

Improved lower bounds on the number of edges in list critical and online list critical graphs

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Abstract

We prove that every k -list-critical graph ($k \geq 8$) on $n \geq k + 2$ vertices has at least $\frac{1}{2} \left(k - 1 + \frac{k-3}{(k-c)(k-1)+k-3} \right) n$ edges where $c = (k-3) \left(\frac{1}{2} - \frac{1}{(k-1)(k-2)} \right)$. This improves the bound established by Kostochka and Stiebitz [15]. The same bound holds for online list-critical graphs, improving the bound established by Riasat and Schauz [20]. Both bounds follow from a more general result stating that either a graph has many edges or it has an Alon-Tarsi orientable induced subgraph satisfying a certain degree condition. We use these lower bounds to prove an Ore-degree version of Brooks' theorem for online list-coloring: every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is online $\lfloor \frac{\theta}{2} \rfloor$ -choosable. Here θ is the Ore-degree. Finally, we give an upper bound on the online choice number (χ_{OL}) of χ_{OL} -critical graphs. For an n -vertex d -regular χ_{OL} -critical graph with independence number α this bound gives $\chi_{OL} < \frac{n-\alpha}{n}d + 2$. For triangle-free graphs this bound gives $\chi_{OL} \leq \Delta + 1 - \lfloor \frac{1}{4} \lg(\Delta) \rfloor$.

1 Introduction

A k -coloring of a graph G is a function $\pi: V(G) \rightarrow [k]$ such that $\pi(x) \neq \pi(y)$ for each $xy \in E(G)$. The least k for which G has a k -coloring is the *chromatic number* $\chi(G)$ of G . We say that G is k -chromatic when $\chi(G) = k$. A graph is k -critical if $\chi(G) = k$ and $\chi(H) < k$ for every proper subgraph H of G . If G is k -chromatic, then any minimal k -chromatic subgraph of G is k -critical. In this way, many questions about k -chromatic graphs can be reduced to questions about k -critical graphs which have more structure. The study of critical graphs was initiated by Dirac [5] in 1951. It is easy to see that a k -critical graph G must have minimum degree at least $k-1$ and hence $2 \|G\| \geq (k-1) |G|$. The problem of determining the minimum number of edges in a k -critical graph has a long history. First, in 1957, Dirac [6] generalized Brooks' theorem [4] by showing that any k -critical graph G with $k \geq 4$ and $|G| \geq k+2$ must satisfy

$$2 \|G\| \geq (k-1) |G| + k - 3.$$

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In 1963, this bound was improved for large $|G|$ by Gallai [8]. Put

$$g_k(n, c) := \left(k - 1 + \frac{k - 3}{(k - c)(k - 1) + k - 3} \right) n.$$

Gallai showed that every k -critical graph G with $k \geq 4$ and $|G| \geq k + 2$ satisfies $2 \|G\| \geq g_k(|G|, 0)$. In 1997, Krivelevich [16] improved Gallai's bound by replacing $g_k(|G|, 0)$ with $g_k(|G|, 2)$. Then, in 2003, Kostochka and Stiebitz [15] improved this by showing that a k -critical graph with $k \geq 6$ and $|G| \geq k + 2$ must satisfy $2 \|G\| \geq g_k(|G|, (k - 5)\alpha_k)$ where

$$\alpha_k := \frac{1}{2} - \frac{1}{(k - 1)(k - 2)}.$$

In 2012, Kostochka and Yancey [12] drastically improved these bounds by showing that every k -critical graph G with $k \geq 4$ must satisfy

$$\|G\| \geq \left\lceil \frac{(k + 1)(k - 2)|G| - k(k - 3)}{2(k - 1)} \right\rceil.$$

Moreover, they show that this bound is tight for $k = 4$ and $n \geq 6$ as well as for infinitely many values of $|G|$ for any $k \geq 5$. This bound has many interesting coloring applications such as a very short proof of Grötsch's theorem on the 3-colorability of triangle-free planar graphs [11] and short proofs of the results on coloring with respect to Ore degree in [10, 17, 14]. Given the applications to coloring theory, it makes sense to investigate the same problem for more general types of coloring. In this article, we obtain improved lower bounds on the number of edges for both the list coloring and online list coloring problems. To state our results we need some definitions.

List coloring was introduced by Vizing [23] and independently Erdős, Rubin and Taylor [7]. Let G be a graph. A list assignment on G is a function L from $V(G)$ to the subsets of \mathbb{N} . A graph G is L -colorable if there is $\pi: V(G) \rightarrow \mathbb{N}$ such that $\pi(v) \in L(v)$ for each $v \in V(G)$ and $\pi(x) \neq \pi(y)$ for each $xy \in E(G)$. A graph G is L -critical if G is not L -colorable, but every proper subgraph H of G is $L|_{V(H)}$ -colorable. For $f: V(G) \rightarrow \mathbb{N}$, a list assignment L is an f -assignment if $|L(v)| = f(v)$ for each $v \in V(G)$. If $f(v) = k$ for all $v \in V(G)$, then we also call an f -assignment a k -assignment. We say that G is f -choosable if G is L -colorable for every f -assignment L . The best known lower bound on the number of edges in an L -critical graph where L is a $(k - 1)$ -assignment was given by Kostochka and Stiebitz [15] in 2003 and shows that for $k \geq 9$ and $G \neq K_k$ an L -critical graph where L is a $(k - 1)$ -assignment on G , we have $2 \|G\| \geq g_k(|G|, \frac{1}{3}(k - 4)\alpha_k)$. We improve the bound to $2 \|G\| \geq g_k(|G|, (k - 3)\alpha_k)$.

Online list coloring was independently introduced by Zhu [24] and Schauz [21] (Schauz called it *paintability*). Let G be a graph and $f: V(G) \rightarrow \mathbb{N}$. We say that G is *online f -choosable* if $f(v) \geq 1$ for all $v \in V(G)$ and for every $S \subseteq V(G)$ there is an independent set $I \subseteq S$ such that $G - I$ is online f' -choosable where $f'(v) := f(v)$ for $v \in V(G) - S$ and $f'(v) := f(v) - 1$ for $v \in S - I$. We say that G is *online f -critical* if G is not online f -choosable, but every proper subgraph H of G is online $f|_{V(H)}$ -choosable. In 2012, Riasat and Schauz [20] showed that Gallai's bound $2 \|G\| \geq g_k(|G|, 0)$ holds for online f -critical graphs where $f(v) := k - 1$ for all $v \in V(G)$. We improve this for $k \geq 8$ by proving the same bound as we have for list coloring: $2 \|G\| \geq g_k(|G|, (k - 3)\alpha_k)$.

Our main theorem shows that a graph either has many edges or an induced subgraph which has a certain kind of good orientation. To describe these good orientations we need a few definitions. A subdigraph H of a directed multigraph D is called *eulerian* if $d_H^-(v) = d_H^+(v)$ for every $v \in V(H)$. We call H *even* if $\|H\|$ is even and *odd* otherwise. We write $EE(D)$ (resp. $EO(D)$) for the number of even (resp. odd) spanning subdigraphs of D . Note that the edgeless subgraph of D is even and hence we always have $EE(D) > 0$.

