

DEGREE, GIRTH AND CHROMATIC NUMBER

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§1.

The following notations will be used throughout: $\chi(G)$ is the chromatic number of a graph G ; \mathcal{L}_σ is the class of graphs with the maximal degree of vertices not exceeding σ ; \mathcal{L}^g is the class of graphs whose girth is at least g ; $\mathcal{L}_\sigma^g = \mathcal{L}^g \cap \mathcal{L}_\sigma$.

$\lfloor x \rfloor$ and $\lceil x \rceil$ denote respectively the lower and upper integers of x (i.e. $x - 1 < \lfloor x \rfloor \leq x$ and $x \leq \lceil x \rceil < x + 1$).

It is evident that for any σ and g

$$\max_{G \in \mathcal{L}_\sigma^g} \chi(G) \geq \max_{G \in \mathcal{L}_\sigma^{g+1}} \chi(G),$$

hence the sequence of these maxima becomes constant (depending on σ) $\Psi(\sigma) = \min_g \max_{G \in \mathcal{L}_\sigma^g} \chi(G)$ after a finite number of steps.

In 1968 Vizing [7] set up the problem: *Determine the maximal chromatic number of the graphs, contained in \mathcal{L}_σ^4 .*

The following question is also of interest: *How large is the number $\Psi(\sigma)$?*

Grünbaum [5] has formulated the conjecture, suggested by the papers [9], [3], [4], that $\Psi(\sigma) = \sigma$, if $\sigma \geq 3$.

In [1] and [2] it has been independently shown that for any $G \in \mathcal{L}_\sigma^4$ ($\sigma \leq 4$)

$$\chi(G) \leq \left\lfloor \frac{3(\sigma + 2)}{4} \right\rfloor,$$

and, consequently

$$\Psi(\sigma) \leq \left\lfloor \frac{3(\sigma + 2)}{4} \right\rfloor.$$

Thus, Grünbaum's conjecture does not hold for $\sigma \geq 7$. The present paper is devoted to proving the following fact.

Theorem. *Let $\sigma \geq 5$, $g \geq 7$, $G \in \mathcal{L}_\sigma^g$. Let, further, $q = \left\lceil \frac{g}{2} \right\rceil$. If some natural number χ satisfies the inequalities*

$$(1) \quad \chi \geq \left\lfloor \frac{\sigma}{2} \right\rfloor + 2,$$

$$(2) \quad \left(\frac{\chi - 1}{\sigma - \chi + 1} \right)^{q-1} \geq \frac{e}{2} q \sigma (\sigma + \chi - 2),$$

then $\chi(G) \leq \chi$.

The theorem will be proved in §§2-5.

Corollary 1. *If $\sigma \geq 5$ then $\Psi(\sigma) \leq \left\lfloor \frac{\sigma}{2} \right\rfloor + 2$.*

Proof. It is easy to see that for any $\sigma \geq 5$, and for any natural number $g \geq 4(\sigma + 2) \ln \sigma$, $\chi = \left\lfloor \frac{\sigma}{2} \right\rfloor + 2$ satisfies all the conditions of the theorem.

Corollary 2. $\max_{G \in \mathcal{L}_5^{35}} \chi(G) \leq 4$.

For the proof it suffices to verify that the conditions of the theorem are satisfied for $\sigma = 5$, $g = 35$, $\chi = 4$.

According to Corollary 1, Grünbaum's conjecture is not true for $\sigma \geq 5$.

§ 2.

We only consider $\chi \leq \sigma - 1$, for at $\chi \geq \sigma$ the statement of the theorem is the weakening of Brooks' theorem which asserts $\chi(G) \leq \sigma$ for any graph $G (\in \mathcal{L}_\sigma)$ not containing a complete subgraph with $\sigma + 1$ vertices.

Assume that the statement of the theorem is not true. Then there exists a $(\chi + 1)$ -critical graph $G \in \mathcal{L}_\sigma^g$. Let $v_0 \in V(G)$ be chosen and some colouring f of the vertices of $G \setminus \{v_0\}$ with χ colours be given.

Let $f(A)$, where $A \subseteq V(G)$, denote the set $\{f(v) \mid v \in A \setminus \{v_0\}\}$, and let $I(v)$ denote the set $\{w \in V(G) \mid (v, w) \in E(G)\}$. We shall call $w \in I(v_0)$ an $O_1(v_0)$ -vertex, if $f(w) \notin f(I(v_0) \setminus \{w\})$. The set of all $O_1(v_0)$ -vertices is denoted by $O_1(v_0)$.

As G is critical, $|O_1(v_0)| \geq 2\chi - \sigma \geq 3$.

We propose an algorithm for determining a subset Γ of the set $V(G) \cup E(G)$ in G . The set Γ will play the main role in the proof of the theorem. The edges of Γ will be oriented, some of them in both directions. While constructing Γ , the edges and vertices, belonging to Γ , will be called Γ -edges and Γ -vertices respectively. Further, the edges in Γ will be divided into Γ_1 -edges and Γ_2 -edges. The algorithm will work in not more than $\sigma |V(G)|$ steps. At the i -th step Γ -edges and Γ -vertices of i -th level will be defined. Simultaneously, with the construction of Γ we shall construct a mapping P , defined on the set of Γ_2 -edges with the values in the set of Γ -vertices.

