Code Rate

Suppose we have an information source producing information at a rate H' bits/second. If we use binary symbols (binits) to represent this source stream then we have a symbol rate:

$$R_B = \frac{H'}{\log 2}$$
 binits/second

When we channel code we *purposefully* add extra *parity* symbols to the information stream to make $R \le C$ and to hopefully lower the error rate. For every k information symbols we form a block of n symbols with $n \ge k$. We demand even with the extra parity symbols added that the information arrive at the destination *at the same rate*:

$$kT_B = nT_S$$

where $T_B = \frac{1}{R_B}$, the original information symbol period and $T_S = \frac{1}{R_S}$, the period of a symbol in the new block of n symbols. Rearranging we define a code rate r:

$$r = \frac{k}{n} = \frac{T_S}{T_B} = \frac{R_B}{R_S}$$

The code rate is a dimensionless quantity and for binary systems $0 < r \le 1$. Using the code rate in signal-to-noise ratio calculations:

$$E_B = ST_B$$
 and $E_S = ST_S$

$$\frac{E_B}{T_B} = S = \frac{E_S}{T_S}$$

$$E_{S} = \frac{T_{S}}{T_{B}} E_{B} = r E_{B}$$

$$\frac{E_S}{N_0} = r \frac{E_B}{N_0}$$

$$\frac{E_{S}}{N_{0(dB)}} = \frac{E_{B}}{N_{0(dB)}} + 10 \log_{10} r$$

To make fair comparisons between various channel coding schemes and code rates the *bit noise* ratio $\frac{E_B}{N_0}$ is used.

The price for demanding that the information rate be unchanged is *bandwidth expansion*. In general, system bandwidth is directly proportional to signal rate:

$$W \propto R$$

$$R_S = \frac{R_B}{r}$$

$$W_S = \frac{1}{r} W_B$$

Example

Suppose we are working with a binary AWGN channel and a $\frac{E_B}{N_0}$ of 6 dB. We are interested in channel coding schemes with a code rate r of about 0.5.

Checking the capacity chart on the next page this is in the "attainable" region meaning that arbitrarily low error rates are possible.

| Uncoded | 2.4e-3 |
|-----------------------|-------------|
| Hamming(8,4) | 3.5e-3 |
| Golay(24,12) | 3.5e-4 |
| Reed-Solomon(255,128) | 1.7e-5 |
| Reed-Solomon(511,256) | 8.6e-6 |
| DSN Concatenated | almost zero |

The price for better error performance is larger block size n and greater coding/decoding complexity.

