

Functional Programming and Interactive Theorem Proving with Isabelle/HOL

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Introduction

Programs may have bugs.

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Hunting bugs:

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Hunting bugs: Testing?

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BUT: when done one paper, its likely to have errors in proof!

This lecture: Using a computer to check proofs

Material

The Theorem Prover Isabelle/HOL:

<https://isabelle.in.tum.de/>

Lecture Material:

https://github.com/lammich/MCR_SS_2019_FunProgProve

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Relax and enjoy! There is no exam on this lecture!

Short Poll

Raise your hand if you

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 - Others?
- Know what quicksort is

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- Know what quicksort is
- Know what (structural) induction means

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 - Others?
- Know what quicksort is
- Know what (structural) induction means
- Have ever used an interactive theorem prover

Lists

$[]$ Empty list

$a \# l$ List with first element a and then list l

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Example: $1\#2\#3\#4\#[]$

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Notation: $[1, 2, 3, 4]$

Lists

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Example: $1\#2\#3\#4\#[]$

Notation: $[1, 2, 3, 4]$

$l_1 @ l_2$: concatenate two lists

Example: $[1, 2, 3] @ [4, 5, 6] = [1, 2, 3, 4, 5, 6]$

Append @

How to define @ ?

Append @

How to define @ ? Using only [] and #?

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

[] @ l₂ =

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

$$(x \# l_1) @ l_2 =$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

$$(x \# l_1) @ l_2 = x \# (l_1 @ l_2)$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

$$(x \# l_1) @ l_2 = x \# (l_1 @ l_2)$$

Example:

$$([1,2] @ [3])$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

$$(x \# l_1) @ l_2 = x \# (l_1 @ l_2)$$

Example:

$$([1,2] @ [3]) = (1 \# 2 \# []) @ (3 \# [])$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

$$(x \# l_1) @ l_2 = x \# (l_1 @ l_2)$$

Example:

$$\begin{aligned}([1,2] @ [3]) &= (1 \# 2 \# []) @ (3 \# []) \\ &= 1 \# ((2 \# []) @ (3 \# []))\end{aligned}$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$\begin{aligned} [] @ l_2 &= l_2 \\ (x \# l_1) @ l_2 &= x \# (l_1 @ l_2) \end{aligned}$$

Example:

$$\begin{aligned} ([1,2] @ [3]) &= (1 \# 2 \# []) @ (3 \# []) \\ &= 1 \# ((2 \# []) @ (3 \# [])) \\ &= 1 \# 2 \# ([] @ (3 \# [])) \end{aligned}$$

Append @

How to define @ ? Using only [] and #?

Case distinction whether first list is empty:

$$[] @ l_2 = l_2$$

$$(x \# l_1) @ l_2 = x \# (l_1 @ l_2)$$

Example:

$$([1,2] @ [3]) = (1 \# 2 \# []) @ (3 \# [])$$

$$= 1 \# ((2 \# []) @ (3 \# []))$$

$$= 1 \# 2 \# ([] @ (3 \# []))$$

$$= 1 \# 2 \# 3 \# []$$

Filter

Erase all elements `not ≤ 4` from a list

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`leq4 [1,42,7,5,2,6,3]`

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`leq4 (x#l) =`

Filter

Erase all elements `not ≤ 4` from a list

`leq4 [1,42,7,5,2,6,3] = [1,2,3]`

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`leq4 (x#l) = if x \leq 4 then x # leq4 l else leq4 l`

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Erase all elements `not ≤ 4` from a list

`leq4 [1,42,7,5,2,6,3] = [1,2,3]`

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`leq4 [1, 42, 7, 5, 2, 6, 3]`

Filter

Erase all elements *not* ≤ 4 from a list

leq4 [1,42,7,5,2,6,3] = [1,2,3]

leq4 [] = []

leq4 (x#l) = if $x \leq 4$ then x # *leq4* l else *leq4* l

leq4 [1, 42, 7, 5, 2, 6, 3]

= 1 # *leq4* [42, 7, 5, 2, 6, 3]

Filter

Erase all elements **not** ≤ 4 from a list

leq4 [1,42,7,5,2,6,3] = [1,2,3]

leq4 [] = []

leq4 (x#l) = if $x \leq 4$ then x # *leq4* l else *leq4* l

leq4 [1, 42, 7, 5, 2, 6, 3]