Let G be a graph and $f: V(G) \rightarrow \mathbb{N}$. We say that G is *f-Alon-Tarsi* (for brevity, *f-AT*) if G has an orientation D where $f(v) \geq d_D^+(v) + 1$ for all $v \in V(D)$ and $EE(D) \neq EO(D)$. Alon and Tarsi [1] showed that such orientations are very useful for list coloring; they proved the following.

Lemma 1.1. *If a graph G is f -AT for $f: V(G) \rightarrow \mathbb{N}$, then G is f -choosable.*

Schauz [22] extended this result to online f -choosability.

Lemma 1.2. *If a graph G is f -AT for $f: V(G) \rightarrow \mathbb{N}$, then G is online f -choosable.*

For a graph G , we define $d_0: V(G) \rightarrow \mathbb{N}$ by $d_0(v) := d_G(v)$. The d_0 -choosable graphs were first characterized by Borodin [3] and independently by Erdős, Rubin and Taylor [7]. The connected graphs which are not d_0 -choosable are precisely the Gallai trees (connected graphs in which every block is complete or an odd cycle). The generalization to a characterization of d_0 -AT graphs was first given in [9] by Hladký, Král and Schauz.

We prove the following general lemma saying that either a graph has many edges or has an induced f_H -AT subgraph H where f_H basically gives the number of colors we would expect the vertices to have left in their lists after $\delta(G)$ -coloring $G - H$.

Theorem 2.13. *Let G be a graph with $\delta := \delta(G) \geq 5$ and $K_{\delta+1} \not\subseteq G$. If G does not have a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$, then $2\|G\| \geq g_{\delta+1}(|G|, c)$ where $c := (\delta - 2)\alpha_{\delta+1}$ when $\delta \geq 7$ and $c := (\delta - 3)\alpha_{\delta+1}$ when $\delta \in \{5, 6\}$.*

The *Alon-Tarsi number* of a graph $AT(G)$ is the least k such that G is f -AT where $f(v) := k$ for all $v \in V(G)$. We have $\chi(G) \leq \text{ch}(G) \leq \text{ch}_{OL}(G) \leq AT(G) \leq \text{col}(G)$. We say that G is k -AT-critical if $AT(G) = k$ and $AT(H) < k$ for all proper induced subgraphs H of G . From Theorem 2.13 we can conclude the following.

Theorem 3.5. *For $k \geq 6$ and $G \neq K_k$ a k -AT-critical graph, we have $2\|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.*

Similarly, applying Lemma 1.1 gives the following.

Theorem 3.1. *For $k \geq 6$ and $G \neq K_k$ an L -critical graph where L is a $(k - 1)$ -assignment on G , we have $2\|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.*

This improves the bound given by Kostochka and Stiebitz in [15]; for L -critical graphs, they have $2\|G\| \geq g_k(|G|, \frac{1}{3}(k - 4)\alpha_k)$ for $k \geq 9$. Now, applying Lemma 1.2 gives the following.

Theorem 3.3. For $k \geq 6$ and $G \neq K_k$ an online f -critical graph where $f(v) := k - 1$ for all $v \in V(G)$, we have $2 \|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.

Definition 1. The *Ore-degree* of an edge xy in a graph G is $\theta(xy) := d(x) + d(y)$. The *Ore-degree* of a graph G is $\theta(G) := \max_{xy \in E(G)} \theta(xy)$.

A bound like Brooks' theorem in terms of the Ore-degree was given by Kierstead and Kostochka [10] and subsequently the required lower bound on Δ was improved in [17, 14, 18]. For example, we have the following.

Theorem 1.3. Every graph with $\theta \geq 10$ and $\omega \leq \frac{\theta}{2}$ is $\lfloor \frac{\theta}{2} \rfloor$ -colorable.

Another method for achieving the tightest of these results on Ore-degree was given by Kostochka and Yancey [12]. Their proof combined their new lower bound on the number of edges in a color critical graph together with a list coloring lemma derived via the kernel lemma. In Section 4 we improve this latter lemma and, in a similar way, use it in combination with Theorem 3.1 to prove an Ore-degree version of Brooks' theorem for list coloring. The improved lemma can be seen to be giving another lower bound on the number of edges in G . Let $\text{mic}(G)$ be the maximum of $\sum_{v \in I} d_G(v)$ over all independent sets I of G .

Theorem 4.5. For any graph G we have either:

1. G has a nonempty induced subgraph H which is online f_H -choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or
2. $2 \|G\| \geq (\delta(G) - 1) |G| + \text{mic}(G) + 1$.

This quickly gives the aforementioned Ore degree version of Brooks' theorem for list coloring.

Theorem 5.9. Every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is $\lfloor \frac{\theta}{2} \rfloor$ -choosable.

Note that using Kostochka and Stiebitz's above lower bound on the number of edges in a list critical graph gives Theorem 5.9 with $\theta \geq 54$. Similarly, we get the online version.

Theorem 5.8. Every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is online $\lfloor \frac{\theta}{2} \rfloor$ -choosable.

We expect that Theorems 5.8 and 5.9 actually hold for $\theta \geq 10$. In the regular coloring case, it was shown in [14] that the only exception when $\theta \geq 8$ is the graph O_5 ; again, the expectation is that the same result will hold for Theorems 5.8 and 5.9.

Finally, we give a general upper bound on the online choice number that follows immediately from the results in Section 4. The *online choice number* of a graph G is the least k for which G is online f -choosable where $f(v) = k$ for all $v \in V(G)$. We say that G is χ_{OL} -critical if $\chi_{OL}(G - v) < \chi_{OL}(G)$ for all $v \in V(G)$. The following also hold with χ or χ_L in place of χ_{OL} .

Corollary 1.4. For any χ_{OL} -critical graph G , we have $\chi_{OL}(G) < \Delta(G) + 2 - \frac{\text{mic}(G)}{|G|}$.

Since we always have $\text{mic}(G) \geq \frac{2\|G\|}{\chi(G)}$, this gives the following (where $d(G)$ is the average degree of G).

Corollary 1.5. *For any χ_{OL} -critical graph G , we have $\chi_{OL}(G) < \Delta(G) + 2 - \frac{d(G)}{\chi(G)}$.*

The regular case of Corollary 1.4 simplifies somewhat.

Corollary 1.6. *For any regular χ_{OL} -critical graph G , we have $\chi_{OL}(G) < \frac{|G| - \alpha(G)}{|G|} \Delta(G) + 2$.*

Note that Corollary 1.6 reduces proving Brooks' theorem to showing that $\alpha(G) \geq \frac{|G|}{\Delta(G)}$ when $K_{\Delta(G)+1} \not\subseteq G$.