THE ALGORITHM OF CONSTRUCTING Γ

Step 0. v_0 is called a Γ -vertex of level 0.

Step 1. Direct each edge (v_0, w) , where w is a $O_1(v_0)$ -vertex, towards w . We call these directed edges Γ_1 -edges of level 1, and we refer

to the $O_1(v_0)$ -vertices as Γ -vertices of level 1. There are no Γ_2 -edges of level 1. Go to Step 2.

Definition. For any natural number i and for each Γ -vertex $v \neq v_0$ we denote by $T_i(v)$ the set of those Γ -vertices, which belong to $I(v)$, and from which Γ_1 -edges of level $\leq i$ go to v .

Example. If v is an $O_1(v_0)$ -vertex then $T_1(v) = \{v_0\}$.

Definition. For any natural number $i \geq 2$ and each Γ -vertex $v \neq v_0$ we denote

$$O_i(v) = \left\{ w \in I(v) \setminus (T_{i-1}(v) \cup \{v_0\}) \mid f(w) \notin f(I(v) \setminus (T_{i-1}(v) \cup \{w\})) \text{ \& } w \notin \bigcup_{j=2}^{i-1} O_j(v) \right\}.$$

Example. If $v \in O_1(v)$, then

$$O_2(v) = \{w \in I(v) \setminus \{v_0\} \mid f(w) \notin f(I(v) \setminus \{w, v_0\})\}.$$

Definition. Let $v \neq v_0$ be a Γ -vertex. We shall say that the $D_i(v)$ -situation takes place in G , if $|f(I(v) \setminus T_{i-1}(v))| < \chi - 1$.

Step k ($k \geq 2$).

(a) If for at least one Γ -vertex $v \neq v_0$ the $D_{k-1}(v)$ -situation takes place, then the algorithm terminates. Otherwise, the algorithm terminates if no Γ -edge of the $(k-1)$ -th level has been constructed in the $(k-1)$ -th step. In all other cases go to item (b).

(b) For each ordered pair of vertices $\{v, w\}$, where $v \neq v_0$ is a Γ -vertex of level 1 or 2 or 3... or $k-1$, and $w \in O_k(v)$, we direct the edge (v, w) towards w . We call all such edges the Γ -edges of the k -th level. Go to item (c).

Remark 1. It may happen that some edge is directed in both senses.

(c) Let us consider an arbitrary Γ -edge $(\overrightarrow{v, w})$ of the k -th level. $f(v) = \alpha$, $f(w) = \beta$. If there is a Γ -vertex $u (\in I(v) \cup I(w))$ of level not exceeding $k-q$, then we call $(\overrightarrow{v, w})$ a Γ_2 -edge of the k -th level. Be-

sides, we call $(\overrightarrow{v, w})$ a Γ_2 -edge of the k -th level if for some $s (\geq 2)$ there exists in G a directed chain

$$(\overrightarrow{v_1, v_2}), (\overrightarrow{v_2, v_3}), \dots, (\overrightarrow{v_{s-1}, v_s})$$

of Γ_1 -edges such that $v_s = v$, $f(v_j) \in \{\alpha, \beta\}$ (where $1 \leq j \leq s$) and at least one of the vertices v_1, v_2, \dots, v_s is adjacent to a Γ -vertex u' the level of which does not exceed $k - q$. The vertex u (or u'), because of which $(\overrightarrow{v, w})$ became a Γ_2 -edge, will be called the image of $(\overrightarrow{v, w})$ in the mapping P . (If there exist more than one such vertices u or u' , then we choose $P(\overrightarrow{v, w})$ arbitrarily from among them.) We check each Γ -edge of level k whether or not it is a Γ_2 -edge; if it is not, we call it a Γ_1 -edge of level k . Go to item (d).

(d) A vertex $v \in V(G)$ will be called a Γ -vertex of the k -th level, if at least one Γ_1 -edge of the k -th level enters it, but no Γ_1 -edge of lower level. Go to Step $k + 1$.

Remark 2. If an edge of G is directed in both directions, it may be a Γ_1 -edge in one direction and a Γ_2 -edge in the other.

If no Γ_1 -edge of the k -th level appears in Step k of the algorithm, then for any Γ -vertex v

$$T_{k-1}(v) = T_k(v), \quad O_{k+1}(v) = \phi,$$

and at Step $(k + 1)$ there will not appear any Γ -edge of the $(k + 1)$ -th level. That is, the algorithm terminates not later than at Step $(k + 1)$. Consequently, the algorithm works in at most $2 \cdot |E(G)|$ steps.

