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Lab 1 – Homework Solutions

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A Homework

Homework 1.1. Find the minimum positive number representable in MATLAB/Octave by implementing an ad hoc procedure. Compare with realmin.

Solution Homework 1.1.

```
k = 0;
zero = 1/2;
while (0 + zero) > 0
  zero_old = zero;
  zero = zero / 2;
  k = k + 1;
end

* Remember: realmin is the minimum !normalized! positive floating point number
realmin
* The minimum positive floating point number is instead a !denormalized! number
zero_old % < realmin
k</pre>
```

Homework 1.2. a. Use Taylor polynomial approximation to avoid the loss of significance errors in the following function when x approaches 0

$$f(x) = \frac{1 - \cos(x)}{x^2}$$

b. Reformulate the following function g(x) to avoid the loss of significance error in its evaluation for increasing values of x towards $+\infty$

$$g(x) = x\left(\sqrt{x+1} - \sqrt{x}\right).$$

Solution Homework 1.2.

format long

```
% Use Taylor polynomial approximation to avoid the loss of significance errors in the following function when x approaches 0 % f(x) = \{1-\cos(x)\}/\{x^2\}% The limit of f(x) = \{1-\cos(x)\}/\{x^2\} as x \to 0 is well known: % \lim_{x \to 0} \{x \to 0\} = \{1-\cos(x)\}/\{x^2\} = \{1\}/\{2\} clear all, close all, clc k = [1:30]';
```

```
x = 2.^(-k);
f = @(x) (1 - cos(x))./(x.^2);
[k \times f(x)]
% If we try to evaluate f(x) directly with a sequence xk = 2^{-k} that approaches 0, we
notice anomalous behaviour at k = 13, k = 27, k = ...
% The reason is numerical cancellation of the terms on the numerator. We can use Taylor
expansion of \cos(x) for x \to 0 to obtain a better approximation. Indeed
% \cos(x) = 1 - \{1\}/\{2\} x^2 + \{1\}/\{24\} x^4 - \{1\}/\{720\} x^6 + o(x^8)
% so that
% f(x) = \{1\}/\{2\} - \{1\}/\{24\} x^2 + \{1\}/\{720\} x^4 + o(x^6)
f_{taylor_4} = @(x) 1/2 - x.^2/24 + x.^4/720;
[k \times f(x) f_{taylor_4}(x)]
% The obtained formula is more stable that the direct evaluation of f(x), and the computed
values approach {1}/{2}.
% Reformulate the following function g(x) to avoid the loss of
% significance error in its evaluation for increasing values of x towards +infinity.
% g(x) = x (sqrt\{x+1\} - sqrt\{x\}).
% It is well known that
% \lim\{x \rightarrow +infinity\} \ g(x) = +infinity
clear all, close all, clc
format short e
k = [1:20]';
x = 10.^(k);
g = @(x) x.*(sqrt(x+1) - sqrt(x));
[k x g(x)]
% If we try to evaluate g(x) for a sequence xk = 10^{k} that approaches 0, we notice
anomalous behaviour at k = 16, k = \dots, for which the computed value is 0. This happens since,
 in floating point representation, x16 + 1 = x16!
% Rationalization of the expression in bracket solves this numerical cancellation: indeed,
multiplying the numerator and the denominator by sqrt\{x+1\} + sqrt\{x\}
 \begin{tabular}{ll} \$ & g(x) = \{x \; (sqrt\{x+1\} \; - \; sqrt\{x\}) \; (sqrt\{x+1\} \; + \; sqrt\{x\}) \; ) \; / \; (sqrt\{x+1\} \; + \; \{x\}) \; ) \; = \; \{x \; \} \; / \; (sqrt\{x+1\} \; + \; sqrt\{x\}) \; ) \; / \; (sqrt\{x+1\} \; + \; sqrt\{x\}) \; / \; (sqrt\{x+1\} \; + \; sqrt\{x\}) \; ) \; / \; (sqrt\{x+1\} \; + \; sqrt\{x\}) \; / \; (s
  + sqrt{x})}.
g2 = @(x) x./(sqrt(x+1) + sqrt(x));
[k \times g(x) g2(x)]
% This formula is more stable and does not suffer of numerical cancellation.
```

Homework 1.3. We can compute e^{-x} around x=0 using Taylor polynomials in two ways, either using

$$e^{-x} \approx 1 - x + \frac{1}{2}x^2 - \frac{1}{6}x^3 + \cdots$$

or using

$$e^{-x} = \frac{1}{e^x} \approx \frac{1}{1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \cdots}.$$

Which approach is the most accurate?

