<https://leetcode.com/problems/subarray-sum-equals-k/>

Given an array of integers nums and an integer k, return *the total number of subarrays whose sum equals to* k.

A subarray is a contiguous **non-empty** sequence of elements within an array.

**Example 1:**

Input: nums = [1,1,1], k = 2

Output: 2

**Example 2:**

Input: nums = [1,2,3], k = 3

Output: 2

**Constraints:**

1 <= nums.length <= 2 \* 10^4

-1000 <= nums[i] <= 1000

-10^7 <= k <= 10^7

**Attempt 1: 2023-02-04**

**Solution 1:  Native for loop with 2 passes (30 min, first pass is create preSum array, second pass is calculate interval with nested for loop)**

class Solution {

    public int subarraySum(int[] nums, int k) {

        int[] preSum = new int[nums.length + 1];

        preSum[0] = 0;

        for(int i = 1; i < preSum.length; i++) {

            preSum[i] = preSum[i - 1] + nums[i - 1];

        }

        int count = 0;

        for(int start = 0; start < preSum.length; start++) {

            for(int end = start + 1; end < preSum.length; end++) {

                if(preSum[end] - preSum[start] == k) {

                    count++;

                }

            }

        }

        return count;

    }

}

Time Complexity: O(n^2)

Space Complexity: O(n)

**Solution 2:  Hash Table + Auxiliary array (30 min)**

class Solution {

    public int subarraySum(int[] nums, int k) {

        Map<Integer, Integer> map = new HashMap<Integer, Integer>();

        //map.put(0, 0);

        map.put(0, 1);

        int[] preSum = new int[nums.length + 1];

        int count = 0;

        for(int i = 1; i <= nums.length; i++) {

            preSum[i] = preSum[i - 1] + nums[i - 1];

            if(map.containsKey(preSum[i] - k)) {

                count += map.get(preSum[i] - k);

            }

            map.put(preSum[i], map.getOrDefault(preSum[i], 0) + 1);

        }

        return count;

    }

}

Time Complexity: O(n)

Space Complexity: O(n)

**Refer to**

<https://leetcode.com/problems/subarray-sum-equals-k/solutions/803317/java-solution-with-detailed-explanation/>

**Thinking**

1. Use an array to store the sum accumulated from the beginning to a certain position.

Example:

nums = [1,  2,  3  ]

sum  = [1, 1+2, 1+2+3]

2. How to create array "sum" ?

sum[i] = sum[i - 1] + nums[i]

Q : If i == 0, the index is out of range. How to solve this problem ?

A : Set the first element of the array "sum" to 0, and initialize the array "sum" **from index 1 rather than 0**.

nums = [1,  2,  3  ]

sum  = [0,  1,  1+2, 1+2+3] // Also, the length of "sum" is one more than "nums"

sum[i] = sum[i - 1] + nums[i - 1]

// Java Version

int[] sum = new int[nums.length + 1];

sum[0] = 0;

for (int i = 1; i < (nums.length + 1); i++)

  sum[i] = sum[i - 1] + nums[i - 1];

3. Using array "sum" to calculate the sum of a subarray

sumOfSubarray = sum[end] - sum[start];

For example : Calculate the sum of "nums" means using the last element of "sum" minus the first element of "sum" which is 0.

nums[0] + nums[1] + nums[2] = sum[3] - sum[0] = 6 - 0

4.Using array "sum"to caculate all possibilities .

**Code**

public class Solution {

    public int subarraySum(int[] nums, int k) {

        int count = 0;

        int[] sum = new int[nums.length + 1];

        sum[0] = 0;

        for (int i = 1; i <= nums.length; i++)

            sum[i] = sum[i - 1] + nums[i - 1];

        for (int start = 0; start < sum.length; start++) {

            for (int end = start + 1; end < sum.length; end++) {

                if (sum[end] - sum[start] == k)

                    count++;

            }

        }

        return count;

    }

}

**Complexity Analysis**

Time complexity : *O(n^2)*.

Space complexity : *O(n)*.

**Optimization by Hashmap**

**Thinking**

1.In the previous method

Step 1. The "nums" array is traversed to calculate all the elements of the sum array

Step 2. Use the nested loop to judge.

key : Can we judge when the array is traversed(Step 1) ?

