

Lecture 2

Container

Set : Stores only unique elements in sorted manner

Declare : `set<int> st;`

Function

`st.insert(2);` // $O(\log n)$ time

X No random access `st[0]` X

access by iterators = ~~`*st.begin()`~~ `*st.begin()`

Function :-

`st.erase(st.begin())` // $O(\log n)$

→ iterator to delete

// $O(\log n)$ `st.erase(st.begin(), st.begin() + 2);` []

✓ deletes first 2 elements.
~~length you deleted~~

`st.erase(5);` // $O(\log n)$

Declare :- `set<int> st = {1, 5, 7, 8};`

`set<int> st1(st.begin(), st.end());`

Function :-

`auto it = st.find(7);` // $O(\log n)$

If 7 exists in st, ~~it~~ it points to 7

If 7 doesn't exist in st, ~~it~~ `it = st.end();`

Function

`st.emplace(6);` // emplace faster than insert.

Function

`st.size();` // gives size of set.

Function

`st.erase(st.begin(), st.end());` // delete all elements.

Container

`unordered_set<int> st;`

- 1) No guaranteed ordering of elements.
 - 2) Average time complexity is $O(1)$
 - 3) Worst case time complexity is $O(\text{set size})$ ($O(n)$)
- Worst case depends on hashes generated.

→ Thumb Rule:-

If no need of storing elements in ascending order
go for unordered-set(); if you get TLE → switch
to set()

Container

multiset<int> ms; stores duplicate also
stores all elements in sorted manner.

function

Function

ms.erase(2) // erase all instances of 2 in multiset

~~auto~~ auto it = ms.find(2); // returns iterator to the first
instance of 2.

$O(\log n)$ operations

function

~~ms~~ ms.erase(ms.find(2)). ← only delete 1 instance
ms.erase(ms.find(2), ms.find(2)+2) ← delete 2
instances of 2.

function

ms.count(2); // no. of instances of 2

× × × × × × × ×

Container

Map:- Key value pairs.

Only unique keys; $O(\log n)$ operations

Dedare: map<string, int> mpp;

mpp["John"] = 100;

Function

mpp.emplace("Tiger", 40);

function

mpp.erase("John");

mpp.erase(mpp.begin());

mpp.clear();

mpp.erase(mpp.begin(), mpp.begin()+2);

auto it = mpp.find("John");

mpp.empty() → true if map empty; false otherwise

function: `mpp.count("raj")` // One in case map if "raj" key exists, zero otherwise

Pair:
`pair<int, int> pr;`
`pr.first = 1;`
`pr.second = 2;`

it → pointer to the pair
 it → first ~~or~~ (*it).first

Printing map:

for (auto it : mpp)
 {
 cout << it.first << " " << it.second << endl;
 }
 Pair

for (auto it = mpp.begin(); it != mpp.end(); it++)
 {
 cout << it->first << " " << it->second << "\n";
 }

Container :- `unordered_map<int, int> mpp;`
 // Random order; O(1) in almost all cases
 // O(n) in the worst case.

~~Container~~
PAIR:
`pair<int, int> pr = {1, 2};`
~~pr~~ `pr = make_pair(1, 2);`
`pair<pair<int, int>, int> pr = {{1, 2}, 3};`
`pr.first.first = 1; pr.first.second = 2;`
`pr.second = 3;`

unordered_set < pair<int, int> > ✗ Not Allowed.
unordered_map < pair<int, int>, int > ✗ Not allowed.
~~map~~ `map<pair<int, int>, int>` ✓ Allowed
`set<pair<int, int>>` ✓ Allowed.
`vector<pair<int, int>>` ✓ Allowed

Container :- `multimap<string, int> mpp;`
 # stores duplicate keys also.
 # Sorted fashion.

No iterator concept in Stack - Queue

Data Structures

Stack and Queue

stack <int> st; Last In First Out D.S.

// pop, top, size, empty, push and emplace // $O(1)$ operation

No clear function No clear() function

Delete :- To delete entire stack \Rightarrow while (!st.empty()) // linear time
 $\{$
 3. st.pop();
 $\}$

stack <int> st;

"cout << st.top(); or st.pop() on empty stack \Rightarrow throws error.

Usually runtime errors in online judges

Imp :-

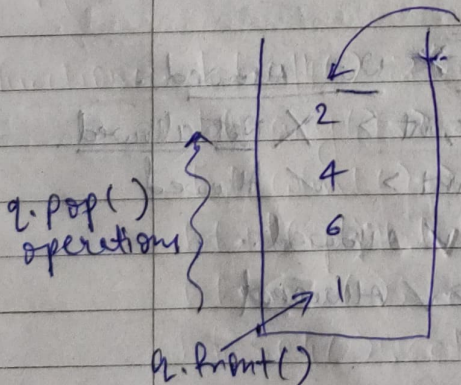
Always use top() & pop() with a check condition
 if (!st.empty())
 $\{$
 cout << st.top(); or st.pop();
 $\}$

Data Structures

Queue First In First Out D.S

queue <int> q;

// push, emplace, front, pop, size, empty // $O(1)$ operations



q.push(1)

" " (6)

" " (4)

" " (2)

clear

To clear whole queue use while loop.

while (!q.empty()) // linear time

$\{$
 q.pop();
 $\}$

No iterator concept

Data Structure

Priority-queue (Most important). // Use concept of Heap Sort
 → stores all elements in sorted order.

→ $O(\log n)$ operations.

// push, size, top, pop, empty

priority-queue <int> pq; → MAXIMUM PRIORITY QUEUE :-

pq.push(6)	6	← top	} pq.top() operation
" (1)	4		
" (2)	2		
" (4)	1		

→ To store element like minimum priority queue, negate elements while inserting.

pq.push(-6)	-1	← top	# while fetching do negation eg again:- cout << (-1) * pq.top();
" (-1)	-2		
" (-2)	-4		
" (-4)	-6		

Data Structure

* Minimum priority Queue :-

Means Ultra

priority-queue <int, ~~vector~~ vector<int>, greater<int>> pq;

pq.push(6)	1	← top()	} top operations
pq.push(1)	2		
pq.push(2)	4		
pq.push(4)	6		

* for min. priority queue of pairs, replace int with pair.
 priority-queue <pair<int, int>, vector<pair<int, int>>, greater<pair<int, int>>>

Containers

Deque: Doubly ended Queue

`deque<int> dq;`

`O(1)` push-front # pop-back

// push-back // begin, end, rbegin, rend

// pop-front // size // empty

// clear // at

All functions of vector

Container

List

`list<int> ls;`

All operations of deque

// remove() → O(1) operation

eg. `ls.push-back(1)`

" " (2)

" " (0)

`ls.remove(2);` // O(1) operation

Containers

Bitset → 1 bit

int: ~~32~~ 32 bits

char: 8 bits

Declare

`bitset<10> bt;` → 10 bits

Input

`cin >> bt;` // 1011011100

function

`bt.all()` → returns true if all bits are set, false otherwise

`bt.any()` → " " " at least one bit set, false "

`bt.count()` → returns no. of set bits

`bt.flip()` `bt.flip(index)`; 0 → 1 or 1 → 0

`bt.none()`; returns true if none bits set, false otherwise

`bt.set()`; sets all bits; `bt.set(index)`; `bt.set(index, 0/1)`

`bt.reset()`; make all bits 0; `bt.reset(index)`

`bt.size()`; returns size of bitset

`bt.test(index)`; true if bit at index is set, false otherwise