

CSIT 113 (Problem Solving)

Assignment 1

Name: Mohamed Haneefa Jiyavudeen

Tutorial group: T03F

UOW number: 8496407

Question 1)

- a) $4! = 4 \times 3 \times 2 \times 1 = 24$
- b) $5! = 5 \times 4 \times 3 \times 2 \times 1 = 120$
- c) $6! = 6 \times 5 \times 4 \times 3 \times 2 \times 1 = 720$
- d) $n! = n \times (n-1) \times (n-2) \times \dots \times 3 \times 2 \times 1$

Question 2)

- a) Number of total possible outcomes (without constraint)
 $= 4^4$
 $= 256$
- b) Number of total possible outcomes (with constraint)
 $= 4!$
 $= 4 \times 3 \times 2 \times 1$
 24
- c) Legend:
 Freestyle → 1
 Breaststroke → 2
 Butterfly → 3
 Backstroke → 4

Andrew	1	1	1	1	1	1	2	2	2	2	2	2	3	3	3	3	3	3	4	4	4	4	4	4
Bernard	2	3	2	4	4	3	1	3	1	4	4	3	1	2	1	4	4	2	1	2	1	3	3	2
Chester	3	2	4	2	3	4	3	1	4	1	3	4	2	1	4	1	2	4	2	1	3	1	2	3
Daniel	4	4	3	3	2	2	4	4	3	3	1	1	4	4	2	2	1	1	3	3	2	2	1	1
Outcomes	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24

d) All combinations

1 – 183	7 – 177	13 – 173	19 – 172
2 – 177	8 – 174	14 – 176	20 – 175
3 – 182	9 – 176	15 – 172	21 – 172
4 – 174	10 – 171	16 – 167	22 – 169
5 – 174	11 – 176	17 – 172	23 – 174
6 – 176	12 – 178	18 – 180	24 – 180

Best combination for optimal time is:

Andrew – Butterfly (3)

Bernard – Backstroke (4)

Chester – Freestyle (1)

Daniel – Breaststroke (2)

Question 3)

1. $\{A|B|C|D\}\{C|D|A|B, \rightarrow | \}\{C|D| |A|B\}$: Initially all at left bank, then, **Alice** and **Bob** row over to right bank, now, **Carol** and **Dave** at left bank and **Alice** and **Bob** at right bank
2. $\{C|D| |A|B\}\{C|D|A, \leftarrow |B\}\{A|C|D| |B\}$: **Carol** and **Dave** at left bank, **Alice** row from right bank to left bank, and now **Alice**, **Carol** and **Dave** at left bank and **Bob** at right bank
3. $\{A|C|D| |B\}\{D|A|C, \rightarrow |B\}\{D| |A|B|C\}$: **Alice**, **Carol** and **Dave** at left bank, **Alice** and **Carol** row to right bank, and now **Alice**, **Bob** and **Carol** at right bank while **Dave** is at left bank.
4. $\{D| |A|B|C\}\{D|A, \leftarrow |B|C\}\{A|D| |B|C\}$: **Dave** is at left bank, **Alice** row back to left bank, and now, **Alice** and **Dave** at left bank, while **Bob** and **Carol** at right bank.
5. $\{A|D| |B|C\}\{A|D, \rightarrow |B|C\}\{A| |B|C|D\}$: **Alice** at **Dave** at left bank, then **Dave** row to right bank leaving **Alice** at left bank and **Bob**, **Carol** and **Dave** at right bank.
6. $\{A| |B|C|D\}\{A|C, \leftarrow |B|D\}\{A|C| |B|D\}$: Since **Alice** is at left bank, **Carol** has to row back from right bank to left bank. Now **Alice** and **Carol** at left bank, **Bob** and **Dave** at right bank.
7. $\{A|C| |B|D\}\{A|C, \rightarrow |B|D\}\{ | |A|B|C|D\}$: **Alice** and **Carol** row from left bank to right bank and now everyone has crossed over to right bank safely. **Alice**, **Bob**, **Carol** and **Dave** are at right bank and no one is at left bank.

Minimum number of rows is 7.

