

- 1. 背景介绍
- 2. 赛题详解
- 3. 解题思路和案例
- 4. 评分标准

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背景介绍 - 验证方法

- Formal verification v.s. simulation based verification
- Directed test, constrained random test and UVM

- Solver简介
- 类型: SAT, SMT, CSP and BDD-based
- 特性:正确性,随机性和可重现,概率分布,性能要求

赛题详解 – 约束问题定义

约束问题,由变量、常量和表达式组成,其中

- 变量是无符号(unsigned)的、具有固定位宽(bit-width)的位向量(bit-vector)
 - 比如bit[7:0] x;定义了一个unsigned的8bit位宽的变量。
- 常量是有符号或无符号、具有固定位宽的整型数
- 表达式由操作数和运算符构成,操作数可以常量或变量,运算符有以下6类:
 - 逻辑运算: &&, ||,!
 - 位运算: &, |, ^, ~ ...
 - 算术运算: +, -, *, /, %
 - 关系运算: >, <, >=, <=, ==,!=

赛题详解 – 约束问题示例

bit[1:0] x,y,z;

x > y;

y > z;

随机解1:

x = 2'b11

y = 2'b10

z = 2'b01

随机解2:

x = 2'b10

y = 2'b01

z = 2'b00

. . .

赛题详解 - BDD & CUDD

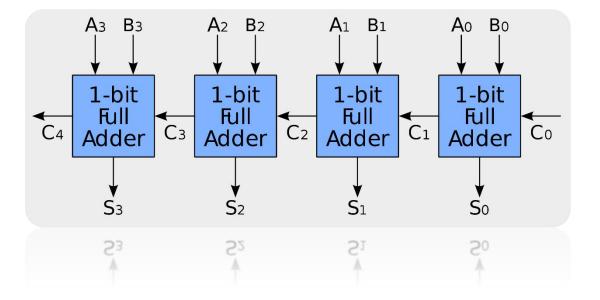
BDD^[9,10]是一种可以用来表达布尔函数的数据结构。在形式化验证、电路验证中,许多任务都涉及大型命题逻辑公式的运算,BDD在这些领域中有着广泛的应用。在约束问题求解中,因为约束本身就是一系列逻辑、位运算和算术运算组成的表达式,所以也可以通过基于BDD的运算来实现。

CUDD是一个用于计算和处理BDD/ADD/ZDD的开源库。使用CUDD提供的API可以比较方便地 从表达式计算得到BDD

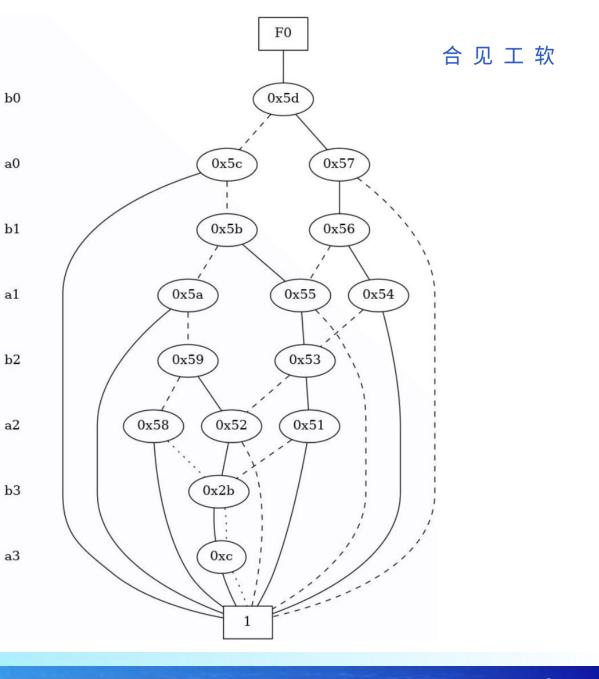
赛题详解 - BDD & CUDD

bit[3:0] A,B;

$$A + B > 4'b0000$$

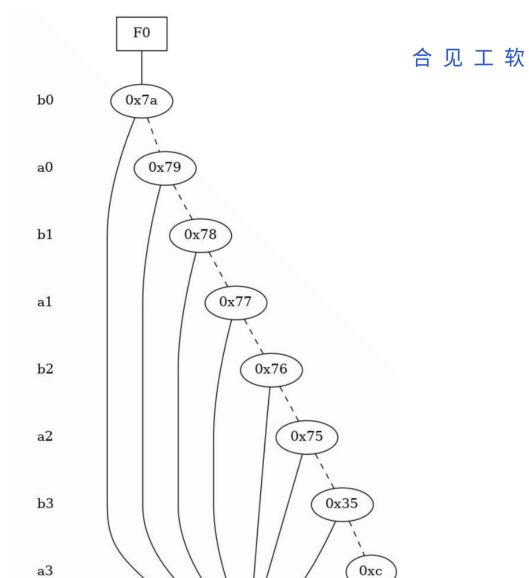


```
#include "cuddObj.hh"
BDD a0, a1, a2, a3, b0, b1, b2, b3; // bddvar
BDD s0, s1, s2, s3;
BDD result;
a0 = bddvar();
s0 = a0 \land b0;
c0 = (a0 \& b0);
s1 = a1 \wedge b1 \wedge c0;
c1 = (a1 \& b1) | (c0 \& (a1 | b1))
s2 = a2 \wedge b2 \wedge c1;
c2 = (a2 \& b2) | (c1 \& (a2 | b2))
s3 = a3 \wedge b3 \wedge c2;
c3 = (a3 \& b3) | (c2 \& (a3 | b3))
result = s0.Xor(s1.Xor(s2.Xor(s3)))
```



