

CS483 Review 02
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Algorithm Design: Chapter 1 ~ Chapter 8

1 Overview

This document is written to help you prepare the coming exam. Do not use this document as your cheat sheet for the the final exam. So far we have covered Chapters 1 ~ 8 of our textbook.)

1.1 Algorithm Design Strategies

- Theoretical analysis of algorithm: To analyze an algorithm, we should ask:
 - What is the input size?
 - What is the basic operation?
 - What is the efficiency class when the input size is infinitely large?
- Asymptotic notations for a function $f(n)$ (which normally represents the number of basic step of an algorithm with input size n)
 - $O(f(n))$ means $f(n)$ is an upper bound
 - $\Omega(f(n))$ means $f(n)$ is a lower bound
 - $\Theta(f(n))$ means $f(n)$ is a tight bound
- Compare order of growth: $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$
 - $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$ means $f(n) = O(g(n))$
 - $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \text{constant}$ means $f(n) = \Theta(g(n))$
 - $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$ means $f(n) = \Omega(g(n))$
- Greedy algorithm
 - Decision is made based on current (or local) state only
 - usually done iteratively (with partial solution)
 - used for solving optimization problems
 - Example: making change, minimum spanning tree, single-source shortest paths, Huffman coding
- Divide-n-Conquer (divide into sub-problems and solve all sub-problems)
 - merge sort, integer and matrix multiplications
 - Selection (Medians, k-the smallest element, binary search)
 - FFT
- Graph algorithms
 - Depth-/Breadth-first search and their applications
 - * what data structures are used?
 - Directed graphs (cycles, types of edges)
 - Directed acyclic graph (dag) and linearization (topological sort)

- (Strongly) connected components
 - * undirected graph
 - * directed graph
- Single-source shortest paths
 - * graphs without weight
 - * weight graphs with positive weights
 - what data structures are used?
 - * weight graphs with negative weights
 - * DAG
- Dynamic programming
 - solve and **store** results from sub-problems
 - Reuse the results from the sub-problems
 - Do not confuse dynamic programming with divide and conquer
 - “programming” means optimization. Therefore, dynamic programming is usually used to solve optimization problems (problems that ask you to find maximum value or minimum cost).
 - When you are asked to solve to problem using dynamic programming, the first thing you do is to think and define how such a problem can be solved using **sub-problems**. The second thing you should think is how to re-use the solutions of the sub-problems.
 - Example: Fibonacci number, transitive closure, all-pairs shortest-paths, knapsack, matrix sequence multiplication
- Linear programming
 - Representing problems using linear programs
 - Network flow
 - * Max flow problem
 - * Min cut problem
 - * Bipartite matching
 - Duality
 - * Game theory (min-max theorem)
 - Simplex
 - * Start with a (competent) solution
 - * Improve the quality of the solution in each iteration
 - * Until no better solution can be obtained
 - * used for solving optimization problems
- NP completeness
 - Definition of search problems
 - Classical search problems (some hard, some easy)
 - Complexity classes
 - * P (solvable in polynomial time), NP (checkable in polynomial time), NP-hard (all problems in NP can be reduced to any problem in NP hard), NP-complete (NP and NP-hard)
 - * Problems in P (MST), in NP (SAT), in NP hard (halting problem), in NP-complete (SAT), in NP but neither NP-hard nor P (graph isomorphism)
 - Reduction
 - * Used in solving problem, establishing lower bound, proving NP-hardness
 - * Problem $A \rightarrow$ Problem B (reduce A to B), the complexity of problem A is known, we want to show the complexity of B
 - * Given an instance I_A of problem A modify I_A as an input of problem S
 - Example: Vertex Cover, Independent Set, Clique, TSP, Knapsack, Rudrata cycle, SAT, ...

2 Problems and Strategies We Learned

The following table listed a set of problems that we have learned from Chapter 6 to Chapter 8.

| Category | Problem | Complexity | Strategy |
|---------------------------|-----------------------------|------------|----------|
| Graph | transitive closure | | |
| | all-pairs shortest-paths | | |
| | maximum flow | | |
| | bipartite matching | | |
| | Rudrata circuit/path | | |
| | traveling salesmen | | |
| | vertex cover | | |
| | 3-coloring | | |
| | independent set | | |
| Other optimization | knapsack | | |
| | edit distance | | |
| | longest increasing sequence | | |
| | matrix sequence mult | | |
| | bandwidth allocation | | |
| | task assignment | | |
| search | SAT | | |
| | 3SAT | | |
| | circulation with demand | | |
| | survey design problem | | |
| | subset sum | | |
| | Integer LP | | |
| | ZOE | | |
| other | Halting problem | | |
| | Factoring | | |

3 Sample Questions

1. (10 points) Answer and provide explanations to the following questions.

- a. (5 pts) Compare Dijkstra's algorithm and Prim's algorithm.
 - What are the problems that these two algorithms try to solve?
 - What kind of strategies are used by these two algorithms (e.g., divide&conquer, space-time, iterative, or ...)?
 - What functions are used to order unvisited vertiecs?
 - What kind of data structure can be used to maintain unvisited vertiecs?
 - What are their time complexities?

2. (10 points) Problem reduction.

a. (7 pts) Reduce the following 3 SAT problem to a 3-coloring problem.

$$(x \vee y \vee \bar{z}) \wedge (\bar{w} \vee z) \wedge (\bar{y} \vee w) \wedge (\bar{x} \vee \bar{z})$$

b. (3 pts) Solve the 3 SAT problem above using the solution of the 3-coloring problem. (**Hint:** Find a clique with 4 vertices.)