Lab No. 10 Synchronization: Two-Process Solutions, MUTEX, and Semaphore

Objective

This lab is designed to implement the solutions to the critical-solution problem.

Activity Outcomes:

On completion of this lab students will be able to

- Implement two process solutions to critical-section problem
- Solve the CS problem using MUTEX and Semaphore

Instructor Notes

As pre-lab activity, read the content from the following (or some other) internet source: https://www.geeksforgeeks.org

1) Useful Concepts

The Critical-Section Problem

Cooperating processes or threads share some data with each other. If two or more threads or processes access and manipulate the shared data concurrently then this may result in data inconsistency. To avoid such data inconsistencies, we need to make it sure that threads/processes must be synchronized and if one thread/process is manipulating the shared data then no other thread/process should be allowed to access that data.

Each cooperating thread/process has some segments of critical code that is the segment of code where shared data is accessed and manipulated. These segment of codes are called critical-sections. We need to make it sure that if one thread/process is executing its critical section then no other process should be allowed to execute its critical section. Designing solutions to ensure this; is called the CS problem. We can define the CS problem as:

The critical section problem is used to design a protocol followed by a group of processes, so that when one process has entered its critical section, no other process is allowed to execute in its critical section.

The general structure of a solution to critical-section problem is:

```
do {

    entry section

    critical section

    exit section

    remainder section
} while (true);
```

Two Process Solutions to CS problem

A simple solution:

First, we discuss a simple solution. This solution uses a variable turn. The value of turn decides, whose turn it is to enter in CS. The solution is given below:

```
Code for process i
                                    Code for process j
do {
                                    do {
     while (turn == j); // Entry
                                         while (turn == i); // Entry
Code
                                    Code
     critical section
                                          critical section
      turn = j; //Exit code
                                          turn = i; //Exit code
     remainder section
                                         remainder section
    while (true);
                                         while (true);
```

This simple solution ensure mutual exclusion and bounded waiting conditions but it fail to ensure progress condition.

Peterson's Solution:

Peterson's solution satisfies all the conditions for a good solution. It is also a two process solution. The two processes share two variables: int turn and Boolean flag[2]. The variable turn indicates whose turn it is to enter the critical section while the flag array is used to indicate if a process is ready to enter the critical section. flag[i] = true implies that process Pi is ready. The pseudo code of the Peterson's solution is given below:

```
Code for process i
                                     Code for process j
do
                                     do
       flag[ i ] = true;
                                            flag[ j] = true;
       turn = j;
                                            turn = i;
                                            while (flag[ i ] && turn =
       while (flag[ j ] && turn =
                                     = i);
= j);
     critical section
                                          critical section
       flag[ i ] = false;
                                            flag[ j ] = false;
     remainder section
                                          remainder section
 while (true);
                                      while (true); while (true);
```

MUTEX Lock

MUTEX lock is software based solution to CS problem and is applicable on n threads/processes. MUTEX is a shortened form of the words "mutual exclusion". MUTEX variables are one of the primary means of implementing thread synchronization. A MUTEX variable acts like a "lock" protecting access to a shared data resource. The basic concept of a MUTEX; as used in pthreads is that only one thread can lock (or own) a MUTEX variable at any given time. Thus, even if several threads try to lock a MUTEX only one thread will be successful. No other thread can own that MUTEX until the owning thread unlocks that MUTEX.

A typical sequence in the use of a MUTEX is as follows:

Create and initialize a MUTEX variable

- Several threads attempt to lock the MUTEX
- Only one succeeds and that thread owns the MUTEX
- The owner thread performs some set of actions
- The owner unlocks the MUTEX
- Another thread acquires the MUTEX and repeats the process
- Finally the MUTEX is destroyed.

The routines to perform these tasks are given below:

```
pthread_mutex_init(pthread_mutex_t var, pthread_mutexattr_t attr)
// to initialize MUTEX variable
pthread_mutex_lock (pthread_mutex_t var) // to lock CS
pthread_mutex_unlock (pthread_mutex_t var) // to unlock CS
pthread_mutex_destroy(pthread_mutex_t var) // to destroy MUTEX
variable
```

Semaphore

Semaphore is another synchronization tool that can be used to solve several synchronization problems. Semaphore is an integer variable but it can be accessed only through two functions which are wait() and signal().

```
Pseudo code for wait function
                                     Pseudo code for solution to CS
                                    problem using semaphore
wait(S) {
                                    do {
    while (S \le 0)
                                                  //entry code
                                        wait(s)
       ; // busy wait
                                            critical section
    s - -;
                                         signal(S) // exit code
}
                                           remainder section
                                     } while (true);
Pseudo code for signal function
signal (S) {
    S++;
}
```

The following routines are used to implement POSIX semaphore

```
#include <semaphore.h> // header-file
sem t // semaphore data type
int sem_init(sem_t *sem, int pshared, unsigned value); // to
initialize of semaphore
```

sem_wait()	// wait function
sem_post()	// signal function
int sem dest	roy(sem t *sem); // to destroy semaphore

