**BigDecimal(double val, MathContext mc)**

Translates a double into a `BigDecimal`, with rounding according to the context settings.

**BigDecimal(int val)**

Translates an int into a `BigDecimal`.

**BigDecimal(int val, MathContext mc)**

Translates an int into a `BigDecimal`, with rounding according to the context settings.

**BigDecimal(long val)**

Translates a long into a `BigDecimal`.

**BigDecimal(long val, MathContext mc)**

Translates a long into a `BigDecimal`, with rounding according to the context settings.

**BigDecimal(String val)**

Translates the string representation of a `BigDecimal` into a `BigDecimal`.

**BigDecimal(String val, MathContext mc)**

Translates the string representation of a `BigDecimal` into a `BigDecimal`, accepting the same strings as the **BigDecimal(String)** constructor, with rounding according to the context settings.

## Method Summary

### All Methods    Static Methods    Instance Methods    Concrete Methods

Modifier and Type	Method and Description
<b>BigDecimal</b>	<b>abs()</b> Returns a <code>BigDecimal</code> whose value is the absolute value of this <code>BigDecimal</code> , and whose scale is <code>this.scale()</code> .
<b>BigDecimal</b>	<b>abs(MathContext mc)</b> Returns a <code>BigDecimal</code> whose value is the absolute value of this <code>BigDecimal</code> , with rounding according to the context settings.
<b>BigDecimal</b>	<b>add(BigDecimal augend)</b> Returns a <code>BigDecimal</code> whose value is <code>(this + augend)</code> , and whose scale is <code>max(this.scale(), augend.scale())</code> .
<b>BigDecimal</b>	<b>add(BigDecimal augend, MathContext mc)</b> Returns a <code>BigDecimal</code> whose value is <code>(this + augend)</code> , with rounding according to the context settings.

byte	<b>byteValueExact()</b> Converts this <code>BigDecimal</code> to a byte, checking for lost information.
int	<b>compareTo(BigDecimal val)</b> Compares this <code>BigDecimal</code> with the specified <code>BigDecimal</code> .
BigDecimal	<b>divide(BigDecimal divisor)</b> Returns a <code>BigDecimal</code> whose value is $(\text{this} / \text{divisor})$ , and whose preferred scale is $(\text{this}.\text{scale}() - \text{divisor}.\text{scale}())$ ; if the exact quotient cannot be represented (because it has a non-terminating decimal expansion) an <code>ArithmeticException</code> is thrown.
BigDecimal	<b>divide(BigDecimal divisor, int roundingMode)</b> Returns a <code>BigDecimal</code> whose value is $(\text{this} / \text{divisor})$ , and whose scale is <code>this.scale()</code> .
BigDecimal	<b>divide(BigDecimal divisor, int scale, int roundingMode)</b> Returns a <code>BigDecimal</code> whose value is $(\text{this} / \text{divisor})$ , and whose scale is as specified.
BigDecimal	<b>divide(BigDecimal divisor, int scale, RoundingMode roundingMode)</b> Returns a <code>BigDecimal</code> whose value is $(\text{this} / \text{divisor})$ , and whose scale is as specified.
BigDecimal	<b>divide(BigDecimal divisor, MathContext mc)</b> Returns a <code>BigDecimal</code> whose value is $(\text{this} / \text{divisor})$ , with rounding according to the context settings.
BigDecimal	<b>divide(BigDecimal divisor, RoundingMode roundingMode)</b> Returns a <code>BigDecimal</code> whose value is $(\text{this} / \text{divisor})$ , and whose scale is <code>this.scale()</code> .
BigDecimal[]	<b>divideAndRemainder(BigDecimal divisor)</b> Returns a two-element <code>BigDecimal</code> array containing the result of <code>divideToIntegerValue</code> followed by the result of remainder on the two operands.
BigDecimal[]	<b>divideAndRemainder(BigDecimal divisor, MathContext mc)</b> Returns a two-element <code>BigDecimal</code> array containing the result of <code>divideToIntegerValue</code> followed by the result of remainder on the two operands calculated with rounding according to the context settings.
BigDecimal	<b>divideToIntegerValue(BigDecimal divisor)</b> Returns a <code>BigDecimal</code> whose value is the integer part of the quotient $(\text{this} / \text{divisor})$ rounded down.
BigDecimal	<b>divideToIntegerValue(BigDecimal divisor, MathContext mc)</b> Returns a <code>BigDecimal</code> whose value is the integer part of $(\text{this} / \text{divisor})$ .
double	<b>doubleValue()</b>

Converts this `BigDecimal` to a `double`.

`boolean`

**`equals(Object x)`**

Compares this `BigDecimal` with the specified `Object` for equality.

`float`

**`floatValue()`**

Converts this `BigDecimal` to a `float`.

