SQL Aggregate Functions

Solve the following SQL queries using aggregate functions to calculate the following from the Employee table:

- a. Find the total salary of all employees.
- b. Find the average salary of employees in the IT department.
- c. Find the department with the highest total salary.
- d. Find the total salary paid in each department.

Here are the SQL queries using **aggregate functions** to answer each part based on an Employee table that includes columns like Salary, Department, etc.:

a. Find the total salary of all employees:

```
SELECT SUM(Salary) AS TotalSalary
FROM Employee;
```

b. Find the average salary of employees in the IT department:

```
sql

SELECT AVG(Salary) AS AverageITSalary
FROM Employee
WHERE Department = 'IT';
```

c. Find the department with the highest total salary:

```
sql
```

```
SELECT Department, SUM(Salary) AS TotalDeptSalary
FROM Employee
GROUP BY Department
ORDER BY TotalDeptSalary DESC
LIMIT 1;
```

d. Find the total salary paid in each department:

```
SELECT Department, SUM(Salary) AS TotalSalary
FROM Employee
GROUP BY Department;
```

Let me know if you want to include employee names, handle nulls, or filter further!

Solve the following SQL queries using date functions to display the following:

- a. Get the current system date.
- b. Display employees' names and the number of years they have worked.
- c. Find employees hired on a weekend (Saturday or Sunday).
- d. Find the employee who has been working the longest

Here are the SQL queries using **date functions** to solve each part from the Employee table (assuming there are columns like Name and HireDate):

a. Get the current system date:

```
sql
```

```
SELECT CURRENT_DATE AS SystemDate;
```

In some systems, you can also use GETDATE() or SYSDATE.

b. Display employees' names and the number of years they have worked:

```
SELECT Name,

TIMESTAMPDIFF(YEAR, HireDate, CURDATE()) AS YearsWorked

FROM Employee;
```

If you're using Oracle, use:

```
SELECT Name,

FLOOR(MONTHS_BETWEEN(SYSDATE, HireDate) / 12) AS YearsWorked

FROM Employee;
```

c. Find employees hired on a weekend (Saturday or Sunday):

```
SELECT Name, HireDate
FROM Employee
WHERE DAYOFWEEK(HireDate) IN (1, 7);

1 = Sunday, 7 = Saturday in MySQL.
In Oracle, use:
```

```
WHERE TO_CHAR(HireDate, 'DY') IN ('SAT', 'SUN');
```

d. Find the employee who has been working the longest:

```
SELECT Name, HireDate
FROM Employee
ORDER BY HireDate
LIMIT 1;
```

In SQL Server, use TOP 1, and in Oracle, use FETCH FIRST 1 ROWS ONLY.

Let me know your SQL dialect (MySQL, Oracle, SQL Server, etc.) if you want exact compatibility!

Solve the following SQL queries on Employee & Department Schemas

- a. To Perform a LEFT JOIN between the Employee and Department tables to display: Employee name, Department name & Salary. All employees should be displayed, with NULL for those without a department.
- b. To select Only employees who are assigned a department should be displayed(Inner Join)
- c. Convert all employee names to uppercase
- d. Find employees whose name contains 'o' as the second letter

Here are the SQL queries for each part based on **Employee** and **Department** tables:

a. LEFT JOIN: Show all employees with department name (NULL if no department)

```
SELECT E.Name AS EmployeeName,
D.DeptName AS DepartmentName,
E.Salary
FROM Employee E
LEFT JOIN Department D
ON E.DeptID = D.DeptID;
```

b. INNER JOIN: Show only employees who are assigned a department

```
SELECT E.Name AS EmployeeName,
D.DeptName AS DepartmentName,
E.Salary
FROM Employee E
INNER JOIN Department D
ON E.DeptID = D.DeptID;
```

c. Convert all employee names to uppercase

```
SELECT UPPER(Name) AS UpperCaseName
FROM Employee;
```

d. Find employees whose name contains 'o' as the second letter

```
SELECT Name
FROM Employee
WHERE Name LIKE '_o%';

_ represents one character, and % means any sequence of characters.
```

Let me know if you want to add sorting, filtering, or use aliases for better readability!