= 1 # *leq4* [42, 7, 5, 2, 6, 3]

= 1 # *leq4* [7, 5, 2, 6, 3]

Filter

Erase all elements `not ≤ 4` from a list

`leq4 [1,42,7,5,2,6,3] = [1,2,3]`

`leq4 [] = []`

`leq4 (x#l) = if x \leq 4 then x # leq4 l else leq4 l`

`leq4 [1, 42, 7, 5, 2, 6, 3]`

`= 1 # leq4 [42, 7, 5, 2, 6, 3]`

`= 1 # leq4 [7, 5, 2, 6, 3]`

`= 1 # leq4 [5, 2, 6, 3]`

Filter

Erase all elements $\text{not } \leq 4$ from a list

leq4 [1,42,7,5,2,6,3] = [1,2,3]

leq4 [] = []

leq4 (x#l) = if $x \leq 4$ then $x \# \text{leq4 } l$ else *leq4* l

leq4 [1, 42, 7, 5, 2, 6, 3]

= 1 # *leq4* [42, 7, 5, 2, 6, 3]

= 1 # *leq4* [7, 5, 2, 6, 3]

= 1 # *leq4* [5, 2, 6, 3]

= 1 # *leq4* [2, 6, 3]

Filter

Erase all elements $\text{not } \leq 4$ from a list

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3] = [1, 2, 3]$

$\text{leq4 } [] = []$

$\text{leq4 } (x\#l) = \text{if } x \leq 4 \text{ then } x \# \text{leq4 } l \text{ else leq4 } l$

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [5, 2, 6, 3]$

$= 1 \# \text{leq4 } [2, 6, 3]$

$= 1 \# 2 \# \text{leq4 } [6, 3]$

Filter

Erase all elements $\text{not } \leq 4$ from a list

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3] = [1, 2, 3]$

$\text{leq4 } [] = []$

$\text{leq4 } (x\#l) = \text{if } x \leq 4 \text{ then } x \# \text{leq4 } l \text{ else leq4 } l$

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [5, 2, 6, 3]$

$= 1 \# \text{leq4 } [2, 6, 3]$

$= 1 \# 2 \# \text{leq4 } [6, 3]$

$= 1 \# 2 \# \text{leq4 } [3]$

Filter

Erase all elements $\text{not } \leq 4$ from a list

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3] = [1, 2, 3]$

$\text{leq4 } [] = []$

$\text{leq4 } (x\#l) = \text{if } x \leq 4 \text{ then } x \# \text{leq4 } l \text{ else leq4 } l$

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [5, 2, 6, 3]$

$= 1 \# \text{leq4 } [2, 6, 3]$

$= 1 \# 2 \# \text{leq4 } [6, 3]$

$= 1 \# 2 \# \text{leq4 } [3]$

$= 1 \# 2 \# 3 \# \text{leq4 } []$

Filter

Erase all elements $\text{not } \leq 4$ from a list

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3] = [1, 2, 3]$

$\text{leq4 } [] = []$

$\text{leq4 } (x \# l) = \text{if } x \leq 4 \text{ then } x \# \text{leq4 } l \text{ else leq4 } l$

$\text{leq4 } [1, 42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [42, 7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [7, 5, 2, 6, 3]$

$= 1 \# \text{leq4 } [5, 2, 6, 3]$

$= 1 \# \text{leq4 } [2, 6, 3]$

$= 1 \# 2 \# \text{leq4 } [6, 3]$

$= 1 \# 2 \# \text{leq4 } [3]$

$= 1 \# 2 \# 3 \# \text{leq4 } []$

$= 1 \# 2 \# 3 \# []$

Filter

Condition as parameter to function

filter P l = l

filter P $(x \# l)$ = (**if** P x **then** $x \# \text{filter } P \ l$ **else** $\text{filter } P \ l$)

Filter

Condition as parameter to function

filter P $[] = []$

filter P $(x \# l) = (\text{if } P\ x \text{ then } x \# \text{filter } P\ l \text{ else filter } P\ l)$

filter $(\lambda x. x \leq 4)$ l : Same as *leq4*

Filter

Condition as parameter to function

filter P $[] = []$

filter P $(x \# l) = (\text{if } P\ x \text{ then } x \# \text{filter } P\ l \text{ else filter } P\ l)$

filter $(\lambda x. x \leq 4)$ l : Same as *leq4*

$\lambda x. x \leq 4$ is anonymous function, with parameter x .

