

#ADS05 Binomial Queue

二项队列

左式堆将合并、删除最小元的操作的时间控制在 $O(\log N)$,二项队列进一步减少了这个时间。

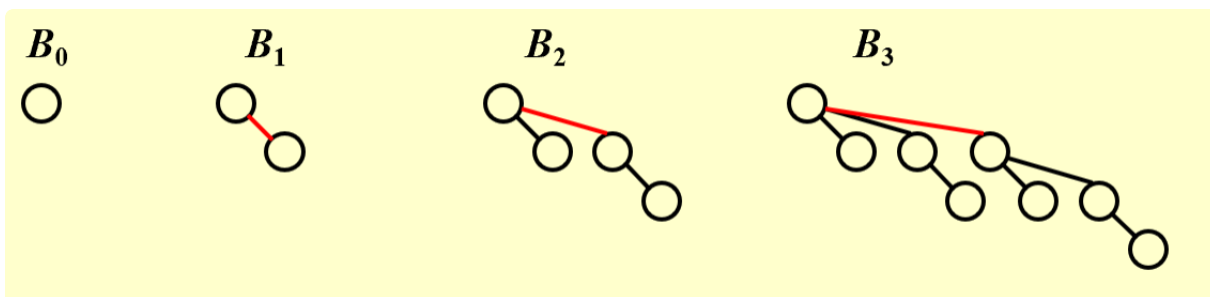
Binomial以最坏时间 $O(\log N)$ 支持上述操作, **插入操作平均花费常数时间**

Structure:

A binomial queue is not a heap-ordered tree, but rather **a collection of heap-ordered trees**, known as a forest. Each heap-ordered tree is a binomial tree.

A binomial tree of height 0 is a one-node tree.

A binomial tree, B_k , of height k is formed by attaching a binomial tree, B_{k-1} , to the root of another binomial tree, B_{k-1} .



B_k consists of a root with k children, which are B_0 to B_{k-1} .

B_k has exactly 2^k nodes. The number of nodes at depth d is Ckd

B_k structure + heap order + one binomial tree for each height

→ A priority queue of any size can be uniquely represented by a collection of binomial trees.

如果我们把堆序添加在二项树上，并允许任意高度上最多有一棵二项树，那么我们能由二项树的集合唯一标识任意大小的优先队列。

【Example】 Represent a priority queue of size **13** by a collection of binomial trees.

Solution: $13 = 2^0 + 0 \times 2^1 + 2^2 + 2^3 = 1101_2$

B_0



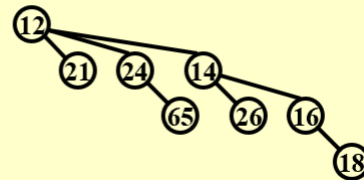
B_1



B_2



B_3



FindMin

The minimum is in one of the roots.

There are at most $\lceil \log N \rceil$ roots, hence $T_p = O(\log N)$.

We can remember the minimum and update whenever it is changed. Then this operation will take $O(1)$

⊕ Merge:

H_1



H_2

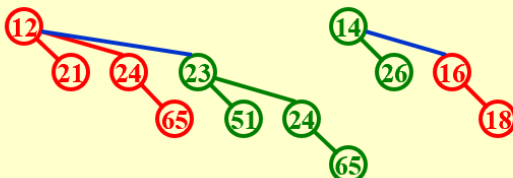


13

$$110_2 = 6$$

$$+ 111_2 = 7$$

H_1



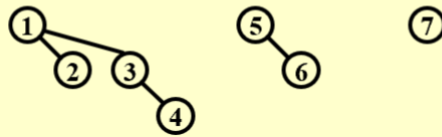
13

$$1101$$

$$T_p = O(\log N)$$

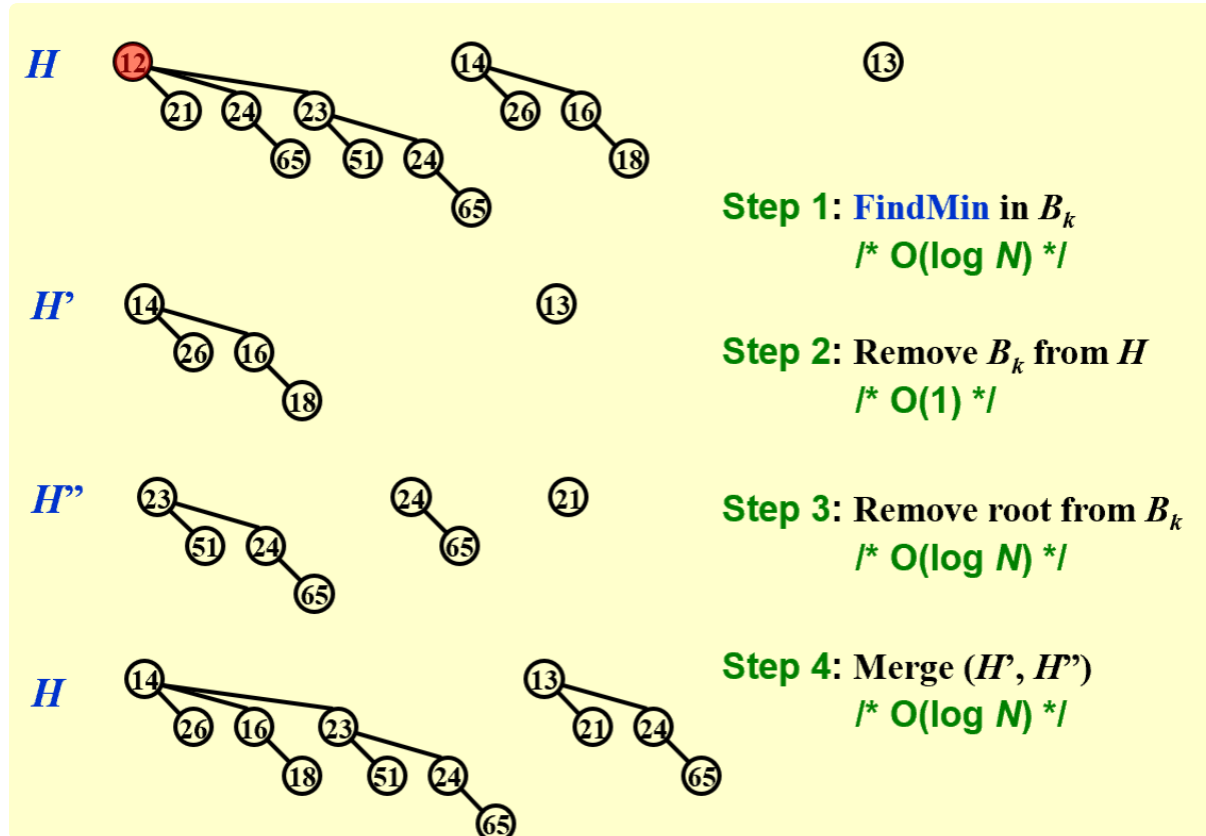
Insert is a special case for merging

【Example】 Insert 1, 2, 3, 4, 5, 6, 7 into an initially empty queue.



Performing N inserts on a initially empty binomial queue will take $O(N)$ worst-case time. Hence the average time is constant.

DeleteMin(H)



Implementation

Binomial queue = array of binomial trees

Operation	Property	Solution
DeleteMin	Find all the subtrees quickly	
Merge	The children are ordered by their sizes	

Decreasing.