A simple probabilistic argument gives a reasonable bound on $\text{mic}(G)$ for triangle-free graphs and we get the following.

Corollary 4.10. *Triangle-free graphs are online $(\Delta + 1 - \lfloor \frac{1}{4} \lg(\Delta) \rfloor)$ -choosable.*

2 Either many edges or an Alon-Tarsi orientable subgraph

Let \mathcal{T}_k be the Gallai trees with maximum degree at most $k - 1$, excepting K_k . For a graph G , let $W^k(G)$ be the $v \in V(G)$ contained in some K_{k-1} in G .

2.1 Extending Alon-Tarsi orientations

In [15] Kostochka and Stiebitz gave a method for extending list colorings into Gallai trees. We generalize these ideas in terms of extension of orientations.

Lemma 2.1. *Let G be a multigraph without loops and $f: V(G) \rightarrow \mathbb{N}$. If there are $F \subseteq G$ and $Y \subseteq V(G)$ such that:*

1. *any multiple edges in G are contained in $G[Y]$; and*
2. *$f(v) \geq d_G(v)$ for all $v \in V(G) - Y$; and*
3. *$f(v) \geq d_{G[Y]}(v) + d_F(v) + 1$ for all $v \in Y$; and*
4. *For each component T of $G - Y$ there are different $x_1, x_2 \in V(T)$ where $N_T[x_1] = N_T[x_2]$ and $T - \{x_1, x_2\}$ is connected such that either:*
 - (a) *there are $x_1y_1, x_2y_2 \in E(F)$ where $y_1 \neq y_2$ and $N(x_i) \cap Y = \{y_i\}$ for $i \in [2]$; or*
 - (b) *$|N(x_2) \cap Y| = 0$ and there is $x_1y_1 \in E(F)$ where $N(x_1) \cap Y = \{y_1\}$,*

then G is f -AT.

Proof. Suppose not and choose a counterexample (G, f, F, Y) minimizing $|G - Y|$. If $|G - Y| = 0$, then $f(v) \geq d_G(v) + 1$ for all $v \in V(G)$, so any acyclic orientation of G will do. Hence $|G - Y| > 0$.

Pick a component T of $G - Y$ and pick $x_1, x_2 \in V(T)$ as guaranteed by (4). First, suppose (4a) holds. Put $G' := (G - T) + y_1y_2$, $F' := F - T$, $Y' := Y - V(T)$ and let f' be f restricted to $V(G')$. Then G' has an orientation D' where $f'(v) \geq d_{D'}^+(v) + 1$ for all $v \in V(D')$ and $EE(D') \neq EO(D')$, for otherwise (G', f', F', Y') would contradict minimality. By symmetry we may assume that the new edge y_1y_2 is directed toward y_2 . Now we use the orientation

of D' to construct the desired orientation of D . First, we use the orientation on $D' - y_1y_2$ on $G - T$. Now, order the vertices of T as $x_1, x_2, z_1, z_2, \dots$ so that every vertex has at least one neighbor to the right. Orient the edges of T left-to-right in this ordering. Finally, we use y_1x_1 and x_2y_2 and orient all other edges between T and $G - T$ away from T . Plainly, $f(v) \geq d_D^+(v) + 1$ for all $v \in V(D)$. Now, there is a bijection between eulerian subgraphs of D' containing y_1y_2 and eulerian subgraphs of $D - (T - \{x_1, x_2\})$ containing $y_1x_1x_2y_2$, moreover this bijection preserves the parity of the eulerian subgraph. Any other eulerian subgraph A of D must contain y_1x_1 and not x_2y_2 . Since x_1 must have in-degree 1 in A , it must also have out-degree 1 in A . We show that A has a mate A' of opposite parity. Suppose $x_2 \notin A$ and $x_1z_1 \in A$; then we make A' by removing x_1z_1 from A and adding $x_1x_2z_1$. If $x_2 \in A$ and $x_1x_2z_1 \in A$, we make A' by removing $x_1x_2z_1$ and adding x_1z_1 . We conclude that $EE(D) \neq EO(D)$, a contradiction.

Now suppose (4b) holds. Put $G' := G - T$, $F' := F - T$, $Y' := Y - V(T)$ and define f' by $f'(v) = f(v)$ for all $v \in V(G' - y_1)$ and $f'(y_1) = f(y_1) - 1$. Then G' has an orientation D' where $f'(v) \geq d_{D'}^+(v) + 1$ for all $v \in V(D')$ and $EE(D') \neq EO(D')$, for otherwise (G', f', F', Y') would contradict minimality. We orient $G - T$ according to D , orient T as in the previous case, again use y_1x_1 and orient all other edges between T and $G - T$ away from T . Since we decreased $f'(y_1)$ by 1, the extra out edge of y_1 is accounted for and we have $f(v) \geq d_D^+(v) + 1$ for all $v \in V(D)$. Again any additional eulerian subgraph must contain y_1x_1 and since x_2 has no neighbor in $G - T$ we can use x_2 as before to build a mate of opposite parity for any additional eulerian subgraph. Hence $EE(D) \neq EO(D)$ giving our final contradiction. \square

Lemma 2.2. *Let $r \geq 0$, $k \geq r + 4$ and $G \neq K_k$ be a graph with $x \in V(G)$ such that:*

1. $G - x$ is connected; and
2. $d_G(x) \geq r + 2$; and
3. $|N(x) \cap W^k(G - x)| \geq 1$; and
4. $d_G(v) \leq k - 1$ for all $v \in V(G - x)$.

Then G is f -AT where $f(x) = d_G(x) - r$ and $f(v) = d_G(v)$ for all $v \in V(G - x)$.

Proof. Suppose not and choose a counterexample minimizing $|G|$. Let Q be the non-separating vertices in $G - x$. Suppose we have $y \in Q$ such that $G - y$ satisfies all the hypotheses of the theorem. Then minimality of $|G|$ shows that $G - y$ is f' -AT where $f'(v) := f(v) + d_{G-y}(v) - d_G(v)$ for $v \in V(G)$. Create an orientation D of G from the orientation of $G - y$ by directing all edges incident to y into y . These new edges are on no cycle and thus the eulerian subgraph counts did not change. Also, we have increased the out degree of any vertex v by at most $d_G(v) - d_{G-y}(v)$. Hence G is f -AT, a contradiction. Therefore $G - y$ must fail some hypothesis for each $y \in Q$; note that it is only possible for $G - y$ to fail (2) or (3).

We show that $Q \subseteq N(x)$. Suppose otherwise that we have $y \in Q - N(x)$. Since (2) is satisfied for $G - y$, (3) must fail and hence y is contained in a K_{k-1} , call it B , in $G - x$ such that $N(x) \cap B \neq \emptyset$. Pick $z \in N(x) \cap B$. Since $d_G(z) \leq k - 1$ we must have

$N_{G-x}(z) \subseteq B$ and hence $z \in Q$. Since $d_G(y) \leq k-1$ and $y \in Q$ we must have $N_{G-x}(y) \subseteq B$. But then the conditions of Lemma 2.1 are satisfied with $F := G[x, z]$ and $Y := \{x\}$ since $f(x) \geq d_G(x) - r \geq 2 = d_{G[Y]}(x) + d_F(x) + 1$. This is a contradiction and hence we must have $Q \subseteq N(x)$.

