

Problem Set II

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Introduction to Analysis of Algorithms

Problem 1

$\frac{2}{N}, 12, \sqrt{N}, N, N \log \log N, N \log N, N \log N^2,$
 $N \log^2 N, N^{1.5}, N^2, N^2 \log N, N^3, 2^{\frac{N}{2}}, 2^N$

$N^{1.5}, N^2$, and N^3 grow polynomially.

$2^{\frac{N}{2}}$ and 2^N grow exponentially.

$N \log N$ and $N \log N^2$ grow linearithmically

Problem 2

A. $O(2^{2^N})$

B. $O(\log_2(\log_2 D))$

C. $O(\log_2 \log_2 D)$

Problem 3

$$3 \text{ min} = 1.8 \times 10^8 \mu s$$

A. linear $O(N)$

$$C(100) = 700 \rightarrow C = 7$$

$$2.7 \times 10^7 \text{ items}$$

B. linearithmic $O(N \log N)$

$$C(100) \log_2(100) = 700 \rightarrow C = \frac{7}{2 \log_2 10}$$

$$\frac{7}{2 \log_2 10} x \log_2 x = 1.8 \times 10^8$$

Using Wolfram-Alpha to solve:

$$7.48 \times 10^6 \text{ items}$$

C. quadratic $O(N^2)$

$$C(100^2) = 700 \rightarrow C = 0.07$$

$$= 5.07 \times 10^4 \text{ items}$$

D. cubic $O(N^3)$

$$C(100^3) = 700 \rightarrow C = 0.0007$$

$$6.40 \times 10^3 \text{ items}$$

E. exponential $O(2^N)$

$$C(2^{100}) = 1.8 \times 10^8 \rightarrow C = 5.52 \times 10^{-28}$$

$$1.18 \times 10^1 \text{ items}$$

Problem 4

A. `Bar` is more guaranteed to run faster for values of $N < 100$

B. `Foo` is more guaranteed to run faster for values of $N > 100000$

C. $221N \log_2 N = 3N^2 \Rightarrow N = 1$

D. `Bar` can run faster than `Foo` if the input size is sufficiently small enough

Run Time Calculations

Problem 1

1. Fragment 1
 $O(N)$

```
void frag1(size_t n) {  
    size_t sum = 0;  
    for (size_t i = 0; i < n; ++i)  
        sum = sum + 1;  
}
```

Figure 1: Fragment 1

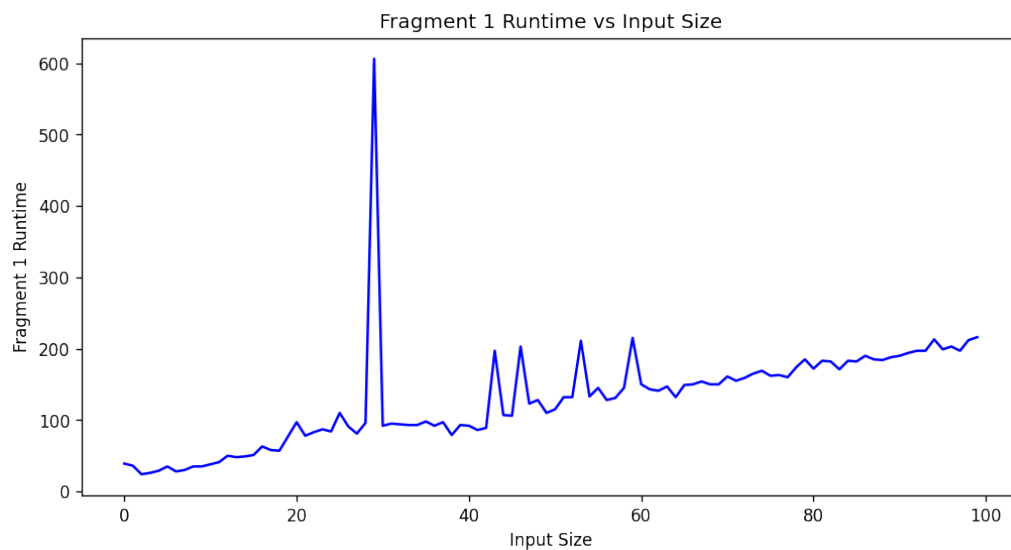


Figure 2: Fragment 1 Runtime

2. Fragment 2

$O(N^2)$

```
void frag2(size_t n) {  
    size_t sum = 0;  
    for (size_t i = 0; i < n; ++i)  
        for (size_t j = 0; j < i; ++j)  
            sum = sum + 1;  
}
```