2) Solved Lab Activities

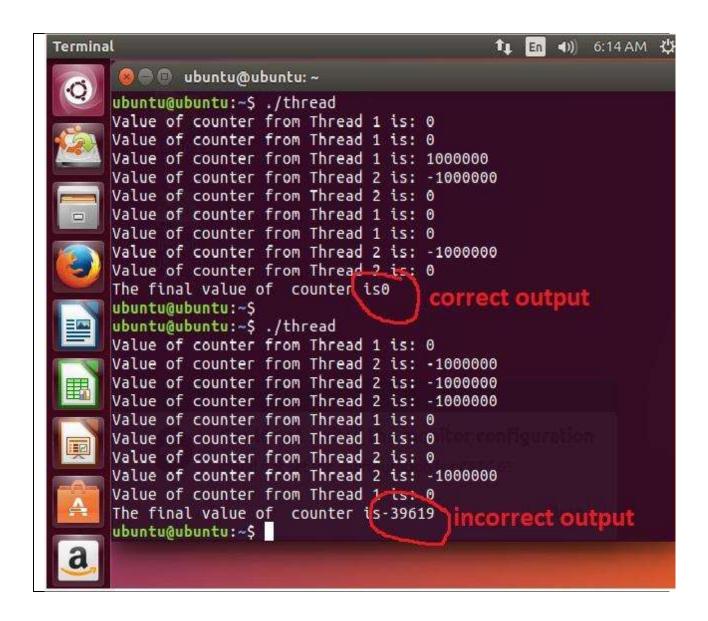
Sr.No	Allocated Time	Level of Complexity	CLO Mapping
1	20	Medium	CLO-7
2	10	Medium	CLO-7
3	15	Medium	CLO-7
4	15	Medium	CLO-7

Activity 1:

In this activity, we identify how the concurrent access to shared data may result in data inconsistency. We create two threads that manipulate a shared variable counter concurrently. As a result, sometimes we get correct output while sometimes the output is not c Solution:

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```
Text Editor
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         🔞 🖨 📵 thread.cpp (~/) - gedit
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         #include<iostream>
         #include<pthread.h>
         using namespace std:
        int counter=0;// shared data
        void *Thread1(void *args) // routine executed by thread 1
         for(int i=0; i<=5000000; i++)
         counter++;
         if(counter%1000000==0)
         cout<<"Value of counter from Thread 1 is: "<<counter<<endl;
         void *Thread2(void *args) // routine executed by thread 2
         for(int i=0; i<=50000000; i++)</pre>
   counter--;
        if(counter%1000000==0)
         cout<<"Value of counter from Thread 2 is: "<<counter<<endl;
         int main()
        pthread t t1,t2;
        pthread_create(&t1,NULL,Thread1,NULL);
         pthread_create(&t2,NULL,Thread2,NULL);
         pthread join(t1, NULL);
         pthread_join(t2, NULL);
         cout<<"The final value of counter is"<<counter<<endl;
         return 0;
Out-put
```



Activity 2:

Out-put

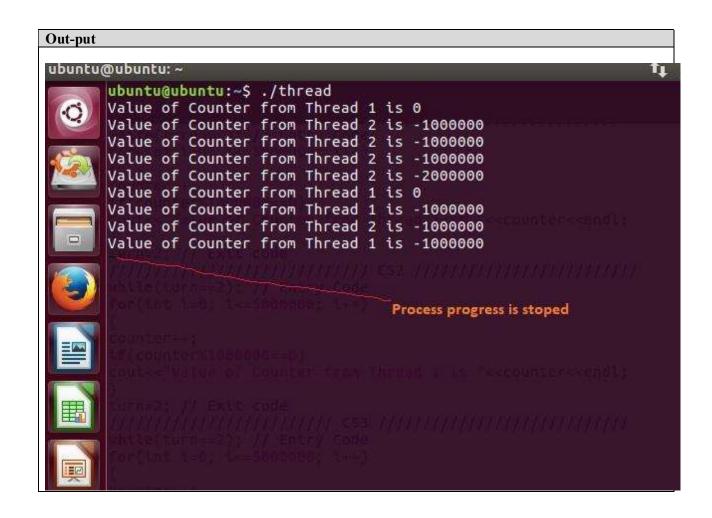
Now, we implement the simple solution in the code written in Activity 1 to protect the CS Solution:

```
Code
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       #include<iostream>
       #include<pthread.h>
       using namespace std;
       int turn=0; // turn variable
       int counter=0;
       void *thread1(void* args)
       while(turn==2); // Entry Code
       for(int i=0; i<=50000000; i++)
       counter++;
       if(counter%1000000==0)
       cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
       turn=2; // Exit code
       void *thread2(void* args)
       while(turn==1); //Entry code
       for(int i=0; i<=50000000; i++)
       counter --:
       if(counter%1000000==0)
       cout<<"Value of Counter from Thread 2 is "<<counter<<endl;
       turn=1; //Exit code
       int main()
       pthread_t t1,t2;
       pthread create(&t1,NULL,thread1,NULL);
       pthread create(&t2,NULL,thread2,NULL);
       pthread_join(t1,NULL);
       pthread_join(t2,NULL);
       cout<<"Final value of counter is: "<<counter<<endl;
       return 0;
```



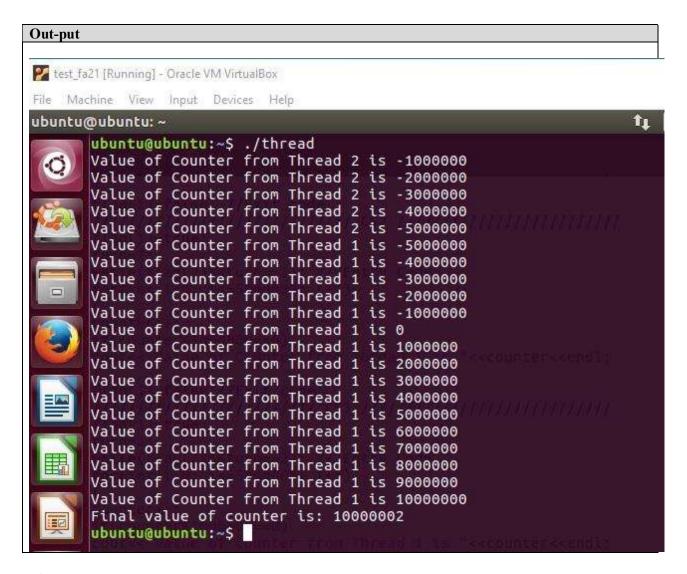
Activity 3: In this activity, we show that how the progress in not satisfied in simple solution Solution:

```
Code
   thread.cpp (~/) - gedit
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                  a
         #include<iostream>
         #include<pthread.h>
         using namespace std;
         int turn=0; // turn variable
         int counter=0;
         void *thread1(void* args)
         for(int i=0; i<=50000000; i++)
         counter++;
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
         turn=2; // Exit code
         while(turn==2); // Entry Code
         for(int i=0; i<=5000000; i++)</pre>
    counter++;
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
         turn=2; // Exit code
         for(int i=0; i<=50000000; i++)
         counter++;
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
         turn=2; // Exit code
         void *thread2(void* args)
         while(turn==1); //Entry code
         for(int i=0; i<=5000000; i++)
         counter--;
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 2 is "<<counter<<endl;
         turn=1; //Exit code
         int main()
         pthread_t t1,t2;
pthread_create(&t1,NULL,thread1,NULL);
         pthread_create(&t2,NULL,thread2,NULL);
         pthread_join(t1,NULL);
         pthread join(t2, NULL);
         cout<<"Final value of counter is: "<<counter<<endl;
         return 0;
         3
```



Activity 4: In this activity, we implement the Peterson's solution and show that progress condition is satisfied. Solution:

```
Code
  thread.cpp (~/) - gedit
                                                                           11
         #include<iostream>
         #include<pthread.h>
         using namespace std;
bool flag[3]={false,false,false}; // flag array
         int turn=0; // turn variable
int counter=0;
         void *thread1(void* args)
         turn=2
         while(flag[2] && turn==2); // Entry Code
for(int i=0; i<=5000000; i++)</pre>
         counter++:
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
   flag[1]=true;
         turn=2
         while(flag[2]&&turn==2); // Entry Code
for(int i=0; i<=5000000; i++)</pre>
         counter++:
         if(counter%1000000==0)
    a
         cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
         flag[1]=false;//Exit code
         turn=2;
         while(flag[2]&&turn==2); // Entry Code
for(int i=0; i<=5000000; i++)</pre>
         counter++:
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 1 is "<<counter<<endl;
         flag[1]=false;//Exit code
         void *thread2(void* args)
    a
         flag[2]=true;
         turn=1
         while(flag[1]&&turn==1); // Entry Code
for(int i=0; i<=5000000; i++)</pre>
         counter --:
         if(counter%1000000==0)
         cout<<"Value of Counter from Thread 2 is "<<counter<<endl;
         flag[2]=false; //Exit code
         int main()
         pthread_t t1,t2;
pthread_create(&t1,NULL,thread1,NULL);
         pthread_create(&t2,NULL,thread2,NULL);
pthread_join(t1,NULL);
         pthread_join(t2,NULL);
cout<<"Final value of counter is: "<<counter<<endl;</pre>
         return 0;
```



3) Graded Lab Tasks

Note: The instructor can design graded lab activities according to the level of difficult and complexity of the solved lab activities. The lab tasks assigned by the instructor should be evaluated in the same lab.

Task 1:

Write the code for the problem given in Activity 3 and protect the CS's using the MUTEX lock

Task 2:

Write the code for the problem given in Activity 3 and protect the CS's using the Semaphore