

DEPARTMENT OF MECHANICAL ENGINEERING



Milestone II for Ph.D. Program

Discrete Path Planning For Platonic Solids

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1 ABSTRACT

Background: Rolling geometry between two regular rigid surfaces of objects in continuous space can generate paths which is a well-known nonholonomic system. However, obtaining the path planning of a compact-closed surface on a general surface is challenging. From the rolling contact viewpoint, path-finding of convex polyhedrons has not received enough attention, and their paths on discrete surfaces have remained unexplored.

Aims: The aim of this project is to develop path planning algorithms for rolling polyhedron on a surface.

Approaches: The approach presented in this paper is based on using tree exploration technique to solve the closed-path planning disregarding orientation. The graph search starts from a root node then expands to different branches which indicates multiple levels of rolling. We also focus on close-path planning for platonic solids on discrete surfaces.

Conclusions: The results show multi-suboptimal closed paths based on rolling for the platonic solids to achieve the desired configuration on the known environments. The results also show successful in reaching the closed-path after combining with the point-to-point path planning technique in terms of changing polyhedrons' orientation. Matlab simulations are carried out to demonstrate the effectiveness of the proposed path-finding algorithm for the regular polyhedra.

Implications: its implementation can be extended to any convex polyhedrons.

2 INTRODUCTION

Path planning algorithm is one of the challenging problems in nonholonomic systems which, when resolved achieves the dexterous manipulation of objects in an unknown or part-known environment. This problem is mainly applied to the fields of robotics, artificial intelligence and autonomous vehicles. In robotics, the motivation of path planning is to find a possible path from an initial configuration, avoiding the obstacles and achieving the goal configuration [1]. Based on the task of robot performance, there are mainly two kinds of planning including feasibility and optimality. The former is to find a plan for only achieving the path while the latter is to find an optimal path. In the artificial intelligence fields, searching for actions to attain the desired goal state with receiving reward is employed, including decision-theoretic methods. Each specific path planning algorithm is usually implemented in a parameter space, such as configuration space or free space, which generates the feasible path connecting the two given points. Defining the state space is also one of the important steps for planning purpose. The configuration space or C-space which includes all possible configurations in a physical system is applied for solving path planning problems in n-dimensional. Examples of solving the path planning problems from Lavelle [2] and Kavraki [3] presented the feasible paths avoiding obstacles in the high dimensional configuration space.

Rolling contact between rigid bodies has been considered as nonholonomic systems in order to solve the problem of dexterous manipulation of industrial parts. The goal of rolling manipulation is to roll the part from an initial configuration to the goal configuration. It can be divided into three types of rolling contacts including point contact [4], [5], line contact [6] and surface contact [7]. A simple experiment of a rolling polyhedral part on a table, mentioned in [8], showed that object manipulation with polyhedral surfaces without sliding can be executed by nonholonomic constraints through rolling. Some cases of rolling polyhedral objects through graspless manipulation have been studied in the robotics field [9], [10]. Due to the lack of complete research of contact kinematics and rolling manipulation with discretized objects, planning for rolling polyhedral parts under reorientation with smooth and non-smooth systems still attracts attention from the research community.

Planning techniques are categorized into different aspects. The basic idea of discrete path planning in most cases is that state-space models will be used to demonstrate the distinct situation in which the task of a planning algorithm solves the sequence actions, transforming from an initial state to other states [11]. For example, Thomas [12] applied Delaunay triangulations to discretize the environment, and cubic spline representations are proposed to meet robot kinematic constraints. Considering other research on geometry, the paper [13] studied the continuous curvature on smooth curves which has been integrated within the probabilistic approaches in order to compute the piecewise smooth paths for a car-like vehicle as a four-dimensional system. Whereas dealing with nonholonomic constraints, a sampling-based road map technique has been proposed in [14], which determined trajectories and re-entry trajectories for hovercrafts and rigid spacecrafts. Based on decomposing space into cells [15], a potential field without local minima was assigned with polygonal partitions of planar environments to solve the Laplace's equa-

tion problems in each presence cell. Applications for these techniques in discrete space is limited by a grid.

Not much work has been done in path planning considering rolling contact constraint. Some types of moving polyhedral parts have been investigated on a plane such as sliding on a face, tumbling through the edges or pivoting [9] through the vertex. The planning motions of rolling polyhedral parts through the edges were clearly represented in [16]. This paper presented some results about changing an orientation of a polyhedron through its edge's contact on a fixed plane without slipping. Experimental works from the article demonstrated that the manipulation of rolling polyhedron on a plane where the set of configurations has different structures of the polyhedral parts can be reached by rolling through its edges. In the experimental validation, a unit cube will reach the next position by rolling along the edges on a square mesh considering the given tolerance which leads to reaching an orientation closer. The paper also proposed a concept of path planning algorithm with the tolerance to achieve the goal configuration, which was considered as an important condition to generate an accurate path. However, the practical application may not be successful on robot manipulation.

Marigo [17] proposed the path planning for polyhedron in the case of an octahedron with eight faces rolling and translating on a plane. For the octahedron rolling algorithm, a list of faces containing the vertices and edges stored parts of the polyhedron. The defect angles are also computed between two connected faces. The algorithm initialized a start configuration, a desired final configuration, and given a polyhedron with a set of geometrical parameters. The steps of planning include displacing and reorienting the polyhedral part until achieving the final configuration. Nevertheless, the algorithm may not satisfy with the accuracy for more general polyhedron.

Therefore, in this study, we propose a discrete path planning algorithm based on the tree exploration method for the five types of platonic solids including cube, tetrahedron, octahedron, icosahedron, and dodecahedron. The path-finding algorithm focuses on rolling a solid on the associate grid from an original position to its initial position with different orientation. The study is organized as following. Section 3 covers the general properties in geometrical aspect of the five types platonic solids. Section 4 describes path planning algorithm based on tree exploration technique. Finally, section 5 shows the results of the proposed path planning algorithm under considering their different geometrical properties, then comes the conclusion of the paper.

3 MODEL DESCRIPTION

Platonic solids: The platonic solids are also called regular polyhedra have the convex polyhedra properties. There are exactly five regular polyhedra namely cube, tetrahedron, octahedron, dodecahedron and icosahedron (Figure 1). Some of the equivalent statements are used to describe the platonic solids, including all the vertices lie on a sphere, all the dihedral angle are equal, and all solid angles are equivalent. The tetrahedron is folded by 4-sided pyramid, the octahedron has the double-pyramid with 8 faces and 20-sided pyramid for the icosahedron. The cube is constructed by 6 square faces while the dodecahedron is composed of 12-sided of regular pentagons.

Geometrical parameters: The number of faces, edges, and vertices of each type of platonic solids is described in the Table 1. In mathematics, the Euler's formula shows the relationship between total number of vertices (V), edges (E), and faces (F) by the Eq. 1

$$V - E + F = 2 \quad (1)$$

The icosahedron has the largest number of faces with 20 while the tetrahedron has only 4 faces. If each face of platonic solids has i sides and k edges of the polyhedron meet at each vertex. The conditions are i, k greater than 3 ($i, k \geq 3$) because every face have at least three edges and at least three edges of faces meet at each vertex. The faces are equilateral triangle if $i = 3$ and changing the values of k will yield the other three types of polyhedron, including the tetrahedron ($k = 3$), the octahedron ($k = 4$), and the icosahedron ($k = 5$). When $i = 4$, the faces are square and the cube has $k = 3$ edges which meet at each vertex. The last case with $i = 5$ of regular pentagons and only $k = 3$ will generate the dodecahedron.

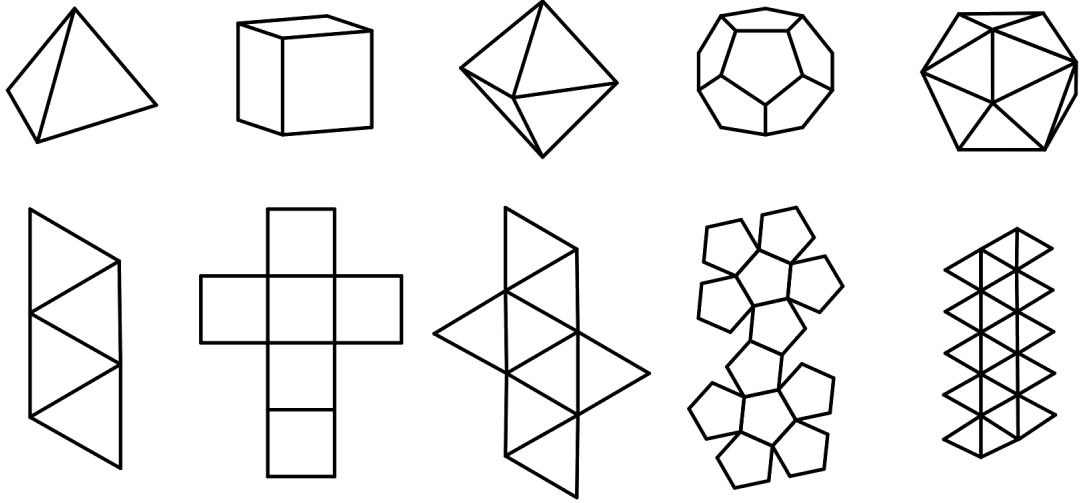


Figure 1: The platonic solids. From left to right with models and unfolding models: the tetrahedron, the cube, the octahedron, the dodecahedron, and the icosahedron