Question 4)

$$p \wedge (q \equiv r) \equiv ((p \wedge q) \equiv (p \wedge r))$$

p	q	r	(q \equiv r)	p \wedge (q \equiv r)	p \wedge q	p \wedge r	((p \wedge q) \equiv (p \wedge r))	p \wedge (q \equiv r) \equiv ((p \wedge q) \equiv (p \wedge r))
T	T	T	T	T	T	T	T	T
T	T	F	F	F	T	F	F	T
T	F	T	F	F	F	T	F	T
F	T	T	T	F	F	F	T	F
F	F	F	T	F	F	F	T	F
F	F	T	F	F	F	F	T	F
F	T	F	F	F	F	F	T	F
T	F	F	T	T	F	F	T	T

In this case, $p \wedge (q \equiv r) \equiv ((p \wedge q) \equiv (p \wedge r))$ is only true under 4 scenarios and is false under 4 scenarios. Hence $p \wedge (q \equiv r) \equiv ((p \wedge q) \equiv (p \wedge r))$ is not always true.

Question 5)

$$(p \Rightarrow q) \wedge (q \Rightarrow r) \Rightarrow (p \Rightarrow r)$$

p	q	r	p \Rightarrow q	q \Rightarrow r	p \Rightarrow r	(q \Rightarrow r) \Rightarrow (p \Rightarrow r)	(p \Rightarrow q) \wedge (q \Rightarrow r) \Rightarrow (p \Rightarrow r)
T	T	T	T	T	T	T	T
T	T	F	T	F	F	T	T
T	F	T	F	T	T	T	T
F	T	T	T	T	T	T	T
F	F	F	T	T	T	T	T
F	F	T	T	T	T	T	T
F	T	F	T	F	T	T	T
T	F	F	F	T	F	F	T

Hence it is proven that the formula is always true for $(p \Rightarrow q) \wedge (q \Rightarrow r) \Rightarrow (p \Rightarrow r)$

Question 6)

- a) Given, $A \equiv (B \equiv C)$ and $B \equiv \sim C$
Substitute $B \equiv \sim C$ into $A \equiv (B \equiv C)$

$$A \equiv (\sim C \equiv C)$$

Since $(\sim C \equiv C)$, C is not true.

If C is not true, then C is a knave, hence proves that Bob is a knight.

And Alice says that Bob and Carol are the same type.

Since we know Bob is a knight and Carol is a knave, we can conclude that **Alice is a knave.**

b)

1. Let J represent the proposition "John is a knight".
2. Let S represent the proposition "Steven is a knight".
3. Let Q represent the proposition "The question to ask Steven to identify exactly what John is"

$$(S \equiv Q) \equiv J$$

$$Q \equiv (S \equiv J)$$

Therefore, the question to ask Steven is "Is John a knight?"

If Steven say yes, it proves that John is a knight as we have assumed Steven is saying the truth. If Steven say no, then we know that John is a knave assuming that Steven is saying the truth.

Question 7)

Recursive Algorithm for

```
SumOfR(A,n) {  
    for n ≥ 1  
    if n = 1  
        sum = 1/A[0]  
    else  
        sum = SumOfR(A,n-1) + 1/A[n-1]
```

Question 8)

$$8) 1^3 + 2^3 + 3^3 + \dots + n^3 = \frac{[n^2(n+1)^2]}{4}$$

$$\text{Base step: } 1^3 = \frac{(1^2(1+1)^2)}{4}$$

$$1 = 1$$

LHS = RHS, hence the formula holds for base step,

Inductive step

Assume formula holds for $n = k$

$$1^3 + 2^3 + 3^3 + \dots + k^3 = \frac{[k^2(k+1)^2]}{4}$$

for $n = k+1$,

$$\text{LHS} = 1^3 + 2^3 + 3^3 + \dots + k^3 + (k+1)^3$$

$$= \frac{k^2(k+1)^2}{4} + (k+1)^3$$

$$= \frac{k^2(k+1)^2 + 4(k+1)^3}{4} = \frac{(k+1)^2(k^2 + 4(k+1))}{4}$$

$$= \frac{(k+1)^2(k^2 + 4k + 4)}{4}$$

$$= \frac{[(k+1)^2(k+2)^2]}{4}$$

for $n = k+1$,

$$\text{RHS} = \frac{[(k+1)^2(k+1+1)^2]}{4}$$

$$= \frac{[(k+1)^2(k+2)^2]}{4}$$

\therefore LHS = RHS as proven above, for all $n \geq 1$
the formula holds.