Figure 3: Fragment 2

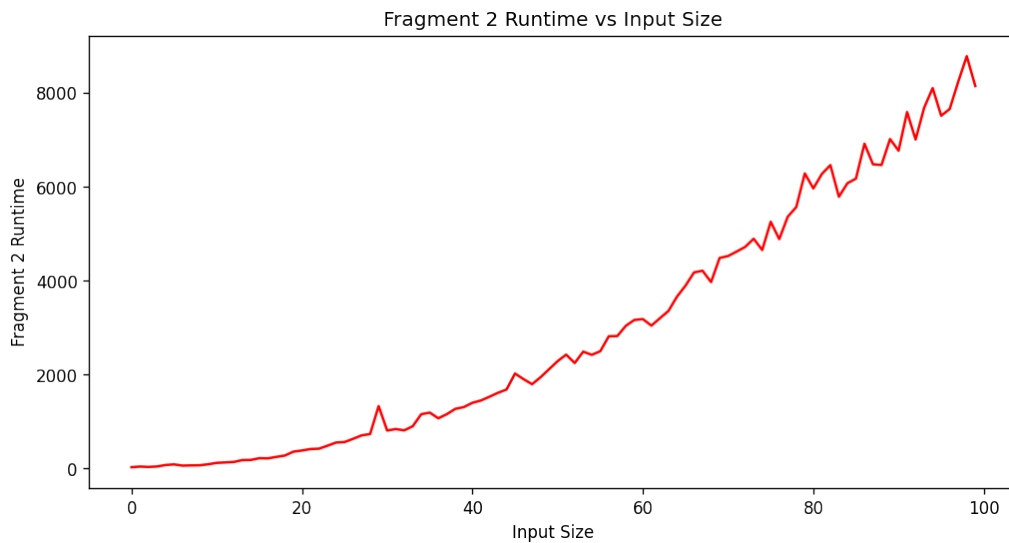


Figure 4: Fragment 2 Runtime

3. Fragment 3

$O(N^5)$

```
void frag3(size_t n) {  
    size_t sum = 0;  
    for (size_t i = 0; i < n; ++i)  
        for (size_t j = 0; j < i * i; ++j)  
            for (size_t k = 0; k < j; ++k)  
                sum = sum + 1;  
}
```

Figure 5: Fragment 3

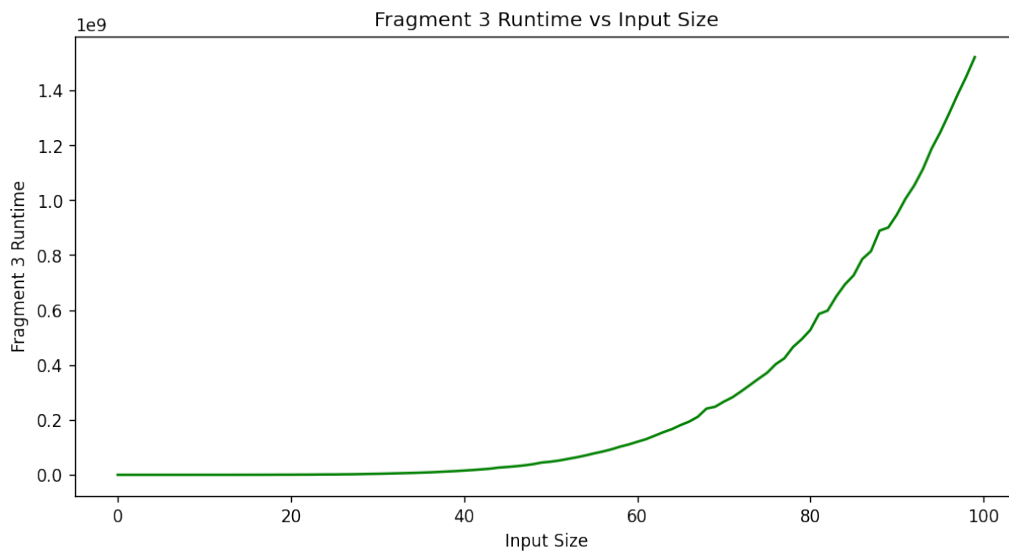


Figure 6: Fragment 3 Runtime

Problem 2

For each of the following operations, write out (e.g. in pseudocode) the algorithm you typically use for hand-calculations and give a big-O estimate of the time complexity of that algorithm.

A. Add two N -digit integers

```
def add(int a, int b):
    Let x = min(a, b)
    Let y = max(a, b)
    Let carry = 0
    Let digits = []

    while x > 0:
        Let dx = x % 10
        Let dy = y % 10
        Let sum = dx + dy + carry
        if sum > 10:
            sum %= 10
            carry = 1
        else:
            carry = 0
        end

        digits.append(sum)
        x //= 10
        y //= 10
    end

    print digits in reverse
end
```

$O(n)$

B. Multiply two N -digit integers

```
def mult(a, b):  
    Let product = 0;  
    while (b > 0):  
        product = add(product, a)  
        b = b - 1  
    end  
    return product  
end
```

 $O(N^2)$ **C. Divide two N -digit integers**

```
def divide(a, b):  
    Let quotient = 0;  
    Let neg_b = mult(-1, b)  
    if (b < a):  
        return 0  
    end  
    while (a > 0):  
        quotient = quotient + 1  
        a = add(a, neg_b)  
    end  
    return quotient  
end
```

 $O(N^2)$

Problem 3

Use big-O to estimate how much time is required to compute the following function:

$$f(x) = \sum_{i=0}^N a_i x^i$$

- A. $O(N)$
- B. $O(\log N)$
- C. $O(N^2)$

Puzzle Problem [optional]

Show that X^{62} can be computed with only 8 multiplications

$$x^2 \rightarrow x^4 \rightarrow x^6 \rightarrow x^{12} \rightarrow x^{24} \rightarrow x^{30} \rightarrow x^{31} \rightarrow x^{62}$$

Extra Challenge

What is the smallest possible power k such that X^k required at least eight multiplications.