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Class: COMP1D-Y

Comp Arch

Q4

Ans

A	B	C	D	X
0	0	0	0	1
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

Q1. D-latch

D	Q
0	0
1	1

Q3

expression: $\bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + A\bar{B}\bar{C} + ABC$

	$\bar{A}\bar{B}$	$A\bar{B}$	AB
\bar{C}	0	1	0
C	1	1	1

$$\bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + A\bar{B}\bar{C}$$

$$\bar{A}\bar{B}\bar{C} + AC$$

Q4:

$19/2 = 9$	R
$9/2 = 4$	1
$4/2 = 2$	1
$2/2 = 1$	0
$1/2 = 0$	0
	1

$$131/2 = 65 \quad R1$$

$$65/2 = 32 \quad R1$$

$$32/2 = 16 \quad R0$$

$$16/2 = 8 \quad R0$$

$$8/2 = 4 \quad R0$$

$$4/2 = 2 \quad R0$$

$$2/2 = 1 \quad R0$$

$$1/2 = 0 \quad R1$$

ii $19_{(10)} = 10011_{(2)} \Rightarrow 0.25$ fraction

iii $0.25 \times 2 = 0, 0.5$

$0.5 \times 2 = 1, 0$

$$0.25_{(10)} = 0.01_{(2)}$$

iv $19.25_{(10)} = 10011.01 \times 2^0 \Rightarrow 1.001101 \times 2^4$

v $4 + 2^{(8-1)} - 1$
 $= 131_{(10)}$

1000 0011

8 expression = $\bar{A}\bar{B}\bar{C}D + \bar{A}BCD + A\bar{B}\bar{C}D + A\bar{B}CD$
 $+ ABC\bar{D} + ABCD$

	CD 00	01	11	10
00 AB	x	1	1	x
01	x	0	0	x
11	1	0	0	1
10	0	1	1	0

= $A\bar{B}D + \bar{A}\bar{B}D + ABC\bar{D} + ABCD$

Q10 Expression = $\bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}\bar{C}D + \bar{A}\bar{B}C\bar{D} + \bar{A}\bar{B}CD$
 $+ A\bar{B}\bar{C}\bar{D} + A\bar{B}C\bar{D} +$

	CD 00	01	11	10
00 AB	1	1	1	1
01	0	0	0	0
11	0	0	0	0
10	1	0	0	1

= $AB + \bar{A}\bar{B}\bar{C}\bar{D} + A\bar{B}C\bar{D}$

Q4 mantissa:- 001 1010 0000 0000 0000 0000
 sign:- 1
 exponent:- 1000 0011

1-10000011-0011010 0000 0000 0000 0000.

Q5 ^{1's} complement

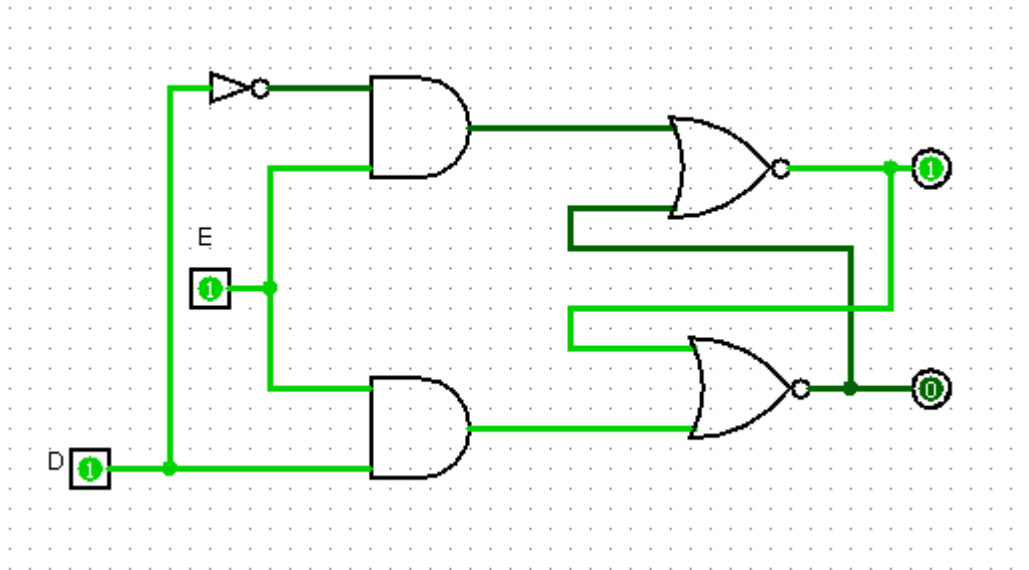
	R		R
$36/2 = 18$	0	$28/2 = 14$	0
$18/2 = 9$	0	$14/2 = 7$	0
$9/2 = 4$	1	$7/2 = 3$	1
$4/2 = 2$	0	$3/2 = 1$	1
$2/2 = 1$	0	$1/2 = 0$	1
$1/2 = 0$	1		