Now, by (3), $G - x$ has at least one K_{k-1} , call it B , such that $N(x) \cap V(B) \neq \emptyset$. If $V(G-x) = B$, then $B = Q \subseteq N(x)$ and $G = K_k$, impossible. Hence we may pick $y \in Q - B$. Then $G - y$ satisfies (3) and hence must not satisfy (2). We conclude that $d_G(x) = r+2$ and hence $|Q| \leq r+2$. But $|Q| \geq \Delta(G-x) = k-1$ and hence $k \leq r+3$, a contradiction. \square

We will need to know what happens when we patch two d_0 -choosable graphs together at a vertex. To determine this we first need to understand the structure of d_0 -choosable graphs. The d_0 -choosable graphs were first characterized by Borodin [3] and independently by Erdős, Rubin and Taylor [7]. The generalization to a characterization of d_0 -AT graphs was first given in [9] by Hladký, Král and Schauz. This generalization follows quickly using the following lemma from [7] (this is often referred to as “Rubin’s Block Theorem”).

Lemma 2.3 (Rubin [7]). *A 2-connected graph is either complete, an odd cycle or contains an induced even cycle with at most one chord.*

Lemma 2.4. *For a connected graph G , the following are equivalent:*

1. G is not a Gallai tree,
2. G contains an induced even cycle with at most one chord,
3. G is d_0 -choosable,
4. G is d_0 -AT,
5. G has an orientation D where $d_G(v) \geq d_D^+(v) + 1$ for all $v \in V(D)$, $EE(D) \in \{2, 3\}$ and $EO(D) \in \{0, 1\}$.

Proof. That (1), (2) and (3) are equivalent is the characterization of d_0 -choosable graphs in [3] and [7]. Since (5) implies (4) and (4) implies (3) it will suffice to show that (2) implies (5). The proof we give of (5) is the same as in [9]. Suppose (2) holds and let H be an induced even cycle with at most one chord in G . Orient the even cycle in H clockwise and the (possible) other edge arbitrarily. Collapse H to a single vertex x_H to form H' and take a spanning tree T of H' with root x_H . Orient the remaining edges in G away from the root in this tree to get D . Then every vertex has in degree at least 1 in D and hence $d_G(v) \geq d_D^+(v) + 1$ for all $v \in V(D)$. Also, since the orientation of $D - H$ is acyclic, the only spanning eulerian subgraphs of D are the edgeless graph, the graph with just the edges from the even cycle in H and possibly one other using the chord in H . Hence $EE(D) \in \{2, 3\}$ and $EO(D) \in \{0, 1\}$, thus (5) holds. \square

Lemma 2.5. *If $\{A, B\}$ is a separation of G such that $G[A]$ and $G[B]$ are connected d_0 -AT graphs and $A \cap B = \{x\}$, then G is f -AT where $f(v) = d_G(v)$ for all $v \in V(G) - x$ and $f(x) = d_G(x) - 1$.*

Proof. By Lemma 2.4 we may choose an orientation D_A of A with $d^+(v) < d(v)$ for all $v \in V(D_A)$ and $EE(D_A) \neq EO(D_A)$ and an orientation D_B of B with $d^+(v) < d(v)$ for all $v \in V(D_B)$ and $EE(D_B) \neq EO(D_B)$. Together these give the desired orientation D of G since no cycle has vertices in both $A - x$ and $B - x$ and thus $EE(D) - EO(D) = EE(D_A)EE(D_B) + EO(D_A)EO(D_B) - (EE(D_A)EO(D_B) + EO(D_A)EE(D_B)) = (EE(D_A) - EO(D_A))(EE(D_B) - EO(D_B)) \neq 0$. \square

Lemma 2.2 restricts the interaction of a high vertex and a single low component. Similarly to [15] we'll use the following lemma to restrict a high vertex's interaction with two low components.

Lemma 2.6. *Let $k \geq 4$ and let G be a graph with $x \in V(G)$ such that:*

1. $G - x$ has two components $H_1, H_2 \in \mathcal{T}_k$; and
2. $|N(x) \cap V(H_i)| = 2$ for $i \in [2]$; and
3. $|N(x) \cap W^k(H_i)| \in [2]$ for $i \in [2]$; and

Then G is f -AT where $f(x) = d_G(x) - 1$ and $f(v) = d_G(v)$ for all $v \in V(G - x)$.

Proof. Using Lemma 2.5, we just need to show that $Q_i := G[\{x\} \cup V(H_i)]$ is d_0 -AT for $i \in [2]$; that is show that Q_i is not a Gallai tree. If Q_i is a Gallai tree, then x 's two neighbors in H_i must be in the same block in H_i and this block must be a K_{k-1} , but this creates a diamond since $k \geq 4$, impossible. \square

Combining Lemma 2.2 and Lemma 2.6 gives the following.

Lemma 2.7. *Let $k \geq 5$ and let G be a graph with $x \in V(G)$ such that:*

1. $K_k \not\subseteq G$; and
2. $G - x$ has components $H_1, H_2, \dots, H_t \in \mathcal{T}_k$; and
3. $d_G(v) \leq k - 1$ for all $v \in V(G - x)$; and
4. $|N(x) \cap W^k(H_i)| \geq 1$ for $i \in [t]$; and
5. $d_G(x) \geq t + 2$.

Then G is f -AT where $f(x) = d_G(x) - 1$ and $f(v) = d_G(v)$ for all $v \in V(G - x)$.

Proof. Since $d_G(x) \geq t + 2$, either x has 3 neighbors in some H_i or x has two neighbors in each of H_i, H_j . In either case, let C_1, \dots, C_q be the other components of $G - x$. For each $i \in [q]$, pick $z_i \in N(x) \cap V(C_i)$. Then order the vertices of C_i with z_i first and orient all the edges in C_i to the left with respect to this ordering. Now orient all edges between C_i and $G - C_i$ into C_i . Note that each vertex in C_i has in-degree at least one and no cycle passes through C_i . Hence we can complete the orientation using one of Lemma 2.2 or Lemma 2.6 to get our desired orientation D of G . \square

To deal with more than one high vertex we need to define the following auxiliary bipartite graph. For a graph G , $\{X, Y\}$ a partition of $V(G)$ and $k \geq 4$, let $\mathcal{B}_k(X, Y)$ be the bipartite graph with one part Y and the other part the components of $G[X]$. Put an edge between $y \in Y$ and a component T of $G[X]$ iff $N(y) \cap W^k(T) \neq \emptyset$. This lemma gives the substantive improvement over [15] on the lower bound on the number of edges in a list critical graph.