`int`

**`hashCode()`**

Returns the hash code for this `BigDecimal`.

`int`

**`intValue()`**

Converts this `BigDecimal` to an `int`.

`int`

**`intValueExact()`**

Converts this `BigDecimal` to an `int`, checking for lost information.

`long`

**`longValue()`**

Converts this `BigDecimal` to a `long`.

`long`

**`longValueExact()`**

Converts this `BigDecimal` to a `long`, checking for lost information.

**`BigDecimal`**

**`max(BigDecimal val)`**

Returns the maximum of this `BigDecimal` and `val`.

**`BigDecimal`**

**`min(BigDecimal val)`**

Returns the minimum of this `BigDecimal` and `val`.

**`BigDecimal`**

**`movePointLeft(int n)`**

Returns a `BigDecimal` which is equivalent to this one with the decimal point moved `n` places to the left.

**`BigDecimal`**

**`movePointRight(int n)`**

Returns a `BigDecimal` which is equivalent to this one with the decimal point moved `n` places to the right.

**`BigDecimal`**

**`multiply(BigDecimal multiplicand)`**

Returns a `BigDecimal` whose value is  $(\text{this} \times \text{multiplicand})$ , and whose scale is  $(\text{this.scale}() + \text{multiplicand.scale}())$ .

**`BigDecimal`**

**`multiply(BigDecimal multiplicand, MathContext mc)`**

Returns a `BigDecimal` whose value is  $(\text{this} \times \text{multiplicand})$ , with rounding according to the context settings.

**`BigDecimal`**

**`negate()`**

Returns a `BigDecimal` whose value is  $(-\text{this})$ , and whose scale is `this.scale()`.

**`BigDecimal`**

**`negate(MathContext mc)`**

Returns a `BigDecimal` whose value is  $(-\text{this})$ , with rounding according to the context settings.

**`BigDecimal`**

**`plus()`**

Returns a `BigDecimal` whose value is `(+this)`, and whose scale is `this.scale()`.

**BigDecimal****plus(MathContext mc)**

Returns a `BigDecimal` whose value is `(+this)`, with rounding according to the context settings.

**BigDecimal****pow(int n)**

Returns a `BigDecimal` whose value is `(thisn)`. The power is computed exactly, to unlimited precision.

**BigDecimal****pow(int n, MathContext mc)**

Returns a `BigDecimal` whose value is `(thisn)`.

int

**precision()**

Returns the *precision* of this `BigDecimal`.

**BigDecimal****remainder(BigDecimal divisor)**

Returns a `BigDecimal` whose value is `(this % divisor)`.

**BigDecimal****remainder(BigDecimal divisor, MathContext mc)**

Returns a `BigDecimal` whose value is `(this % divisor)`, with rounding according to the context settings.

**BigDecimal****round(MathContext mc)**

Returns a `BigDecimal` rounded according to the `MathContext` settings.

int

**scale()**

Returns the *scale* of this `BigDecimal`.

**BigDecimal****scaleByPowerOfTen(int n)**

Returns a `BigDecimal` whose numerical value is equal to `(this * 10n)`.

**BigDecimal****setScale(int newScale)**

Returns a `BigDecimal` whose scale is the specified value, and whose value is numerically equal to this `BigDecimal`'s.

**BigDecimal****setScale(int newScale, int roundingMode)**

Returns a `BigDecimal` whose scale is the specified value, and whose unscaled value is determined by multiplying or dividing this `BigDecimal`'s unscaled value by the appropriate power of ten to maintain its overall value.

**BigDecimal****setScale(int newScale, RoundingMode roundingMode)**

Returns a `BigDecimal` whose scale is the specified value, and whose unscaled value is determined by multiplying or dividing this `BigDecimal`'s unscaled value by the appropriate power of ten to maintain its overall value.

short

**shortValueExact()**

Converts this `BigDecimal` to a `short`, checking for lost information.

int

**signum()**

[PREV CLASS](#) [NEXT CLASS](#) [FRAMES](#) [NO FRAMES](#) [ALL CLASSES](#)[SUMMARY: NESTED](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#) [DETAIL: FIELD](#) | [CONSTR](#) | [METHOD](#)

compact1, compact2, compact3

java.util

## Class `TreeMap<K,V>`

java.lang.Object

java.util.AbstractMap&lt;K,V&gt;

java.util.TreeMap&lt;K,V&gt;

### Type Parameters:

K - the type of keys maintained by this map

V - the type of mapped values

### All Implemented Interfaces:

`Serializable`, `Cloneable`, `Map<K,V>`, `NavigableMap<K,V>`, `SortedMap<K,V>`

```
public class TreeMap<K,V>
extends AbstractMap<K,V>
implements NavigableMap<K,V>, Cloneable, Serializable
```

A Red-Black tree based `NavigableMap` implementation. The map is sorted according to the **natural ordering** of its keys, or by a `Comparator` provided at map creation time, depending on which constructor is used.

