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Report

The analysis of the space and time complexity of the four non-recursive tree:

The space complexity for the postorder, preorder and inorder non-recursive tree traversals would all be $O(h)$ (h is the depth of the binary tree), since I am only using one stack to implement the function. And the time complexity would be $O(m + n)$ which can be finalized into $O(n)$, too.

The time complexity for the Breadth first search would be $O(n^2)$ since each node will donate two child node and n is the number of total number of nodes in the tree. The space complexity for the breadth first search would be $O(n)$ since the function save the whole tree in the stack and pop one by one.

The tabulation of the number of comparisons for three different search case:

The best case for all three case would be the key is the head of the tree or list, it will only one comparison to find it, time complexity = $O(1)$. The worst for the list and the unbalanced tree would be all the elements are identical and searching for a different element. Then the function would search all the nodes to prove that the element doesn't exist with time complexity $O(n)$. The worst case for the balanced tree is that the element at the bottom of the tree which will take $O(\log n)$ to find it.

Summary: It will take $O(n)$ to make the linked list, $O(n \log n)$ to build the tree from the linked list, and $O(n \log n)$ to build the balanced tree, too. Because instead of balance the unbalanced tree I made two new pointers and built a new balanced tree from list.