Child – Sibling

Structure:

▼ Declaration

C | 复制代码

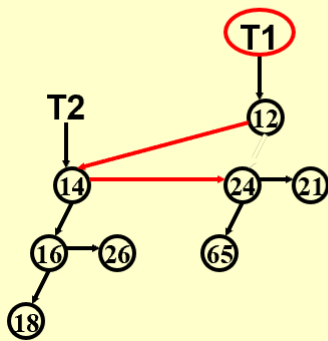
```

1  typedef struct BInNode *Position
2  typedef struct Collection *BinQueue
3  typedef struct BinNode *BinTree
4
5  struct BinNode
6  {
7      ElementType element;
8      Position LeftChild;
9      Position NextSibling;
10 }
11
12 struct Collection
13 {
14     int CurrentSize;
15     BinTree TheTrees[MaxTrees];
16 }

```

```
1  BinTree CombineTrees(BinTree T1,BinTree T2)
2  {
3      //merge equal-sized T1 and T2
4      if(T1->Element > T2->Element)
5          //attach the larger one to the smaller one
6          return CombineTrees(T2,T1)
7      //insert T2 to the front of the children list of T1
8      T2->NextSibling = T1->LeftChild;
9      T1->LeftChild = T2;
10     return T1;
11 }
```

$$T_p = O(1)$$



```
1  BinQueue Merge(BinQueue H1, BinQueue H2)
2  {
3      BinTree T1,T2,Carry = NULL;
4      int i,j;
5      if(H1->CurrentSize + H2->CurrentSize > Capacity)
6          ErrorMessage();
7
8      H1->CurrentSize += H2->CurrentSize;
9      for(i=0,j=1;j<=H1->CurrentSize;i++,j+=2)
10     {
11         //current trees
12         T1 = H1->TheTrees[i];
13         T2 = H2->TheTrees[i];
14         switch(4*!!carry + 2*!!T2+!!T1)
15         {
16             case 0: /*000*/
17             case 1: /*001*/ break;
18             case 2: /*010*/
19                 H1->TheTree[i] = T2;
20                 H2->TheTree[i] = NULL;
21                 break;
22             case 4: /*100*/
23                 H1->TheTree[i] = Carry;
24                 Carry = NULL;
25                 break;
26             case 3: /*011*/
27                 Carry = BombineTree(T1,T2);
28                 H1->TheTree[i] = H2->TheTree[i] = NULL;
29                 break;
30             case 5: /*101*/
31                 Carry = BombineTree(T1,Carry);
32                 H1->TheTree[i] = NULL;
33                 break;
34             case 6: /*110*/
35                 Carry = BombineTree(T2,Carry);
36                 H2->TheTree[i] = NULL;
37                 break;
38             case 7: /*111*/
39                 H1->TheTrees[i] = Carry;
40                 Carry = BombineTree(T1,T2);
41                 H2->TheTree[i] = NULL;
42                 break;
43         }
44     }
45     return H1;
```

▼ DeleteMin

C | 复制代码

```

1  ElementType DeleteMin(BinQueue H)
2  ▼ {
3      BinQueue DeleteQueue;
4      Position DeleteTree, OldRoot;
5      Elementtype MinItem = Infinity;
6      int i, j, MinTree;
7
8      if(IsEmpty(H)) {PrintErrorMessage();return -Infinity;}
9
10     for(i=0; i<MaxTrees; i++) // Find the minimum item
11     ▼ {
12         if(H->TheTrees[i] && H->TheTrees[i] ->Element < MinItem)
13     ▼     {
14         MinItem = H->TheTrees[i] ->Element;
15         MinTree = i;
16     }
17     }
18
19     DeleteTree = H->TheTrees[MinTree];
20     H->TheTree[MinTree] = NULL; //remove the MinTree
21     OldRoot = DeleteTree; // remove the root;
22     DeletedTree = DeletedThee->LeftChild;
23     Free(OldRoot);
24     DeleteQueue = Initialize();
25     DeleteQueue->CurrentSize = (1<<MinTree) - 1; //2^MinTree - 1
26     for(j = MInTree - 1; j>=0; j--)
27     ▼ {
28         DeleteQueue->TheTrees[j] = DeletedTree;
29         DeletedTree = DeletedTree->Nextsibling;
30         DeletedQueue->TheTrees[j]->NextSibling = NULL;
31     }
32     H->CurrentSize -= DeleteQueue->CurrentSize + 1;
33     H = Merge(H, DeleteQueue);
34     return MinItem;
35 }

```

【Chaim】

A binomial queue of N elements can be built by N successive insertions in $O(N)$ time.