Table 1: Properties of polyhedron

	Faces	Edges	Vertices	Edges on each face	Edges meeting at each vertex
Tetrahedron	4	6	4	3	3
Cube	6	12	8	4	3
Octahedron	8	12	6	3	4
Dodecahedron	12	30	20	5	3
Icosahedron	20	30	12	3	5

The path planning for platonic solids focuses on the rolling of the models through edge contact. Each edge is shared by two faces and each face may have e edges. Assume that Δ is the quantity faces which contact at an edge or $E = \frac{1}{2}\Delta$. However, each vertex will be shared by f faces. Then, $V = \frac{\Delta}{f}$. After substituting these two quantities into the Eq. 1, the result is:

$$\begin{aligned}
 V - E + F &= 2 \\
 \frac{\Delta}{f} - \frac{1}{2}\Delta + F &= 2 \\
 \rightarrow F &= \frac{4f}{2e - fe + 2f}
 \end{aligned} \tag{2}$$

From the Eq. 2, the condition is that $2e - fe + 2f$ should be positive because F and $4f$ are both positive. As above requirements with $(e \geq 3)$ and $(f \geq 3)$, there are only five pairs of integers (e, f) satisfy this condition. So, the results of these pairs are shown in the last two columns in the Figure 1.

The important condition in the path planning through rolling contact is rotation angle. In the context of regular convex polyhedra, a rotation angle is supplementary with a dihedral angle which is the angle between two faces inside the polyhedra. The Table 2 shows the radii of each solid with the inradius (r_i), the midradius (ρ), the circumradius (R) and the dihedral angles (β).

Table 2: Geometrical parameters of platonic solids

	r_d	ρ	R	dihedral angles (β)
Tetrahedron	$\frac{1}{12}\sqrt{6}$	$\frac{1}{4}\sqrt{2}$	$\frac{1}{4}\sqrt{6}$	$\cos^{-1}(\frac{1}{3})$
Cube	$\frac{1}{2}$	$\frac{1}{2}\sqrt{2}$	$\frac{1}{2}\sqrt{3}$	$\frac{1}{2}\pi$
Octahedron	$\frac{1}{6}\sqrt{6}$	$\frac{1}{2}$	$\frac{1}{2}\sqrt{2}$	$\cos^{-1}(-\frac{1}{3})$
Dodecahedron	$\frac{1}{20}\sqrt{250 + 110\sqrt{5}}$	$\frac{1}{4}(3 + \sqrt{5})$	$\frac{1}{4}(\sqrt{15} + \sqrt{3})$	$\cos^{-1}(-\frac{1}{5}\sqrt{5})$
Icosahedron	$\frac{1}{12}(3\sqrt{3} + \sqrt{15})$	$\frac{1}{4}(1 + \sqrt{5})$	$\frac{1}{4}\sqrt{10 + 2\sqrt{5}}$	$\cos^{-1}(-\frac{1}{3}\sqrt{5})$

4 ALGORITHM\METHODOLOGY

4.1 Path Planning Based Rolling Contact

Rolling on discretized surfaces: The surface contacts between platonic solids and the plane can be categorized into three types as shown in Figure 2 including square polygon for a cube, triangle polygon for the tetrahedron, the octahedron, and the dodecahedron, pentagon polygon for the dodecahedron. The bottom surface of a cube occupies each square on the grid when the path planning process is executed. This property is applied for triangular grid with different rotation angle. In physics, the rotation angle of the cube is $\pi/2(\text{rad})$ while the rotation angles of tetrahedron, octahedron, icosahedron and dodecahedron are $\pi - \arctan(2\sqrt{2})$, $\pi - 2\arctan\sqrt{2}$, $\arccos(-\sqrt{5}/3)$, and $\pi - \arccos(-\sqrt{5}/5)$ in radian respectively.

The square grid has $\pi/2$ at all corners while the triangular grid has $\pi/3$ between two arbitrary edges at a vertex. In the case of dodecahedron rolling contact, the Figure 2d shows the two types of connections between pentagons where the first case has a gap (Figure 2c) and the other has overlap pentagon connection. A regular pentagon has five interior angles of 108° which generate a gap between three pentagons surrounding because of $3 * 108^\circ = 324^\circ$, which is different 360° of the full circle. Another case of four overlap pentagons with $4 * 108^\circ = 432^\circ$ is greater than the circle of 360° . The path planning through rolling of the dodecahedron solid can be categorized into two these cases. It would be found the possible paths in the first case of dodecahedron without overlap rolling while the second case with overlap rolling cannot guarantee the paths.

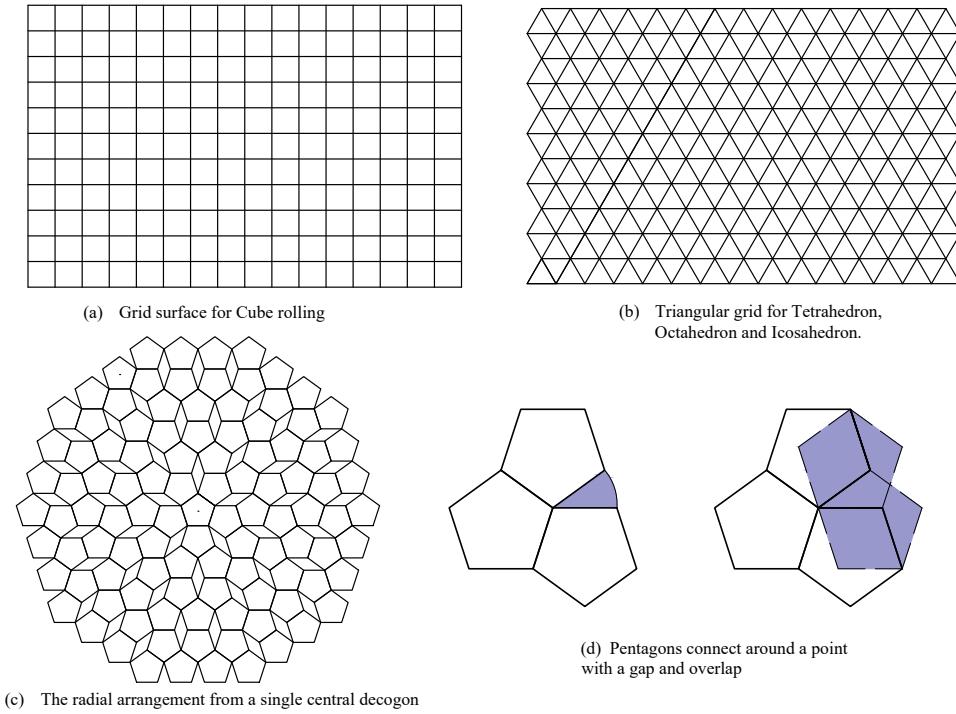


Figure 2: Grid of platonic solids

Rodrigues' roatation: It is assumed that the motion of rolling the platonic solids is a pure rolling without slipping or spinning at the line contact. Rolling the platonic solids means rolling all its vertices in 3D environment indicated by changing the coordinates. There were several algorithms to transform vertices stored in a matrix in 3D space. In this case, the Rodrigues's rotation method in [18] is used to represent the rotation matrix in the path planning algorithm.

Algorithms: Due to the different surface contacts, there are three types of direction for the rolling of platonic solids. As shown in the Figure 3, the cube has four directions with the square surface contact while the tetrahedron, the octahedron and the icosahedron have three rolling directions with the triangular surface contact. The dodecahedron with pentagon surface contact has five rolling directions. In the case of rolling cube, the surface contact is surrounded by four edges which means there are four possible directions through the edges. In this work, the proposed path planning algorithm deals with rolling from initial configuration within the position and orientation to the goal within the same position but different orientation. While rolling on the smooth plane, the platonic solid models will contact to the plane through their edges.

The Algorithm 1 shows that path planning for cube rolling based on tree graph search has some important steps. The first step is to initial the coordinates and the orientations of the initial cube and the target cube which is stored as the initial path. The same as tree expansion, cube will roll in four different directions including the right, left, up and down is the next step. From these new positions and orientations, the cubes will continue expand with only three directions to avoid return the previous positions. An example for this step is that from the initial coordinate the cube achieves a new position after doing rolling for right direction, the new three positions of the cube by rolling through right, up, and down direction. After implementing the expansion steps through rolling, the function of checking whether updated models reach the goal is called through the loop. By that means, the loop will stop when reaching the goal, whereas the loop will continue to execute and store new models to the initial path. While the searching algorithm is executing, the data structure is used to store the positions and orientations from the start to the current. This process runs in time $O(|E|^3)$ (where $|E|$ is the number of updated cubes) which causes the longer the running time of the searching technique.