Transposition

int[] sum = new int[nums.length + 1];

sum[0] = 0;

for (int end = 1; end < (nums.length + 1); end++)

  sum[end] = sum[end - 1] + nums[end - 1];

a. Put each element of "sum" array into hashmap according to this format : (*sumi*, *number of occurence*)

b. **When constructing the "sum" array, we take the currently constructed element as sum[end], then all the elements before "end" which have been calculated can be regarded as all sum[start] for this "end".**

Transform the judgment condition

Obviously, when sum[end] is calculated, all its possible sum[start] are already in the map.

sum[end] - sum[start] == k

sum[end] - k == sum[start]

c. When *sumend* is calculated, we only need to determine whether there is key == *sumend*- k in the hashmap and add the number of occurrence to the answer.

**Attention : In the previous method, we set the first element of sum to 0. Similarly, we put it in the hashmap, which is (0, 1).**

**Code**

public class Solution {

    public int subarraySum(int[] nums, int k) {

        int count = 0, sum = 0;

        HashMap < Integer, Integer > map = new HashMap < > ();

        map.put(0, 1);

        for (int i = 0; i < nums.length; i++) {

            sum += nums[i];

            if (map.containsKey(sum - k))

                count += map.get(sum - k);

            map.put(sum, map.getOrDefault(sum, 0) + 1);

        }

        return count;

    }

}

**Complexity Anaysis**

Time complexity : *O(n)*.

Space complexity : *O(n).*

**Solution 3:  Hash Table + Single Varaiable (30 min)**

**Actually no need preSum array, we can replace its functionality with a single variable 'presSum'**

class Solution {

    public int subarraySum(int[] nums, int k) {

        int count = 0;

        int preSum = 0;

        Map<Integer, Integer> map = new HashMap<>();

        map.put(0, 1);

        for(int num : nums) {

            preSum += num;

            count += map.getOrDefault(preSum - k, 0);

            map.put(preSum, map.getOrDefault(preSum, 0) + 1);

        }

        return count;

    }

}

Time Complexity: O(n)

Space Complexity: O(1)

**Refer to**

<https://leetcode.com/problems/subarray-sum-equals-k/solutions/102106/java-solution-presum-hashmap/>

Solution 1. Brute force. We just need two loops (i, j) and test if SUM[i, j] = k. Time complexity O(n^2), Space complexity O(1). I bet this solution will TLE.

Solution 2. From solution 1, we know the key to solve this problem is SUM[i, j]. So if we know SUM[0, i - 1] and SUM[0, j], then we can easily get SUM[i, j]. To achieve this, we just need to go through the array, calculate the current sum and save number of all seen PreSum to a HashMap. Time complexity O(n), Space complexity O(n).

public class Solution {

public int subarraySum(int[] nums, int k) {

int sum = 0, result = 0;

Map<Integer, Integer> preSum = new HashMap<>();

preSum.put(0, 1);

for (int i = 0; i < nums.length; i++) {

sum += nums[i];

if (preSum.containsKey(sum - k)) {

result += preSum.get(sum - k);

}

preSum.put(sum, preSum.getOrDefault(sum, 0) + 1);

}

return result;

}

}

**A harder problem**

<https://leetcode.com/problems/number-of-submatrices-that-sum-to-target/>

<https://leetcode.com/problems/number-of-submatrices-that-sum-to-target/discuss/803353/java-solution-with-detailed-explanation>

**Refer to**

[L713.Subarray Product Less Than K](note://4B9EF6834F24414AB80E2515E98D02C4)

[L724.Find Pivot Index](note://WEB5ebe7989ed9b56c712df3e501b165608)

[L1074.Number of Submatrices That Sum to Target (Ref.L560)](note://WEB47e31ed01fd0cb1f8aa407face798345)

[L1658.Minimum Operations to Reduce X to Zero (Ref.L1423)](note://WEB2d3951b594e0ba1abab46b9e91d0033c)

[L2090.K Radius Subarray Averages (Ref.L560)](note://WEBe84a8ca2031715d4b4b367e29560077f)

[L2219.Maximum Sum Score of Array (Ref.L560)](note://WEB2af62d40f8ba85a933c982e8a5479ed2)