Solve the following SQL queries
a. Write a SQL query using the IN operator to Find
employees who work in departments with IDs 10, 20, or 30.
b. Write a SQL query using the NOT IN operator to Find
Employees Who Are NOT in HR and IT Departments.

Here are the SQL queries for your requirements using the IN and NOT IN operators:

a. Find employees who work in departments with IDs 10, 20, or 30:

```
SELECT *
FROM Employee
WHERE DeptID IN (10, 20, 30);
```

b. Find employees who are NOT in HR and IT departments

(Assuming you have to check based on department **names**, and the **Department** table has columns like **DeptID** and **DeptName**):

```
sql
```

```
SELECT *
FROM Employee
WHERE DeptID NOT IN (
    SELECT DeptID
    FROM Department
    WHERE DeptName IN ('HR', 'IT')
);
```

If DeptName is directly in the Employee table, then simply:

```
SELECT *
FROM Employee
WHERE DeptName NOT IN ('HR', 'IT');
```

Let me know if you'd like to see this joined with department names for clarity!

Solve the following SQL queries
a. Write a SQL query using the EXISTS operator to Find
employees who are not assigned to any department.
b. Write a SQL query using the ANY operator to Find
employees earning more than at least one employee in
department 10.

Here are the SQL queries using the EXISTS and ANY operators:

a. Find employees who are not assigned to any department using EXISTS:

```
SELECT *
FROM Employee E
WHERE NOT EXISTS (
```

```
SELECT 1
FROM Department D
WHERE E.DeptID = D.DeptID
);
```

Alternatively, if you just want to check for unassigned employees (i.e., <code>DeptID IS NULL</code>), this is simpler:

```
SELECT *
FROM Employee
WHERE DeptID IS NULL;
```

b. Find employees earning more than at least one employee in department 10 using ANY:

```
SELECT *
FROM Employee
WHERE Salary > ANY (
    SELECT Salary
    FROM Employee
    WHERE DeptID = 10
);
```

This will return employees whose salary is greater than the lowest salary in Dept 10.

Let me know if you want to adjust these queries to use JOINs or filter further!

Consider the following un normalized table:
StudentID StudentName Course Instructor
InstructorDept
101 Alice DBMS, OS Dr. Smith, Dr. John CS, IT

102 Bob DBMS Dr. Smith CS 103 Charlie OS Dr. John IT

- a. Identify the anomalies in the table.
- b. Normalize the table to 1NF, 2NF, 3NF, and BCNF by creating proper relations.
- c. Explain how each normalization step resolves anomalies.

Let's walk through the normalization process for the given **unnormalized table** and answer each part clearly:

Given Unnormalized Table:

| StudentID | StudentName | Course | Instructor | InstructorDept |
|-----------|-------------|----------|---------------------|----------------|
| 101 | Alice | DBMS, OS | Dr. Smith, Dr. John | CS, IT |
| 102 | Bob | DBMS | Dr. Smith | CS |
| 103 | Charlie | OS | Dr. John | IT |

a. Anomalies in the Table:

1. Insertion Anomaly:

Can't insert a new instructor or course unless a student takes it.

2. Update Anomaly:

• If Dr. John changes department from IT to AI, we must update all rows where he appears—risk of inconsistent data.

3. **Deletion Anomaly**:

• If Alice is deleted, and she's the only one taking "OS" with Dr. John, we lose all info about that course and instructor.

4. Redundancy:

• Instructor names and departments repeat for each student-course combo.

5. Multivalued Attributes:

Course , Instructor , and InstructorDept contain multiple values in a single field
 —this violates 1NF.

b. Normalization Steps

1NF (First Normal Form)

➡ Eliminate multivalued attributes by creating one row per atomic value:

| StudentID | StudentName | Course | Instructor | InstructorDept |
|-----------|-------------|--------|------------|----------------|
| 101 | Alice | DBMS | Dr. Smith | CS |
| 101 | Alice | OS | Dr. John | IT |
| 102 | Bob | DBMS | Dr. Smith | CS |
| 103 | Charlie | OS | Dr. John | IT |

2NF (Second Normal Form)

- → Remove partial dependencies (i.e., non-prime attributes depending only on part of the composite key).
- → Assume the composite primary key is (StudentID, Course).

