

# 10. Logic Programming With Prolog

## Overview

- Logic Programming
- Prolog

Note: Study Section 11.3 of the textbook, excluding 11.3.2.

## Logic Programming

- Logic programming is a form of declarative programming
- A program is a collection of *axioms*
  - Each axiom is a *Horn clause* of the form:  

$$H \text{ :- } B_1, B_2, \dots, B_n$$
where  $H$  is the *head term* and  $B_i$  are the *body terms*
  - Meaning  $H$  is true if all  $B_i$  are true
- A user of the program states a *goal* (a theorem) to be proven
  - The logic programming system attempts to find axioms using *inference steps* that imply the goal (theorem) is true

## Resolution

- To deduce a goal (theorem), the logic programming system searches axioms and combines sub-goals
- For example, given the axioms:  

$$C \text{ :- } A, B.$$

$$D \text{ :- } C.$$
- To deduce goal  $D$  given that  $A$  and  $B$  are true:
  - *Forward chaining* deduces that  $C$  is true:  

$$C \text{ :- } A, B$$
and then that  $D$  is true:  

$$D \text{ :- } C$$
  - *Backward chaining* finds that  $D$  can be proven if sub-goal  $C$  is true:  

$$D \text{ :- } C$$
the system then deduces that the sub-goal is  $C$  is true:  

$$C \text{ :- } A, B$$
Since the system could prove  $C$  it has proven  $D$

## Prolog

- Uses backward chaining
  - More efficient than forward chaining for larger collections of axioms
- Interactive (hybrid compiled/interpreted)
- Applications: expert systems, artificial intelligence, natural language understanding, logical puzzles and games
- Popular system: SWI-Prolog
  - Login: **program.cs.fsu.edu**
  - Type: **pl** to start SWI-Prolog
  - Type: **halt.** to halt Prolog (note that a period is used as a command terminator)

## Prolog Terms

- Terms are symbolic expressions that form the building blocks of Prolog
  - A Prolog program consists of *terms*
  - Data structures processed by a Prolog program are *terms*
- A *term* is either
  - a *variable*: a name beginning with an upper case letter
  - a *constant*: a number or string
  - an *atom*: a symbol or a name beginning with a lower case letter
  - a *structure* of the form:  
 $functor(arg_1, arg_2, ..., arg_n)$   
where *functor* is an atom and  $arg_i$  are terms
- Examples:
  - **X, Y, ABC**, and **Alice** are variables
  - **7, 3.14**, and **"hello"** are constants
  - **foo, bAR**, and **+** are atoms
  - **bin\_tree(foo, bin\_tree(bar, glarch))** and **+(3,4)** are structures

## Prolog Clauses

- A program consists of a database of Horn *clauses*
- Each *clause* consists of a *head predicate* and *body predicates*:  
 $H :- B_1, B_2, ..., B_n$ 
  - A clause is either a *rule*, e.g.  
**snowy(X) :- rainy(X), cold(X).**  
Meaning "If X is rainy and X is cold then this implies that X is snowy"
  - Or a clause is a *fact*, e.g.  
**rainy(rochester).**  
Meaning "Rochester is rainy."  
This fact is identical to the rule with **true** as the body predicate:  
**rainy(rochester) :- true.**
- A *predicate* is a *term* (must be an atom or a structure)
  - **rainy(rochester)**
  - **member(X,Y)**
  - **true**

## Queries and Goals

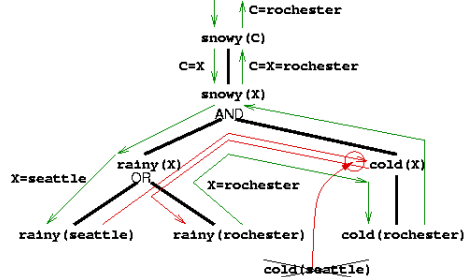
- *Queries* are used to "execute" *goals*
  - A *query* is interactively entered by a user after a program is loaded and stored in the database
  - A *query* has the form  
 $?- G_1, G_2, ..., G_n$   
where  $G_i$  are *goals*
- A *goal* is a predicate to be proven true by the programming system
  - Example program with two facts:  
**rainy(seattle).**  
**rainy(rochester).**
  - Query with one goal to find which city C is rainy (if any):  
**?- rainy(C).**
  - Response by the interpreter:  
**C = seattle**
  - Type a semicolon ; to get next solution:  
**C = rochester**
  - Type another semicolon ;:  
**no**  
(no more solutions)

## Example

- Program with three facts and one rule:  
**rainy(seattle).**  
**rainy(rochester).**  
**cold(rochester).**  
**snowy(X) :- rainy(X), cold(X).**
- Query and response:  
**?- snowy(rochester).**  
**yes**
- Query and response:  
**?- snowy(seattle).**  
**no**
- Query and response:  
**?- snowy(Paris).**  
**no**