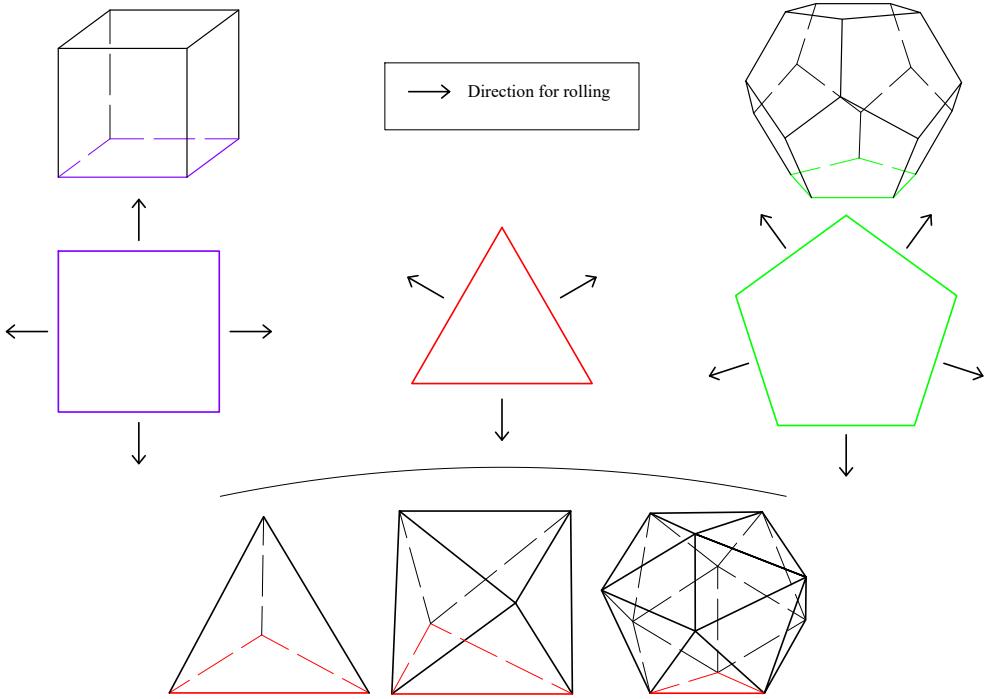


Figure 3: Rolling direction for each type of platonic solids

4.2 Tree Exploration Algorithm

The node tree exploration for searching algorithm described in Algorithm 1 is similar to non-recursive depth-first-search algorithm [19]. The graph search in the Figure 4 shows the expansion from the *root* with node *S* to multi-level from 1^{st} level ... n^{th} level for the case study of a cube solid. Each node indicates the position of the cube's center and the orientation of the cube. The node *S* means Start-Point while *R*, *L*, *U*, *D* are labelled for four different directions including right, left, up and down respectively. For each iteration, a tree with a node including 3D coordinate and orientation is stored in each level. At the same time, the algorithm of checking the goal configuration will be called to check whether the current executing level achieves the target.

In other cases of tetrahedron, octahedron and icosahedron with the triangular grid (Figure 2b), there are three directions at the first rolling and only two directions for the rest of path-finding process. Only the case of dodecahedron has the different approach from the algorithm. The path planning algorithm depends on the environment including gaps or overlaps between two pentagon connections as can be seen in the Figure 2b.

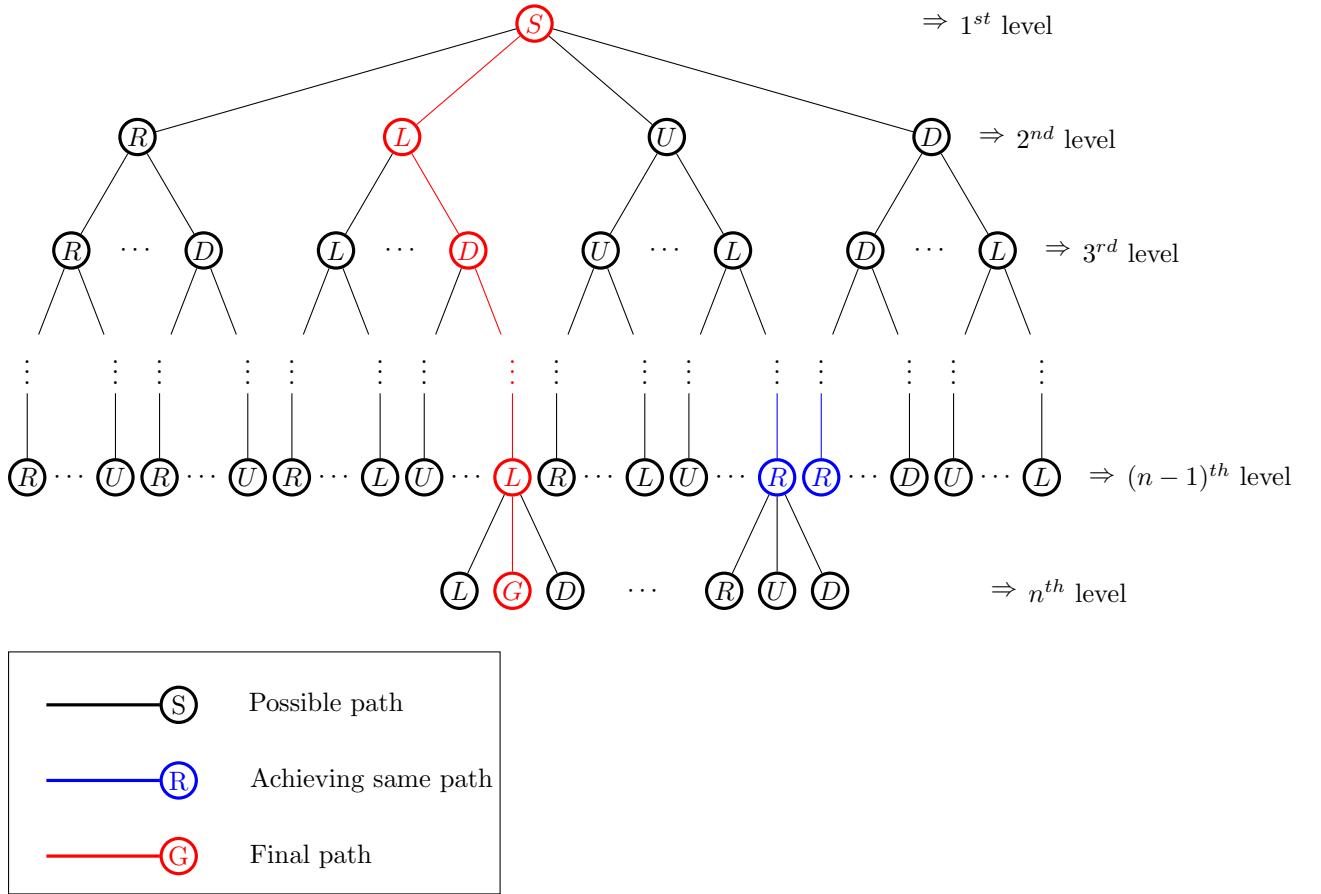


Figure 4: Tree Exploration of Cube Rolling

Starting from the root *S*, path planning based rolling of the cube model at the first level of expansion will generate to four different direction *R*, *L*, *U*, *D*. In the next level, the cube can only roll with three directions without rolling back to the previous position. An example of the second level is that node *R* will roll to right, up, and down directions indicated by node *R*, *U*, *D* respectively.

To eliminate the processing time in the proposed algorithms, whenever any updated points achieved the same position and orientation, these nodes will merge at that *level*. An example from Figure 4 shows two updated nodes *R* (blue node) at $(n-1)^{th}$ level have achieved the same position. The next path is generated from this merged nodes. From the tree exploration algorithm, the result can show only one path or various paths which depends on the initial and goal configuration. The first path is the shortest path because the executing time is shortest based on the condition of achieving goal configuration.