This implementation provides guaranteed  $\log(n)$  time cost for the `containsKey`, `get`, `put` and `remove` operations. Algorithms are adaptations of those in Cormen, Leiserson, and Rivest's *Introduction to Algorithms*.

Note that the ordering maintained by a tree map, like any sorted map, and whether or not an explicit comparator is provided, must be *consistent with equals* if this sorted map is to correctly implement the `Map` interface. (See `Comparable` or `Comparator` for a precise definition of *consistent with equals*.) This is so because the `Map` interface is defined in terms of the `equals` operation, but a sorted map performs all key comparisons using its `compareTo` (or `compare`) method, so two keys that are deemed equal by this method are, from the standpoint of the sorted map, equal. The behavior of a sorted map is well-defined even if its ordering is inconsistent with `equals`; it just fails to obey the general contract of the `Map` interface.

**Note that this implementation is not synchronized.** If multiple threads access a map concurrently, and at least one of the threads modifies the map structurally, it *must* be synchronized externally. (A structural modification is any operation that adds or deletes one or more mappings; merely changing the value associated with an existing key is not a structural modification.) This is typically accomplished by synchronizing on some object that naturally encapsulates the map. If no such object exists, the map should be "wrapped" using the `Collections.synchronizedSortedMap` method. This is best done at creation time, to prevent accidental unsynchronized access to the map:

```
SortedMap m = Collections.synchronizedSortedMap(new TreeMap(...));
```

The iterators returned by the `iterator` method of the collections returned by all of this



class's "collection view methods" are *fail-fast*: if the map is structurally modified at any time after the iterator is created, in any way except through the iterator's own `remove` method, the iterator will throw a `ConcurrentModificationException`. Thus, in the face of concurrent modification, the iterator fails quickly and cleanly, rather than risking arbitrary, non-deterministic behavior at an undetermined time in the future.

Note that the fail-fast behavior of an iterator cannot be guaranteed as it is, generally speaking, impossible to make any hard guarantees in the presence of unsynchronized concurrent modification. Fail-fast iterators throw `ConcurrentModificationException` on a best-effort basis. Therefore, it would be wrong to write a program that depended on this exception for its correctness: *the fail-fast behavior of iterators should be used only to detect bugs*.

All `Map.Entry` pairs returned by methods in this class and its views represent snapshots of mappings at the time they were produced. They do **not** support the `Entry.setValue` method. (Note however that it is possible to change mappings in the associated map using `put`.)

This class is a member of the [Java Collections Framework](#).

**Since:**

1.2

**See Also:**

[Map](#), [HashMap](#), [Hashtable](#), [Comparable](#), [Comparator](#), [Collection](#), [Serialized Form](#)

## ***Nested Class Summary***

### **Nested classes/interfaces inherited from class `java.util.AbstractMap`**

`AbstractMap.SimpleEntry<K,V>`, `AbstractMap.SimpleImmutableEntry<K,V>`

## ***Constructor Summary***

### **Constructors**

#### **Constructor and Description**

##### **`TreeMap()`**

Constructs a new, empty tree map, using the natural ordering of its keys.

##### **`TreeMap(Comparator<? super K> comparator)`**

Constructs a new, empty tree map, ordered according to the given comparator.

##### **`TreeMap(Map<? extends K,? extends V> m)`**

Constructs a new tree map containing the same mappings as the given map, ordered according to the *natural ordering* of its keys.

##### **`TreeMap(SortedMap<K,? extends V> m)`**

Constructs a new tree map containing the same mappings and using the same ordering as the specified sorted map.

## ***Method Summary***



Modifier and Type	Method and Description
<b>Map.Entry&lt;K, V&gt;</b>	<b>ceilingEntry(K key)</b> Returns a key-value mapping associated with the least key greater than or equal to the given key, or null if there is no such key.
<b>K</b>	<b>ceilingKey(K key)</b> Returns the least key greater than or equal to the given key, or null if there is no such key.
<b>void</b>	<b>clear()</b> Removes all of the mappings from this map.
<b>Object</b>	<b>clone()</b> Returns a shallow copy of this TreeMap instance.
<b>Comparator&lt;? super K&gt;</b>	<b>comparator()</b> Returns the comparator used to order the keys in this map, or null if this map uses the <b>natural ordering</b> of its keys.
<b>boolean</b>	<b>containsKey(Object key)</b> Returns true if this map contains a mapping for the specified key.
<b>boolean</b>	<b>containsValue(Object value)</b> Returns true if this map maps one or more keys to the specified value.
<b>NavigableSet&lt;K&gt;</b>	<b>descendingKeySet()</b> Returns a reverse order <b>NavigableSet</b> view of the keys contained in this map.
<b>NavigableMap&lt;K, V&gt;</b>	<b>descendingMap()</b> Returns a reverse order view of the mappings contained in this map.
<b>Set&lt;Map.Entry&lt;K, V&gt;&gt;</b>	<b>entrySet()</b> Returns a <b>Set</b> view of the mappings contained in this map.
<b>Map.Entry&lt;K, V&gt;</b>	<b>firstEntry()</b> Returns a key-value mapping associated with the least key in this map, or null if the map is empty.
<b>K</b>	<b>firstKey()</b> Returns the first (lowest) key currently in this map.
<b>Map.Entry&lt;K, V&gt;</b>	<b>floorEntry(K key)</b> Returns a key-value mapping associated with the greatest key less than or equal to the given key, or null if there is no such key.
<b>K</b>	<b>floorKey(K key)</b> Returns the greatest key less than or equal to the given key, or null if there is no such key.

