

## Laboratory practice No. 5: Divide to Conquer and Dynamic Programming

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### 3) Practice for final project defense presentation

#### 3.2 order $O(n^2 \cdot 2^n)$ with 50 data ( $50^2 \cdot 2^{50}$ )

**3.3** The data structure used to solve this problem was an arraylist composed by pairs. The algorithm uses its own method to compare doubles, which is in charge of taking the starting point, it means, the spot where the robot is located, and then it compares which is the closest point according to the data related to the radioactive garbage, then when the algorithm knows which one is the fewer, it returns it, and it moves one square, it could be in a horizontal way or a vertical way and each movement has a cost of 1. To be able to find the minimum distance, we use an algorithm to find the hypotenuse, it's known as Pythagoras..

**3.4** The question is repeated in the document.

**3.5** the order of the algorithm (without taking into account when looking for the closest current point)  $O(N \cdot K)$

#### 3.6

n is the number of maps

K is the number of waste

### 4) Practice for midterms

#### 4.1

	$\epsilon$	c	a	l	l	e
$\epsilon$	0	1	2	3	4	5
c	1	0	1	2	3	4
a	2	1	0	1	2	3
s	3	2	1	1	2	3
a	4	3	2	2	2	3

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**ESTRUCTURA DE DATOS 2**  
**Código ST0247**

**4.2**

	$\epsilon$	m	a	d	r	e
$\epsilon$	0	1	2	3	4	5
m	1	0	1	2	3	4
a	2	1	0	1	2	3
m	3	2	1	1	2	3
a	4	3	2	2	2	3

**4.2**

4.2.1  $O(\text{length } x * \text{length } y)$

4.2.2 return table [lenx][leny];

**4.3**

4.3.1 a)  $O(n)$

4.3.2 c)  $T(n) = c_1 2n + c_2$

**4.4** c)**4.5**

4.5.1 c)  $T(n) = T(n/2) + C$  que es  $O(\log n)$

4.5.2 a[mitad]

4.5.3 a, mitad + 1, de, z

**4.6**

4.6.1 scm[i]=0

4.6.2 scm[i]=arr[j];

4.6.3 max++;

4.6.4 c)  $O(n^2)$

**4.7**

4.7.1 int ni = g[i][j];

4.7.2 int nj = g[k][j];

4.7.3 int nk = g[i][k];

4.7.4  $O(n^3)$

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