## Example (cont'd)

- Program:  
**rainy(seattle).**  
**rainy(rochester).**  
**cold(rochester).**  
**snowy(X) :- rainy(X), cold(X).**
- ?- snowy(C).  
**C = rochester**  
because **rainy(rochester)** and **cold(rochester)** are sub-goals that are both true facts in the database
- snowy(X)** with **X=seattle** is a goal that fails, because **cold(X)** fails, triggering *backtracking*



## Backtracking

- For every successful match of a (sub-)goal with a head predicate of a clause, the system keeps this execution point in memory together with the current variable bindings to enable *backtracking*
- An unsuccessful match later forces *backtracking* in which alternative clauses are searched that match (sub-)goals
- Backtracking* unwinds variable bindings to enable establishing new bindings

## Example: Family Relationships

- Facts:  
**male(albert).**  
**male(edward).**  
**female(alice).**  
**female(victoria).**
- Rules:  
**parents(edward, victoria, albert).**  
**parents(alice, victoria, albert).**  
**sister(X,Y) :- female(X), parents(X,M,F), parents(Y,M,F).**
- Query:  
?- sister(alice, X1).
- The system applies backward chaining to find the answer:
  - sister(alice, X1)** matches 2nd rule: **X = alice, Y = X1**
  - New goals: **female(alice), parents(alice, M, F), parents(X1, M, F)**
  - female(alice)** matches 3rd fact
  - parents(alice, M, F)** matches 2nd rule: **M = victoria, F = albert**
  - parents(X1, victoria, albert)** matches 1st rule: **X1 = edward**

## Example: Murder Mystery

% the murderer had brown hair:  
**murderer(X) :- hair(X, brown).**

% mr\_holman had a ring:  
**attire(mr\_holman, ring).**

% mr\_pope had a watch:

**attire(mr\_pope, watch).**

% If sir\_raymond had tattered cuffs then mr\_woodley had the pincenez:

**attire(mr\_woodley, pincenez) :-**  
**attire(sir\_raymond, tattered\_cuffs).**

% and vice versa:

**attire(sir\_raymond, pincenez) :-**  
**attire(mr\_woodley, tattered\_cuffs).**

% A person has tattered cuffs if he is in room 16:

**attire(X, tattered\_cuffs) :- room(X, 16).**

% A person has black hair if he is in room 14, etc:

**hair(X, black) :- room(X, 14).**

**hair(X, grey) :- room(X, 12).**

**hair(X, brown) :- attire(X, pincenez).**

**hair(X, red) :- attire(X, tattered\_cuffs).**

% mr\_holman was in room 12, etc:

**room(mr\_holman, 12).**

**room(sir\_raymond, 10).**

**room(mr\_woodley, 16).**

**room(X, 14) :- attire(X, watch).**

## Example: Murder Mystery (cont'd)

- Question: who is the murderer?
- ?- **murderer(X)**.
- Trace (indentation showing nesting depth):  
**murderer(X)**  
  **hair(X, brown)**  
    **attire(X, pincenez)**  
      **X = mr\_woodley**  
      **attire(sir\_raymond, tattered\_cuffs)**  
        **room(sir\_raymond, 16)**  
          **FAIL (no facts or rules)**  
        **FAIL (no alternative rules)**  
      **REDO (found one alternative rule)**  
      **attire(X, pincenez)**  
        **X = sir\_raymond**  
        **attire(mr\_woodley, tattered\_cuffs)**  
          **room(mr\_woodley, 16)**  
            **SUCCESS**  
            **SUCCESS: X = sir\_raymond**  
          **SUCCESS: X = sir\_raymond**  
        **SUCCESS: X = sir\_raymond**  
      **SUCCESS: X = sir\_raymond**

## Unification and Variables

- In the previous examples we saw the use of variables, e.g. **C** and **X**
- A variable is *instantiated* to a term as a result of *unification*
- *Unification* takes place when goals are matched to head predicates of rules and facts
  - Goal in query: **rainy(C)**
  - Fact in database: **rainy(seattle)**
  - Unification is the result of the goal-fact match: **C = seattle**
- Unification is recursive:
  - An *uninstantiated* variable unifies with anything, even with other variables which makes them identical (aliases)
  - An atom unifies with an identical atom
  - A constant unifies with an identical constant
  - A structure unifies with another structure if the functor and number of arguments are the same and the corresponding arguments unify recursively
- Once a variable is instantiated to a non-variable term, it cannot be changed and cannot be instantiated with a term that has a different structure