2-1 分数 1

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Which of the following binomial trees can represent a binomial queue of size 42?

- ☐ A. $B_0 B_1 B_2 B_3 B_4 B_5$
- ☒ B. $B_1 B_3 B_5$
- ☐ C. $B_1 B_5$
- ☐ D. $B_2 B_4$

答案正确: 1 分

💡 创建提问

2-2 分数 1

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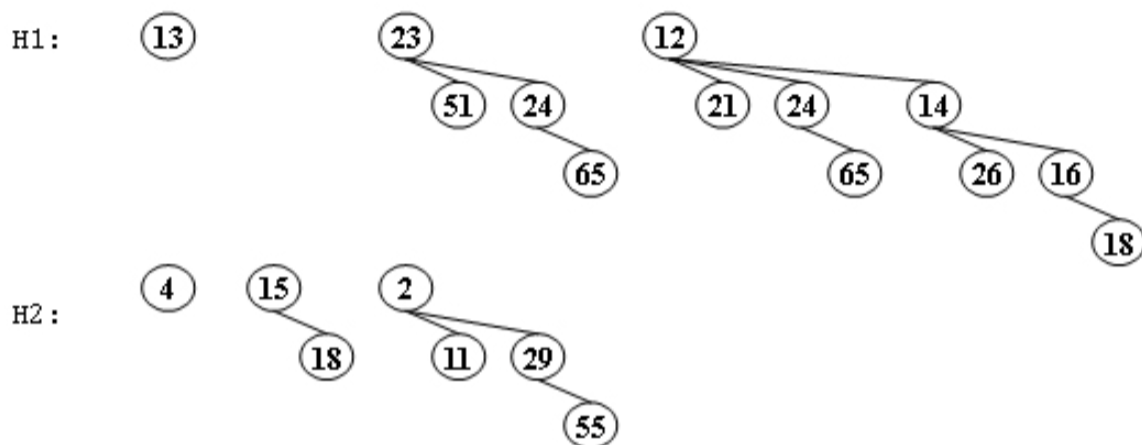
For a binomial queue, __ takes a constant time on average.

- ☐ A. merging
- ☐ B. find-max
- ☐ C. delete-min
- ☒ D. insertion

答案正确: 1 分

💡 创建提问

Merge the two binomial queues in the following figure. Which one of the following statements must be FALSE?

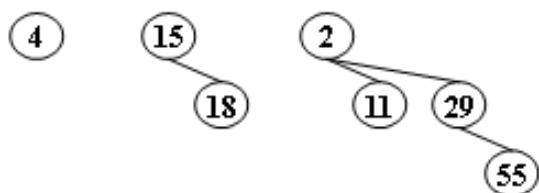


- ☐ A. there are two binomial trees after merging, which are B_2 and B_4
- ☐ B. 13 and 15 are the children of 4
- ☐ C. if 23 is a child of 2, then 12 must be another child of 2
- ☒ D. if 4 is a child of 2, then 23 must be another child of 2

答案正确: 2 分

💡 创建提问

Delete the minimum number from the given binomial queues in the following figure. Which one of the following statements must be FALSE?



- ☐ A. there are two binomial trees after deletion, which are B_1 and B_2
- ☐ B. 11 and 15 can be the children of 4
- ☒ C. 29 can never be the root of any resulting binomial tree
- ☐ D. if 29 is a child of 4, then 15 must be the root of B_1

答案正确: 2 分

💡 创建提问