Lemma 2.8. *Let $k \geq 8$ and let G be a graph with $Y \subseteq V(G)$ such that:*

1. $K_k \not\subseteq G$; and
2. the components of $G - Y$ are in \mathcal{T}_k ; and
3. $d_G(v) \leq k - 1$ for all $v \in V(G - Y)$; and
4. with $\mathcal{B} := \mathcal{B}_k(V(G - Y), Y)$ we have $\delta(\mathcal{B}) \geq 3$.

Then G is f -AT where $f(y) = d_G(y) - 1$ for $y \in Y$ and $f(v) = d_G(v)$ for all $v \in V(G - Y)$.

Proof. Suppose not and choose a counterexample G minimizing $|G|$. Suppose a component of $G - Y$ has an endblock $B \neq K_{k-1}$. Let X_B be the non-separating vertices of B . Then $G' := G - X_B$ still satisfies the hypotheses of the theorem since the degrees in \mathcal{B} are not affected. Hence, by minimality of $|G|$, we have that G' is f' -AT where $f'(v) := f(v) + d_{G'}(v) - d_G(v)$ for all $v \in V(G')$. But we can extend the orientation of G' to G by ordering the vertices of B with the separating vertices before the non-separating vertices, orienting all edges of B to the right and then orienting all edges from X_B to $G - B$ into X_B . In doing so, we don't change the eulerian subgraph counts since no cycle passes through X_B . Moreover, we don't break the out-degree condition since the out-degree of any $v \in V(G)$ is increased by at most $d_{G'}(v) - d_G(v)$. Therefore G is f -AT, a contradiction. Hence every endblock in $G - Y$ is a K_{k-1} . We can do the same thing with a K_{k-1} endblock in which the non-separating vertices have no neighbors in Y . Thus every endblock B has a non-separating vertex x_B with a neighbor $y_B \in Y$.

To each component T of $G - Y$ we associate a set of edges $u(T)$ from $W^k(T)$ to Y . We'll also classify the components as either *type 1*, *type 2a*, *type 2b*, *type 2c* or *type 3*. Call T *saturated* if it has an endblock B_T such that every non-separating $z \in V(B_T)$ has a neighbor in Y .

Suppose T is saturated. If T has one block, we let $u(T)$ be all edges from T to Y . Such a T is of type 2a. Otherwise, T has an endblock B' other than B_T . First, suppose every non-separating $z \in V(B')$ has a neighbor in Y . Then we let $u(T)$ be all edges from non-separating vertices in B_T to Y together with all edges from non-separating vertices in B' to Y . Such a T is of type 3. Otherwise some non-separating $v \in V(B')$ has no edge to Y and we let $u(T)$ be all edges from non-separating vertices in B_T to Y together with $x_{B'}y_{B'}$. Such a T is of type 2b.

Now suppose T is not saturated. If T has at least 3 endblocks B_1, B_2, B_3 , then put $u(T) := \{x_{B_i}y_{B_i} \mid i \in [3]\}$. Such a T is of type 1. If T has only one block, then pick 3 edges from T with different endpoints in Y . Such a T is also of type 1. Now suppose T has exactly two endblocks B_1 and B_2 . If, for some $i \in [2]$, there are two edges incident to non-separating vertices in B_i with different endpoints in Y , then let $u(T)$ be these two edges together with $x_{B_{3-i}}y_{B_{3-i}}$. Such a T is of type 1.

If we haven't defined $u(T)$ yet, then T must contain a block $B = K_{k-1}$ that is not an endblock since $d_{\mathcal{B}}(T) \geq 3$. Moreover, B must contain a non-separating x_B which is adjacent to $y_B \in Y - \{y_{B_1}, y_{B_2}\}$. If some non-separating $v \in V(B)$ has no edge to Y , then put $u(T) := \{x_{B_1}y_{B_1}, x_{B_2}y_{B_2}, x_By_B\}$. In this case, T is of type 1. Otherwise, let $u(T)$ be $\{x_{B_1}y_{B_1}, x_{B_2}y_{B_2}\}$ together with all edges from non-separating vertices of B to Y . Note that as T has exactly 2 endblocks, B has $k - 3$ non-separating vertices. This final kind of T is of type 2c.

We need to handle one more issue before proceeding with the application of Lemma 2.1. When T is type 2a, 2b, 2c or 3, it is possible that $u(T)$ contains edges of the form yx_1 and yx_2 . To see why this could be a problem, consider a component T of type 2a. We'd like to get the degree of T in \mathcal{B} as large as possible, but it could be that for some $y \in Y$ we have yx_1 and yx_2 for different $x_1, x_2 \in V(T)$ and hence T has degree at most $k - 2$ in \mathcal{B} . If this situation happens more than once for T , we need to handle T differently. By Lemma 2.7 each $y \in Y$ has at most 2 neighbors in any given T , so we don't need to worry about triples of edges. If T is of type 2a, 2b, or 2c, we call T *tricky* if there are different $y_1, y_2 \in Y$ and different $x_1, x_2, x_3, x_4 \in V(T)$ such that $y_1x_1, y_1x_2, y_2x_3, y_2x_4 \in u(T)$. If T is of type 3, we call T *tricky* if there are different $y_1, y_2, y_3 \in Y$ and different $x_1, x_2, x_3, x_4, x_5, x_6 \in V(T)$ such that $y_1x_1, y_1x_2, y_2x_3, y_2x_4, y_3x_5, y_3x_6 \in u(T)$.

Now we will create a bipartite graph \mathcal{B}' from \mathcal{B} by splitting vertices corresponding to components T of $G - Y$ as follows. Split each nontricky T into T and T' by divvying up T 's edges in \mathcal{B} so that T is adjacent to all and only the endpoints in Y of the edges in $u(T)$ and T' is adjacent to the rest. We call T an *original* and T' a *copy*. We have the nontricky vertex degrees under control since if T is nontricky type 2a, we have $d_{\mathcal{B}'}(T) \geq k - 1 - 1 = k - 2$. If T is nontricky type 2b, then $d_{\mathcal{B}'}(T) \geq k - 2 - 1 + 1 = k - 2$. If T is nontricky type 2c, then $d_{\mathcal{B}'}(T) \geq k - 3 - 1 + 2 = k - 2$. Finally, if T is nontricky type 3, then $d_{\mathcal{B}'}(T) \geq 2(k - 2) - 2 = 2k - 6$.