Algorithm 1 Path planning based rolling contact for Cube model.

```
1: procedure CUBE PATH PLANNING( $S_p, G_p$ )  $\triangleright$  Find the shortest path from start to goal position with
   different orientation
2:    $flag \leftarrow false$ 
3:    $Path[S_p] \leftarrow S_p$ 
4:    $newPoints \leftarrow \text{ROLLING4DIRECTIONS}(S_p)$   $\triangleright$  Generate first four updated points
5:   while  $newPoints \neq G_p$  do
6:      $updatedPoints \leftarrow \text{TREEEXPLORATION}(newPoints)$   $\triangleright$  Update new three right rolling models
7:      $n \leftarrow \text{size}(updatedPoints)$ 
8:     for  $i \leftarrow 0, n$  do
9:       for  $j \leftarrow 1, n$  do
10:        if  $updatedPoints[i] = updatedPoints[j]$  then
11:           $\text{remove}(updatedPoints[i])$ 
12:        end if
13:      end for
14:    end for
15:     $flag \leftarrow \text{CHECKINGTARGETPOINT}(updatedPoints)$   $\triangleright$  Compare updated points with goal point
16:    if  $flag = true$  then
17:      return  $Path[S_p, G_p]$   $\triangleright$  Store new point to  $Path$ 
18:    end if
19:     $newPoints = updatedPoints$ 
20:  end while
21:  return "no path found"
22: end procedure
23: procedure ROLLING4DIRECTIONS( $S_p$ )  $\triangleright$  Generate new points in different direction of rolling
24:    $(newRightPoint, newLeftPoint, newUpPoint, newDownPoint) \leftarrow \text{ROLLINGCONTACT}(S_p)$ 
25:   return  $newPoints \leftarrow (newRightPoint, newLeftPoint, newUpPoint, newDownPoint)$ 
26: end procedure
27: procedure TREEEXPLORATION( $newPoints$ )
28:   if  $dir = right$  then
29:      $updatedPoints \leftarrow (newRightPoint, newUpPoint, newDownPoint)$ 
30:   else if  $dir = left$  then
31:      $updatedPoints \leftarrow (newLeftPoint, newUpPoint, newDownPoint)$ 
32:   else if  $dir = up$  then
33:      $updatedPoints \leftarrow (newRightPoint, newLeftPoint, newUpPoint)$ 
34:   else
35:      $updatedPoints \leftarrow (newRightPoint, newLeftPoint, newDownPoint)$ 
36:   end if
37:   return  $updatedPoints$ 
38: end procedure
39: procedure CHECKINGTARGETPOINTS( $updatedPoints, G_p$ )
40:   if  $updatedPoints = G_p$  then  $\triangleright$  Consider both position and orientation
41:      $flag \leftarrow true$ 
42:   end if
43:   return  $flag$ 
44: end procedure
```

5 EVALUATION

The proposed algorithm for platonic solids path planning by rolling through edge contact was implemented in MATLAB environment. In general of path planning, there are three case studies including same location and different orientation between initial configuration and goal configuration, long distance between two configurations, and bi-direction path finding. To validate the proposed algorithm, this study only considers the first case study of path planning that both initial and goal configuration have the same positions and different orientations.

Assume that the platonic solids' edges has the same length with a ($a = 1$) and one of the faces of the platonic solids contact to OXY . The environment for each type of platonic solids is discretized from the smooth surface. For example, cube solid will roll on the square grid while tetrahedron, octahedron, and icosahedron solids roll on the triangle gird. Only dodecahedron solid rolls on the pentagon grid with the two specific cases including the gaps and overlaps between pentagons. The path planning based on rolling for the five cases of platonic solids is represented as the following.

5.1 Cube solid

Properties: The cube has a length a which is the same as the length of side of each grid square. Path planning for rolling a cube is not complex as other cases of platonic solids. The rolling angle β has the value of $\frac{\pi}{2}$ as shown in the Figure 5 ($\beta = \angle I K I_1 = \frac{\pi}{2}$). The center of cube will move along the curve II_2I_1 following the Ox axis. The proposed algorithms will focus on rolling along Ox or Oy axis. However, all the vertices are stored in a matrix will change their coordinates in 3D space. An example from the Figure 5, eight vertices of the cube $ABCDMNOP$ and the center I are stored in a matrix $M(9, 3)$ with 9 rows for nine 3D points, and 3 columns for 3D coordinate Ox, Oy, Oz of each point. After finishing the rotation angle β based on the rotation axis MN , the coordinates of the cube's vertices are updated within a new cube $MA'D'PN'B'C'O'$. This is the basic step of rolling a cube through a line contact.

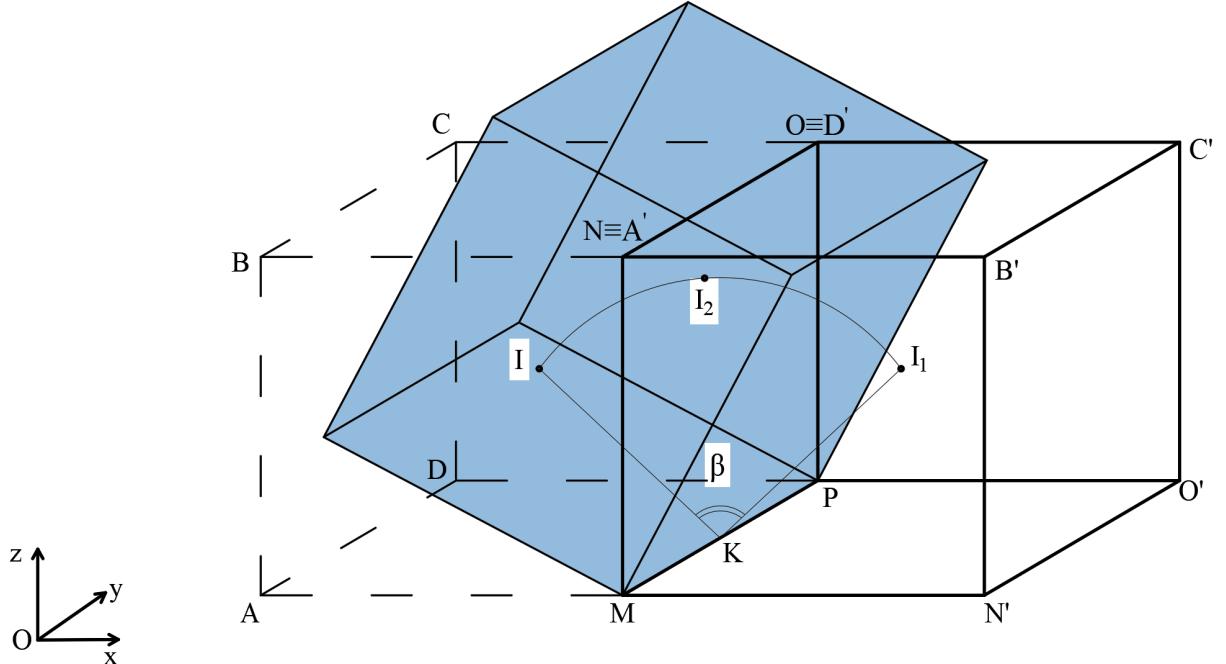


Figure 5: Rotation angle of a cube on a plane

Path planning: The only way to move from initial position to goal position of the cube solid is by rolling from square to square on the grid without moving diagonal. The Figure 6 shows the first three layers of cube path finding on the grid. The algorithm implements in $O(|E|^3)$ running time from the second layer called the expansion three branches from a tree. All 3D coordinates of the cube are stored in a matrix which can affect to the computer's storage capacity. When the updated cubes achieved one of the same previous configurations, the execution time can be reduced by releasing this configuration or stopping the expansion of this tree branch. The * position in the grid (Figure 6) is occupied by the two updated cubes with different orientations in the Layer 2.

