

Chapter 1

Library eo

```
Require Import Arith.
Require Import Ascii.
Require Import Bool.
Require Import Coq.Strings.Byte.
Import IFNOTATIONS.

Inductive string : Set :=
| EmptyString : string
| String : ascii → string → string.

Inductive EOExpr : Type :=
| EOUnit : EOExpr
| EOInt : nat → EOExpr
| EOBool : bool → EOExpr
| EOVar : string → EOExpr
| EOAssign : EOExpr → EOExpr → EOExpr
| EOPlus : EOExpr → EOExpr → EOExpr
| EOMinus : EOExpr → EOExpr → EOExpr
| EOMult : EOExpr → EOExpr → EOExpr
| EODiv : EOExpr → EOExpr → EOExpr
| EOIf : EOExpr → EOExpr → EOExpr → EOExpr
| EOLambda : string → EOExpr → EOExpr
| EOClosure : string → EOExpr → EOEnv → EOExpr
| EOApply : EOExpr → EOExpr → EOExpr

with EOEnv : Type :=
| EOEmptyEnv : EOEnv
| EOBind : string → EOExpr → EOEnv → EOEnv.
```

Chapter 2

Library main

Require Import Bool.

Inductive **bool** :=

| true

| false.

Definition negb ($b : \mathbf{bool}$) : **bool** :=

match b with

| true \Rightarrow false

| false \Rightarrow true

end.

Theorem negb_negb: $\forall (b : \mathbf{bool}), \text{negb} (\text{negb } b) = b$.

Proof.

intros b .

destruct b .

+ simpl. reflexivity.

+ simpl. reflexivity.

Qed.

Theorem negb_trice : $\forall (b : \mathbf{bool}), \text{negb} (\text{negb} (\text{negb } b)) = \text{negb } b$.

Proof.

intros b .

destruct b .

+ simpl. reflexivity.

+ simpl. reflexivity.

Qed.

Definition andb ($a \ b : \mathbf{bool}$) : **bool** :=

match a, b with

| true, true \Rightarrow true

| -, - \Rightarrow false

end.

Theorem and_true_both_arg_true : $\forall (a\ b : \mathbf{bool}),$
 $a = \mathbf{true} \rightarrow b = \mathbf{true} \rightarrow \mathbf{andb}\ a\ b = \mathbf{true}.$

Proof.

```
intros a b Ha Hb.
rewrite Ha.
rewrite Hb.
simpl. reflexivity.
```

Qed.

Theorem and_true_otherside : $\forall (a\ b : \mathbf{bool}),$
 $\mathbf{andb}\ a\ b = \mathbf{true} \rightarrow a = \mathbf{true} \wedge b = \mathbf{true}.$

Proof.

```
intros a b Ha.
destruct a.
destruct b.
split. reflexivity.
reflexivity.
split. reflexivity.
simpl in Ha.
inversion Ha.
destruct b.
simpl in Ha.
inversion Ha.
simpl in Ha.
inversion Ha.
```

Qed.

Theorem andb_associative : $\forall (a\ b\ c : \mathbf{bool}),$
 $\mathbf{andb}\ a\ (\mathbf{andb}\ b\ c) = \mathbf{andb}\ (\mathbf{andb}\ a\ b)\ c.$

Proof.

```
intros [] [] [];
reflexivity.
```

Qed.

Theorem andb_comutative : $\forall a\ b, \mathbf{andb}\ a\ b = \mathbf{andb}\ b\ a.$

Proof.

```
intros [] []; simpl; reflexivity.
```

Qed.

Definition orb($a\ b : \mathbf{bool}$) : $\mathbf{bool} :=$
 match a, b with
 | false, false \Rightarrow false
 | -, - \Rightarrow true
 end.

Theorem andb_negb_orb : $\forall (a\ b : \mathbf{bool}),$

negb (andb a b) = orb (negb a) (negb b).
Proof.

intros [] []; simpl; reflexivity.
Qed.

Definition material(*a b* : **bool**) : **bool** :=
match *a,b* with
| false, true ⇒ false
| -, - ⇒ true
end.

Theorem positive_material : $\forall a b : \mathbf{bool}$,
 $a = \text{true} \rightarrow \text{material } a b = \text{true}$.

Proof.
intros *a b* *Ha*.
destruct *a eqn:E*.
- simpl. destruct *b*; reflexivity.
- discriminate *Ha*.
Show Proof.
Qed.

Chapter 3

Library phi

Require Import Relations.Relation_Definitions.

Inductive Process :=

| Nil : Process
| Input : nat → (nat → Process) → Process
| Output : nat → Process → Process
| Sum : Process → Process → Process
| Res : nat → Process → Process.

Definition small_step : relation Process :=

fun p1 p2 ⇒
 match p1 with
 | Input x f ⇒
 ∃ v, p2 = f v
 | Output x p ⇒
 p2 = p
 | Sum p1' p2' ⇒
 p2 = Sum p1' p2' ∨ p2 = Sum p1 p2'
 | Res x p ⇒
 ∃ y, p2 = Res y p ∧ x ≠ y
 | _ ⇒ False
end.

Inductive multi_step : relation Process :=

| refl : ∀ p, multi_step p p
| step : ∀ p1 p2 p3, small_step p1 p2 → multi_step p2 p3 → multi_step p1 p3.

Inductive bisim : relation Process :=

| bisim_refl : ∀ p, bisim p p
| bisim_trans : ∀ p1 p2 p3, bisim p1 p2 → small_step p2 p3 → bisim p1 p3
| bisim_symm : ∀ p1 p2, bisim p1 p2 → bisim p2 p1.

Definition bisimulation (R : relation Process) :=

∀ p q, R p q → (∀ p', small_step p p' → ∃ q', R p' q' ∧ small_step q q') ∧

$$(\forall q', \text{small_step } q \ q' \rightarrow \exists p', R \ p' \ q' \wedge \text{small_step } p \ p').$$

Definition bisimilarity ($R : \text{relation Process}$) :=

$$\forall p \ q, R \ p \ q \rightarrow \exists q', \text{bisim } q \ q' \wedge \text{multi_step } p \ q'.$$