Later we shall denote the level of the Γ -vertex v or that of the Γ -edge \overrightarrow{e} by $Y(v)$ or $Y(\overrightarrow{e})$ respectively. It is clear that $Y((\overrightarrow{v, w}))$ and $Y((\overrightarrow{w, v}))$ may be different.

§3.

In this section we consider some properties of Γ .

(I). If $(\overrightarrow{v, w})$ is a Γ -edge, then $Y((\overrightarrow{v, w})) \geq Y(v) + 1$. If $(\overrightarrow{v, w})$ is a Γ_1 -edge, then $Y((\overrightarrow{v, w})) \geq Y(w)$.

Proof. The first inequality follows from the definition of Γ -edges of the k -th level. By the definition of Γ -vertices of level k , the level of any Γ -vertex v is equal to the minimum of levels of Γ_1 -edges entering into v . This implies the second inequality.

(II). If $(\overrightarrow{v, w})$ and $(\overrightarrow{w, u})$ are Γ -edges, $v \neq u$ and $f(v) = f(u)$, then $(\overrightarrow{v, w})$ is a Γ_1 -edge, and $Y((\overrightarrow{v, w})) \leq Y((\overrightarrow{w, u})) - 1$.

Proof. Let $Y((\overrightarrow{w, u})) = i$. If $v \notin T_{i-1}(v)$, then $u \notin O_i(v)$, and $(\overrightarrow{w, u})$ would not be a Γ -edge of the i -th level. Consequently, $v \in T_{i-1}(v)$. That is, $(\overrightarrow{v, w})$ is a Γ_1 -edge, and $Y((\overrightarrow{v, w})) \leq i - 1$.

The next statement is obvious.

(III). For any Γ -vertex $v \neq v_0$ there exists a Γ_1 -edge $(\overrightarrow{w, v})$ such that $Y(v) = Y((\overrightarrow{w, v})) \geq Y(w) + 1$.

As an immediate corollary of (I) and (III) we state:

(IV). For any Γ -vertex v the length of the shortest directed chain, consisting of Γ_1 -edges and leading from v_0 to v , does not exceed $Y(v)$.

(V). There is no Γ_1 -edge of level > 1 , which terminates at a vertex adjacent to v_0 .

Proof. Each edge, whose level exceeds 1, and whose end vertex belongs to $I(v_0)$ is a Γ_2 -edge. By (IV) and since $G \in \mathcal{L}^g$, its level is at least $g - 2$.

(VI). For any directed two-coloured chain $(\overrightarrow{v_1, v_2}), (\overrightarrow{v_2, v_3}), \dots, (\overrightarrow{v_{s-1}, v_s})$, consisting of Γ_1 -edges, the following is true:

$$Y((\overrightarrow{v_{s-1}, v_s})) \leq Y(v_1) + q - 2.$$

Proof. Since the chain is two-coloured, v_1 has a colour, and $v_1 \neq v_0$. Then, according to (III), there exists a $w \in I(v_1)$ with $Y(w) \leq Y(v_1) - 1$. But, taking into account the definition of Γ_2 -edges, if $Y((\overrightarrow{v_{s-1}, v_s})) \geq Y(w) + q$ then $(\overrightarrow{v_{s-1}, v_s})$ will be a Γ_2 -edge.

The following statement results from the definition of the sets $O_i(v)$ and the Γ -edges of the k -th level.

(VII). Let $v \in \Gamma \setminus \{v_0\}$, $f(v) = \alpha$. Then for any $\beta \neq \alpha$ there exists at most one Γ -edge, going from v to a vertex of colour β . Moreover, if there exists a vertex $w \in I(v)$ such that $f(w) = \beta$, $(\overrightarrow{v, w})$ is a Γ -edge and $Y(\overrightarrow{v, w}) = k$, then any vertex $u \in I(v) \setminus \{w\}$ with $f(u) = \beta$ belongs to $T_{k-1}(v)$.

Definition. Let α and β be arbitrary colours. We denote by $G_{\alpha\beta}$ the subgraph of the graph G , spanned by the vertices whose colour is α and β .

Definition. Let $(\overrightarrow{u, v})$ be Γ_1 -edge, $f(u) = \alpha$, $f(v) = \beta$. We denote by $G_{\alpha\beta}((\overrightarrow{u, v}))$ the connected component of the graph $G_{\alpha\beta} \setminus \{(u, v)\}$, containing the vertex u .

From (II), (VI) and (VII) we obtain:

(VIII). Each component $G_{\alpha\beta}((\overrightarrow{u, v}))$ is a rooted tree* with the root u , and its height does not exceed $q - 3$; furthermore, any edge of the tree is a Γ_1 -edge, and is directed towards u .