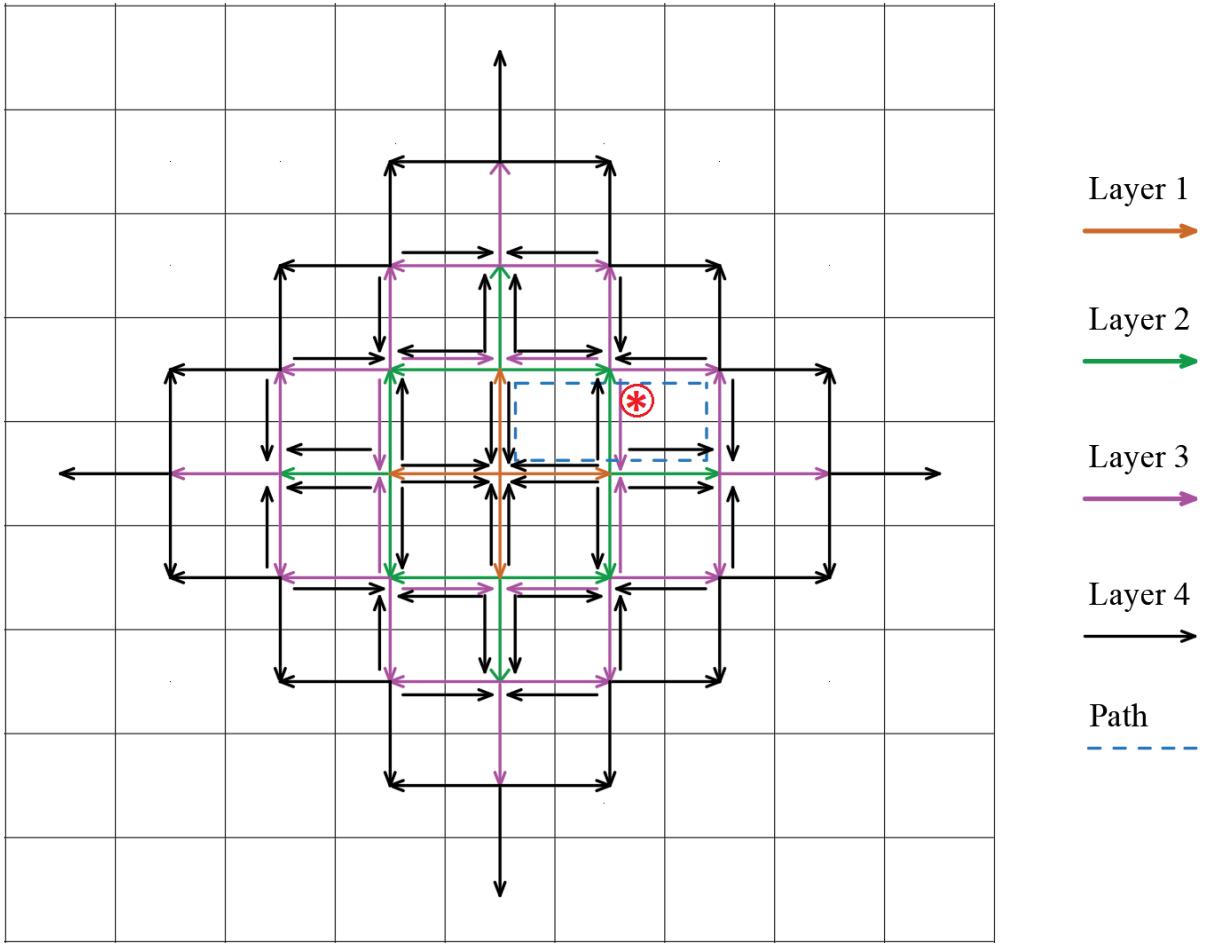
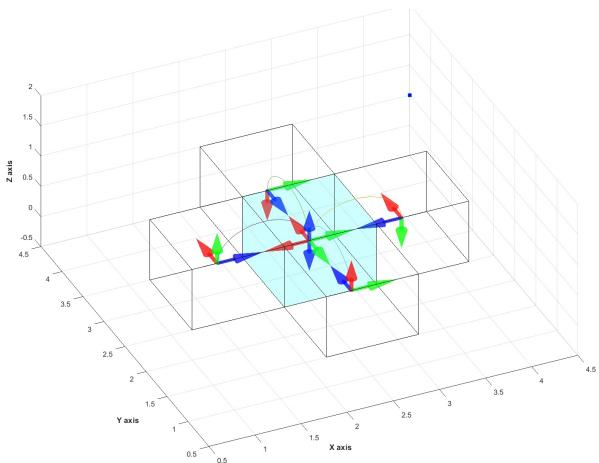
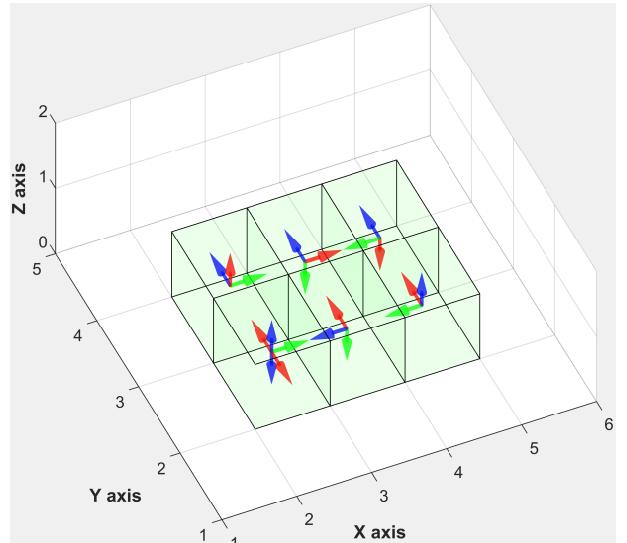


Figure 6: The first four layers of cube rolling

To be more visualized, the Figure 7(a) indicates the first step of rolling of the cube with its coordinates in the red, green and blue arrows. The first shortest path of cube rolling are shown in the Figure 7(b).



(a) The initial configuration is the same position but different orientation with goal configuration



(b) Shortest path of cube rolling

Figure 7: Initial configuration and shortest path of cube rolling

Extension case: The proposed algorithm in this study can be applied for the case study of long distance between the initial and goal configurations. Figure 8 shows an example with two different paths. Shortest path-finding algorithm is added to the original algorithm to find the shortest path from start point to the goal point. After finishing this step, the cube updated to a new orientation with different orientation at the goal configuration. Then, the original algorithm will be implemented from the updated cube to the goal configuration. Assume that the initial configuration is at S_1 and the goal configuration is at S_2 . In the first step, the red line segment S_1S_2 shows the shortest distance from start position to goal position. The cube will roll through this line segment and achieve the updated orientation at the goal position called shortest path for the case of long distance.

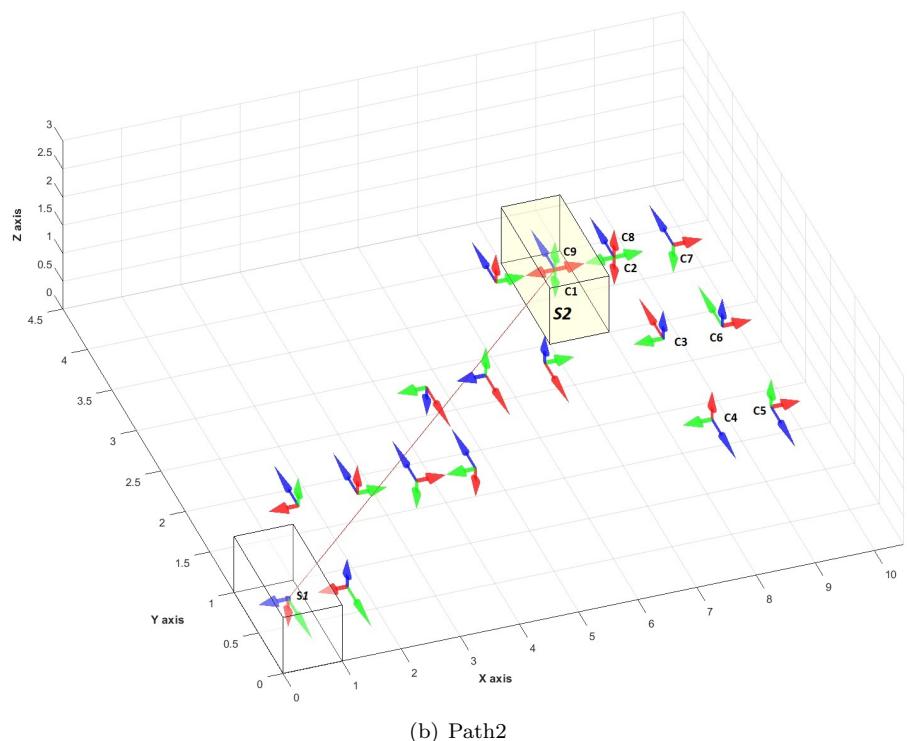
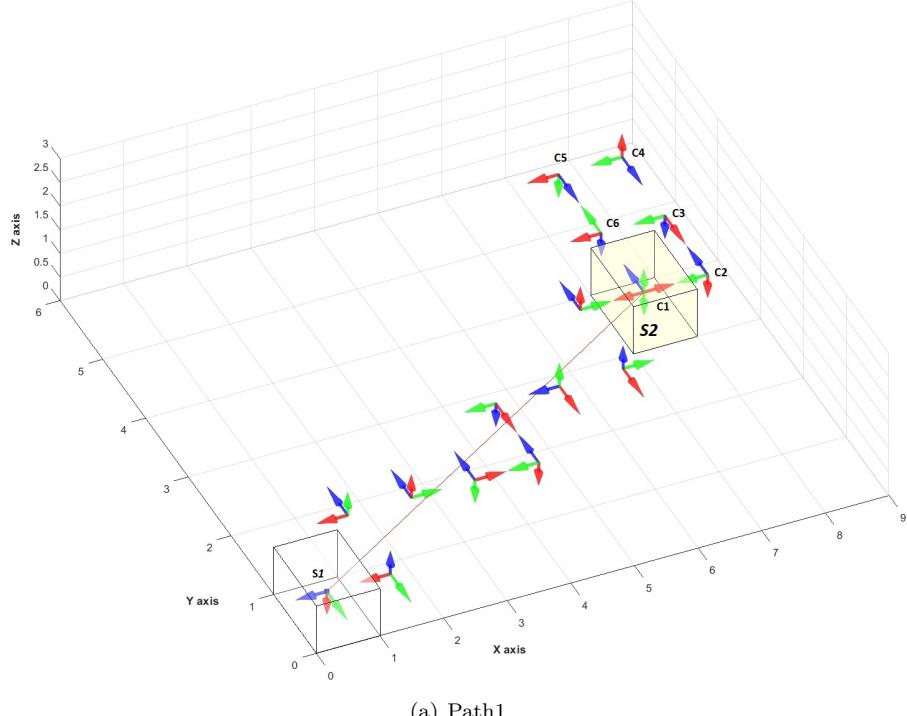


Figure 8: The case study of long distance between the initial and goal configurations

5.2 Tetrahedron solid

Properties As can be seen from the Figure 9, the Tetrahedron has constructed by four faces of the equilateral triangles. Then the height of triangle ABC is AM and $AM = DM = a\sqrt{3}/2$. Because of $r = OH = a\sqrt{6}/12$ (the radius of insphere) and $R = OA = a\sqrt{6}/4$ (the radius of circumsphere), the height of tetrahedron is $AH = OA + OH = a\sqrt{6}/3$. The rotation angle is β determined by supplementary angles $\alpha = \arctan(2\sqrt{2})$ or $\beta = \pi - \alpha = \pi - \arctan(2\sqrt{2})$.



