

Scheduling Algo

Premptive (can interfere)

(No Interference)

Non-Premptive

→ SRTF

→ LRTF

→ Round Robin

→ Priority based

→ FCFS

→ SJF

→ LRF

→ HRRN

→ Multilevel queue

Scheduling Criteria

Arrival Time: At which process enters the ready queue or state

Burst Time: Time required by a process to get execute on (Duration) CPU

Completion time: The time at which process complete its execution

Turn around time: $\{ \frac{\text{Completion Time} - \text{Arrival Time}}{\text{Time}} \}$

Waiting Time: $\{ \frac{\text{Turn around time} - \text{Burst time}}{\text{time}} \}$

Response time: $\{ \frac{\text{The time at which process get CPU First time} - \text{Arrival time}}{\text{time}} \}$

First Come First Serve (FCFS)

Process	AT	BT	CT	TAT	WT	RT
P ₁	0	2	2	2	0	0
P ₂	1	2	4	3	1	1
P ₃	5	3	8	3	0	0
P ₄	6	4	12	6	2	2

Criteria: Arrival Time
 Mode: Non-preemptive

$$\text{Avg TAT} =$$

Grantt chart



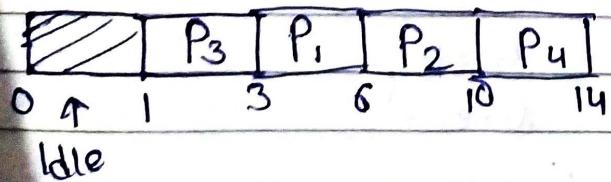
$$\text{Avg WT}$$

Shortest Job First

Process	AT	BT	CT	TAT	WT	RT
P ₁	1	3	6	5	2	2
P ₂	2	4	10	8	4	4
P ₃	1	2	3	2	0	0
P ₄	4	4	14	10	6	6

Criteria: Burst Time
 Mode: Non-preemptive

Grantt chart



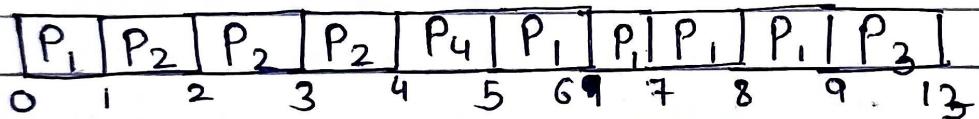
Shortest Remaining Time First

Criteria: Burst Time

Mode: Preemptive

Process	AT	BT	CT	TAT	WT	RT
P ₁	0	5	4	9	4	0
P ₂	1	3	2	4	3	0
P ₃	2	4	12	11	7	7
P ₄	4	1	5	1	0	0

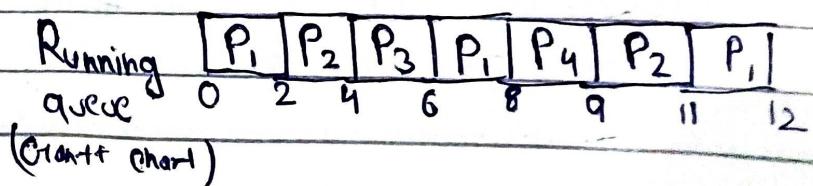
Gantt chart



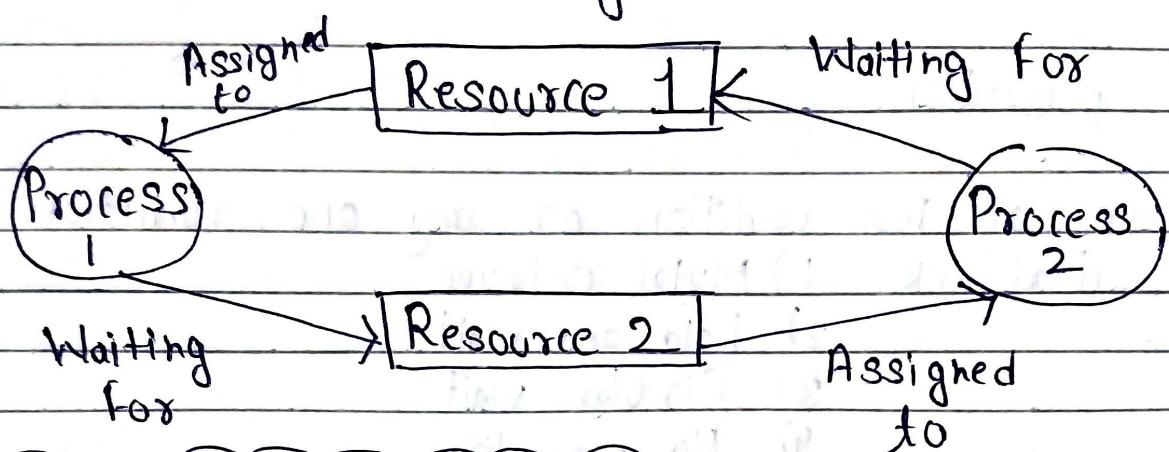
Round Robin

Process	AT	BT	CT	Given TQ = 2
P ₁	0	5	12	
P ₂	1	4	11	
P ₃	2	2	6	
P ₄	4	1	9	

Criteria: Time Quantum
Mod: Preemptive



Deadlock: It is a situation in computer system where multiple process or threads are unable to proceed because each is waiting for a resource that the others are holding.



Characterization of Deadlock

- 1) Mutual Exclusion: Processes compete for exclusive access to resources, meaning that only one process can use a resource at a time.
- 2) Hold and Wait: Processes hold resources while waiting for additional ones, creating the potential for resource contention.
- 3) No Preemption: Resources cannot be forcibly taken from a process; they must be released voluntarily by the process holding them.
- 4) Circular Wait: A circular wait chain of two or more processes exists with each process waiting for a resource held by the next process in chain.

Handling deadlock

1) Deadlock Ignorance (Ostrich method)

'Don't affect performance'

2) Deadlock prevention

- Remove all four condition or any one condition of deadlock
 - 1) Mutual exclusive
 - 2) hold and wait
 - 3) Circular wait
 - 4) No preemption

3) Deadlock Avoidance (Banker's Algo)

Date / /

Baner's Algo (Deadlock Avoidance) Deadlock Detection

Process	Allocation			Max need			Available			Remaining (max - alloc)		
	A	B	C	A	B	C	A	B	C	A	B	C
P ₁	0	1	0	7	5	3	3	3	2	7	4	3 ✓
P ₂	2	0	0	3	2	2	5	3	2	1	2	2 ✓
P ₃	3	0	2	5	0	2				6	0	0 X
P ₄	2	1	1	4	2	2	7	4	3	2	1	1 ✓
P ₅	0	0	2	5	3	3	7	4	5	5	3	1 ✓
	7	2	5				9	5	0			

$$\text{Total } A = 10 \quad B = 5 \quad C = 7 \quad \frac{3}{10} \quad \frac{0}{5} \quad \frac{2}{7}$$

$$\frac{-7}{3} \quad \frac{-2}{3} \quad \frac{-5}{2}$$

Safe \rightarrow No deadlock

$P_2 \rightarrow P_4 \rightarrow P_5 \rightarrow P_1 \rightarrow P_6$, Unsafe \rightarrow Deadlock

Safe sequence