Solution Homework 1.3.

```
clear all, close all, clc
format long
f_{taylor_pos} = @(x) 1 - x + 1/2*x.^2 - 1/6*x.^3 + 1/24*x.^4 - 1/120*x.^5;
f_{taylor_neg} = @(x) 1./(1 + x + 1/2*x.^2 + 1/6*x.^3 + 1/24*x.^4 + 1/120*x.^5);
k = [1:20]';
x_pos = 10.^(-k);
[x_pos f_taylor_pos(x_pos) f_taylor_neg(x_pos)]
x_neg = -10.^(-k);
[x_neg f_taylor_pos(x_neg) f_taylor_neg(x_neg)]
% No bad behaviour is observed since the 1 in both formulas dominates the sum and numerical
cancellation does not occur.
% Therefore both expression can be accepted.
% Instead, if the sum were not dominated by the term 1 (or if you didn't notice that),
  the following combination could have been proposed
  e^{-x} \sim \{1\}/\{1 + x + \{1\}/\{2\}x^2 + \{1\}/\{6\}x^3 + ...\}, x > 0;
       \sim 1 - x + \{1\}/\{2\}x^2 - \{1\}/\{6\}x^3 + \dots, x \le 0.
cancellation because all the terms in the sum are positive. Similarly, the choice for x < 0
grants that the terms -x^{2k+1} are positive, again avoiding any numerical cancellation.
```

Homework 1.4. Consider the following integral

$$I_n(\alpha) = \int_0^1 \frac{x^n}{x+\alpha} dx, \quad \forall n \in \mathbb{N}, \alpha > 0.$$

- a. Give an upper bound for $I_n(\alpha)$, $\forall n \in \mathbb{N}, \alpha > 0$.
- b. Prove the following recursive relation between $I_n(\alpha)$ and $I_{n-1}(\alpha)$:

$$\begin{cases} I_n(\alpha) = -\alpha I_{n-1}(\alpha) + \frac{1}{n} \\ I_0(\alpha) = \ln\left(\frac{\alpha+1}{\alpha}\right) \end{cases}$$

- c. Employing the previous relation, compute $I_{40}(\alpha=8)$ and comment the obtained results.
- d. Write a numerically stable recursive relation for $I_{40}(\alpha = 8)$.

Solution Homework 1.4.

```
clear all;
format short e

n = 40;

% alpha = 1/8 not requested in the homework
for (alpha = [1/8, 8])
        I(1) = log((alpha+1)/alpha);
    for (k = 1:n)
        I(k+1) = -alpha*I(k) + 1/k;
```

```
end
      recursion_integral = I(n+1);
      upper_bound = (1/alpha) * (1/(n+1));
      exact_integral = quadl(@(x) (x.^n)./(x+alpha), 0, 1, 1e-16); % Exact value (not requested)
      [recursion_integral, upper_bound, exact_integral]
end
% The recursive approximation of I\{40\} (alpha = 8) is 1.6389 ... 10^{418}: this result is for
sure incorrect because it violates the upper bound for I_{40}(\alpha=8).
\$ This is due to the finite precision we use and the error propagation: the initial value IO(
alpha) = ln(\{alpha+1\}/\{alpha\}) is represented with a certain error eps; denoting by fl(y) the
floating point representation of the number y, we have
% fl(IO(alpha)) = IO(alpha) + eps,
% f(I(alpha)) = -alpha f(I(alpha)) + 1 = -alpha (I(alpha)) + eps) + 1 = I(alpha) - alpha f(I(alpha)) + 1 = I(alpha) + eps) + eps) + I(alpha) + eps) + eps) + I(alpha) + eps) + eps
 eps,
% fl(Ik(alpha)) = In(alpha) + (-1)^k alpha^k eps,
result in an amplification of the error if alpha > 1. Instead, for alpha < 1, the error is
damped and the recursive relation is stable (e.g. see the previous test for alpha = 1/8).
% Now write a numerically stable recursive relation for I\{40\} (alpha = 8).
% The idea is to transform the factor alpha on the RHS of the recursive relation in a factor
{1}/{alpha}; this can be achieved inverting the recursive relation:
           I\{k-1\}\ (alpha) = -\{1\}/\{alpha\}\ I\{k\}\ (alpha) + \{1\}/\{k\ alpha\}
           lim\{k \rightarrow +infinity\} I\{k\} (alpha) = 0
% The final value \lim\{k \rightarrow +infinity\} \ I\{k\}\ (alpha) is equal to zero because 0 <= I\{k\}\ (alpha) \} 
= \{1\}/\{alpha\} \{1\}/\{k+1\} \rightarrow 0, as k \rightarrow + infinity.
clear all;
n = 40:
big = 1000;
% alpha = 1/8 not requested in the homework
for (alpha = [1/8, 8])
     I(big + 1) = 0;
      for (k = big:-1:n+1)
         I(k) = -1/alpha*I(k+1) + 1/(k*alpha);
      end
      recursion_integral = I(n+1);
      upper_bound = (1/alpha)*(1/(n+1));
      exact_integral = quadl(@(x) (x.^n)./(x+alpha), 0, 1, 1e-16); % Exact value (not requested)
      [recursion_integral, upper_bound, exact_integral]
end
% The recursion is now stable for alpha > 1 (and note that, now, an additional error is
present, because the final condition can be set for an arbitrary large k = k* instead of k = +
infinity).
```