Definition. For any Γ -vertex $v \neq v_0$ we define the notion of the v -tree by induction with respect to the level of the vertices:

1. If $Y(v) = 1$, then a v -tree consists of the vertices v_0, v and of the Γ_1 -edge $(\overrightarrow{v_0, v})$.

2. Let the u -tree be defined for each Γ -vertex $u \neq v_0$ with $Y(u) < k$. We consider a Γ -vertex v with $Y(v) = k$ and $f(v) = \alpha$. According to (III), there exists a $u_0 \in T_k(v)$. Let $f(u_0) = \beta$. We choose among the initial vertices of the graph $G_{\beta\alpha}((\overrightarrow{u_0, v}))$ a vertex w in such a way, that the directed chain, consisting of Γ_1 -edges leading in the graph $G_{\beta\alpha}((\overrightarrow{u_0, v}))$ from w to u_0 , would end in a Γ_1 -edge $(\overrightarrow{w', u_0})$,

* The root is an arbitrary distinguished vertex of the tree. The height of a vertex of a rooted tree is its distance from the root. By (VII), each Γ_1 -edge of the tree in question is directed towards the root u .

having the least level among the edges of $G_{\beta\alpha}(\overrightarrow{(u_0, v)})$, entering u_0 . If $G_{\beta\alpha}(\overrightarrow{(u_0, v)}) = \{u_0\}$, then we take $w \doteq u_0$. From (I) and (II) it follows that $Y(w) < Y(v)$. Then all the vertices and Γ_1 -edges of the w -tree, all the vertices and Γ_1 -edges of $G_{\beta\alpha}(\overrightarrow{(u_0, v)})$, the Γ_1 -edge $\overrightarrow{(u_0, v)}$ and the Γ -vertex v belong to the v -tree (and the v -tree consists of these elements only).

Remark 3. Generally speaking, the v -tree is not unique (since it depends on the choice of the vertices w).

(IX). Let $v \neq v_0$ be some Γ -vertex, and $F(v)$ some arbitrary v -tree. Then

(a) if $\overrightarrow{(u, w)} \in F(v)$, $\overrightarrow{(w, y)} \in F(v)$, $u \neq y$, then $Y(\overrightarrow{(u, w)}) < Y(\overrightarrow{(w, y)})$.

(b) $F(v)$ is the directed tree with root v ; each edge in $F(v)$ is a Γ_1 -edge, its height does not exceed $Y(v)$. Its edges are directed towards v . The vertex v_0 is one of the initial vertices of this tree. Only one Γ_1 -edge, belonging to $F(v)$, goes from any of the vertices $w \in F(v) \setminus \{v\}$.

Proof. Let us prove this statement by the induction on level v . If $Y(v) = 1$, then this statement is obvious.

Suppose that this statement is true for all Γ -vertices of the level not exceeding $k - 1$ and $Y(v) = k$. According to the definition of the v -tree, there exist such Γ -vertices u_0 and w , that $F(v)$ consists of vertices and Γ_1 -edges of a w -tree and $G_{\beta\alpha}(\overrightarrow{(u_0, v)})$, and of the Γ_1 -edge $\overrightarrow{(u_0, v)}$ and the vertex v . Since $Y(w) < k$, for the w -tree (IX) is valid. We show that no vertex of $G_{\beta\alpha}(\overrightarrow{(u_0, v)})$, except w , belongs to the w -tree. Suppose that some vertex $u \neq w$ lies simultaneously in the w -tree and in $G_{\beta\alpha}(\overrightarrow{(u_0, v)})$. Then a directed chain $\overrightarrow{(u, u_1)}, \overrightarrow{(u_1, u_2)}, \dots, \overrightarrow{(u_{s-1}, u_s)}$ leads from u to u_0 such that $u_s = u_0$ and $f(u_j) \in \{\alpha, \beta\}$ ($1 \leq j \leq s$). Besides, another directed chain

$$\overrightarrow{(u, y_1)}, \overrightarrow{(y_1, y_2)}, \dots, \overrightarrow{(y_{r-1}, w)}, \overrightarrow{(w, y_{r+1})}, \dots, \overrightarrow{(y_{l-1}, u_0)}$$

leads from u to u_0 . According to (II) and (VI)

$$(3) \quad Y(\overrightarrow{(u_0, v)}) \leq Y(u) + q - 2;$$

$$Y(\overrightarrow{(u_i, u_{i+1})}) < Y(\overrightarrow{(u_{i+1}, u_{i+2})}), \quad i = 1, 2, \dots, s-2;$$

$$Y(\overrightarrow{(y_i, y_{i+1})}) < Y(\overrightarrow{(y_{i+1}, y_{i+2})}), \quad i = r, r+1, \dots, l-2.$$

Besides, according to the definition of the v -tree $Y(\overrightarrow{(y_{r-1}, w)}) < Y(\overrightarrow{(w, y_{r+1})})$ and in accordance with the induction hypothesis

$$Y(\overrightarrow{(y_i, y_{i+1})}) \leq Y(\overrightarrow{(y_{i+1}, y_{i+2})}), \quad i = 1, 2, \dots, r-2;$$

$$Y(\overrightarrow{(u, y_1)}) \leq Y(\overrightarrow{(y_1, y_2)}),$$

However, by (I),

$$Y(\overrightarrow{(u, u_1)}) \geq Y(u) + 1, \quad Y(\overrightarrow{(u, y_1)}) \geq Y(u) + 1.$$

Consequently,

$$Y(\overrightarrow{(u_{s-1}, u_0)}) \geq Y(u) + s, \quad Y(\overrightarrow{(y_{l-1}, u_0)}) \geq Y(u) + l.$$

Since $Y(\overrightarrow{(u_0, v)}) \geq \max \{Y(\overrightarrow{(y_{l-1}, u_0)}) + 1, Y(\overrightarrow{(u_{s-1}, u_0)}) + 1\}$, it follows from (3) that $\max \{s, l\} \leq q - 3$.