We aim to find an orientation of \mathcal{B}' such that each $y \in Y$ has in degree at least 2, each original T of type 1 has in degree at least 1, each original nontricky T of type 2a, 2b or 2c has in degree at least 2 and each original nontricky T of type 3 has in degree at least 3. Note that we are not placing any conditions on the copies and the trickies. Let's see why we want such an orientation. Suppose we have one. We will construct $F \subseteq G$ needed in Lemma 2.1. For each directed edge yT , pick an edge in G from y to T from $u(T)$ and add it to F . We claim that these edges satisfy (4) of Lemma 2.1. If T is type 1 then the selected edge satisfies (4b). If T is of type 2a, then we have satisfied (4a). If T is type 2b or 2c, then we satisfy (4a) if the two edges end in the same block of T and (4b) otherwise. If T is type 3, then two of the three edges must end in the same block and we satisfy (4a). Also, as each $y \in Y$ has in degree at least 2, we have $f(y) = d_G(y) - 1 \geq d_{G[Y]}(y) + d_F(y) + 1$. It remains to satisfy (4) for the tricky components without breaking (3). Suppose T is a tricky component of type 2a, 2b, or 2c and $y_1x_1, y_1x_2, y_2x_3, y_2x_4$ are in $u(T)$. If any of these edges are in F already, remove them. This only makes (3) better and doesn't change anything else. Now pick one of y_1x_1, y_1x_2 and one of y_2x_3, y_2x_4 so that (4) is satisfied (we can by definition of $u(T)$) and add them to F . We haven't broken (3) because y_1, y_2 had in degree at least two, and T is still contributing its one in edge. The tricky components of type 3 are handled similarly by picking one edge from each pair. Now G is f -AT by Lemma 2.1, a contradiction.

Hence we may assume that \mathcal{B}' has no such orientation. Let Q be the original vertices of

type 1 in \mathcal{B}' , P the original nontricky vertices of type 2a, 2b or 2c in \mathcal{B}' and R the original nontricky vertices of type 3 in \mathcal{B}' . Then each $y \in Y$ has degree at least 3 in \mathcal{B}' , each $T \in Q$ has degree at least 3 in \mathcal{B}' , each $T \in P$ has degree at least $k - 2$ in \mathcal{B}' and each $T \in R$ has degree at least $2k - 6$ in \mathcal{B}' . Now we apply Lemma 4.1 with $S := Q \cup P \cup R$. Let $H \trianglelefteq \mathcal{B}'[S]$. For readability, we'll put $A_H := A \cap V(H)$ for $A \subseteq V(\mathcal{B}')$. To get a contradiction, suppose H doesn't meet the condition in Lemma 4.1; that is, suppose

$$\|H\|^+ < 2|Y_H| + |Q_H| + 2|P_H| + 3|R_H|.$$

We also have

$$\|H\|^+ \geq \max\{3|Y_H|, 3|Q_H| + (k - 2)|P_H| + (2k - 6)|R_H|\}.$$

Hence $|Q_H| + 2|P_H| + 3|R_H| > |Y_H|$ and $2|Y_H| > 2|Q_H| + (k - 4)|P_H| + (2k - 9)|R_H|$. Hence $2|Q_H| + (k - 4)|P_H| + (2k - 9)|R_H| < 2|Q_H| + 4|P_H| + 6|R_H|$ giving $(k - 8)|P_H| + (2k - 15)|R_H| < 0$ a contradiction since $k \geq 8$. \square

With a slightly simpler argument we get the following version with asymmetric degree condition on \mathcal{B} . The point here is that this works for $k \geq 6$. As we'll see in the next section, the consequence is that we trade a bit in our size bound for the proof to go through with $k \in \{6, 7\}$.

Lemma 2.9. *Let $k \geq 6$ and let G be a graph with $Y \subseteq V(G)$ such that:*

1. $K_k \not\subseteq G$; and
2. the components of $G - Y$ are in \mathcal{T}_k ; and
3. $d_G(v) \leq k - 1$ for all $v \in V(G - Y)$; and
4. with $\mathcal{B} := \mathcal{B}_k(V(G - Y), Y)$ we have $d_{\mathcal{B}}(y) \geq 4$ for all $y \in Y$ and $d_{\mathcal{B}}(T) \geq 2$ for all components T of $G - Y$.

Then G is f -AT where $f(y) = d_G(y) - 1$ for $y \in Y$ and $f(v) = d_G(v)$ for all $v \in V(G - Y)$.

Proof. Suppose not and choose a counterexample G minimizing $|G|$. In the exact same way as the proof of Lemma 2.8, we use minimality of $|G|$ to conclude that if T is a component of $G - Y$, then every endblock of T is K_{k-1} and every endblock B of T has a vertex x_B among its non-separating vertices X_B that has a neighbor $y_B \in Y$.

To each component T of $G - Y$ we associate a set of edges $u(T)$ from $W^k(T)$ to Y . We'll also classify the components as either *type 1*, *type 2a*, *type 2b* or *type 3*. Call T *saturated* if it has an endblock B_T such that every non-separating $z \in V(B_T)$ has a neighbor in Y .

Suppose T is saturated. If T has one block, we let $u(T)$ be all edges from T to Y . Such a T is of type 2a. Otherwise, T has an endblock B' other than B_T . First, suppose every non-separating $z \in V(B')$ has a neighbor in Y . Then we let $u(T)$ be all edges from non-separating vertices in B_T to Y together with all edges from non-separating vertices in B' to Y . Such a T is of type 3. Otherwise some non-separating $v \in V(B')$ has no edge to Y and we let $u(T)$ be all edges from non-separating vertices in B_T to Y together with $x_{B'}y_{B'}$. Such a T is of type 2b.

Now suppose T is not saturated. If T has at least 2 endblocks B_1, B_2 , then put $u(T) := \{x_{B_i}y_{B_i} \mid i \in [2]\}$. Such a T is of type 1. Otherwise, T has only one block and we pick 2 edges from T with different endpoints in Y . Such a T is also of type 1.

We need to handle one more issue before proceeding with the application of Lemma 2.1. When T is type 2a, 2b or 3, it is possible that $u(T)$ contains edges of the form yx_1 and yx_2 . To see why this could be a problem, consider a component T of type 2a. We'd like to get the degree of T in \mathcal{B} as large as possible, but it could be that for some $y \in Y$ we have yx_1 and yx_2 for different $x_1, x_2 \in V(T)$ and hence T has degree at most $k - 2$ in \mathcal{B} . If this situation happens more than once for T , we need to handle T differently. By Lemma 2.7 each $y \in Y$ has at most 2 neighbors in any given T , so we don't need to worry about triples of edges. If T is of type 2a or 2b, we call T *tricky* if there are different $y_1, y_2 \in Y$ and different $x_1, x_2, x_3, x_4 \in V(T)$ such that $y_1x_1, y_1x_2, y_2x_3, y_2x_4 \in u(T)$. If T is of type 3, we call T *tricky* if there are different $y_1, y_2, y_3 \in Y$ and different $x_1, x_2, x_3, x_4, x_5, x_6 \in V(T)$ such that $y_1x_1, y_1x_2, y_2x_3, y_2x_4, y_3x_5, y_3x_6 \in u(T)$.

