

## K smallest edit distances

Let  $lev_{a,b}(i, j)$  be the set of  $k$  smallest edit distances between the first  $i$  characters of the string  $a$  and the first  $j$  characters of the string  $b$ .

Let's introduce the following sets:

$I_{a,b}(i, j) := \{d + 1 \mid d \in lev_{a,b}(i - 1, j)\}$	The set of $k$ edit distances obtained through insertions
$D_{a,b}(i, j) := \{d + 1 \mid d \in lev_{a,b}(i, j - 1)\}$	The set of $k$ edit distances obtained through deletions
$S_{a,b}(i, j) := \{d + 1_{(a_i \neq b_j)} \mid d \in lev_{a,b}(i - 1, j - 1)\}$	The set of $k$ edit distances obtained through substitutions
$E_{a,b}(i, j) := I_{a,b}(i, j) \cup D_{a,b}(i, j) \cup S_{a,b}(i, j)$	The set of all $3k$ edit distances

Where  $1_{(a_i \neq b_j)}$  is the indicator function equal to 0 when  $a_i = b_j$  and equal to 1 otherwise.

Let's consider the set  $E_{a,b,k}(i, j) \subseteq E_{a,b}(i, j)$ , such that  $|E_{a,b,k}(i, j)| = k$  (in case if  $|E_{a,b}(i, j)| < k$  then  $E_{a,b,k}(i, j) = E_{a,b}(i, j)$ ), and for every  $x \in E_{a,b,k}(i, j)$  and for every  $y \in E_{a,b}(i, j) \setminus E_{a,b,k}(i, j)$  it follows that  $x < y$ .

Then the set of  $k$  smallest edit distances between the first  $i$  characters of the string  $a$  and the first  $j$  characters of the string  $b$  is defined as follows:  $lev_{a,b}(i, j) := E_{a,b,k}(i, j)$ .

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**Proof of correctness:** (proof by the smallest counterexample)

For the sake of convenience let's introduce an auxiliary notation:  $min_k(X)$  for some set  $X$  with natural numbers, and for some  $k \in \mathbb{N}$ , which means that  $min_k(X) \subseteq X$  and  $\forall a \in min_k(X), \forall b \in X \setminus min_k(X)$  it follows, that  $a < b$ , and  $|min_k(X)| = k$  (in case if  $|X| < k$ , then  $min_k(X) = X$ ).

*Induction Basis:*

*Induction Hypothesis:*

- The set  $lev_{a,b}(i - 1, j)$  contains the  $k$  smallest edit distances between the first  $i - 1$  characters of the string  $a$  and the first  $j$  characters of the string  $b$ .
- The set  $lev_{a,b}(i, j - 1)$  contains the  $k$  smallest edit distances between the first  $i$  characters of the string  $a$  and the first  $j - 1$  characters of the string  $b$ .
- The set  $lev_{a,b}(i - 1, j - 1)$  contains the  $k$  smallest edit distances between the first  $i - 1$  characters of the string  $a$  and the first  $j - 1$  characters of the string  $b$ .

*Inductive Step:*

We want to show, that the *Induction Hypothesis* implies, that the set  $lev_{a,b}(i, j)$  contains the  $k$  smallest edit distances between the first  $i$  characters of the string  $a$  and the first  $j$  characters of the string.

**For the sake of contradictions** let's assume, that there exists an edit distance  $y \notin lev_{a,b}(i, j)$  between the first  $i$  characters of the string  $a$  and the first  $j$  characters of the string, such that  $\exists x \in lev_{a,b}(i, j)$  for which  $y < x$  (thus the set  $lev_{a,b}(i, j)$  doesn't contain the  $k$  smallest edit distances):

$$\exists y \notin lev_{a,b}(i, j) \wedge \exists x \in lev_{a,b}(i, j) : y < x$$

Let's expand the expression in a following way:

$$\begin{aligned}
& \exists y \notin lev_{a,b}(i, j) \wedge \exists x \in lev_{a,b}(i, j) : y < x \Leftrightarrow && \text{By definition of } lev_{a,b}(i, j) \\
& \Leftrightarrow \exists y \notin min_k \left( I_{a,b}(i, j) \cup D_{a,b}(i, j) \cup S_{a,b}(i, j) \right) \wedge \exists x \in lev_{a,b}(i, j) : y < x \Rightarrow && \text{As far as } y < x \\
& \Rightarrow \exists y \notin I_{a,b}(i, j) \cup D_{a,b}(i, j) \cup S_{a,b}(i, j) \wedge \exists x \in lev_{a,b}(i, j) : y < x \Leftrightarrow && \text{By definition of } I_{a,b}(i, j), D_{a,b}(i, j), S_{a,b}(i, j) \\
& \Leftrightarrow \exists y : \left( (y-1) \notin lev_{a,b}(i-1, j) \right) \wedge \left( (y-1) \notin lev_{a,b}(i, j-1) \right) \wedge \\
& \quad \wedge \left( (y-1_{(a_i \neq b_j)}) \notin lev_{a,b}(i-1, j-1) \right) \wedge \exists x \in lev_{a,b}(i, j) : y < x
\end{aligned}$$

According to the *Induction Hypothesis* - the set  $lev_{a,b}(i-1, j)$  is the set of the  $k$  **smallest edit distances**, thus from  $(y-1) \notin lev_{a,b}(i-1, j)$  it follows, that  $\forall a \in lev_{a,b}(i-1, j) \Rightarrow (y-1) > a$  (and the same logic is applicable to the sets  $lev_{a,b}(i, j-1)$  and  $lev_{a,b}(i-1, j-1)$ ):

$$\begin{aligned}
& \exists y : \left( (y-1) \notin lev_{a,b}(i-1, j) \right) \wedge \left( (y-1) \notin lev_{a,b}(i, j-1) \right) \wedge \\
& \quad \wedge \left( (y-1_{(a_i \neq b_j)}) \notin lev_{a,b}(i-1, j-1) \right) \wedge \exists x \in lev_{a,b}(i, j) : y < x \Rightarrow \\
& \Rightarrow \exists y : \left( \forall a \in lev_{a,b}(i-1, j) \Rightarrow y > a+1 \right) \wedge \left( \forall b \in lev_{a,b}(i, j-1) \Rightarrow y > b+1 \right) \wedge \\
& \quad \wedge \left( \forall c \in lev_{a,b}(i-1, j-1) \Rightarrow y > c+1_{(a_i \neq b_j)} \right) \wedge \exists x \in lev_{a,b}(i, j) : y < x \Rightarrow \\
& \Rightarrow \exists y : \left( \forall x \in lev_{a,b}(i, j) \Rightarrow y > x \right) \wedge \exists x \in lev_{a,b}(i, j) : y < x \Rightarrow \\
& \Rightarrow \text{Contradiction.}
\end{aligned}$$

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