# Landlock LSM: toward unprivileged sandboxing

Mickaël Salaün

ANSSI

September 14, 2017

# Secure user-space software

#### How to harden an application?

- secure development
- ▶ follow the least privilege principle
- compartmentalize exposed processes

# Secure user-space software

#### How to harden an application?

- secure development
- follow the least privilege principle
- compartmentalize exposed processes

#### Multiple sandbox uses

- built-in sandboxing (tailored security policy)
- sandbox managers (unprivileged and dynamic compartmentalization)
- container managers (hardened containers)

	Fine-grained control	Embedded policy	Unprivileged use
SELinux	$\checkmark$		

	Fine-grained control	Embedded policy	Unprivileged use
SELinux	✓		
seccomp-bpf		✓	✓
namespaces		✓	~

	Fine-grained control	Embedded policy	Unprivileged use
SELinux	$\checkmark$		
seccomp-bpf		✓	✓
namespaces		✓	~
Landlock	√	✓	<b>√</b>

Tailored access control to match your needs: programmatic access control

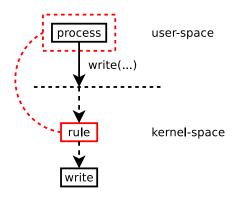
	Fine-grained control	Embedded policy	Unprivileged use
SELinux	$\checkmark$		
seccomp-bpf		✓	✓
namespaces		✓	~
Landlock	√	✓	✓

Tailored access control to match your needs: programmatic access control

#### Example

Run an application allowed to write only on a terminal.

## Landlock overview



# Landlock: patch v7

- ▶ a minimum viable product
- a stackable LSM
- using eBPF
- focused on filesystem access control

# The Linux Security Modules framework (LSM)

#### LSM framework

- allow or deny user-space actions on kernel objects
- policy decision and enforcement points
- kernel API: support various security models
- ▶ 200+ hooks: inode\_permission, inode\_unlink, file\_ioctl...

# The Linux Security Modules framework (LSM)

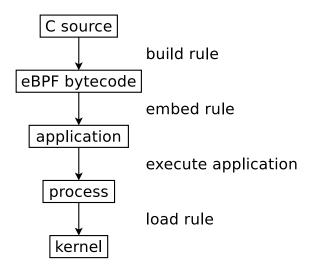
#### LSM framework

- allow or deny user-space actions on kernel objects
- policy decision and enforcement points
- kernel API: support various security models
- ▶ 200+ hooks: inode\_permission, inode\_unlink, file\_ioctl...

#### Landlock

- rule: control an action on an object
- event: use of a kernel object type (e.g. file)
- action: read, write, execute, remove, IOCTL...

# Life cycle of a Landlock rule



- read-only access to the filesystem...
- ...but allowed to write on TTY and pipes
- rule enforced on each filesystem access request

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
            mode = bpf handle fs get mode(ctx->arg1);
10
11
            /* allow write on TTY and pipes */
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1;
15
```

```
SEC("landlock1")
   int landlock fs rule1(struct landlock context *ctx)
3
4
            int mode;
5
6
            /* allow non-write actions */
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
            mode = bpf handle fs get mode(ctx->arg1);
10
11
            /* allow write on TTY and pipes */
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1;
15
```

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
7
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
            mode = bpf handle fs get mode(ctx->arg1);
10
11
            /* allow write on TTY and pipes */
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1;
15
```

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
7
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
10
            mode = bpf handle fs get mode(ctx->arg1);
11
            /* allow write on TTY and pipes */
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1:
15
```

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
7
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
10
            mode = bpf handle fs get mode(ctx->arg1);
11
            /* allow write on TTY and pipes */
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1;
15
```

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
            mode = bpf handle fs get mode(ctx->arg1);
10
            /* allow write on TTY and pipes */
11
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1;
15
```

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
10
            mode = bpf handle fs get mode(ctx->arg1);
            /* allow write on TTY and pipes */
11
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1:
15
```

```
SEC("landlock1")
    int landlock fs rule1(struct landlock context *ctx)
3
    {
4
            int mode;
5
6
            /* allow non-write actions */
7
            if (!(ctx->arg2 & LANDLOCK ACTION FS WRITE))
8
                    return 0;
9
            /* get the file mode */
            mode = bpf handle fs get mode(ctx->arg1);
10
11
            /* allow write on TTY and pipes */
12
            if (S ISCHR(mode) || S ISFIFO(mode))
13
                    return 0;
14
            return 1:
15
```

## extended Berkeley Packet Filter

#### In-kernel virtual machine

- safely execute code in the kernel at run time
- widely used in the kernel: network filtering, seccomp-bpf, tracing...
- can call dedicated functions
- can exchange data through maps between eBPF programs and user-space