Now we will create a bipartite graph \mathcal{B}' from \mathcal{B} by splitting vertices corresponding to components T of $G - Y$ as follows. Split each nontricky T into T and T' by divvying up T 's edges in \mathcal{B} so that T is adjacent to all and only the endpoints in Y of the edges in $u(T)$ and T' is adjacent to the rest. We call T an *original* and T' a *copy*. We have the nontricky vertex degrees under control since if T is nontricky type 2a, we have $d_{\mathcal{B}'}(T) \geq k - 1 - 1 = k - 2$. If T is nontricky type 2b, then $d_{\mathcal{B}'}(T) \geq k - 2 - 1 + 1 = k - 2$. Finally, if T is nontricky type 3, then $d_{\mathcal{B}'}(T) \geq 2(k - 2) - 2 = 2k - 6$.

We aim to find an orientation of \mathcal{B}' such that each $y \in Y$ has in degree at least 2, each original T of type 1 has in degree at least 1, each original nontricky T of type 2a or 2b has in degree at least 2 and each original nontricky T of type 3 has in degree at least 3. Note that we are not placing any conditions on the copies and the trickies. Let's see why we want such an orientation. Suppose we have one. We will construct $F \subseteq G$ needed in Lemma 2.1. For each directed edge yT , pick an edge in G from y to T from $u(T)$ and add it to F . We claim that these edges satisfy (4) of Lemma 2.1. If T is type 1 then the selected edge satisfies (4b). If T is of type 2a, then we have satisfied (4a). If T is type 2b, then we satisfy (4a) if the two edges end in the same block of T and (4b) otherwise. If T is type 3, then two of the three edges must end in the same block and we satisfy (4a). Also, as each $y \in Y$ has in degree at least 2, we have $f(y) = d_G(y) - 1 \geq d_{G[Y]}(y) + d_F(y) + 1$. It remains to satisfy (4) for the tricky components without breaking (3). Suppose T is a tricky component of type 2a or 2b and $y_1x_1, y_1x_2, y_2x_3, y_2x_4$ are in $u(T)$. If any of these edges are in F already, remove them. This only makes (3) better and doesn't change anything else. Now pick one of y_1x_1, y_1x_2 and one of y_2x_3, y_2x_4 so that (4) is satisfied (we can by definition of $u(T)$) and add them to F . We haven't broken (3) because y_1, y_2 had in degree at least two, and T is still contributing its one in edge. The tricky components of type 3 are handled similarly by picking one edge from each pair. Now G is f -AT by Lemma 2.1, a contradiction.

Hence we may assume that \mathcal{B}' has no such orientation. Let Q be the original vertices of type 1 in \mathcal{B}' , P the original nontricky vertices of type 2a or 2b in \mathcal{B}' and R the original nontricky vertices of type 3 in \mathcal{B}' . Then each $y \in Y$ has degree at least 3 in \mathcal{B}' , each $T \in Q$ has degree at least 3 in \mathcal{B}' , each $T \in P$ has degree at least $k - 2$ in \mathcal{B}' and each $T \in R$ has degree at least $2k - 6$ in \mathcal{B}' . Now we apply Lemma 4.1 with $S := Q \cup P \cup R$. Let $H \trianglelefteq \mathcal{B}'[S]$. For readability, we'll put $A_H := A \cap V(H)$ for $A \subseteq V(\mathcal{B}')$. To get a contradiction, suppose

H doesn't meet the condition in Lemma 4.1; that is, suppose

$$\|H\|^+ < 2|Y_H| + |Q_H| + 2|P_H| + 3|R_H|.$$

We also have

$$\|H\|^+ \geq \max\{4|Y_H|, 2|Q_H| + (k-2)|P_H| + (2k-6)|R_H|\}.$$

Hence $|Q_H| + 2|P_H| + 3|R_H| > 2|Y_H|$ and $2|Y_H| > |Q_H| + (k-4)|P_H| + (2k-9)|R_H|$. Hence $|Q_H| + (k-4)|P_H| + (2k-9)|R_H| < |Q_H| + 2|P_H| + 3|R_H|$ giving $(k-6)|P_H| + 2(k-6)|R_H| < 0$ a contradiction since $k \geq 6$. \square

2.2 Many edges

The rest of the proof is basically taken verbatim from [15]. We need the following definitions:

$$\begin{aligned}\mathcal{L}_k(G) &:= G[x \in V(G) \mid d_G(x) = k-1], \\ \mathcal{H}_k(G) &:= G[x \in V(G) \mid d_G(x) \geq k], \\ \sigma_k(G) &:= \left(k-2 + \frac{2}{k-1}\right) |\mathcal{L}_k(G)| - 2\|\mathcal{L}_k(G)\|, \\ \tau_{k,c}(G) &:= 2\|\mathcal{H}_k(G)\| + \left(k-c - \frac{2}{k-1}\right) \sum_{y \in V(\mathcal{H}_k(G))} (d_G(y) - k), \\ \alpha_k &:= \frac{1}{2} - \frac{1}{(k-1)(k-2)}, \\ q_k(G) &:= \alpha_k \sum_{v \in V(G) - W^k(G)} (k-1 - d_G(v)).\end{aligned}$$

As proved in [15], a computation gives the following.

Lemma 2.10. *Let G be a graph with $\delta := \delta(G) \geq 3$ and $0 \leq c \leq \delta + 1 - \frac{2}{\delta}$. If $\sigma_{\delta+1}(G) + \tau_{\delta+1,c}(G) \geq c|\mathcal{H}_{\delta+1}(G)|$, then $2\|G\| \geq g_{\delta+1}(|G|, c)$.*

We need the following degeneracy lemma.

Lemma 2.11. *Let G be a graph and $f: V(G) \rightarrow \mathbb{N}$. If $\|G\| > \sum_{v \in V(G)} f(v)$, then G has an induced subgraph H such that $d_H(v) > f(v)$ for each $v \in V(H)$.*

Proof. Suppose not and choose a counterexample G minimizing $|G|$. Then $|G| \geq 3$ and we have $x \in V(G)$ with $d_G(x) \leq f(x)$. But now $\|G-x\| > \sum_{v \in V(G-x)} f(v)$, contradicting minimality of $|G|$. \square

We'll also need the following consequence of a lemma in [15] giving a lower bound on $\sigma_k(T)$ for $T \in \mathcal{T}_k$.

Lemma 2.12. *Let $k \geq 6$ and $T \in \mathcal{T}_k$. If $K_{k-1} \subseteq T$, then $\sigma_k(T) \geq 2 + q_k(T)$; otherwise $\sigma_k(T) \geq 2 - \alpha_k + q_k(T)$.*

We are now ready to prove the main theorem.