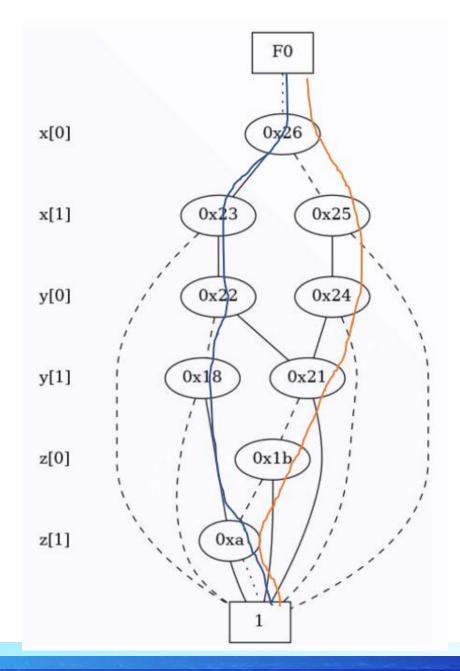
$$A + B > 0$$

$$32'(A) + 32'(B) > unsigned'(0)$$



$$x > y$$
;

$$y > z$$
;



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The CUDD package, BDD, ADD Tutorial and examples | David

<u>Kebo</u>

Henrik Reif Andersen. An Introduction to Binary Decision Diagrams. Fall, 1999

bdd.pdf (utexas.edu)

```
"variable list":
                   // variable
                   "id": <number>,
                   "name": <string>,
                   "signed": <true_of_false>,
                   "bit width": <number>
"constraint list":
                   // expression
                   "op": <string> // operator type
                   "id": <number> // variable id if op is VAR
                   "value": <number> // hexical number is op is CONST
                   "Ihs expression": // if op is another type
                                 // Ihs expression
                                 // same as above expression format
                   "rhs_expression": // if op is another type and rhs exists
                                 // rhs expression
                                 // same as above expression format
                   "if expression": // if op is MUX(If-Then-Else), if: if expression,
then: Ihs expression, else: rhs expression
                                 // pred expression
                                 // same as above expression format
```

```
"assignment list":
                  // assignment 1
                   {"value": <number>} // hexical number, sorted in asending
order of variable ids
                   {"value": <number>}
                  // assignment 2
```

```
Constraint Problem:

bit [31:0] a;

bit [31:0] b;

bit [31:0] c;

bit [31:0] abj_a;

bit [31:0] abj_b;

bit [31:0] abj_c;

constraint cb {

(abj_a < 32'h0000064);

(abj_b < 32'h00000064);

(abj_c < 32'h00000064);

(abj_a == a);

(abj_b == b);

(abj_c == c);

((a * b) * c) == 32'h0000000533);

}
```

```
Constraint Problem in Json:
"variable list": [
     { "id": 0, "name": "a", "signed": false, "bit width": 32 },
     { "id": 1, "name": "b", "signed": false, "bit width": 32 },
     { "id": 2, "name": "c", "signed": false, "bit width": 32 },
     { "id": 3, "name": "abj a", "signed": false, "bit width": 32 },
     { "id": 4, "name": "abj b", "signed": false, "bit width": 32 },
     { "id": 5, "name": "abj c", "signed": false, "bit width": 32 }
"constraint list": [
                    "op": "LT",
                    "lhs expression": { "op": "VAR", "id": 3 },
                    "rhs expression": { "op": "CONST", "value": "000000001234556789012345678901234567890" }
                    "op": "EQ",
                    "lhs expression": {
                                   "op": "MUL",
                                   "lhs expression": {
                                                 "op": "MUL'
                                                 "lhs expression": { "op": "VAR", "id": 0 },
                                                 "rhs expression": { "op": "VAR", "id": 1 }
                                   "rhs expression": { "op": "VAR", "id": 2 }
                    "rhs expression": { "op": "CONST", "value": "000001533" }
```

```
results in Json:
"assignment_list":
                     {"value": "0"},
                     {"value": "1"},
                     {"value": "2"},
                     {"value": "3"},
                     {"value": "4"},
                     {"value": "5"}
                     {"value": "5"},
                     {"value": "4"},
                     {"value": "3"},
                     {"value": "2"},
                     {"value": "1"},
                     {"value": "0"}
     ],
[
                     {"value": "5"},
                     {"value": "4"},
                     {"value": "3"},
                     {"value": "2"},
                     {"value": "1"},
                     {"value": "0"}
```

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赛题详解 - 任务分解

- 1.读取问题文件,存储到内部数据结构
- 2. 精度转换
- 3.遍历表达式,使用CUDD API构造BDD
- 4.基于BDD, 挑选合法路径
- 5.输出随机解到文件

评分标准

本赛题将提供一定数量的测试用例,其中包含了从简单到复杂不同难度的问题。所有的测试用例赛前不对参赛者公开。每个测试用例分配相同的分值,最后根据总分评判参赛者排名。赛前会提供若干测试用例和输出结果的样本,供参赛者练习参考使用。

评分标准包括3个部分:

- 1) 输出结果必须保证正确性;
- 2) 性能通过程序的执行时间来评判;
- 3) 输出结果需要符合均匀联合随机分布;

根据以上标准每个测试用例的评分步骤如下,具体细则另行说明:

- 1) 检查结果正确性,任意一组值错误,本用例得分为0;
- 2) 检查结果分布,如随机值分布偏差过大,本用例得分为0;
- 3) 通过正确性和分布检查后,根据性能排名评分:假设参赛人数为N,名次为M,得分根据以下公式计算

Score = 10* (N-M+1) / N

第一名得10分, 第二名得(N-1)*10/N......

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