So we obtained that there exists in G a cycle $(u, u_1, u_2, \dots, u_s, y_{l-1}, y_{l-2}, \dots, y_1, u)$, the length of which is $s + l$. But

$$s + l \leq 2(q - 3) < g.$$

Hence, $F(v)$ is a tree. Verification of the further parts of the statement does not raise any difficulties.

(X). Let $v \neq v_0$ be a Γ -vertex. Only the Γ_1 -edges, belonging to $F(v)$, are the edges of the subgraph of the graph G , generated by the vertices of an arbitrary v -tree $F(v)$.

Proof. Suppose that the vertices x and y , belonging to $F(v)$, are connected by the edge $(x, y) \in E(G) \setminus E(F(v))$.

Case 1. Let the vertex y lie in the directed chain, leading in $F(v)$ from the vertex x to the vertex v . The necessary condition for $(\overrightarrow{w, y})$

lying in this chain, to be a Γ_1 -edge, is*

$$Y(\overrightarrow{(w, y)}) \leq Y(x) + q - 1.$$

Taking (IX/a) into account we obtain that G contains a cycle of length not exceeding q .

Since the roles of the vertices x and y are symmetric, only the following case remained open.

Case 2. Let u be the first common vertex of the directed chains in $F(v)$, going from x to v , and from y to v , $u \notin \{x, y\}$. Then let $(\overrightarrow{u, u_1}) \in F(v)$, $f(u) = \alpha$, $f(u_1) = \beta$. According to the construction of $F(v)$ at least one of the vertices x and y belongs to $G_{\alpha\beta}(\overrightarrow{(u, u_1)})$. (VIII) implies that at most one of the vertices x and y can lie in $G_{\alpha\beta}(\overrightarrow{(u, u_1)})$. Let, for sake of definiteness, $y \in G_{\alpha\beta}(\overrightarrow{(u, u_1)})$. Let $(\overrightarrow{u_x, x})$ (or $(\overrightarrow{u_y, u})$) denote the last edge of the directed chain in $F(v)$, going from x (from y) respectively to u . Then, according to the construction of $F(v)$,

$$Y(\overrightarrow{(u_y, u)}) \leq Y(\overrightarrow{(u, u_1)}) - 1, \quad Y(\overrightarrow{(u_x, u)}) \leq Y(\overrightarrow{(u, u_1)}) - 1.$$

Since $(x, y) \in E(G)$, $(\overrightarrow{u, u_1})$ is a Γ_1 -edge and the directed chain, going from y to u_1 is two-coloured, therefore

$$Y(\overrightarrow{(u, u_1)}) \leq Y(x) + q - 1$$

and (by (VI))

$$Y(\overrightarrow{(u, u_1)}) \leq Y(y) + q - 2.$$

According to (IX/a) and (I) the length of the directed chain, going from x (or from y) to u in $F(v)$, does not exceed $q - 2$ ($q - 3$, respectively). Thus G contains a cycle with length at most

$$(q - 2) + (q - 3) + 1 = 2q - 4 < g.$$

Hence the proof is complete.

*Suppose the contrary, i.e.

$$Y(\overrightarrow{(w, y)}) \geq Y(x) + q.$$

Then $(\overrightarrow{w, y})$ is a Γ_2 -edge (by definition, such that u is replaced by x).

Definition. Suppose that the $D_k(v)$ -situation arises for some Γ -vertex $v \neq v_0$, and for some natural k in G . Let, further $f(v) = \alpha$. We define the v -trace according to the following rules.

1. If $|f(I(v))| < \chi - 1$, then any v -tree is a v -trace.

2. Suppose that $|f(I(v))| = \chi - 1$. Then, according to the definition of the $D_k(v)$ -situation, there exists such a colour β that $\beta \notin f(I(v) \setminus T_k(v))$. We consider the connected component $G_{\alpha\beta}(v)$ of the graph $G_{\alpha\beta}$, containing the vertex v . Due to (VI) and (VII) $G_{\alpha\beta}(v)$ is a directed tree with root v , whose height does not exceed $q - 2$. Each edge of this tree is a Γ_1 -edge, and is directed towards v . Let v_1 be an initial vertex of the tree $G_{\alpha\beta}(v)$, with the property that the level of the last Γ_1 -edge $(\overrightarrow{v'}, v)$ in the directed chain leading in the graph $G_{\alpha\beta}(v)$ from v_1 to v , is the least in comparison with the levels of the edges from $G_{\alpha\beta}(v)$, entering v . As v -trace we take all the edges and vertices of $G_{\alpha\beta}(v)$ and of arbitrary v_1 -tree.