00011100
 $\Rightarrow 1100011 (-8)$

00100100 ~~00000000~~ (36)

00100100	36
<u>11100011</u>	-8
00000111	
+ 1	
<u>00000110</u>	28 (1's)

Q1. D Latch circuit



Truth table:

E	D	Q(T+1)
0	0	Q _t
0	1	Q _t
1	0	0
1	1	1

How does it operate:

A D latch is like an S-R latch with only one input: the “D” input. Activating the D input sets the circuit, and de-activating the D input resets the circuit. Of course, this is only if the enable input (E) is activated as well. Otherwise, the output(s) will be latched, unresponsive to the state of the D input.

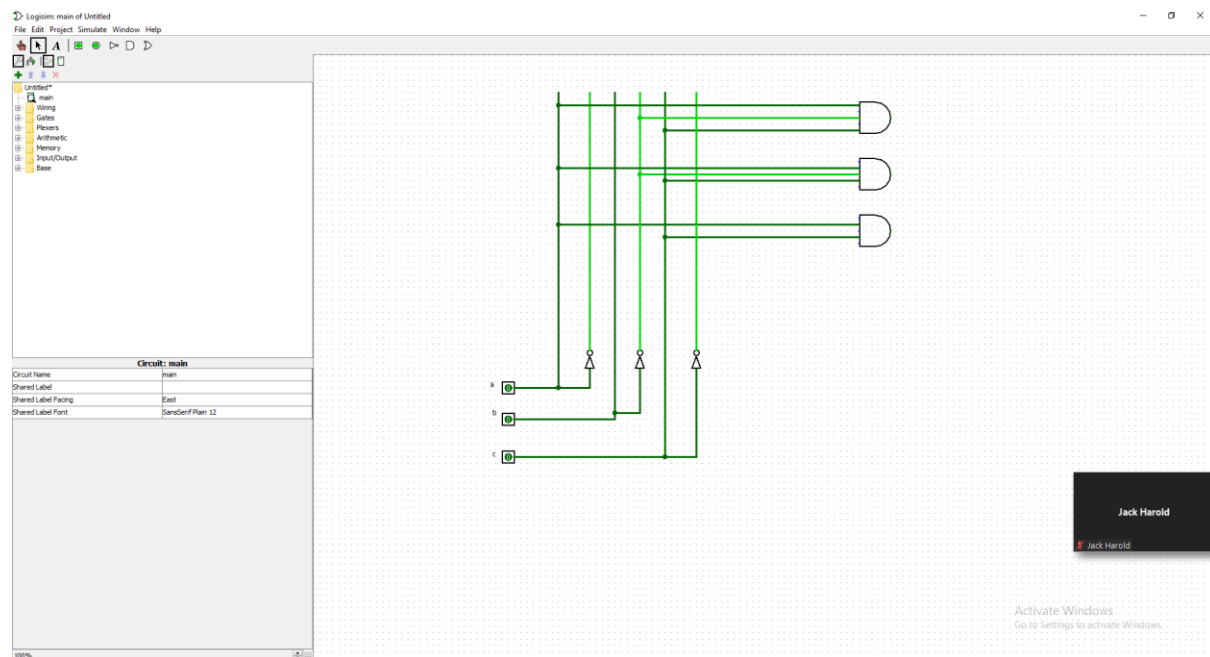
D latches can be used as 1-bit memory circuits, storing either a “high” or a “low” state when disabled, and “reading” new data from the D input when enabled.

Q2.

