CMI ALGEBRA 1 (2021) ASSIGNMENT 1 - T. R. Ramadas

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1. Let $P: V \to V$ be a map such that

$$P = \frac{1}{2}(I_V - i\mathcal{J})$$

Now if $v_1, v_2 \in V$ then

$$P(v_1 + v_2) = \frac{1}{2}((v_1 + v_2) - i\mathcal{J}(v_1 + v_2)) = \frac{1}{2}(v_1 - i\mathcal{J}(v_1)) + \frac{1}{2}(v_2 - i\mathcal{J}(v_2)) = P(v_1) + P(v_2)$$

and

$$P(\lambda \cdot v_1) = \frac{1}{2}(\lambda \cdot v_1 - i\mathcal{J}(\lambda \cdot v_1)) = \frac{1}{2}(\lambda \cdot v_1 - i\lambda \cdot \mathcal{J}(v_1)) = \lambda \cdot \frac{1}{2}(v_1 - i\mathcal{J}(v_1)) = \lambda \cdot P(v_1)$$

Hence P is a linear map. Now

$$\begin{split} P(P(v)) &= P\bigg(\frac{1}{2}(v - i\mathcal{J}(v))\bigg) \\ &= \frac{1}{2}\bigg(\frac{1}{2}(v - i\mathcal{J}(v)) - iJ\bigg(\frac{1}{2}(v - i\mathcal{J}(v))\bigg)\bigg) \\ &= \frac{1}{2}\bigg(\frac{1}{2}(v - i\mathcal{J}(v)) - \bigg(\frac{1}{2}(i\mathcal{J}(v) + \mathcal{J}(\mathcal{J}(v)))\bigg)\bigg) \bigg) \\ &= \frac{1}{2}\bigg(\frac{1}{2}(v - i\mathcal{J}(v)) - \bigg(\frac{1}{2}(i\mathcal{J}(v) - v)\bigg)\bigg) \\ &= \frac{1}{2}(v - i\mathcal{J}(v)) \\ &= P(v) \end{split}$$

Hence for a vector $v \in image(P)$, P(v) = v

Now, if $v \in V_{-i}$ then

$$\mathcal{J}(v) = -iv$$

$$\Longrightarrow i\mathcal{J}(v) = v$$

$$\Longrightarrow v - i\mathcal{J}(v) = 0_V$$

$$\Longrightarrow \frac{1}{2}(v - i\mathcal{J}(v)) = 0_V$$

$$\Longrightarrow P(v) = 0_V$$

Hence if $v \in V_{-i}$ then $v \in ker(P)$ therefore

$$V_{-i} \subseteq ker(P)$$

Now let $v \in ker(P)$ then

$$P(v) = 0_{V}$$

$$\Longrightarrow \frac{1}{2}(v - i\mathcal{J}(v)) = 0_{V}$$

$$\Longrightarrow v - i\mathcal{J}(v) = 0_{V}$$

$$\Longrightarrow v = i\mathcal{J}(v)$$

$$\Longrightarrow iv = -\mathcal{J}(v)$$

$$\Longrightarrow \mathcal{J}(v) = -iv$$

Hence if $v \in ker(P)$ then $v \in V_{-i}$ therefore

$$ker(P) \subseteq V_{-i}$$

Hence

$$V_{-i} = ker(P)$$

Now if $v \in V_i$ then

$$\mathcal{J}(v) = iv$$

$$P(v) = \frac{1}{2}(v - i\mathcal{J}(v))$$
$$= \frac{1}{2}(v - i(iv))$$
$$= \frac{1}{2}(v + v)$$
$$= v$$

Hence if $v \in V_i$ then $v \in image(P)$ therefore

$$V_i \subseteq image(P)$$

Now let $v \in image(P)$ then $\exists u \in V$ such that $P(u) = v = \frac{1}{2}(u - i\mathcal{J}(u))$. Therefore

$$\mathcal{J}(v) = \mathcal{J}(P(u))$$

$$= \mathcal{J}\left(\frac{1}{2}(u - i\mathcal{J}(u))\right)$$

$$= \frac{1}{2}(\mathcal{J}(u) - \mathcal{J}(i\mathcal{J}(u)))$$

$$= \frac{1}{2}(\mathcal{J}(u)) - i\mathcal{J}(\mathcal{J}(u))$$

$$= \frac{1}{2}(\mathcal{J}(u)) + iu)$$

$$= i\frac{1}{2}(-i\mathcal{J}(u) + u)$$

$$= iP(u) = iv$$

Hence if $v \in image(P)$ then $v \in V_i$ therefore

$$image(P) \subseteq V_i$$

Hence

$$V_i = image(P)$$

Hence we need to prove that

$$V = ker(P) \oplus image(P)$$

Let $v \in V$, we can write v = v - P(v) + P(v). Then

$$P(v) = P(v - P(v) + p(v)) = P(vP(v)) + P(P(v))$$

Here

$$P(v - P(v)) = P(v) - P(P(v))$$
$$= P(v) - P(v)$$
$$= 0_V$$

Hence $v - P(v) \in ker(P)$ and $P(v) \in image(P)$. Hence any vector $v \in V$ it can be written as a sum of a vector from ker(P) and a vector from image(P). Hence

$$V \subseteq ker(P) \oplus image(P)$$

Now let $v \in ker(P) \oplus image(P)$. Then $\exists v_1 \in ker(P)$ and $v_2 \in image(P)$ such that

$$v = v_1 + v_2$$

As $v_1 \in ker(P)$, $v_1 \in V$. As the linear map P is $V \to V$, $image(P) \subseteq V$. Hence $v_2 \in V$ also. Therefore $v_1 + v_2 \in V$. Therefore

$$ker(P) \oplus image(P) \subseteq V$$

- . Hence $V = ker(P) \oplus image(P) = V_i \oplus V_{-i}$ [Proved]
- 2. (a) The sets that span (=generate) \mathbb{R}^3 are $\mathcal{S}_2, \mathcal{S}_3, \mathcal{S}_4$.
 - (b) The linearly independent sets are S_1 , S_2 .
 - (c) The bases are S_2 .
- 3. The set $\{(1,0,0),(x,y,0),(x',y',z')\}$ is a basis of \mathbb{R}^3 iff $y \neq 0, z' \neq 0$
- 4. V is a one dimensional vector space. Therefor any non zero vector of V is a basis of V.
- We are considering the characteristic 0 fields.

Let $f \in \mathbb{F}[x_1, x_2, \dots, x_{n-1}, y]$ be an *n*-variate degree d polynomial which can be computed by an arithmetic circuit of size at most s. If g and h are polynomials of degree at least 1 such that f = gh and gcd(g, h) = 1, then g and h have a circuit of size at most poly(s, n, d). Let $f \in \mathbb{F}[x_1, x_2, \dots, x_{n-1}, y]$ be an n-variate degree d polynomial which can be computed by an arithmetic circuit of size at most s. If there is a polynomial s and an integer s such that s and s are circuit of size at most s and s are circuit of size at most s.