

Logical operator

OR

||

a || b

AND

&&

a && b

Not

!

a
0000 0000 outputs
0000 0001

OR | 0010 0000

OR | 0000 0001

Bitwise operator

OR

|

&

~

0010
0000

<u>Hex</u>	<u>ASCII</u>	<code>char a = 'a'; printf ("%c", a);</code>
a ↓ <u>8-bits</u>	- z . A B - y Z , 0 9	<u>256</u> 0 1 2 . . 1

8-bits → 2^8 diff. value

Unicode : 16-bit - 32-bit

255

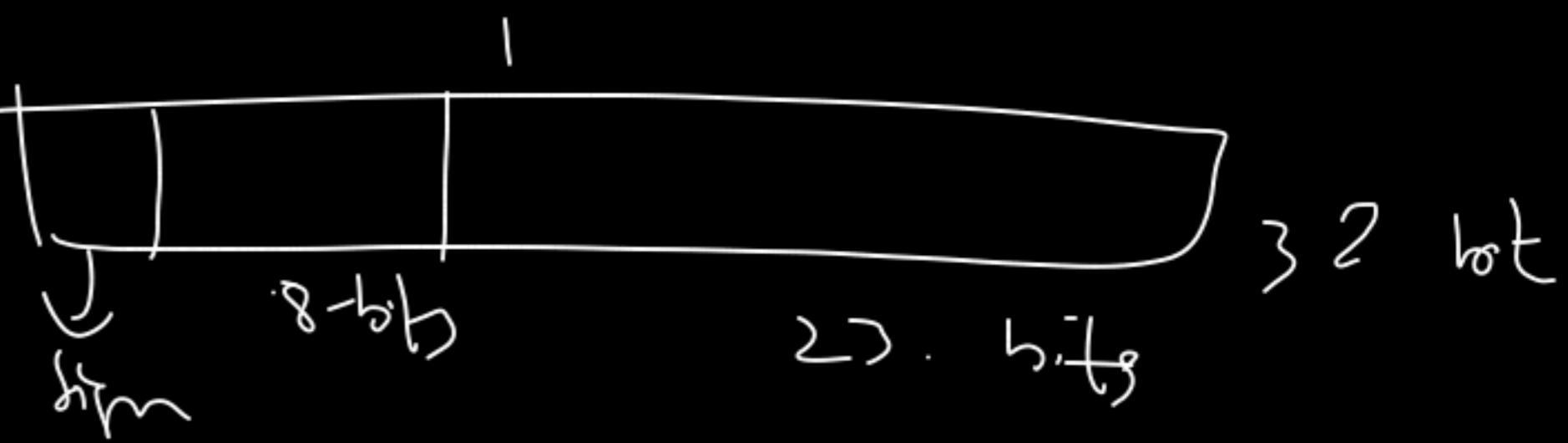
$$\text{Normal form} \leftarrow [1.f \times 2^{\text{exp}}]$$

$f \rightarrow \text{fraction}$

1.f \rightarrow Significant.

for Normal form

$$\text{Subnormal form} \leftarrow \underbrace{\left(\begin{array}{c} 1.\underline{1101} \times 2^{-16} \\ 0.\underline{111101} \end{array} \right)}_{0.f \times 2^{\text{exp}}}$$



- (1) Highest
- 16 lowest
- 4-bit

32-bit Normal

0	0 0 0 0 0 0 0 0 1	0 0 0 0 0 0 0 0
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$0 \cdot f \times 2^{\text{exp}}$