Split into two tables:

StudentCourse Table:

| StudentID | Course |
|-----------|--------|
| 101 | DBMS |
| 101 | OS |
| 102 | DBMS |
| 103 | OS |

Student Table:

| StudentID | StudentName |
|-----------|-------------|
| 101 | Alice |
| 102 | Bob |
| 103 | Charlie |

CourseInstructor Table:

| Course | Instructor | InstructorDept |
|--------|------------|----------------|
| DBMS | Dr. Smith | CS |
| OS | Dr. John | IT |

3NF (Third Normal Form)

→ Remove transitive dependencies (non-prime attribute depending on another non-prime attribute).

Here, InstructorDept depends on Instructor, not directly on the Course.

Split further:

CourseInstructor Table:

| Course | Instructor |
|--------|------------|
| DBMS | Dr. Smith |
| OS | Dr. John |

Instructor Table:

| Instructor | InstructorDept |
|------------|----------------|
| Dr. Smith | CS |
| Dr. John | IT |

BCNF (Boyce-Codd Normal Form)

→ All determinants must be candidate keys.

In **3NF**, if each course has only one instructor (1:1), then **BCNF = 3NF** here.

If the same course can have multiple instructors, we'd need to consider a more complex key, but for now, this design satisfies BCNF.

c. How Each Step Resolves Anomalies:

| Normal Form | Fixes These Anomalies |
|-------------|--|
| 1NF | Removes multivalued fields (atomicity) |
| 2NF | Removes partial dependency (avoids duplication of student names) |
| 3NF | Removes transitive dependency (separates instructor and department info) |
| BCNF | Ensures all dependencies are on candidate keys (eliminates redundancy further) |

Would you like these tables converted into SQL CREATE TABLE and INSERT statements too?

Given a relation R (A, B, C, D) and Functional Dependency set FD = {AB \rightarrow CD, B \rightarrow C}, determine whether the given R is in 2NF? If not convert it into 2 NF.

Let's analyze whether the relation R(A, B, C, D) with $FD = \{AB \rightarrow CD, B \rightarrow C\}$ is in 2NF, and convert it if not.

Step 1: Check for 2NF

What is 2NF?

A relation is in **2NF** if:

- 1. It is in 1NF (i.e., no multi-valued or composite attributes), and
- 2. No partial dependency exists i.e., no non-prime attribute depends on a part of a candidate key.
- Identify candidate keys:
- Given AB → CD (closure of AB includes all attributes), so AB is a candidate key.

Now check for any partial dependencies:

- FD: B → C
 - Here, B is part of candidate key AB, and C is a non-prime attribute.
 - **This is a partial dependency**, which violates 2NF.
- \times Conclusion: R is NOT in 2NF because of partial dependency B \rightarrow C.

Step 2: Convert R to 2NF

We need to remove the partial dependency $B \rightarrow C$.

- Decompose into two relations:
- 1. R1(B, C) to capture the partial dependency $B \rightarrow C$
- 2. **R2(A, B, D)** remove **C** from the original relation

Final 2NF Relations:

R1(B, C)

FD: $B \rightarrow C$

R2(A, B, D)

FD: $AB \rightarrow D$ (after removing C , $AB \rightarrow D$ still holds)

✓ Now, both R1 and R2 are in 2NF — no partial dependencies remain.

Let me know if you'd like to check 3NF or write SQL CREATE TABLE statements for this!

Assume a scenario where salary updates are being made in the Employee table. Solve the following SQL Queries:

- a. BEGIN TRANSACTION to start a transaction.
- b. SAVEPOINT to mark a point in the transaction.
- c. ROLLBACK to undo changes in case of an error.
- d. COMMIT to finalize the changes.