## extended Berkeley Packet Filter

#### In-kernel virtual machine

- safely execute code in the kernel at run time
- widely used in the kernel: network filtering, seccomp-bpf, tracing...
- can call dedicated functions
- can exchange data through maps between eBPF programs and user-space

#### Static program verification at load time

- memory access checks
- register typing and tainting
- pointer leak restrictions
- execution flow restrictions

```
static union bpf prog subtype metadata = {
            .landlock rule = {
3
                     .event = LANDLOCK EVENT FS,
4
                     .ability = LANDLOCK ABILITY DEBUG,
5
6
    };
    union bpf attr attr = {
8
            .insns = bytecode array,
            .prog type = BPF PROG TYPE LANDLOCK RULE,
            .prog subtype = &metadata,
10
11
            // [...]
12
   int rule fd = bpf(BPF PROG LOAD, &attr, sizeof(attr));
13
```

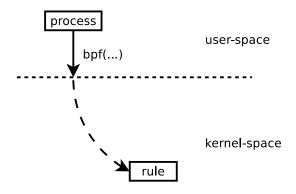
```
static union bpf prog subtype metadata = {
            .landlock rule = {
3
                     .event = LANDLOCK EVENT FS,
4
                     .ability = LANDLOCK ABILITY DEBUG,
5
6
    };
    union bpf attr attr = {
8
            .insns = bytecode array,
            .prog type = BPF PROG TYPE LANDLOCK RULE,
10
            .prog subtype = &metadata,
11
            // [...]
12
    };
13
    int rule fd = bpf(BPF PROG LOAD, &attr, sizeof(attr));
```

```
static union bpf prog subtype metadata = {
            .landlock rule = {
3
                     .event = LANDLOCK EVENT FS,
4
                     .ability = LANDLOCK ABILITY DEBUG,
5
6
    };
    union bpf attr attr = {
8
            .insns = bytecode array,
            .prog type = BPF PROG TYPE LANDLOCK RULE,
10
            .prog subtype = &metadata,
11
            // [...]
12
    };
13
    int rule fd = bpf(BPF PROG LOAD, &attr, sizeof(attr));
```

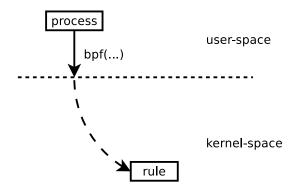
```
static union bpf prog subtype metadata = {
            .landlock rule = {
3
                     .event = LANDLOCK EVENT FS,
4
                     .ability = LANDLOCK ABILITY DEBUG,
5
6
    };
    union bpf attr attr = {
8
            .insns = bytecode array,
            .prog type = BPF PROG TYPE LANDLOCK RULE,
10
            .prog subtype = &metadata,
11
12
    };
13
    int rule fd = bpf(BPF PROG LOAD, &attr, sizeof(attr));
```

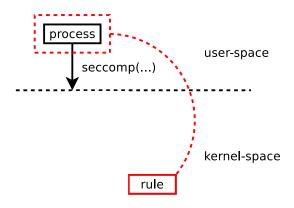
```
static union bpf prog subtype metadata = {
2
            .landlock rule = {
3
                     .event = LANDLOCK EVENT FS,
4
                     .ability = LANDLOCK ABILITY DEBUG,
5
6
    };
    union bpf attr attr = {
8
            .insns = bytecode array,
            .prog type = BPF PROG TYPE LANDLOCK RULE,
10
            .prog subtype = &metadata,
11
            // [...]
12
    };
13
    int rule fd = bpf(BPF PROG LOAD, &attr, sizeof(attr));
```

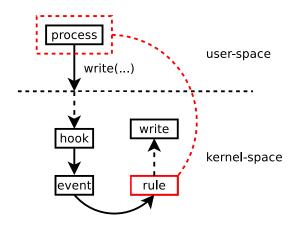
```
static union bpf prog subtype metadata = {
            .landlock rule = {
3
                     .event = LANDLOCK EVENT FS,
4
                     .ability = LANDLOCK ABILITY DEBUG,
5
6
    };
    union bpf attr attr = {
8
            .insns = bytecode array,
            .prog type = BPF PROG TYPE LANDLOCK RULE,
            .prog subtype = &metadata,
10
11
            // [...]
12
   int rule fd = bpf(BPF PROG LOAD, &attr, sizeof(attr));
13
```



```
1 \mid \mathsf{seccomp}(\mathsf{SECCOMP\_PREPEND\_LANDLOCK\_RULE}, \ \emptyset, \ \&\mathsf{rule\_fd});
```



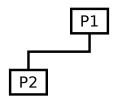


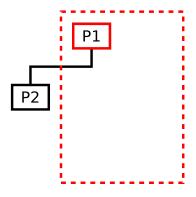


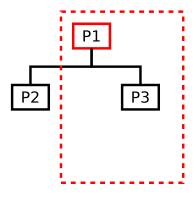
# Rule enforcement on process hierarchy

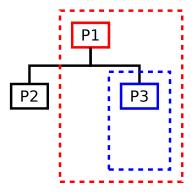


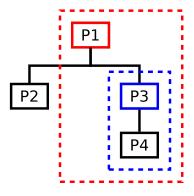
# Rule enforcement on process hierarchy











## Demonstration #1

An (almost) read-only filesystem

# Landlock: pending features

- unprivileged access control
- enforcement through cgroups
- eBPF map fsview
- coming next...