Remark 4. Like the v -tree, the v -trace is not unique either.

(XI). Let $v \neq v_0$ be a Γ -vertex. Then the subgraph of G , generated by the vertices of any v -trace, coincides with this trace, and is the root-orientated tree with root v , the height of which does not exceed $Y(v)$. Each edge of this tree is Γ_1 -edge, and is directed in the direction of v . Only one Γ_1 -edge, belonging to the v -trace, goes from each vertex of this trace except vertex v . Vertex v_0 is one of the initial vertices of this tree.

The proof of (XI) is analogous to that of (X).

Lemma. In the process of constructing Γ we do not get $D_k(v)$ -situations for any pair $v (\in \Gamma)$, k .

Proof. It suffices to show that if for some pair v, k while constructing Γ , $D_k(v)$ -situation arises, then G is χ -colourable.

Let $F(v)$ be some v -trace. We define a function h on the vertices of $F(v)$ according to the following rules. If $|f(I(v))| < \chi - 1$, then a colour $\alpha \notin f(I(v) \cup \{v\})$ will be the image of the vertex v for the map-

ping h . Let $|f(I(v))| = \chi - 1$. Let us recall the definitions of the $D_k(v)$ -situation and the v -trace. All the vertices from $F(v)$, whose Γ_1 -edges enter to v , are coloured with the same colour β , and, besides, this colour has not been used for the colouring of the vertices from $I(v) \setminus F(v)$. Then we assume that $h(v) = \beta$. As about $h(w)$ for each vertex $w \in F(v) \setminus \{v\}$, we take the colour of such a vertex $w' \in F(v)$, that $(\overrightarrow{w, w'}) \in F(v)$. We define the function $f': V(G) \rightarrow \{1, 2, \dots, \chi\}$ such that

$$f'(w) = \begin{cases} f(w), & w \in V(G) \setminus F(v); \\ h(w), & w \in V(G) \cap F(v). \end{cases}$$

It follows from (XI), that f' is a correct colouring of the vertices of G by χ colours. Hence the Lemma is proved.

§4.

Thus, in course of constructing Γ , for each pair $v \in \Gamma$, k

$$(4) \quad |f(I(v) \setminus T_k(v))| = \chi - 1$$

is fulfilled. Thus (in addition to (I)-(VIII)) the following statements are valid for Γ .

(XII). For any Γ -vertex $v \neq v_0$ and for any natural i, j

$$O_j(v) \cap T_i(v) = \emptyset.$$

Proof. Suppose $w \in O_j(v) \cap T_i(v)$, $f(w) = \alpha$. According to (VII) all the vertices of colour α , lying in $I(v) \setminus \{w\}$, belong to $T_{j-1}(v)$. Consequently, $\alpha \notin f(I(v) \setminus T_{\max\{i, j-1\}}(v))$ which contradicts the relation (4).

(XIII). If $|T_i(v)| = a \geq 1$, then $\left| \bigcup_{j=2}^{i+1} O_j(v) \right| \geq 2\chi - \sigma - 2 + a$.

Proof. We have $|I(v) \setminus T_i(v)| \leq \sigma - a$,

$$\bigcup_{j=2}^{i+1} O_j = \{w \in I(v) \setminus T_i(v) \mid f(w) \notin f(I(v) \setminus (T_i(v) \cup \{w\}))\}.$$

By (4), among the $|I(v) \setminus T_i(v)|$ vertices we must come across those of $\chi - 1$ colours. But if not more than $\sigma - a$ elements are coloured by

$\chi - 1$ colours and each colour really occurs, then the number of colours, used only once is not less than

$$(\chi - 1) - ((\sigma - a) - (\chi - 1)) = 2\chi - 2 - \sigma + a.$$

The following statement immediately follows from (4) and (XIII).

(XIV). For any Γ -vertex $v \neq v_0$ and for any $i \geq 1$, $|T_i(v)| \leq \sigma - \chi + 1$ holds and, hence,

$$\begin{aligned} \frac{\left| \bigcup_{j=2}^{i+1} O_j(v) \right|}{|T_i(v)|} &\stackrel{(XIII)}{\geq} \frac{2\chi - 2 - \sigma + |T_i(v)|}{|T_i(v)|} = \\ &= 1 + \frac{2\chi - 2 - \sigma}{|T_i(v)|} \geq \frac{\chi - 1}{\sigma - \chi + 1}. \end{aligned}$$

§5.

For completing the proof of the Theorem it remains to show that if the $D_k(v)$ -situation never arises in course of constructing Γ , then the number of edges in Γ unboundedly increases; this contradicts the finiteness of G .