Demo.thy

Functions

Count

count / *x* How often does element *x* occur in list *l*

Count

count / *x* How often does element *x* occur in list /

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 =

Count

count / x How often does element x occur in list /

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* =

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* =

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

count [2, 2, 1, 2] 2

=

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

count [2, 2, 1, 2] 2

= 1 + *count* [2, 1, 2] 2

=

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

count [2, 2, 1, 2] 2

= 1 + *count* [2, 1, 2] 2

= 1 + 1 + *count* [1, 2] 2

=

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

count [2, 2, 1, 2] 2

= 1 + *count* [2, 1, 2] 2

= 1 + 1 + *count* [1, 2] 2

= 1 + 1 + *count* [2] 2

=

Count

count $l\ x$ How often does element x occur in list l

Example: *count* $[1, 2, 3, 1, 2, 3, 2, 2]\ 2 = 4$

count $[]\ x = 0$

count $(y\ \# l)\ x = \text{if } x=y \text{ then } 1 + \text{count } l\ x \text{ else } \text{count } l\ x$

count $[2, 2, 1, 2]\ 2$

$= 1 + \text{count } [2, 1, 2]\ 2$

$= 1 + 1 + \text{count } [1, 2]\ 2$

$= 1 + 1 + \text{count } [2]\ 2$

$= 1 + 1 + 1 + \text{count } []\ 2$

$=$

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

count [2, 2, 1, 2] 2

= 1 + *count* [2, 1, 2] 2

= 1 + 1 + *count* [1, 2] 2

= 1 + 1 + *count* [2] 2

= 1 + 1 + 1 + *count* [] 2

= 1 + 1 + 1 + 0

=

Count

count *l* *x* How often does element *x* occur in list *l*

Example: *count* [1, 2, 3, 1, 2, 3, 2, 2] 2 = 4

count [] *x* = 0

count (*y* # *l*) *x* = **if** *x*=*y* **then** 1 + *count* *l* *x* **else** *count* *l* *x*

count [2, 2, 1, 2] 2

= 1 + *count* [2, 1, 2] 2

= 1 + 1 + *count* [1, 2] 2

= 1 + 1 + *count* [2] 2

= 1 + 1 + 1 + *count* [] 2

= 1 + 1 + 1 + 0

= 3

Sortedness

Is a list sorted? E.g. *sorted* [1, 2, 4] = *True*, *sorted* [1, 4, 3] = *False*

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sorted [] =

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sorted [] = *True*

sorted [x] =

Sortedness

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sorted [] = *True*

sorted [x] = *True*

Sortedness

Is a list sorted? E.g. *sorted* [1, 2, 4] = *True*, *sorted* [1, 4, 3] = *False*

sorted [] = *True*

sorted [x] = *True*

sorted (x # y # l) =

Sortedness

Is a list sorted? E.g. *sorted* [1, 2, 4] = *True*, *sorted* [1, 4, 3] = *False*

sorted [] = *True*

sorted [x] = *True*

sorted (x # y # l) = $x \leq y \wedge \text{sorted } (y \# l)$

Sortedness

Is a list sorted? E.g. $\text{sorted } [1, 2, 4] = \text{True}$, $\text{sorted } [1, 4, 3] = \text{False}$

$\text{sorted } [] = \text{True}$

$\text{sorted } [x] = \text{True}$

$\text{sorted } (x \# y \# l) = x \leq y \wedge \text{sorted } (y \# l)$

Note \wedge means "and".

Demo.thy

Count and Sortedness

Quicksort

Algorithm to sort a list

Quicksort

Algorithm to sort a list

Idea:

$$qs\ (p\ \# l) = qs\ (\text{elements} \leq p) @ [p] @ qs\ (\text{elements} > p)$$

Quicksort

Algorithm to sort a list

Idea:

$$qs\ (p \# l) = qs\ (\text{elements} \leq p) @ [p] @ qs\ (\text{elements} > p)$$