or null if there is no such key.

void

**forEach**(**BiConsumer**<? super **K**,? super **V**> action)

Performs the given action for each entry in this map until all entries have been processed or the action throws an exception.

**V**

**get**(**Object** key)

Returns the value to which the specified key is mapped, or null if this map contains no mapping for the key.

**SortedMap**<**K**,**V**>

**headMap**(**K** toKey)

Returns a view of the portion of this map whose keys are strictly less than toKey.

**NavigableMap**<**K**,**V**>

**headMap**(**K** toKey, boolean inclusive)

Returns a view of the portion of this map whose keys are less than (or equal to, if inclusive is true) toKey.

**Map.Entry**<**K**,**V**>

**higherEntry**(**K** key)

Returns a key-value mapping associated with the least key strictly greater than the given key, or null if there is no such key.

**K**

**higherKey**(**K** key)

Returns the least key strictly greater than the given key, or null if there is no such key.

**Set**<**K**>

**keySet**()

Returns a **Set** view of the keys contained in this map.

**Map.Entry**<**K**,**V**>

**lastEntry**()

Returns a key-value mapping associated with the greatest key in this map, or null if the map is empty.

**K**

**lastKey**()

Returns the last (highest) key currently in this map.

**Map.Entry**<**K**,**V**>

**lowerEntry**(**K** key)

Returns a key-value mapping associated with the greatest key strictly less than the given key, or null if there is no such key.

**K**

**lowerKey**(**K** key)

Returns the greatest key strictly less than the given key, or null if there is no such key.

**NavigableSet**<**K**>

**navigableKeySet**()

Returns a **NavigableSet** view of the keys contained in this map.

**Map.Entry**<**K**,**V**>

**pollFirstEntry**()

Removes and returns a key-value mapping associated with the least key in this map, or null if the map is empty.

**Map.Entry**<**K**,**V**>

**pollLastEntry**()

Removes and returns a key-value mapping associated with the greatest key in this map, or null if the map is empty.

the greatest key in this map, or null if the map is empty.

<b>V</b>	<b>put(K key, V value)</b> Associates the specified value with the specified key in this map.
void	<b>putAll(Map&lt;? extends K,? extends V&gt; map)</b> Copies all of the mappings from the specified map to this map.
<b>V</b>	<b>remove(Object key)</b> Removes the mapping for this key from this TreeMap if present.
<b>V</b>	<b>replace(K key, V value)</b> Replaces the entry for the specified key only if it is currently mapped to some value.
boolean	<b>replace(K key, V oldValue, V newValue)</b> Replaces the entry for the specified key only if currently mapped to the specified value.
void	<b>replaceAll(BiFunction&lt;? super K,? super V,? extends V&gt; function)</b> Replaces each entry's value with the result of invoking the given function on that entry until all entries have been processed or the function throws an exception.
int	<b>size()</b> Returns the number of key-value mappings in this map.
<b>NavigableMap&lt;K,V&gt;</b>	<b>subMap(K fromKey, boolean fromInclusive, K toKey, boolean toInclusive)</b> Returns a view of the portion of this map whose keys range from fromKey to toKey.
<b>SortedMap&lt;K,V&gt;</b>	<b>subMap(K fromKey, K toKey)</b> Returns a view of the portion of this map whose keys range from fromKey, inclusive, to toKey, exclusive.
<b>SortedMap&lt;K,V&gt;</b>	<b>tailMap(K fromKey)</b> Returns a view of the portion of this map whose keys are greater than or equal to fromKey.
<b>NavigableMap&lt;K,V&gt;</b>	<b>tailMap(K fromKey, boolean inclusive)</b> Returns a view of the portion of this map whose keys are greater than (or equal to, if inclusive is true) fromKey.
<b>Collection&lt;V&gt;</b>	<b>values()</b> Returns a <b>Collection</b> view of the values contained in this map.

### Methods inherited from class java.util.AbstractMap

equals, hashCode, isEmpty, toString

