<https://www.lintcode.com/problem/911/>

Given an array *nums* and a target value *k*, find the maximum length of a subarray that sums to *k*. If there isn't one, return 0 instead.

**Note:** The sum of the entire *nums* array is guaranteed to fit within the 32-bit signed integer range.

**Example 1:**

Input: nums = [1, -1, 5, -2, 3], k = 3

Output: 4

Explanation: The subarray [1, -1, 5, -2] sums to 3 and is the longest.

**Example 2:**

Input: nums = [-2, -1, 2, 1], k = 1

Output: 2

Explanation: The subarray [-1, 2] sums to 1 and is the longest.

**Follow Up:**

Can you do it in O(*n*) time?

**Attempt 1: 2023-02-02**

**Solution 1:  Hash Table (360 min)**

**Style 1: int[] preSum = new int[nums.length + 1], the disadvantage is we also have to specially handle the nums.length == 1 && nums[0] == k case**

public class Solution {

/\*\*

\* @param nums: an array

\* @param k: a target value

\* @return: the maximum length of a subarray that sums to k

\*/

public int maxSubArrayLen(int[] nums, int k) {

// Test out by nums={-1}, k=-1

if(nums.length == 1 && nums[0] == k) {

return 1;

}

// {key, value} = {pre-sum, index}

Map<Integer, Integer> map = new HashMap<Integer, Integer>();

map.put(0, 0);

int[] preSum = new int[nums.length + 1];

int result = 0;

for(int i = 1; i <= nums.length; i++) {

preSum[i] = preSum[i - 1] + nums[i - 1];

if(map.containsKey(preSum[i] - k)) {

result = Math.max(result, i - map.get(preSum[i] - k));

}

// Only store preSum[i] when it happened first time, since we need to find

// maximum length between, and it supposed to get between first happening

// index to current index, all other indexes will only have less length

if(!map.containsKey(preSum[i])) {

map.put(preSum[i], i);

}

}

return result;

}

}

==========================================================================================

OR we can use map.put(k, 0) rather than map.put(0, 0) for initial position

public class Solution {

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\* @param nums: an array

\* @param k: a target value

\* @return: the maximum length of a subarray that sums to k

\*/

public int maxSubArrayLen(int[] nums, int k) {

// {key, value} = {pre-sum, index}

Map<Integer, Integer> map = new HashMap<Integer, Integer>();

// Differ than map.put(0, 0)

map.put(k, 0);

int[] preSum = new int[nums.length + 1];

int result = 0;

for(int i = 1; i <= nums.length; i++) {

preSum[i] = preSum[i - 1] + nums[i - 1];

if(map.containsKey(preSum[i])) {

result = Math.max(result, i - map.get(preSum[i]));

}

// Only store preSum[i] + k when it happened first time, since we need to find

// maximum length between, and it supposed to get between first happening

// index to current index, all other indexes will only have less length

if(!map.containsKey(preSum[i] + k)) {

map.put(preSum[i] + k, i);

}

}

return result;

}

}

Time Complexity: O(n)

Space Complexity: O(n)

**The problem on i < nums.length OR i <= nums.length in for loop**

e.g nums={1, -1, 5, -2, 3}, k=3;

for(int i = 1; i < nums.length; i++) {

preSum[i] = preSum[i - 1] + nums[i - 1];

...

}

preSum={0, 1, 0, 5, 3, 6}

nums[0]= stored on preSum[1], (nums[0] + nums[1]) stored on preSum[2], (nums[0] + nums[1] + nums[2]) stored on preSum[3], (nums[0] + nums[1] + nums[2] + nums[3]) stored on preSum[4], but missing (nums[0] + nums[1]... nums[4]) which suppose to be stored on preSum[5], only set "i <= nums.length" will make it happen, the root cause is we use "nums[i - 1]" which require "i == nums.length" to cover last element in nums

--------------------------------------------------------------------

if(!map.containsKey(preSum[i])) {

map.put(preSum[i], i);

}

map will store below:

i=0 -> {0, 0} -> for preSum initial set as 0, since we didn't store any num at index=0, only start from index=1

i=1 -> {preSum[1]=0+nums[0]=1, 1}

i=2 -> {preSum[2]=preSum[1]+nums[1]=1-1=0, 2} -> ignore since 0 alraedy have {0, 0} as key=0 stored in map, keep index=0

i=3 -> {preSum[3]=preSum[2]+nums[2]=0+5=5, 3}

i=4 -> {preSum[4]=preSum[3]+nums[3]=5-2=3, 4}

i=5 -> {preSum[5]=preSum[4]+nums[4]=3+3=6, 5}

--------------------------------------------------------------------

if(map.containsKey(preSum[i] - k)) {

result = Math.max(result, i - map.get(preSum[i] - k));