0 0 0 0 0 0 0 0	0 0 - - 1 0
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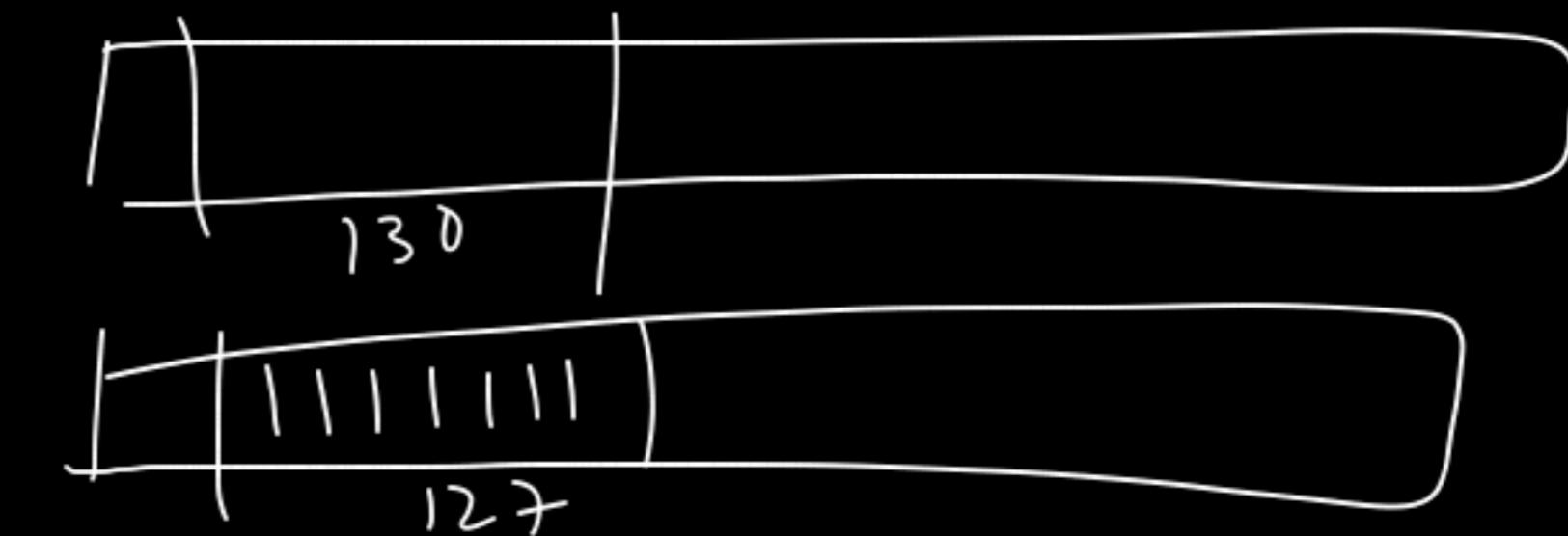
symbol power

if $\text{exp} = 0$
 $\& f \neq 0$, value $\leftarrow 0 \cdot f \times 2^{\text{exp}}$

if $\text{exp} = 0, f = 0$, value = 0

Addition of floats

$$+ \begin{array}{r} 1.11 \times 2^2 \\ 0.0011 \times 2^3 \\ \hline \end{array} \leftarrow 0.11 \times 2^1$$



Step 1: Make exponents equal, make lower \rightarrow larger
Step 2: Calculate diff, $d = \lceil \exp_1 - \exp_2 \rceil$

Step 3: $\text{sig}_2 \gg d$

Step 3 → Add $\text{sig}_1 + \text{sig}_2$ (1.1111)

4.59

4.5

Step 4 → Round off

4.5

Step 5 Normalize \times power
 \downarrow

Round-off
 → Take closest possible

$$\begin{array}{rcl} & \overbrace{}^1 & \\ 6 & 5.56 \rightarrow 5.6 & 6 \\ 5 & 5.39 \rightarrow 5.4 & 5 \\ & 5.46 \rightarrow 5.5 & 6 \\ & \overbrace{}^1 & \end{array}$$

In floating point ops

Guard bit

Floating point Multiplication

$$\begin{array}{r}
 0.11 \\
 0.1 \\
 \hline
 1.1
 \end{array}
 \times
 \begin{array}{r}
 1 \cdot f_1 \times 2^{e_1} \\
 1 \cdot f_L \times 2^{e_L}
 \end{array}$$

$$\begin{array}{c}
 \boxed{1 \cdot sf_1} \\
 \boxed{1 \cdot sf_L}
 \end{array}$$

$$f = \underline{sf_1 + sf_L - 127}$$

1.10

Step 2

Add exponents

$$e_1 + e_L = e$$

Step 2: Multiply $f_{\text{fig. 1}} \times f_{\text{fig. 2}}$

Step 3 → Normalization ($\overset{\text{left}}{\text{Shift}} \text{ Result} \& e = e + \# \text{shift}$)

Circuit for Int multiplication

$$\begin{array}{c}
 x \\
 \downarrow \\
 x_{n-1} \dots x_0
 \end{array}
 \quad
 \begin{array}{c}
 y \\
 \downarrow \\
 y_{n-1} \dots y_0
 \end{array}$$

$$\begin{array}{r}
 x \cdot 0^{\text{b}} \cdot 1^{\text{b}} \\
 \times 1101 \\
 \hline
 \end{array}$$

$s = 0$
for $i = 1$ to 32

$$\begin{aligned}
 \text{if } y[0] &= 1 \\
 s &= s + x
 \end{aligned}$$

$$\begin{aligned}
 y &= y \gg 1 \\
 x &= x \ll 1
 \end{aligned}$$