Let P be the mapping, defined in the course of constructing the Γ_2 -edges. Now we consider an arbitrary Γ -vertex u . Let $y \in I(u)$ and $P_y^{-1}(u)$ be the set of such Γ_2 -edges $(\overrightarrow{v, w})$ from $P^{-1}(u)$, which became Γ_2 -edges because of the fact that they are themselves incident to y , or because of the fact that the two-coloured chain, consisting of Γ_1 -edges of colours $f(v)$ and $f(w)$, passing through y , leads to vertex v . It is obvious that $P^{-1}(u) = \bigcup_{y \in I(u)} P_y^{-1}(u)$. The number of Γ_2 -edges, entering y , does not exceed σ . Besides, according to (VII), not more than $\chi - 1$ directed two-coloured chains come from y . Moreover, for at least one Γ -edge coming from y , it is necessary that $T_{2|E(G)|}(y) \neq \emptyset$.

Consequently, for each Γ -vertex u , and for each vertex $y \in I(u)$

$$|P_y^{-1}(u)| \leq \max \{ \sigma, (\sigma - 1) + (\chi - 1) \} = \sigma + \chi - 2.$$

Thus, for any Γ -vertex u

$$|P^{-1}(u)| = \left| \bigcup_{y \in I(u)} P_y^{-1}(u) \right| \leq \sigma(\sigma + \chi - 2).$$

It is clear that for each Γ_2 -edge $(\overrightarrow{v, w}) \in P^{-1}(u)$

$$Y((\overrightarrow{v, w})) \geq Y(u) + q.$$

Therefore, the number of Γ_2 -edges, the level of which does not exceed k , is bounded from above by $\sigma(\sigma + \chi - 2) \sum_{i=0}^{k-q} |V_i|$, where V_i is the set of the Γ -vertices of the i -th level.

Let E'_i (or E''_i) be the set of Γ_1 -edges (Γ_2 -edges respectively) of the i -th level. The set of all Γ -edges of the i -th level is $E_i = E'_i \cup E''_i$. Then

$$(5) \quad \sum_{i=1}^k |E''_i| \leq \sigma(\sigma + \chi - 2) \sum_{i=0}^{k-q} |V_i|.$$

Since $\chi \geq \left\lfloor \frac{\sigma}{2} \right\rfloor + 2$, and since for any Γ -vertex $v \in V_i$ ($i = 1, 2, \dots$), $|T_i(v)| \geq 1$, according to (XIII), $|O_{i+1}(v)| \geq 2$ ($v \in V_i$). Consequently,

$$(6) \quad |V_i| \leq \frac{1}{2} |E_{i+1}|, \quad i = 0, 1, 2, \dots$$

Thus, for every $k \geq 1$

$$\begin{aligned} (7) \quad \sum_{i=1}^{k+1} |E_i| &= \sum_{v \in \bigcup_{i=0}^k V_i} \left| \bigcup_{j=1}^{k+1} O_j(v) \right| \geq \\ &\geq \sum_{v \in \bigcup_{i=1}^k V_i} \frac{\left| \bigcup_{j=2}^{k+1} O_j(v) \right|}{|T_k(v)|} |T_k(v)| \stackrel{(XIV)}{\geq} \\ &\stackrel{(XIV)}{\geq} \frac{\chi - 1}{\sigma - \chi + 1} \sum_{v \in \bigcup_{i=1}^k V_i} |T_k(v)| \geq \frac{\chi - 1}{\sigma + 1 - \chi} \sum_{i=1}^k |E'_i| = \end{aligned}$$

$$\begin{aligned}
&= \frac{\chi-1}{\sigma-\chi+1} \left(\sum_{i=1}^k |E_i| - \sum_{i=1}^k |E_i''| \right) \stackrel{(5)}{\geq} \\
&\stackrel{(5)}{\geq} \frac{\chi-1}{\sigma-\chi+1} \left(\sum_{i=1}^k |E_i| - \sigma(\sigma+\chi-2) \sum_{i=1}^{k-q} |V_i| \right) \stackrel{(6)}{\geq} \\
&\stackrel{(6)}{\geq} \frac{\chi-1}{\sigma-\chi+1} \left(\sum_{i=1}^k |E_i| - \frac{\sigma(\sigma+\chi-2)}{2} \sum_{i=1}^{k-q+1} |E_i| \right).
\end{aligned}$$

We show that for all $k \geq 1$

$$(8) \quad \sum_{i=1}^{k+1} |E_i| \geq \frac{(\chi-1)(q-1)}{(\sigma-\chi+1)q} \sum_{i=1}^k |E_i|.$$

Since $E_1 \neq \phi$, and, according to (2), $\sqrt[q-1]{\frac{1}{e} \left(\frac{\chi-1}{\sigma-\chi+1} \right)^{q-1}} > 1$,

$$\frac{(\chi-1)(q-1)}{(\sigma-\chi+1)q} > 1.$$

(We have used that

$$\begin{aligned}
\frac{1}{e} \left(\frac{\chi-1}{\sigma-\chi+1} \right)^{q-1} &< \left(\frac{\chi-1}{\sigma-\chi+1} \right)^{q-1}, \\
\frac{1}{\left(1 + \frac{1}{q-1} \right)^{q-1}} &= \left(\frac{(\chi-1)(q-1)}{(\sigma-\chi+1)q} \right)^{q-1}.
\end{aligned}$$

So (8) will imply that the number of Γ -edges increases unboundedly, which contradicts to the finiteness of G .