$qs\ [3, 2, 5, 4, 7]$

$=$

Quicksort

Algorithm to sort a list

Idea:

$$qs\ (p \# l) = qs\ (\text{elements} \leq p) @ [p] @ qs\ (\text{elements} > p)$$

$$\begin{aligned} & qs\ [3, 2, 5, 4, 7] \\ &= qs\ [2] @ [3] @ qs\ [5, 4, 7] \\ &= \end{aligned}$$

Quicksort

Algorithm to sort a list

Idea:

$$qs\ (p \# l) = qs\ (\text{elements} \leq p) @ [p] @ qs\ (\text{elements} > p)$$

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Quicksort

Algorithm to sort a list

Idea:

$$qs\ (p \# l) = qs\ (\text{elements} \leq p) @ [p] @ qs\ (\text{elements} > p)$$

$$\begin{aligned} & qs\ [3, 2, 5, 4, 7] \\ &= qs\ [2] @ [3] @ qs\ [5, 4, 7] \\ &= [2] @ [3] @ qs\ [4] @ [5] @ qs\ [7] \\ &= [2] @ [3] @ [4] @ [5] @ [7] \\ &= [2, 3, 4, 5, 7] \end{aligned}$$

Quicksort

Algorithm to sort a list

Idea:

$$qs\ (p \# l) = qs\ (\text{elements} \leq p) @ [p] @ qs\ (\text{elements} > p)$$

$$\begin{aligned} & qs\ [3, 2, 5, 4, 7] \\ &= qs\ [2] @ [3] @ qs\ [5, 4, 7] \\ &= [2] @ [3] @ qs\ [4] @ [5] @ qs\ [7] \\ &= [2] @ [3] @ [4] @ [5] @ [7] \\ &= [2, 3, 4, 5, 7] \end{aligned}$$

$$\begin{aligned} & qs\ [] = [] \\ & qs\ (p \# l) = qs\ (\text{filter } (\lambda x. x \leq p) l) @ [p] @ qs\ (\text{filter } (\lambda x. x > p) l) \end{aligned}$$

Demo.thy

Quicksort

Correct Sorting

What does it mean that quicksort is *correct*?

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What does it mean that quicksort is *correct*?

- ① The resulting list is sorted: *sorted (qs l)*

Correct Sorting

What does it mean that quicksort is *correct*?

- ① The resulting list is sorted: *sorted (qs l)*
- ② and contains the same elements: $\forall x. \text{count } (qs\ l)\ x = \text{count } l\ x$
 $\forall x. \dots$ means "for all x "

Demo.thy

Correctness of Sorting

Useful Properties

$$\textit{count} (l_1 @ l_2) x =$$

Useful Properties

$$\textit{count } (l_1 @ l_2) x = \textit{count } l_1 x + \textit{count } l_2 x$$

Useful Properties

$$\text{count } (l_1 @ l_2) \ x = \text{count } l_1 \ x + \text{count } l_2 \ x$$

$$\text{count } (\text{filter } P \ l) \ x =$$

Useful Properties

$$\text{count } (l_1 @ l_2) x = \text{count } l_1 x + \text{count } l_2 x$$

$$\text{count } (\text{filter } P l) x = \text{if } P x \text{ then count } l x \text{ else } 0$$

Useful Properties

$$\text{count } (l_1 @ l_2) \ x = \text{count } l_1 \ x + \text{count } l_2 \ x$$

$$\text{count } (\text{filter } P \ l) \ x = \text{if } P \ x \text{ then count } l \ x \text{ else } 0$$

$$\text{count } (\text{filter } (\lambda x. x \leq p) \ l) \ x + \text{count } (\text{filter } (\lambda x. x > p) \ l) \ x =$$

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partitioning preserves elements

Demo.thy

Useful Properties

Induction

To prove correctness of $qs\ l$ for all l :

base show that $qs\ []$ is correct

step show that $qs\ (p\#l)$ is correct,

assuming recursive calls are already correct (IH)

Induction

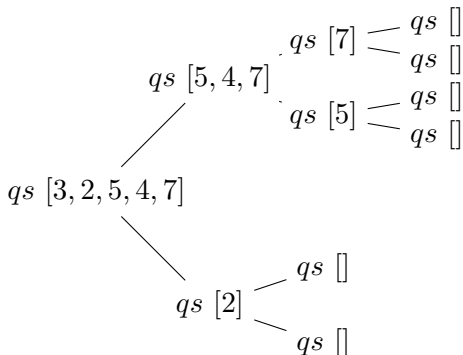
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Idea: Justify correctness, starting at leaves of call tree:



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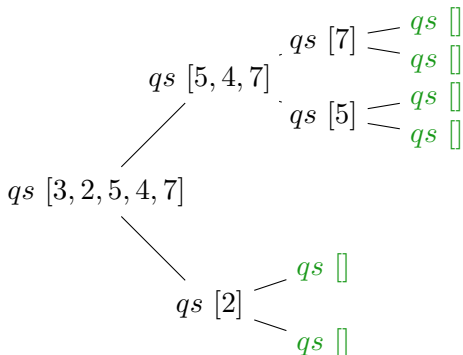
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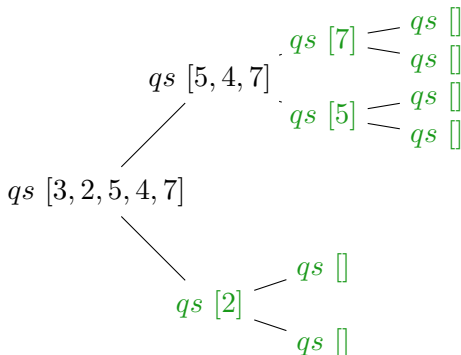
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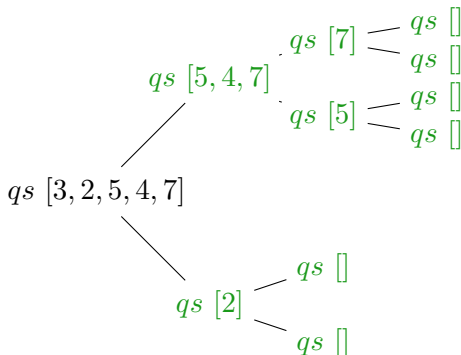
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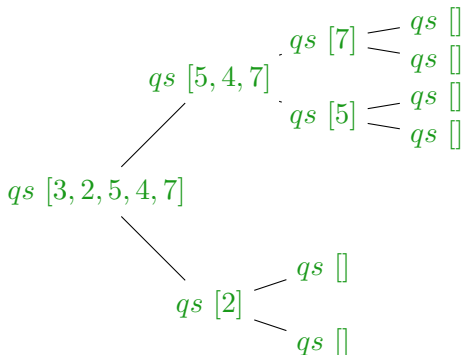
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Element Preservation

count (*qs l*) *x* = *count l x*

Base: *count* (*qs []*) *x* = *count [] x*

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Let $l_1 = \text{filter } (\lambda x. x \leq p) \ l$ and $l_2 = \text{filter } (\lambda x. x > p) \ l$

Element Preservation

$\text{count } (qs \ l) \ x = \text{count } l \ x$

Base: $\text{count } (qs \ []) \ x = \text{count } [] \ x$

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Let $l_1 = \text{filter } (\lambda x. x \leq p) \ l$ and $l_2 = \text{filter } (\lambda x. x > p) \ l$

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Show: $\text{count } (qs \ (p\#l)) \ x = \text{count } (p\#l) \ x$

$$\begin{aligned} & \text{count } (qs \ (p\#l)) \ x \\ &= \text{count } (qs \ l_1 \ @ \ [p] \ @ \ qs \ l_2) \ x \end{aligned}$$

Element Preservation

$count\ (qs\ l)\ x = count\ l\ x$

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$$\begin{aligned} & count\ (qs\ (p\#l))\ x \\ &= count\ (qs\ l_1\ @\ [p]\ @\ qs\ l_2)\ x \\ &= count\ [p]\ x + count\ (qs\ l_1)\ x + count\ (qs\ l_2)\ x \end{aligned}$$

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$$\text{count } (qs \ l) \ x = \text{count } l \ x$$

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Demo.thy

Quicksort preserves Elements

Quicksort Sorts

Set of elements in list $(x \in \text{set } l) = (0 < \text{count } l \ x)$

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What do we know about element x if $x \in \text{set } (\text{filter } P \ l)$?

$x \in \text{set } (\text{filter } P \ l) \implies P \ x$

Demo.thy

More useful Properties and Quicksort Sorts

Conclusions

Proved correct functional implementation of quicksort.

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Proof machine checked, using Isabelle/HOL.

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Thanks!