Figure 9: Tetrahedron geometrical properties

Path Planning: The case study of tetrahedron in this study is the same of the rolling cube path finding which only considering the path planning through rolling from initial configuration to origin coordinate with different orientation. The Figure 10 illustrates two of four cases of the tetrahedron path-finding within rolling (red and cyan arrows are pointing down to plane respectively). Tetrahedron has symmetry properties with indistinguishable for any two faces, edges and vertices. To be more specific, dihedral triangles have same three angles within 60° . In one cycle of rolling a tetrahedron around any vertices, the tetrahedron always achieve the initial configuration due to six times of rolling ($6*60^\circ = 360^\circ$, a full circle).



Figure 10: The two of four cases

5.3 Octahedron solid

Properties: Dennis also went his own way and divided the sides of the triangles into equal-angles (as measured from the center of the geodesic), instead of equal-length pieces. This technique is slightly more effective at evenly distributing the triangles across the surface of the sphere. For example, compare an octahedron subdivided with frequency 20, using the linear technique (as outlined by the quiz) versus the angular technique Dennis used in this picture.



Figure 11: Octahedron geometrical properties

From Figure 11 we have:

$$\begin{aligned} AK &= 2OK = 2\sqrt{(MK^2 - MO^2)} = a\sqrt{2} \\ \Rightarrow OK &= OM = ON = a\frac{\sqrt{2}}{2} \end{aligned}$$

To find the rotation angle, the supplementary angle should be calculated first.

$$\alpha = \angle OPK = \arctan \frac{OK}{OP} = \arctan \sqrt{2}$$

Then, the rotation angle has the result as $\beta = \pi - 2\alpha = \pi - 2 \arctan \sqrt{2}$,

Due to $\angle POK = 90^\circ$, we have

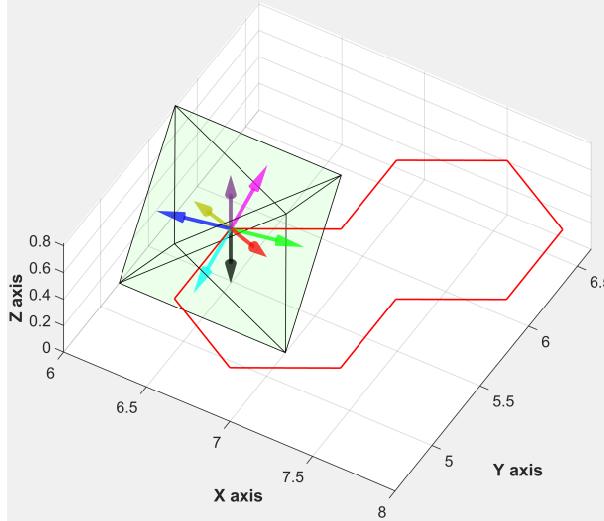
$$\begin{aligned} \frac{1}{OE^2} &= \frac{1}{OP^2} + \frac{1}{OK^2} \\ &= \frac{1}{(\frac{\sqrt{2}}{2})^2} + \frac{1}{(\frac{1}{2})^2} \\ \Rightarrow OE &= \frac{\sqrt{6}}{6} \end{aligned}$$

Applying the theory of the equilateral triangle

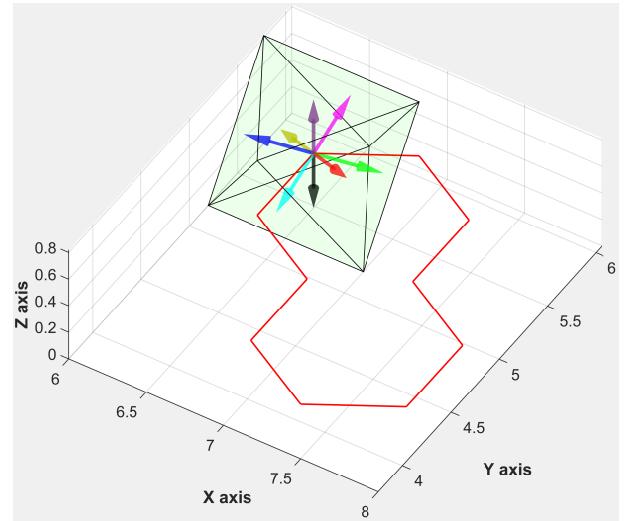
$$\begin{aligned} EO &= OF = a\frac{\sqrt{6}}{6} \\ KE &= FA = a\frac{\sqrt{3}}{3} \\ EP &= FQ = a\frac{\sqrt{3}}{6} \end{aligned}$$

Path planning: Based on the classical path planning which is the movement from point to other points, the octahedron path planning within rolling has three directions to move. In the Figure 11, the bottom layer $\triangle ABC$ which contacts to OXY can roll with the directions FQ, FO_1, FR . The distance between two shortest positions is $2FR = 2a\frac{\sqrt{3}}{6} = a\frac{\sqrt{3}}{3}$.

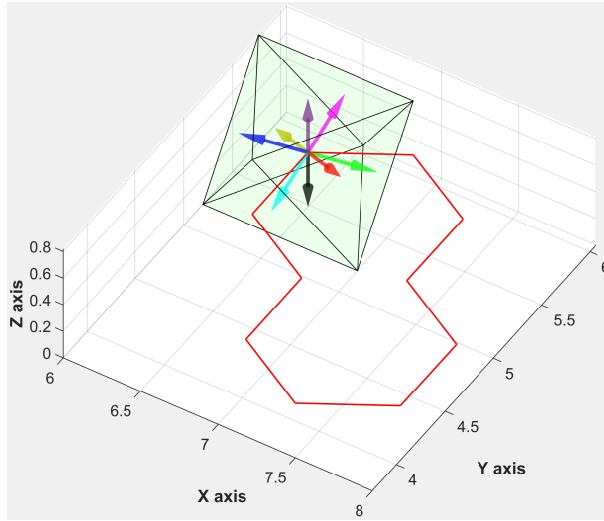
As can be seen from the Figure 12(a), the shortest path within red line includes ten line segments for rolling the octahedron. The total length of the path equals $10a\frac{\sqrt{3}}{3}$ (edge length of the octahedron is a).



(a) The first shortest path of Octahedron path rolling



(b) The second shortest path of Octahedron path rolling



(c) The second shortest path of Octahedron path rolling

Figure 12: Three shortest paths of octahedron based rolling

5.4 Icosahedron solid

Properties: The convex regular icosahedron in the Figure 13 is one of the five regular Platonic solids has 12 vertices, 20 triangular faces, and 30 edges. Assume that the icosahedron has the edge length with a . The crossing surface of the solid at vertex A perpendicular with OD will generate a pentagon $ABCEF$.



Figure 13: Icosahedron geometrical properties

At the right side of the Figure 13, the pentagon $ABCEF$ has circumcircle radius R , inscribed circle r_i , and the height d_{CS} ($R + r_i$). The golden ratio Φ (irrational number) has the value of $\frac{1+\sqrt{5}}{2}$ can be found by:

$$\Phi = \frac{OX}{XB} = \frac{OX}{\frac{1}{2}a}$$

Same as the case of cube's path planning algorithm, the initial step of rolling icosahedron is to find the rotation angle. The rotation angle is a supplementary of the $\angle QRS$. The line OR is perpendicular to QS and $OS = OQ = OM = r_i = \frac{\Phi^2}{2\sqrt{3}}a$. BI is a diagonal of pentagon $BCLIJ$ and $\angle QRS = 2\angle ORS$.