## Unification Examples

- The built-in predicate **=(A,B)** succeeds if and only if **A** and **B** can be unified
- The goal **=(A,B)** may be written as **A = B**
  - ?- **a = a.**  
  **yes**
  - ?- **a = 5.**  
  **no**
  - ?- **5 = 5.0.**  
  **no**
  - ?- **a = X.**  
  **X = a**
  - ?- **foo(a,b) = foo(a,b).**  
  **yes**
  - ?- **foo(a,b) = foo(X,b).**  
  **X = a**
  - ?- **foo(X,b) = Y.**  
  **Y = foo(X,b)**
  - ?- **foo(Z,Z) = foo(a,b).**  
  **no**

## Lists

- A *list* is of the form:  
   $[elt_1, elt_2, \dots, elt_n]$   
  where  $elt_i$  are terms
- The special list form  
   $[elt_1, elt_2, \dots, elt_n \mid tail]$   
  denotes a list whose tail list is *tail*
  - ?- **[a,b,c] = [a|T].**  
      **T = [b,c]**
  - ?- **[a,b,c] = [a,b|T].**  
      **T = [c]**
  - ?- **[a,b,c] = [a,b,c|T].**  
      **T = []**

## List Membership

- List membership is tested with the **member** predicate, defined by  
**member(X, [X|T]).**  
**member(X, [H|T]) :- member(X, T).**
- ?- member(b, [a,b,c]).**
- Execution:
  - member(b, [a,b,c])** does not match predicate **member(X<sub>1</sub>, [X<sub>1</sub>|T<sub>1</sub>])**
  - member(b, [a,b,c])** matches predicate **member(X<sub>1</sub>, [H<sub>1</sub>|T<sub>1</sub>])**  
with **X<sub>1</sub> = b**, **H<sub>1</sub> = a**, and **T<sub>1</sub> = [b,c]**
  - Sub-goal to prove: **member(X<sub>1</sub>, T<sub>1</sub>)** with **X<sub>1</sub> = b** and **T<sub>1</sub> = [b,c]**
  - member(b, [b,c])** matches predicate **member(X<sub>2</sub>, [X<sub>2</sub>|T<sub>2</sub>])**  
with **X<sub>2</sub> = b** and **T<sub>2</sub> = [c]**
  - The sub-goal is proven, so **member(b, [a,b,c])** is proven (deduced)
- Note: variables are "local" to a clause (just like the formal arguments of a function)
  - Local variables such as **X<sub>1</sub>** and **X<sub>2</sub>** are used to indicate a match of a (sub)-goal and a head predicate of a clause

## Predicates are Relations

- Predicates are *not* functions with distinct inputs and outputs
- Predicates are more general and define *relationships* between objects (terms)
  - member(b, [a,b,c])** relates term **b** to the list that contains **b**
  - ?- member(X, [a,b,c]).**  
**X = a ;**    % type ';' to try to find more solutions  
**X = b ;**    % ... try to find more solutions  
**X = c ;**    % ... try to find more solutions  
**no**
  - ?- member(b, [a,Y,c]).**  
**Y = b**
  - ?- member(b, L).**  
**L = [b|\_G255]**  
therefore, **L** is a list with **b** as head and **\_G255** as tail, where **\_G255** is a new variable
- List appending predicate:
  - append([], A, A).**  
**append([H|T], A, [H|L]) :- append(T, A, L).**
  - ?- append([a,b,c], [d,e], X).**  
**X = [a,b,c,d,e]**  
**?- append(Y, [d,e], [a,b,c,d,e]).**  
**Y = [a,b,c]**  
**?- append([a,b,c], Z, [a,b,c,d,e]).**  
**Z = [d,e]**

## Example: Bubble Sort

- bubble(List, Sorted) :-**

```
append(InitList, [B,A|Tail], List),
A < B,
append(InitList, [A,B|Tail], NewList),
bubble(NewList, Sorted).
bubble(List, List).
?- bubble([2,3,1], L).
  append([], [2,3,1], [2,3,1]),
  3 < 2, fails: backtrack
  append([2], [3,1], [2,3,1]),
  1 < 3,
  append([2], [1,3], NewList1), this makes: NewList1=[2,1,3]

  bubble([2,1,3], L).
  append([], [2,1,3], [2,1,3]),
  1 < 2,
  append([], [1,2,3], NewList2), this makes:
NewList2=[1,2,3]
  bubble([1,2,3], L).
  append([], [1,2,3], [1,2,3]),
  2 < 1, fails: backtrack
  append([1], [2,3], [1,2,3]),
  3 < 2, fails: backtrack
  append([1,2], [3], [1,2,3]), does not unify: backtrack

  bubble([1,2,3], L). try second bubble-clause which makes
L=[1,2,3]
  bubble([2,1,3], [1,2,3]).
  bubble([2,3,1], [1,2,3]).
```

