

We may grade a **subset of the assigned questions**, to be determined after the deadline, so that we can provide better feedback on the graded questions.

Unless otherwise stated, each question requires sufficient justification to convince the reader of the correctness of your answer.

For bonus questions, we will not provide any insight during office hours or Piazza, and we do not guarantee anything about the difficulty of these questions.

We strongly encourage you to typeset your solutions in L<sup>A</sup>T<sub>E</sub>X.

If you collaborated with someone, you must state their name(s). You must write your own solution for all problems and may not look at any other student's write-up.

0. If applicable, state the name(s) and uniusername(s) of your collaborator(s).

**Solution:**

1. **Extra credit:** *You do not have to do this question to receive full credit on this assignment.* To receive the bonus points, you must typeset this **entire** assignment in L<sup>A</sup>T<sub>E</sub>X and draw a table with two columns that includes the *name* (e.g., “fraction”) and an *example* of each of the following:

- fraction (using `\frac`)
- less than or equal to
- union of two sets
- an expression involving a sum ( $\sum$ ) and product ( $\prod$ ).
- write a quantified formula about the reals ( $\mathbb{R}$ ) using both quantifiers ( $\exists, \forall$ ).

Unlike most extra credit questions, we will help with this in office hours.

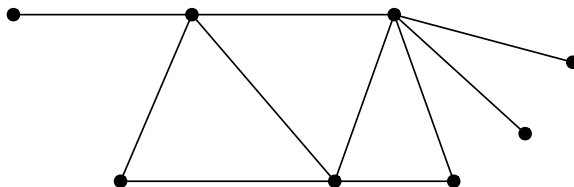
**Solution:**

2. Use induction to prove the following statements for all integers  $n \geq 1$ .

(a)  $\sum_{i=0}^n x^i = \frac{x^{n+1}-1}{x-1}$ , for all  $x \neq 1$ .

(b) The number of binary strings of length  $n$  is  $2^n$ .

(c) In any drawing of a planar graph  $G$  having  $v$  vertices,  $e$  edges,  $f$  faces, and  $c$  connected components,  $v + f - e - c = 1$ . (A planar graph is one that can be drawn in the plane without crossing edges. A *face* is a connected region of the plane after removing vertices and edges. If you remove a vertex or edge from a planar graph, it is still planar. Below is a planar graph with  $v = 8, e = 10, f = 4, c = 1$ .)



**Solution:**

3. A rational number may be written as  $a/b$  for some integers  $a, b$ , where  $b \neq 0$ . Prove that  $\sqrt{p}$  is not a rational number, for any prime  $p$ . *Hint:* Use a proof by contradiction.

**Solution:**

4. Recall that for positive functions  $f(n), g(n)$ , we say that  $f(n) = O(g(n))$  if there exist constants  $c, n_0 > 0$  such that  $f(n) \leq c \cdot g(n)$  for every  $n \geq n_0$ . Alternatively, a *sufficient* condition is that  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$  is finite. (Note, however, that this condition is not *necessary*; it is possible that  $f(n) = O(g(n))$  even when the limit does not exist.)

For the following pairs of  $f(n)$  and  $g(n)$ , is it true that  $f(n) = O(g(n))$ ? Justify your answer. *Hint:* You might find L'Hôpital's Rule useful for some questions: it says that  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)}$  (when the latter limit exists), where  $f'$  and  $g'$  are the derivatives of  $f$  and  $g$ , respectively.

- (a)  $f(n) = n + \log_2(n^4)$ ,  $g(n) = \frac{1}{9}n + 5$ .

**Solution:**

- (b)  $f(n) = (\ln n)^3$ ,  $g(n) = 3^{\log_2 n}$

**Solution:**

- (c)  $f(n) = 2^{1.3n}$ ,  $g(n) = \frac{1}{2}e^n$

**Solution:**

- (d)  $f(n) = \log_2(n^{100})$ ,  $g(n) = \log_2(n)$

**Solution:**

- (e)  $f(n) = n^3$ ,  $g(n) = \binom{n}{3}$ .

**Solution:**

5. Suppose you want to count the number of primes less than or equal to  $t$ . You have a function `is-prime( $i$ )` that takes 1 time-step and returns true if  $i$  is prime.

```
function count-prime(t): // t >= 2 is a non-negative integer
    count = 0
    for i from t down to 2
        if (is-prime(i))
            count++
    return count
```

What is the running time of count-prime, as a function of its input size? (First determine what the input size is.) Give an asymptotic upper bound (big-O) on its running time.

**Solution:**

6. Suppose you want to unambiguously represent all of the elements of  $\{0, 1, \dots, n-1\}$  as strings in  $\Sigma^k$ . Here  $\Sigma$  is a set of symbols (the *alphabet*) and  $\Sigma^k$  means the set of all strings over  $\Sigma$  with length exactly  $k$ . Given  $|\Sigma| \geq 2$  and  $n$ , what is the smallest possible  $k$  that allows this?

**Solution:**