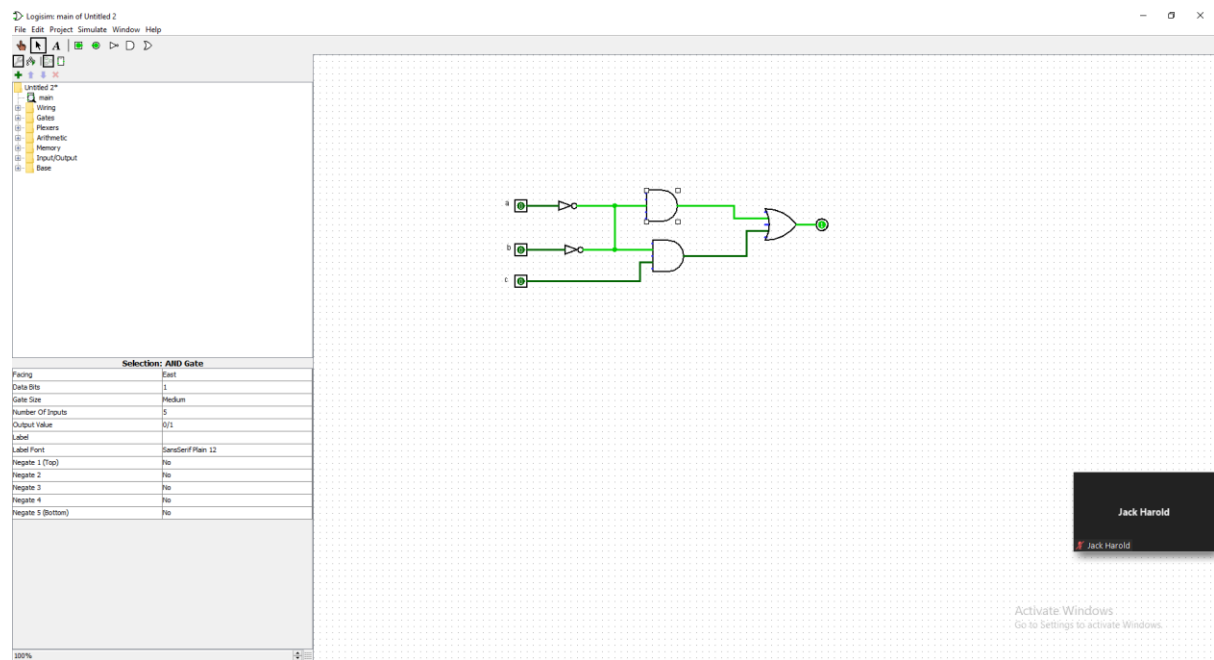
Hamming codes may detect and repair one-bit and two-bit errors while leaving untreated faults undetected. The simple parity code, on the other hand, cannot rectify faults and can only identify an odd number of bits in error. The Hamming code is a set of error-correction codes that can be used to detect and fix faults that can arise when data is transferred or stored from

one source to another. R.W. Hamming invented this approach for error correction. Only an odd number of faults in a codeword can be detected by a single parity check code. The difference is that error detection is a technique for examining data to see if it has been corrupted during storage or transmission. Error repair is a step up from error detection; when it identifies an error, it attempts to restore the data to its original state. Because a single extra bit is added to each byte of data and assigned a value of 1 or 0, often depending on whether there are an even or odd number of "1" bits, parity is the optimum method for error detection.

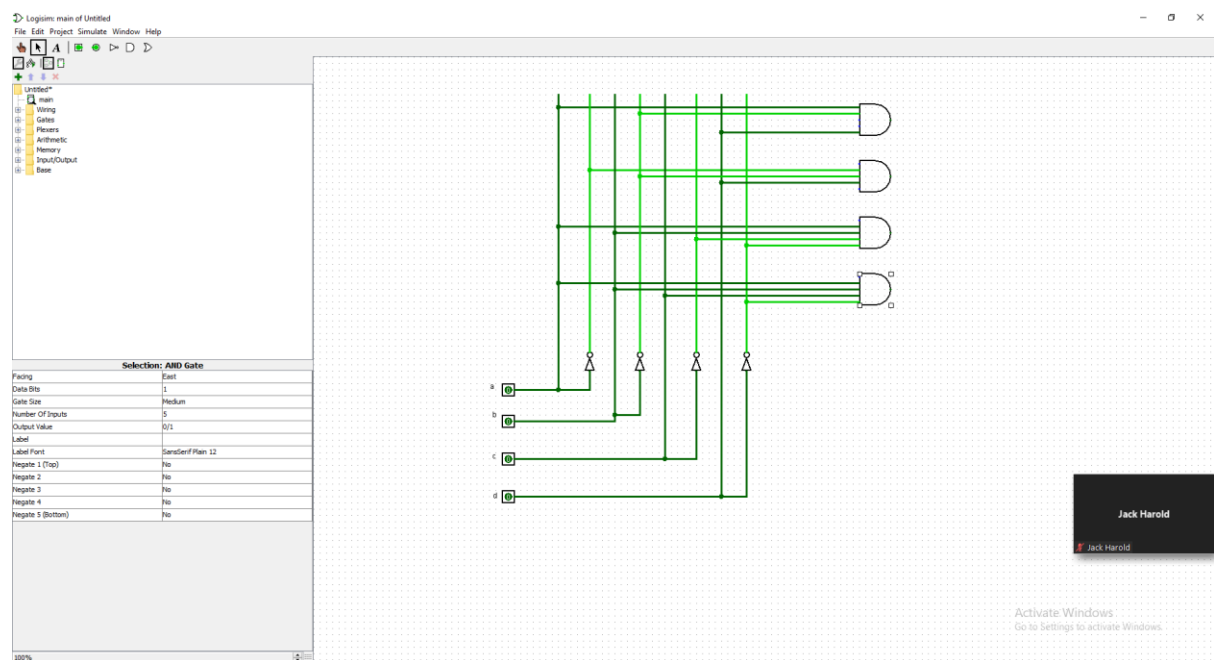
Q3.Circuit



Q7.circuit



Q8.Circuit



Q10.Circuit