Homework 1.5. Given the following sequence:

$$\begin{cases} x_{n+1} = 2^{n+1} \left[\sqrt{1 + \frac{x_n}{2^n}} - 1 \right] \\ x_0 > -1 \end{cases}$$

for which $\lim_{n \to +\infty} x_n = \ln(1 + x_0)$.

- a. Set $x_0 = 1$, compute x_1, x_2, \ldots, x_{71} and explain the obtained results.
- b. Transform the sequence in an equivalent one that converges to the theoretical limit.

Solution Homework 1.5.

```
% Set x0 = 1, compute x1, x2, ..., x\{71\} and explain the obtained results.
clear all, close all, clc
format long
n_max = 71;
x = zeros(n max+1, 1);
x(1) = 1; % x0 set to 1
for n = 0:n_{max-1}
 x(n+2) = 2^{(n+1)} * (sqrt(1 + x(n+1)/2^{(n)}) - 1); % Pay attention to the indexing!!!
x (end)
x_{\lim} = \log(1 + x(1))
figure
hold on, box on
plot([0:71], x, 'bx-', 'Linewidth',3)
\verb"plot([0:71], x_lim*ones(n_max+1,1)", "r-", "Linewidth",3)"
axis([-1 72 -0.05 1.05])
set(gca,'LineWidth',2)
set (gca, 'FontSize', 16)
% When computing the sequence with Octave/MATLAB we find that x{71} differs a lot from ln(1+
x0) = ln(2). The plots also shows that xn = 0, for all n \ge n* = 52: this is due to numerical
cancellation effects, because at n = n* it holds \{x\{n*\}\}/\{2^nn*\}\} < eps, hence in finite
precision arithmetic x\{n*\} = 0!
x (53)
% The value of n* can also be computed with the following argument: since (hopefully)
numerical cancellation will occur for large n
% \quad x\{n\} \sim \ln(1+x0)
  therefore
% \{x\{n\}\}/\{2^{n}\}\} < eps is approximately equivalent to \{\ln(1+x0)\}/\{2^{n}\}\} < eps
% which can be solved for n, finding
n > log2({ln(1+x0)})/{eps}
x 0 = 1;
n = log(log(1 + x_0)/eps)/log(2)
% Transform the sequence in an equivalent one that converges to the theoretical limit.
% After a razionalization, the recurrence can be written as
  x\{n+1\} = 2^{n+1} [ sqrt\{1+\{xn\}/\{2^n\}\} - 1 ] = \{2 xn\}\{ sqrt\{1+\{xn\}/\{2^n\}\} + 1\}
clc, clear all
n_max = 71;
x = zeros(n_max+1, 1);
x(1) = 1; % x0 set to 1
for n = 0:n_max-1
 x(n+2) = 2*x(n+1)/(sqrt(1 + x(n+1)/2^n(n)) + 1);
x (end)
% The error in this case is
x(end) - log(1 + x(1))
figure
hold on, box on
plot([0:n_max], x,'bx-','Linewidth',3)
```