

Example 2: Greatest Common Divisor

Theoretical analysis of time efficiency

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Example 1: Sequential Search

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Asymptotic Notation

□ Best case:

1. When does the best case happen?
when $\gcd(a, b) = \min(a, b)$
2. What is $C_{best}(n)$? $C_{best}(n) = 1$

□ Average case:

1. Assumptions:

- Assume that a and b are two randomly chosen integers
- Assume that all integers have the same probability of being chosen
- **hint:** The probability that an integer d is a and b 's greatest common divisor is $P_{a,b}(d) = \frac{6}{\pi^2 d^2}$

2. When does the average case happen? when K and A satisfy our assumptions

3. What is $C_{avg}(n)$?

Let us denote $n = \min(a, b)$

$$C_{avg}(n) = 1 \cdot P_{a,b}(n) + 2 \cdot P_{a,b}(n-1) + \cdots + (n) \cdot P_{a,b}(1) =$$

$$\frac{6}{\pi^2} \left(\frac{1}{n^2} + \frac{2}{(n-1)^2} + \cdots + \frac{n}{1^2} \right)$$

(when $n = 10$,

$$\left(\frac{1}{100} + \frac{2}{81} + \frac{3}{64} + \frac{4}{49} + \frac{5}{36} + \frac{6}{25} + \frac{7}{16} + \frac{8}{9} + \frac{9}{4} + 10 \right) \cdot (6/\pi^2) = 8.58300468)$$

Question: what is the tightest upper bound of $C_{avg}(n)$? is it $O(n)$, $O(n^2)$ or $O(n \log n)$ and why?

Useful Property

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Asymptotic Notation

O -notation

Ω -notation

Θ -notation

▷ Useful Property

Comparing Orders of Growth

Orders of growth of some important functions

1. $f(n) \in O(f(n))$

Proof.

□

2. $f(n) \in O(g(n))$ if and only if $g(n) \in \Omega(f(n))$

Proof.

□

3. $f(n) \in O(g(n))$ and $g(n) \in O(h(n))$, then $f(n) \in O(h(n))$

Proof.

□

4. $f_1(n) \in O(g_1(n))$ and $f_2(n) \in O(g_2(n))$, then
 $f_1(n) + f_2(n) \in O(\max\{g_1(n), g_2(n)\})$

Proof.

□

Orders of growth of some important functions

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1. All logarithmic functions $\log_a n$ belong to the same class $\Theta(\log n)$ no matter what the logarithm's base $a > 1$ is

Proof.

□

2. All polynomials of the same degree k belong to the same class:
 $a_k n^k + a_{k-1} n^{k-1} + \dots + a_0 \in \Theta(n^k)$

Proof.

□

3. Exponential functions a^n have different orders of growth for different a 's, i.e.,
 $2^n \notin \Theta(3^n)$

Proof.

□

4. $\text{order } \log n < \text{order } n^{a>0} < \text{order } a^n < \text{order } n! < \text{order } n^n$