Here's how you can use transaction control in SQL for salary updates in the Employee table:

a. Start the transaction

```
BEGIN TRANSACTION;
-- or simply:
START TRANSACTION;
```

• b. Set a savepoint

```
sql
SAVEPOINT BeforeSalaryUpdate;
```

• c. Rollback to savepoint (if needed)

```
sql

ROLLBACK TO BeforeSalaryUpdate;
```

This undoes changes after the savepoint, but keeps the transaction open.

• d. Commit the transaction

```
sql
COMMIT;
```

This finalizes all changes made since the transaction began.

Example Scenario

```
BEGIN TRANSACTION;

UPDATE Employee
SET Salary = Salary + 1000
WHERE DeptID = 10;

SAVEPOINT BeforeSalaryUpdate;

-- Suppose an error occurs in the next update:
UPDATE Employee
SET Salary = Salary + 500
WHERE DeptID = 'XYZ'; -- Invalid DeptID

-- Rollback the second update
ROLLBACK TO BeforeSalaryUpdate;

-- Finalize the first update
COMMIT;
```

Let me know if you want this with error handling in stored procedures or PL/SQL!

- a. Create a stored Procedure to Retrieve Student Details by ID.
- b. Create a Procedure to Retrieve Highest and Lowest Marks.

Sure! Here's how you can write stored procedures in SQL (assuming you're using **MySQL** syntax; let me know if you're using Oracle, SQL Server, etc.).

a. Stored Procedure to Retrieve Student Details by ID

```
sql

DELIMITER //

CREATE PROCEDURE GetStudentByID(IN student_id INT)
```

```
BEGIN
    SELECT *
    FROM Student
    WHERE StudentID = student_id;
END //
DELIMITER;
```

Call the procedure like:

```
sql

CALL GetStudentByID(101);
```

b. Stored Procedure to Retrieve Highest and Lowest Marks

```
sql

DELIMITER //

CREATE PROCEDURE GetHighestAndLowestMarks()
BEGIN
    SELECT
        MAX(Marks) AS HighestMark,
        MIN(Marks) AS LowestMark
    FROM Student;
END //
DELIMITER ;
```

Call the procedure like:

```
sql

CALL GetHighestAndLowestMarks();
```

Let me know if your table has a different structure (like a separate Marks table), or if you want to return the student names along with the highest/lowest marks!

- a. On schema Employees(Empid,Name,Department,salary)
 Create a View which displays IT department
 employees whose salary is greater than 60000
 b. Create a View which displays IT department
 information
- c. Create trigger to ensure that employee records are not permanently lost by copying them to a backup table before deletion.

Here's how you can write the required **SQL view and trigger statements** based on the schema:

a. Create a view to display IT department employees with salary > 60000

```
sql

CREATE VIEW IT_HighEarners AS

SELECT Empid, Name, Department, Salary
FROM Employees
WHERE Department = 'IT' AND Salary > 60000;

✓ Use:

sql

SELECT * FROM IT_HighEarners;
```

b. Create a view to display only IT department information

```
sql
```

```
CREATE VIEW IT_Department AS

SELECT *

FROM Employees

WHERE Department = 'IT';

V Use:

sql

SELECT * FROM IT_Department;
```

c. Create a trigger to back up employee records before deletion

First, create a backup table:

```
CREATE TABLE Employees_Backup (
    Empid INT,
    Name VARCHAR(100),
    Department VARCHAR(100),
    Salary DECIMAL(10,2),
    DeletedAt TIMESTAMP DEFAULT CURRENT_TIMESTAMP
);
```

Then, create the trigger:

```
DELIMITER //

CREATE TRIGGER BeforeEmployeeDelete

BEFORE DELETE ON Employees

FOR EACH ROW

BEGIN

INSERT INTO Employees_Backup (Empid, Name, Department, Salary)

VALUES (OLD.Empid, OLD.Name, OLD.Department, OLD.Salary);

END //
```

DELIMITER ;

This trigger automatically copies the deleted record to Employees_Backup before it's removed from Employees.

Let me know if you want similar triggers for updates or other departments too!