## Unprivileged access control

#### Why?

embed a security policy in any application, following the least privilege principle

## Unprivileged access control

#### Why?

embed a security policy in any application, following the least privilege principle

#### Challenges

- applying a security policy requires privileges
- unlike SUID, Landlock should only reduce accesses
- prevent accesses through other processes: ptrace restrictions
- protect the kernel: eBPF static analysis
- prevent information leak: an eBPF program shall not have more access rights than the process which loaded it

## Enforcement through cgroups

## Why?

user/admin security policy (e.g. container): manage groups of processes

## Enforcement through cgroups

#### Why?

user/admin security policy (e.g. container): manage groups of processes

#### Challenges

- complementary to the process hierarchy rules (via seccomp(2))
- processes moving in or out of a cgroup
- unprivileged use with cgroups delegation (e.g. user session)

# eBPF map fsview

Why?

restrict access to a subset of the filesystem

## eBPF map fsview

#### Why?

restrict access to a subset of the filesystem

#### Challenges

- efficient
- updatable from user-space
- unprivileged use (i.e. no xattr)

## eBPF map fsview

#### Why?

restrict access to a subset of the filesystem

#### Challenges

- efficient
- updatable from user-space
- unprivileged use (i.e. no xattr)

#### Proposal

- new eBPF map to identify a filesystem view: mount point hierarchies at a given time
- new eBPF function to compare a file access to such a view

## Demonstration #2

What might a filesystem access control looks like?

## Current roadmap

#### Incremental upstream integration

- 1. minimum viable product
- 2. cgroups handling
- 3. new eBPF map type for filesystem-related checks
- 4. unprivileged mode

## Landlock: wrap-up

## User-space hardening

- programmatic access control
- designed for unprivileged use

## Landlock: wrap-up

## User-space hardening

- programmatic access control
- designed for unprivileged use

#### Current status: patch v7

- autonomous patches merged in net, security and kselftest trees
- ▶ security/landlock/\*: ~1K SLOC
- ongoing patch series: LKML, @l0kod
- growing interest for containers, secure OS and service managers

# https://landlock.io

#### Landlock events

► LANDLOCK\_EVENT\_**FS** 

#### Landlock events

► LANDLOCK\_EVENT\_**FS** 

#### Landlock events

► LANDLOCK\_EVENT\_**FS** 

#### Landlock actions for an FS event

- ► LANDLOCK\_ACTION\_FS\_EXEC
- LANDLOCK\_ACTION\_FS\_WRITE
- ► LANDLOCK\_ACTION\_FS\_READ
- ► LANDLOCK\_ACTION\_FS\_NEW
- ► LANDLOCK\_ACTION\_FS\_GET
- ► LANDLOCK\_ACTION\_FS\_REMOVE
- ► LANDLOCK\_ACTION\_FS\_IOCTL
- ► LANDLOCK\_ACTION\_FS\_LOCK
- ► LANDLOCK\_ACTION\_FS\_FCNTL

#### Landlock events

- LANDLOCK\_EVENT\_FS
- ► LANDLOCK\_EVENT\_FS\_IOCTL
- ► LANDLOCK EVENT **FS LOCK**
- LANDLOCK\_EVENT\_FS\_FCNTL

#### Landlock actions for an FS event

- ► LANDLOCK\_ACTION\_FS\_EXEC
- LANDLOCK\_ACTION\_FS\_WRITE
- ► LANDLOCK\_ACTION\_FS\_READ
- ► LANDLOCK\_ACTION\_FS\_NEW
- LANDLOCK\_ACTION\_FS\_GET
- ► LANDLOCK\_ACTION\_FS\_REMOVE
- ► LANDLOCK\_ACTION\_FS\_IOCTL
- ► LANDLOCK\_ACTION\_FS\_LOCK
- ► LANDLOCK\_ACTION\_FS\_FCNTL

## Available eBPF functions for Landlock rules

## Any rule

bpf\_handle\_fs\_get\_mode

#### Available eBPF functions for Landlock rules

## Any rule

bpf\_handle\_fs\_get\_mode

## Debug mode: need CAP\_SYS\_ADMIN

- bpf\_get\_current\_comm
- bpf\_get\_current\_pid\_tgid
- bpf\_get\_current\_uid\_gid
- bpf\_get\_trace\_printk