So, for $1 \leq k < q$ inequality (8) immediately follows from (7). Let now (8) hold for all $k < k_0$. Then

$$\begin{aligned}
&\sum_{i=1}^{k_0+1} |E_i| \stackrel{(7)}{\geq} \\
&\stackrel{(7)}{\geq} \frac{\chi-1}{\sigma-\chi+1} \left(\sum_{i=1}^{k_0} |E_i| - \frac{\sigma(\sigma+\chi-2)}{2} \sum_{i=1}^{k_0-q+1} |E_i| \right) = \\
&= \frac{(\chi-1)(q-1)}{(\sigma-\chi+1)q} \sum_{i=1}^{k_0} |E_i| +
\end{aligned}$$

$$+ \frac{(\chi - 1)}{(\sigma - \chi + 1)} \left(\frac{1}{q} \sum_{i=1}^{k_0} |E_i| - \frac{\sigma(\sigma + \chi - 2)}{2} \sum_{i=1}^{k_0 - q + 1} |E_i| \right).$$

By the induction hypothesis

$$\begin{aligned} \sum_{i=1}^{k_0} |E_i| &\geq \frac{(\chi - 1)(q - 1)}{(\sigma - \chi + 1)q} \sum_{i=1}^{k_0 - 1} |E_i| \geq \dots \\ &\dots \geq \left(\frac{(\chi - 1)(q - 1)}{(\sigma - \chi + 1)q} \right)^{q-1} \sum_{i=1}^{k_0 - q + 1} |E_i|. \end{aligned}$$

Consequently,

$$\begin{aligned} \frac{1}{q} \sum_{i=1}^{k_0} |E_i| - \frac{\sigma(\sigma + \chi - 2)}{2} \sum_{i=1}^{k_0 - q + 1} |E_i| &\geq \\ &\geq \left(\frac{1}{q} \left(\frac{(\chi - 1)(q - 1)}{(\sigma - \chi + 1)q} \right)^{q-1} - \frac{\sigma(\sigma + \chi - 2)}{2} \right) \sum_{i=1}^{k_0 - q + 1} |E_i|, \end{aligned}$$

$$\text{and, in accordance with (2) } \frac{1}{q} \sum_{i=1}^{k_0} |E_i| - \frac{\sigma(\sigma + \chi - 2)}{2} \sum_{i=1}^{k_0 - q + 1} |E_i| \geq 0.$$

Thus,

$$\sum_{i=1}^{k_0 + 1} |E_i| \geq \frac{(\chi - 1)(q - 1)}{(\sigma - \chi + 1)q} \sum_{i=1}^{k_0} |E_i|.$$

which proves the theorem.

§6.

Remark 5. Using a theorem of Lovász [6] it is easy to prove that for any real $\alpha > \frac{1}{2}$ there exist natural numbers $g(\alpha)$ and $\sigma(\alpha)$ such that for any $g \geq g(\alpha)$, $\sigma \geq \sigma(\alpha)$ and $G \in \mathcal{L}_\sigma^g$

$$\chi(G) \leq \alpha \sigma.$$

The theorem of Lovász states: if $G \in \mathcal{L}_\sigma$ and $\sigma_1, \dots, \sigma_n$ are non-nega-

tive integers such that $\sigma + 1 = \sum_{i=1}^n (\sigma_i + 1)$, then the vertices of G admit a covering by subgraphs G_1, G_2, \dots, G_n such that, for any $1 \leq i \leq n$,

$$G_i \in \mathcal{L}_{\sigma_i}.$$

We remind now the notion, (introduced by V.G. Vizing [8]), of the prescribed colouring and the upper chromatic number.

Definition. By a prescription for the vertices of a graph $G(V, E)$ we understand the mapping Φ of the set of vertices of G to the set of subsets of natural numbers.

Definition. We say that the colouring f of the vertices of a graph G satisfies the prescription Φ , if $f(v) \in \Phi(v)$ for each vertex $v \in V(G)$.

Definition. The smallest natural number k with the following properties is called the upper chromatic number $W(G)$ of a graph G : for each prescription Φ , satisfying

$$(\forall v)(v \in V(G) \Rightarrow |\Phi(v)| \geq k)$$

there exists a colouring f_Φ of vertices of G such that f_Φ satisfies Φ . It is clear that $W(G) \geq \chi(G)$. V.G. Vizing [8] has constructed a graph G_k for each $k \geq 2$ such that $\chi(G_k) = 2$, $W(G_k) \geq k$.

Remark 6. The proof of our Theorem can be applied, practically without any alterations for the prescribed colourings.

In conclusion I would like to call the reader's attention to the difficult and important problem:

To find the best upper estimate for the chromatic number of the graph in terms of the maximal degree and density or girth.

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