Otherwise, $OR = OX = \frac{1}{2}a\Phi$, and

$$\begin{aligned} \angle QRS &= 2\angle ORS \\ \sin \angle ORS &= \frac{OS}{OR} = \frac{\frac{\Phi^2}{2\sqrt{3}}a}{\frac{1}{2}a\Phi} = \frac{\Phi}{\sqrt{3}} \\ \Rightarrow \angle QRS &= 2 \arcsin \frac{\Phi}{\sqrt{3}} \end{aligned}$$

To be more detail in the Figure 14, the rotation angle of icosahedron can be determined from the result of $\angle QRS$, as the following.

$$\begin{aligned} \beta &= \pi - \angle QRS \\ &= \pi - 2 \arcsin \frac{\Phi}{\sqrt{3}} \\ &= \pi - \arccos -\frac{\sqrt{5}}{3} \end{aligned}$$

To calculate the rotation axis for all the cases of the triangular surface contact, there are three axis such as IJ , JK , and IK with the 3D coordinates as following.

$$JK = [a \ 0 \ 0]$$

$$IJ = [-\frac{1}{2}a \ (\frac{\sqrt{3}}{6} + \frac{\sqrt{3}}{3})a \ 0]$$

$$IK = [-\frac{1}{2}a \ -(\frac{\sqrt{3}}{6} + \frac{\sqrt{3}}{3})a \ 0]$$

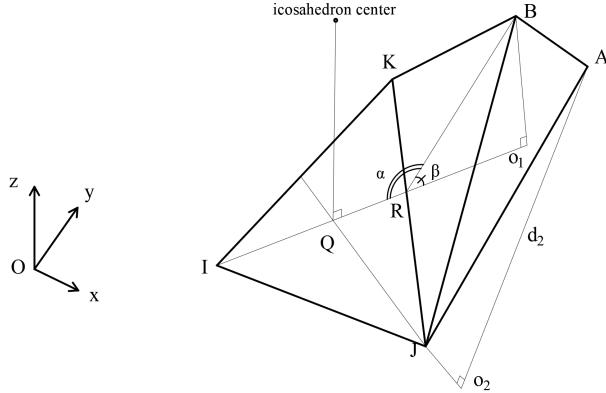


Figure 14: Rotation angle and rotation axis

Path planning Path planning of the regular icosahedron through rolling in known environment has initial coordinate at $[6.5, 5.5, 0]$ and the surface contact with red arrow points to bottom. The goal configuration with same as the initial coordinate has the contact surface with black arrow as shown in the Figure 15(a). This figure shows the shortest path with the red line including 14 line segments. It means that the icosahedron rolled 14 times from the initial configuration to the goal configuration. The path can be seen more precisely from the top view in the Figure 15(b). The total length of this icosahedron path equals $14a\frac{\sqrt{3}}{3}$ (edge length of the octahedron is a).



Figure 15: Path of icosahedron rolling though edges

5.5 Dodecahedron solid

Properties: An dodecahedron is constructed from twelve self-intersecting faces with the same pentagon shape. These faces meet at each vertex and there are total of twenty vertices in a dodecahedron as shown in Figure 16. From the Eq. 1, the great dodecahedron satisfy the Euler's formula ($V = 20, F = 12, E = 30$). It will be assumed that the coordinates $Oxyz$ lie on $ABCDE$ surface within Oy through A and Oz perpendicular to $ABCDE$. The 30 edges have the same length as a . It should be determined all the vertices' coordinates in the three dimensional system. The Figure 16 indicates the lengths of each vertex from l_1 to l_4 and the angles α_1 to α_4 which correspond to the five sides of a pentagon.

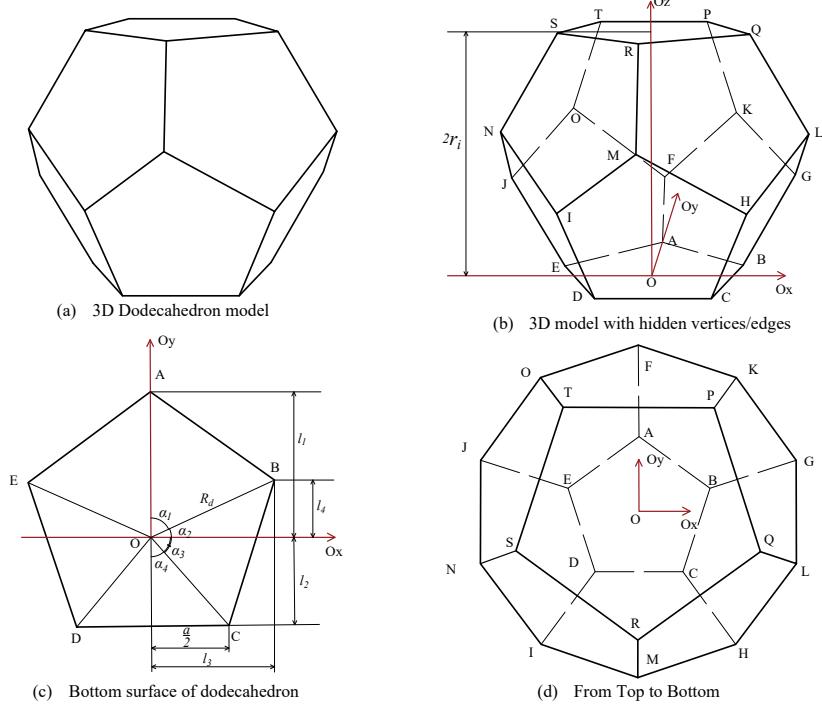


Figure 16: Dodecahedron's vertices.

The path planning will implement on a surface but it will be considered in 3D spaces. Then, each vertex will be determined on 3D coordinates such as the vertex A has coordinate with $[A_x A_y A_z]$. Based on the properties of pentagon, the angle $\alpha_1 = \frac{2\pi}{5}$ and $\alpha_4 = \frac{\pi}{5}$, because the angle between Ox and Oy is $\frac{\pi}{2}$, the sum of α_1 and α_2 is $\alpha_1 + \alpha_2 = \frac{\pi}{2}$. Then the other two angles α_2 and α_3 can be founded as below.

$$\begin{aligned}\alpha_2 &= \frac{\pi}{2} - \alpha_1 = \frac{\pi}{2} - \frac{2\pi}{5} = \frac{\pi}{10} \\ \alpha_3 &= \alpha_1 - \alpha_2 = \frac{2\pi}{5} - \frac{\pi}{10} = \frac{3\pi}{10}\end{aligned}$$

From the Figure 16(c), these labelled dimensions can be calculated as $l_1 = R_d = \frac{a}{2 \sin \alpha_4}$ with R_d is the circumradius of dodecahedron, $l_2 = l_1 \cos \alpha_4$, $l_3 = l_1 \cos \alpha_2$, and $l_4 = l_1 \sin \alpha_2$. The golden ratio Φ with the value of $\frac{1+\sqrt{5}}{2}$ which is the length of the diagonal of a square with 1 unit length of sides is used to calculate the circumscribed sphere radius of dodecahedron. Assume that the dodecahedron has the length a , the radius of an inscribed sphere is $r_i = \frac{a}{20} \sqrt{10(25 + 11\sqrt{5})}$ and the circumscribed sphere radius is $r = a \frac{\sqrt{3}}{2} \frac{1+\sqrt{5}}{2}$.

In the path-finding through rolling, the proposed algorithm focuses on the transformation and translation of 20 vertices of a dodecahedron. Although rolling dodecahedron's faces on 2D surface, the bottom surface which integrates to the Oxy will contact to the surface in the three dimensional space. This condition

expresses the Oz dimension of all the (A, B, C, D, E) vertices which equal to 0 or $A_z = B_z = C_z = D_z = E_z = 0$. Then,

$$\begin{aligned} P_z &= Q_z = R_z = S_z = T_z = 2.r_i \\ &= \frac{a}{10}\sqrt{10(25 + 11\sqrt{5})} \end{aligned}$$

It can be seen that the distance $|AF|$ is a and the distance $|BF|$ is $2l_3$. Using the distance properties and squaring the results give:

$$\begin{aligned} AF^2 &= a^2 = (A_x - F_x)^2 + (A_y - F_y)^2 + (A_z - F_z)^2 \\ BF^2 &= (2l_3)^2 = (B_x - F_x)^2 + (B_y - F_y)^2 + (B_z - F_z)^2 \end{aligned}$$

Figure 16(d) shows that $A_x = F_x = 0$, $A_y = l_1$, $B_y = l_4$, $B_x = l_3$, $A_z = B_z$. Define $l_5 = A_z - F_z$, the relations of these equations are:

$$\begin{aligned} a^2 &= (F_y - l_1)^2 + l_5^2 \\ (2l_3)^2 &= (F_y - l_4)^2 + l_5^2 + l_3^2 \end{aligned}$$

Solving F_y and l_5 gives:

$$\begin{aligned} F_y &= \frac{a^2 - (2l_3)^2 - (l_1^2 - l_3^2 - l_4^2)}{2(l_4 - l_1)} \\ l_5 &= \frac{1}{\sqrt{2}}\sqrt{a^2 + (2l_3)^2 - (F_y - l_1)^2 - (F_y - l_4)^2 - l_3^2} \end{aligned}$$

From these above equations, all the vertices can be determined and stored in a 3D matrix. Path planning based on rolling contact is the rotation of this matrix through edges' contact.