}

length calculate below:

i=0 -> map not contains (0-3=-3)

i=1 -> map not contains (1-3=-2)

i=2 -> map not contains (0-3=-3)

i=3 -> map not contains (5-3=2)

i=4 -> map contains (3-3=0) -> i-map.get(preSum[i]-k)=4-map.get(0)=4-0=4 -> in preSum={0, 1, 0, 5, 3, 6} the initial preSum=0 get from index=0, the current preSum=3 get from index=4, length between calculate directly as 4-0=4, NOT including compensation +1 (e.g 4-0+1=5), because we introduce preSum[0]=0 and store nums beginning only from preSum[1], now comes back to raw input nums={1, -1, 5, -2, 3}, if we based on raw input, since it shift all elements index one step to left than preSum, maximum length to get between index=0 to index=3 needs to be calculated as 3-0+1=4, the compensation +1 is required then

i=5 -> map contains (6-3=3) -> i-map.get(preSum[i]-k)=5-map.get(3)=5-4=1

**Style 2: Set map.put(0, -1) with setting int[] preSum array on accumulate preSum[i + 1] = preSum[i] + nums[i] to avoid problem when i = 0, preSum[0] is nothing, nums[i] only accumulate on preSum[i + 1], e.g nums[0] only stored on preSum[1], nums[1] only stored on preSum[2]**

public class Solution {

/\*\*

\* @param nums: an array

\* @param k: a target value

\* @return: the maximum length of a subarray that sums to k

\*/

public int maxSubArrayLen(int[] nums, int k) {

// {key, value} = {pre-sum, index}

Map<Integer, Integer> map = new HashMap<Integer, Integer>();

// Subarray sum from index i to index j (e.g array[i:j]) = sum(array[0:j]) - sum(array[0:i-1])

// So this problem is another two-sum like problem, we can calculate prefix-sum of

// the given array and use a hashmap to look up answer.

map.put(0, -1);

int[] preSum = new int[nums.length + 1];

int result = 0;

for(int i = 0; i < nums.length; i++) {

preSum[i + 1] = preSum[i] + nums[i];

if(map.containsKey(preSum[i + 1] - k)) {

result = Math.max(result, i - map.get(preSum[i + 1] - k));

}

// Only store preSum[i] when it happened first time, since we need to find

// maximum length between, and it supposed to get between first happening

// index to current index, all other indexes will only have less length

if(!map.containsKey(preSum[i + 1])) {

map.put(preSum[i + 1], i);

}

}

return result;

}

}

Time Complexity: O(n)

Space Complexity: O(n)

**Style 3: Set map.put(0, -1) with change int[] preSum array to int preSum only to avoid problem on accumulate preSum[i] = preSum[i - 1] + nums[i] issue when i = 0, we can NOT do preSum[i] = preSum[i] + nums[i], its wrong because of NO accumulate at all**

public class Solution {

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\* @param nums: an array

\* @param k: a target value

\* @return: the maximum length of a subarray that sums to k

\*/

public int maxSubArrayLen(int[] nums, int k) {

// {key, value} = {pre-sum, index}

Map<Integer, Integer> map = new HashMap<Integer, Integer>();

// Subarray sum from index i to index j (e.g array[i:j]) = sum(array[0:j]) - sum(array[0:i-1])

// So this problem is another two-sum like problem, we can calculate prefix-sum of

// the given array and use a hashmap to look up answer.

map.put(0, -1);

//int[] preSum = new int[nums.length + 1];

int preSum = 0;

int result = 0;

for(int i = 0; i < nums.length; i++) {

preSum += nums[i];

if(map.containsKey(preSum - k)) {

result = Math.max(result, i - map.get(preSum - k));

}

// Only store preSum[i] when it happened first time, since we need to find

// maximum length between, and it supposed to get between first happening

// index to current index, all other indexes will only have less length

if(!map.containsKey(preSum)) {

map.put(preSum, i);

}

}

return result;

}

}

Time Complexity: O(n)

Space Complexity: O(1)

**Refer to**

<https://www.lintcode.com/problem/911/solution/59403>

解题思路

保存从头开始以i结尾的sum，以及i的位置，有相同的sum时，只保留最先出现的那个，key就是sum，value是对应的数组下标.

有了sum之后，令target=sum-k，如果target之前出现过，那从target出现的那个点到当前这个点的和就是k。

我们要找的就是target出现的坐标和当前的坐标之差。

特殊的，没有元素时默认和为0，所以要put一个-1进去，比如恰好第一个数num=k，有个-1可以使我们正好得到答案。

题解代码

public class Solution {

/\*\*

\* @param nums: an array

\* @param k: a target value

\* @return: the maximum length of a subarray that sums to k

\*/

public int maxSubArrayLen(int[] nums, int k) {

// Write your code here

// 保存从头开始以i结尾的sum，以及i的位置，有相同的sum时，只保留最先出现的那个

// key就是sum，value是对应的数组下标，即i

// 例如 [1, -1, 5, -2, 3]

// [0, 1, 0, 5, 3, 6]

// 有了sum之后，令target=sum-k，如果target之前出现过，那从target出现的那个点到当前这个点的和就是k

// 我们要找的就是target出现的坐标和当前的坐标之差

// 特殊的，没有元素时默认和为0，所以要put一个-1进去，比如恰好第一个数num=k，有个-1可以使我们正好得到答案

int ans = 0;

Map<Integer, Integer> map = new HashMap<>();

map.put(0, -1);

int sum = 0;

for(int i=0;i<nums.length;i++){

sum = sum + nums[i];

int target = sum-k;

// 如果target之前出现过，则target出现的坐标和当前的坐标之差就是答案

if(null!=map.get(target)){

ans = Math.max(ans, (i-map.get(target)));

}

// sum有相同的时候只要最先出现的，只有以前没put过的才需要put

if(null==map.get(sum)){

map.put(sum, i);

}

}

return ans;

}

}