- Prolog offers a few built-in constructs to support a form of control-flow
  - \+ G** negates a (sub-)goal **G**
  - !** (cut) terminates backtracking *for a predicate and within the body of the clause of that predicate*
  - fail** always fails
- Examples
  - ?- \+ member(b, [a,b,c]).**  
**no**
  - ?- \+ member(b, []).**  
**yes**
  - Define:  
**if(Cond, Then, Else) :- Cond, !, Then.**  
**if(Cond, Then, Else) :- Else.**
  - ?- if(true, X=a, X=b).**  
**X = a ;**    % type ';' to try to find more solutions  
**no**
  - ?- if(fail, X=a, X=b).**  
**X = b ;**    % type ';' to try to find more solutions  
**no**
  - ?- if(true, a=b, X=b).**  
**no**
  - The cut makes sure that the **Cond** is not executed again upon backtracking and that the second **if**-clause is not executed when **Cond** is true when backtracking
  - Therefore, this example would not work without the cut when backtracking

## Imperative Control Flow

## Example: Tic-Tac-Toe

- Board layout:

1	2	3
4	5	6
7	8	9

- Facts:

```
ordered_line(1,2,3).
ordered_line(4,5,6).
ordered_line(7,8,9).
ordered_line(1,4,7).
ordered_line(2,5,8).
ordered_line(3,6,9).
ordered_line(1,5,9).
ordered_line(3,5,7).
```

Note: You can download the program from here  
(instructions are included in the source).

## Example: Tic-Tac-Toe (cont'd)

- Rules to find line of three (permuted) cells:

```
line(A,B,C) :- ordered_line(A,B,C).
line(A,B,C) :- ordered_line(A,C,B).
line(A,B,C) :- ordered_line(B,A,C).
line(A,B,C) :- ordered_line(B,C,A).
line(A,B,C) :- ordered_line(C,A,B).
line(A,B,C) :- ordered_line(C,B,A).
```

- How to make a good move to a cell:  
**move(A) :- good(A), empty(A).**
- Which cell is empty?  
**empty(A) :- \+ full(A).**
- Which cell is full?  
**full(A) :- x(A).**  
**full(A) :- o(A).**
- Which cell is best to move to? (check this in this order)  
**good(A) :- win(A).**      *% a cell where we win*  
**good(A) :- block\_win(A).**      *% a cell where we block the opponent from a win*  
**good(A) :- split(A).**      *% a cell where we can make a split to win*  
**good(A) :- block\_split(A).**      *% a cell where we block the opponent from a split*  
**good(A) :- build(A).**      *% choose a cell to get a line*  
**good(5).**      *% choose a cell in a good location*  
**good(1).**  
**good(3).**  
**good(7).**  
**good(9).**  
**good(2).**  
**good(4).**  
**good(6).**  
**good(8).**

## Example: Tic-Tac-Toe (cont'd)

- How to find a winning cell:  
**win(A) :- x(B), x(C), line(A,B,C).**
- Choose a cell to block the opponent from choosing a winning cell:  
**block\_win(A) :- o(B), o(C), line(A,B,C).**
- Choose a cell to split for a win later:  
**split(A) :- x(B), x(C), \+ (B = C), line(A,B,D), line(A,C,E), empty(D), empty(E).**

O		
	X	O
	X	X

- Choose a cell to block the opponent from making a split:  
**block\_split(A) :- o(B), o(C), \+ (B = C), line(A,B,D), line(A,C,E), empty(D), empty(E).**
- Choose a cell to get a line:  
**build(A) :- x(B), line(A,B,C), empty(C).**

## Example: Tic-Tac-Toe (cont'd)

- Board positions:

O		
X	O	
X		

- Are stored as facts in the database:  
**x(7).**  
**o(5).**  
**x(4).**  
**o(1).**
- Move query:  
**?- move(A).**  
**A = 9**

## Arithmetic

- Arithmetic is useful for many computations in Prolog
- The **is** predicate evaluates an arithmetic expression and instantiates a variable with the result
  - For example  
**X is 2\*sin(1)+1**  
instantiates **X** with the results of **2\*sin(1)+1**

## Arithmetic Examples

- A predicate to compute the length of a list:  
**length([], 0).**  
**length([H|T], N) :- length(T, K), N is K + 1.**  
where the first argument of **length** is a list and the second is the computed length
- Example query:  
**?- length([1,2,3], X).**  
**X = 3**
- A predicate to compute GCD:  
**gcd(A, A, A).**  
**gcd(A, B, G) :- A > B, N is A-B, gcd(N, B, G).**  
**gcd(A, B, G) :- A < B, N is B-A, gcd(A, N, G).**