Path Planning: As mentioned in the Section 4, the dodecahedron can roll on the two grid types. The simple case has a gap at the connection between three regular pentagons as shown in Figure 17. In this paper, we do not apply the proposed algorithm to find the path of dodecahedron with rolling on the case of overlaps between four pentagons. This overlap grid is complicated environment which can not guarantee to generate paths with proposed algorithm. In the case of grid with gaps (Figure 17), attaching other five pentagons generates five 36° gaps from a regular pentagon. The six pentagons has a shape of large pentagon and the cycle of attaching more five larger pentagons will shape a grid of pentagons.

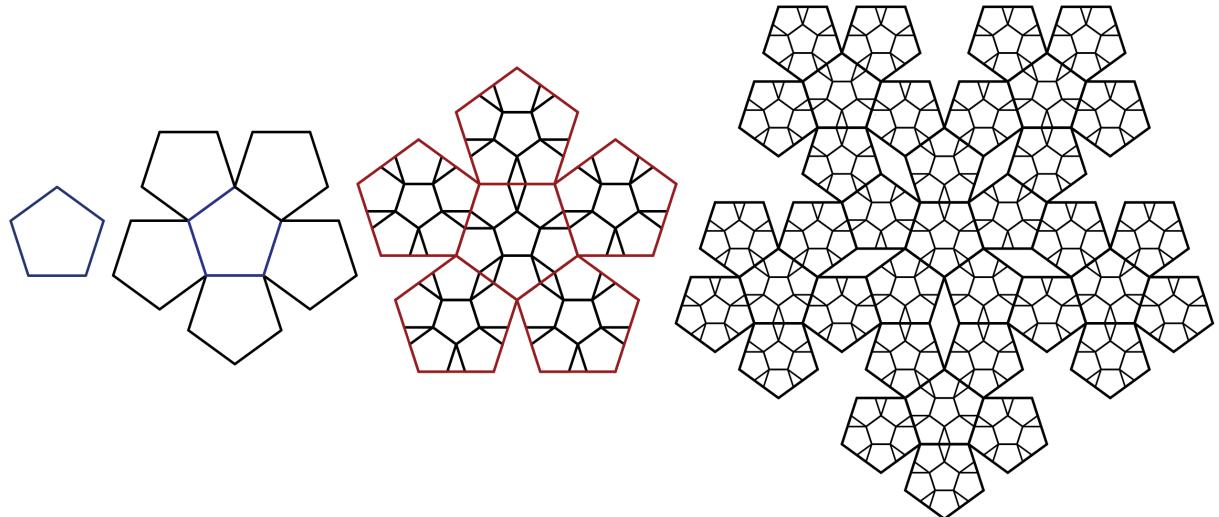


Figure 17: Initial dodecahedron grid with gaps

The rotation angle of dodecahedron on the above grid equals $(\pi - \arccos(\frac{-\sqrt{5}}{5}))$. One of the shortest paths is shown in the Figure 18 with [3.0 2.25 0.0] for both the initial and goal coordinates. One time

rolling of dodecahedron generates a new dodecahedron which has a distance $d = 2l_2$ from the previous coordinate. From the Figure 16b, the distance d_d can be determined as following.

$$\begin{aligned} l_2 &= l_1 \cos \alpha_4 \\ &= \frac{a}{2 \sin \alpha_4} \cos \alpha_4 \\ &= \frac{1}{2} a \cot \alpha_4 \\ \rightarrow d_d &= a \cot \alpha_4 \end{aligned}$$

Then, the total length of the shortest path equals $15a \cot \alpha_4$ as shown in the Figure 18 with the red line.

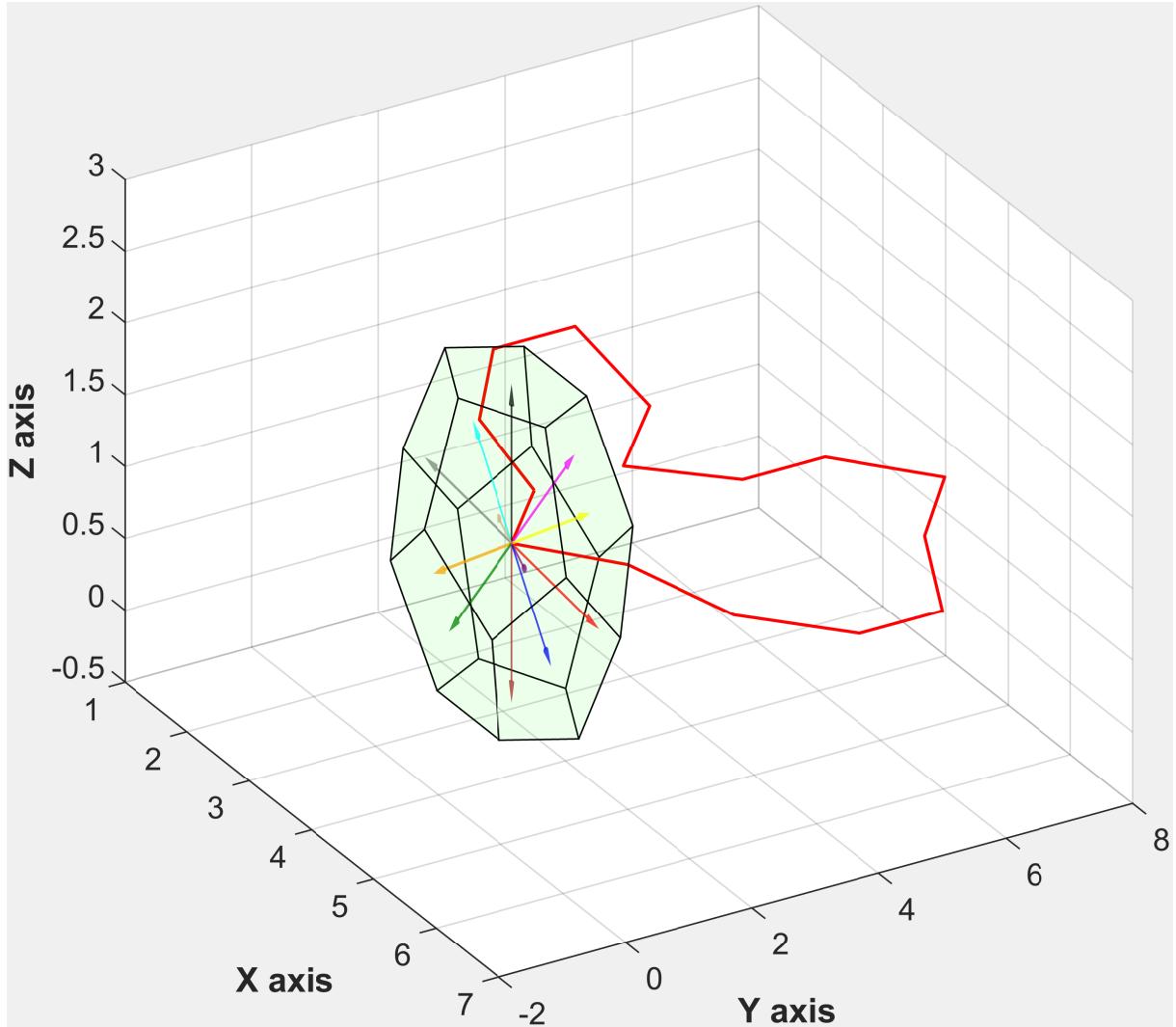


Figure 18: The rolling path of dodecahedron at [3.0 2.25 0.0] coordinate

Path planning through rolling contact for the case of dodecahedron is complex. According to the simulation results, implementing the proposed algorithm for other four types of platonic solids (cube, tetrahedron, octahedron, and icosahedron) is more effective than for the dodecahedron case. Executing the algorithm is no guarantee of path-finding results with the case of overlap grid.

6 CONCLUSIONS

In this paper, we have proposed a method of path-finding for platonic solids - regular convex polyhedra - that is able to produce paths and a shortest path under rolling constraint without sliding. In particular, using tree exploration algorithm can find the path for cube, tetrahedron, octahedron and icosahedron solids. In the first case of dodecahedron with a gap between two pentagons in the initial environment, the algorithm is applied to find shortest paths while the other case with overlaps is not guaranteed to find a path.

The results of this study concern path planning in discrete environment through rolling for only platonic solids with free-obstacles. In future work some more case studies will be included such as rolling platonic solids and the general convex polyhedra with obstacles in the environment. Finally, extensions of the proposed algorithm for rolling convex polyhedra on 3D convex surface and optimization of the algorithm in storage capacity and executing time will be considered.

7 TIME LINE

Year	2018	2019	2020	2021						
Activity	Sep	Oct	Dec	Feb	May	Aug	Nov	Feb	May	Aug
Candidacy proposal										
Literature review										
Mathematical model										
Simulated approach										
Result & Validation										
Thesis preparation										

Refer to the Data Management Plan in Appendix A.

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A APPENDIX