writeup.md 10/11/2020

[Binary] mBRrrr Writeup

We are given a file named bootloader.bin. Lets run some basic command:

```
lotus@Kalin:~/Documents/CTF/Reply2020/pwn/mbrrrr$ file bootloader.bin
bootloader.bin: DOS/MBR boot sector
lotus@Kalin:~/Documents/CTF/Reply2020/pwn/mbrrrr$ strings bootloader.bin
f^f]f
fSgf
f[f]f
fUf1
(3B3G
{FLG:This_is_not_the_real_flag_sorry_keep_rev}
```

So we have a boot sector file, and it has a fake flag inside it as well. Let's try to execute it with qemu and see what happens:

• qemu-system-i386 -fda bootloader.bin

It's printing "Ops", let's disassemble it and look at the code:

• ndisasm -b 16 bootloader.bin

Let's also see the structure of the file with xxd:

• xxd bootloader.bin

Executing these commands we can see that this boot sector is divided in two parts, the first is basically loading the second at address 0x8000 in memory and then jumping to instructions at address 0x802a.

We know the second part of our boot sector starts at address 0×200 in our disassemble so effective execution will start at $0 \times 22a$, let's see what instructions are there:

writeup.md 10/11/2020

```
00000200
         6655
                            push ebp
00000202
         6689E5
                           mov ebp,esp
00000205
         6653
                            push ebx
00000207 67668B5508
                           mov edx,[ebp+0×8]
0000020C 66BB07000000
                           mov ebx,0×7
00000212 67660FBE02
                           movsx eax, byte [edx]
00000217
         84C0
                            test al, al
00000219
         7409
                            jz 0×224
0000021B
         80CC0E
                           or ah,0×e
0000021E
         CD10
                            int 0×10
00000220 6642
                            inc edx
00000222 EBEE
                            jmp short 0×212
00000224 665B
                            pop ebx
00000226 665D
                            pop ebp
00000228 66C3
                            retd
0000022A
         6655
                            push ebp
0000022Csh6631D2
                           xor edx,edx
0000022F 6689E5
                           mov ebp,esp
00000232 6683EC08
                          sub esp,byte +0×8
00000236 678A8240810000 mov al,[edx+0×8140]
0000023D 67328200810000
                           xor al,[edx+0×8100]
00000244
                            add al,dl
         00D0
00000246
         3413
                            xor al,0×13
00000248
         00D0
                            add al,dl
0000024A 3419
                           xor al,0×19
0000024C
         00D0
                           add al,dl
                           cmp al,[edx+0×8180]
0000024E
         673A8280810000
00000255
         740C
                            jz 0×263
00000257
         6683EC0C
                            sub esp,byte +0×c
         6668AE810000
0000025B
                            push dword 0×81ae
00000261
         EB12
                            jmp short 0×275
00000263 6642
                            inc edx
                            cmp edx,byte +0×2e
00000265 6683FA2E
00000269
         75CB
                            jnz 0×236
0000026B 6683EC0C
                            sub esp,byte +0×c
0000026F
         666840810000
                            push dword 0×8140
00000275
         66E885FFFFFF
                            call dword 0×200
```

This is the interesting part of this challenge, here we can see that a byte from an address (edx+0x8140) is loaded in all register and then it's xored with some other bytes, one taken from another address (edx+0x8100) and the others are some constants 0x13 and 0x19, then, it's going to be compared to the value stored at edx+0x8180. What is going on here? This is simply a loop that will execute for 0x2e times and after that, if all the compares are successful, we are pushing a value which is the same used before 0x8140 and then calling code at offset 0x200. If compares will fail the value pushed will be 0x81ae but still the call will happen. Looking at these addresses we see that 0x8140 is where the fake flag is stored at while 0x81ae contains the "Ops" string. Now we know what to do, we know that the execution will be successful if we can pass all the compares, so I took the bytes strings from addresses 0x8100 and 0x8140 which are used for xor and cmp and reverse the for loop in a simple python script:

```
#!/bin/python3

flag = b""
string2=bytes.fromhex("c7bb877c20f7f36583b12370ed0283a9c31fd98fe402faf939fe
ddbd0de943489bee622c89102833423347c59555")
test =
bytearray(bytes.fromhex("b6f8fb2c0cc9d752b6c37a852a9ffed040312b08288c11126e
2cdcfd973986d208e1c06c0e8778a4e8b8ca4c1b97"))
```

writeup.md 10/11/2020

```
for i in range(46):
    test[i] = (test[i]-i) % 256
    test[i] = test[i]^0x19
    test[i] = (test[i]-i) % 256
    test[i] = test[i]^0x13
    test[i] = (test[i]-i) % 256
    test[i] = test[i]-string2[i]
print(test.decode())
```

We run the script and get the flag {FLG:18471a01b9b9528273857ee47a19d6710848f568}

Tips

To debug this dynamically we can use qemu with gdb in the following way:

• qemu-system-i386 -s -S -fda bootloader.bin

Inside gdb:

• target remote localhost:1234