

Problem 1. Proof by resolution. Solution:

Translate the following into first-order logic.

Every kid likes Snoopy	$\forall x(K(x) \rightarrow L(x, Snoopy))$
Everyone who likes Snoopy likes all beagles	$\forall x(L(x, Snoopy) \rightarrow (\forall y B(y) \rightarrow L(x, y)))$
Underdog is a beagle, and Underdog is a superhero	$B(Underdog) \wedge S(Underdog)$
Every superhero can fly or time travel	$\forall x(S(x) \rightarrow (F(x) \vee T(x)))$
No beagle can time travel	$\neg \exists x(B(x) \wedge T(x))$
Garfield does not like all things that can fly	$\neg \forall x(F(x) \rightarrow L(Garfield, x))$
Garfield is not a kid. (goal that will be negated latter)	$\neg K(Garfield)$

Convert the resulting formulas into a canonical form

- $\forall x(K(x) \rightarrow L(x, Snoopy))$
 $\forall x(\neg K(x) \vee L(x, Snoopy))$
- $\forall x(L(x, Snoopy) \rightarrow \forall y(B(y) \rightarrow L(x, y)))$
 $\forall x \forall y(L(x, Snoopy) \rightarrow (B(y) \rightarrow L(x, y)))$
 $\forall x \forall y(L(x, Snoopy) \rightarrow (\neg B(y) \vee L(x, y)))$
 $\forall x \forall y(\neg L(x, Snoopy) \vee (\neg B(y) \vee L(x, y)))$
 $\forall x \forall y(\neg L(x, Snoopy) \vee \neg B(y) \vee L(x, y))$
- $B(Underdog) \wedge S(Underdog)$
 $B(Underdog)$
 $S(Underdog)$
- $\forall x(S(x) \rightarrow (F(x) \vee T(x)))$
 $\forall x(\neg S(x) \vee (F(x) \vee T(x)))$
 $\forall x(\neg S(x) \vee F(x) \vee T(x))$
- $\neg \exists x(B(x) \wedge T(x))$
 $\forall x \neg(B(x) \wedge T(x))$
 $\forall x(\neg B(x) \vee \neg T(x))$

6. $\neg\forall x(F(x) \rightarrow L(\text{Garfield}, x))$
 $\neg\forall x(\neg F(x) \vee L(\text{Garfield}, x))$
 $\exists x(F(x) \wedge \neg L(\text{Garfield}, x))$
 $F(\alpha) \wedge \neg L(\text{Garfield}, \alpha)$
 $F(\alpha)$
 $\neg L(\text{Garfield}, \alpha)$

7. Negated goal: $K(\text{Garfield})$

Resolving negated goal clause with $\neg K(x) \vee L(x, \text{Snoopy})$

$\sigma = \{x/\text{Garfield}\}$

$L(\text{Garfield}, \text{Snoopy})$

Resolving $L(\text{Garfield}, \text{Snoopy})$ with $\forall x\forall y(\neg L(x, \text{Snoopy}) \vee \neg B(y) \vee L(x, y))$

$\sigma = \{x/\text{Garfield}\}$

$\neg B(y) \vee L(\text{Garfield}, y)$

Resolving $\neg B(y) \vee L(\text{Garfield}, y)$ with $B(\text{Underdog})$

$\sigma = \{y/\text{Underdog}\}$

$L(\text{Garfield}, \text{Underdog})$

Resolving $L(\text{Garfield}, \text{Underdog})$ with $\neg L(\text{Garfield}, \alpha)$

$\sigma = \{\alpha/\text{Underdog}\}$

NULL Clause

The statement is proven \square

Problem 2. The resolution theorem prover.

Solution:

The resolution theorem prover is included in the submission folder.

Problem 3. Prolog

Solution:

The knowledge base is included in the submission folder as requested.

The queries are translated to prolog as follows:

```
grandparent(charles, X) .  
brother(william, X), spouse(X, Y) .  
uncle(X, charlotte) .  
aunt(X, archie) .  
nephew(X, harry) .  
stepmother(X, william) .  
sibling(X, Y) .  
cousin(X, Y), female(X) .  
niece(X, Y) .
```

The answers as determined by the Knowledge Base are as follows:

```
X = george  
X = harry, Y = meghan  
X = harry  
X = catherine  
X = george  
X = camila  
X = william, Y = harry  
X = charlotte, Y = archie  
X = charlotte, Y = harry
```