Theorem 2.13. *Let G be a graph with $\delta := \delta(G) \geq 5$ and $K_{\delta+1} \not\subseteq G$. If G does not have a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$, then $2\|G\| \geq g_{\delta+1}(|G|, c)$ where $c := (\delta - 2)\alpha_{\delta+1}$ when $\delta \geq 7$ and $c := (\delta - 3)\alpha_{\delta+1}$ when $\delta \in \{5, 6\}$.*

Proof. Suppose G does not have a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$. Put $k := \delta + 1$, $\mathcal{L} := \mathcal{L}_k(G)$ and $\mathcal{H} := \mathcal{H}_k(G)$. Using Lemma 2.10, we just need to show that $\sigma_k(G) + \tau_{k,c}(G) \geq c|\mathcal{H}|$. Put $W := W^k(\mathcal{L})$, $L' := V(\mathcal{L}) - W$ and $H' := \{v \in V(\mathcal{H}) \mid d_G(v) = k\}$. For $y \in V(\mathcal{H})$, put $\tau_{k,c}(y) := d_{\mathcal{H}}(y) + (k - c + \frac{2}{k-1})(d_G(y) - k)$. We have $\tau_{k,c}(G) = \sum_{y \in V(\mathcal{H})} \tau_{k,c}(y) \geq \sum_{y \in H'} d_{\mathcal{H}}(y) + (k - c - \frac{2}{k-1})|\mathcal{H} - H'| \geq \sum_{y \in H'} d_{\mathcal{H}}(y) + c|\mathcal{H} - H'|$. Thus it will be sufficient to prove that $S := \sigma_k(G) + \sum_{y \in H'} d_{\mathcal{H}}(y) \geq c|H'|$.

Let \mathcal{D} be the components of \mathcal{L} containing K_{k-1} and \mathcal{C} the components of \mathcal{L} not containing K_{k-1} . Then $\mathcal{D} \cup \mathcal{C} \subseteq \mathcal{T}_k$ for otherwise some $T \in \mathcal{D} \cup \mathcal{C}$ is d_0 -AT and hence f_T -AT and (1) holds. We have $\sigma_k(G) = \sum_{T \in \mathcal{D}} \sigma_k(T) + \sum_{T \in \mathcal{C}} \sigma_k(T) \geq 2|\mathcal{D}| + (2 - \alpha_k)|\mathcal{C}| + \alpha_k \sum_{v \in L'} (k - 1 - d_{\mathcal{L}}(v))$.

Now we define an auxiliary bipartite graph F with parts A and B where:

1. $B = H'$ and A is the disjoint union of the following sets A_1, A_2 and A_3 ,
2. $A_1 = \mathcal{D}$ and each $T \in \mathcal{D}$ is adjacent to all $y \in H'$ where $N(y) \cap W^k(T) \neq \emptyset$,
3. For each $v \in L'$, let $A_2(v)$ be a set of $|N(v) \cap H'|$ vertices connected to $N(v) \cap H'$ by a matching in F . Let A_2 be the disjoint union of the $A_2(v)$ for $v \in L'$,
4. For each $y \in H'$, let $A_3(y)$ be a set of $d_{\mathcal{H}}(y)$ vertices which are all joined to y in F . Let A_3 be the disjoint union of the $A_3(y)$ for $y \in H'$.

Case 1. $\delta \geq 7$.

Define $f: V(F) \rightarrow \mathbb{N}$ by $f(v) = 1$ for all $v \in A_2 \cup A_3$ and $f(v) = 2$ for all $v \in B \cup A_1$. First, suppose $\|F\| > \sum_{v \in V(F)} f(v)$. Then by Lemma 2.11, F has an induced subgraph Q such that $d_Q(v) > f(v)$ for each $v \in V(Q)$. In particular, $V(Q) \subseteq B \cup A_1$ and $\delta(Q) \geq 3$. Put $Y := B \cap V(Q)$ and let X be $\bigcup_{T \in V(Q) \cap A_1} V(T)$. Now $H := G[X \cup Y]$ satisfies the hypotheses of Lemma 2.8, a contradiction.

Therefore $\|F\| \leq \sum_{v \in V(F)} f(v) = 2(|H'| + |\mathcal{D}|) + |A_2| + |A_3|$. By Lemma 2.7, for each $y \in B$ we have $d_F(y) \geq k - 1$. Hence $\|F\| \geq (k - 1)|H'|$. This gives $(k - 3)|H'| \leq 2|\mathcal{D}| + |A_2| + |A_3|$. By our above estimate we have $S \geq 2|\mathcal{D}| + \alpha_k \sum_{v \in L'} (k - 1 - d_{\mathcal{L}}(v)) + \sum_{y \in H'} d_{\mathcal{H}}(y) = 2|\mathcal{D}| + \alpha_k |A_2| + |A_3| \geq \alpha_k(2|\mathcal{D}| + |A_2| + |A_3|)$. Hence $S \geq \alpha_k(k - 3)|H'|$. Thus our desired bound holds by Lemma 2.10.

Case 2. $\delta \in \{5, 6\}$.

Define $f: V(F) \rightarrow \mathbb{N}$ by $f(v) = 1$ for all $v \in A_1 \cup A_2 \cup A_3$ and $f(v) = 3$ for all $v \in B$. First, suppose $\|F\| > \sum_{v \in V(F)} f(v)$. Then by Lemma 2.11, F has an induced subgraph Q such that $d_Q(v) > f(v)$ for each $v \in V(Q)$. In particular, $V(Q) \subseteq B \cup A_1$ and $d_Q(v) \geq 4$ for $v \in B \cap V(Q)$ and $d_Q(v) \geq 2$ for $v \in A_1 \cap V(Q)$. Put $Y := B \cap V(Q)$ and

let X be $\bigcup_{T \in V(Q) \cap A_1} V(T)$. Now $H := G[X \cup Y]$ satisfies the hypotheses of Lemma 2.9, a contradiction.

Therefore $\|F\| \leq \sum_{v \in V(F)} f(v) = 3|H'| + |\mathcal{D}| + |A_2| + |A_3|$. By Lemma 2.7, for each $y \in B$ we have $d_F(y) \geq k - 1$. Hence $\|F\| \geq (k - 1)|H'|$. This gives $(k - 4)|H'| \leq |\mathcal{D}| + |A_2| + |A_3|$. By our above estimate we have $S \geq 2|\mathcal{D}| + \alpha_k \sum_{v \in L'} (k - 1 - d_{\mathcal{L}}(v)) + \sum_{y \in H'} d_{\mathcal{H}}(y) = 2|\mathcal{D}| + \alpha_k |A_2| + |A_3| \geq \alpha_k(|\mathcal{D}| + |A_2| + |A_3|)$. Hence $S \geq \alpha_k(k - 4)|H'|$. Thus our desired bound holds by Lemma 2.10. \square

We note a corollary of the above proof that will be useful in our computations in the section on Ore degrees. When $\mathcal{H}_k(G)$ is edgeless, A_3 is empty and $S = \sigma_k(G)$. Also from the proof, we have $\sigma_k(G) \geq 2|\mathcal{D}| + (2 - \alpha_k)|\mathcal{C}| + \alpha_k \sum_{v \in L'} (k - 1 - d_{\mathcal{L}}(v)) \geq \alpha_k(2|\mathcal{D}| + |A_2|) + 2(1 - \alpha_k)|\mathcal{D}| + (2 - \alpha_k)|\mathcal{C}|$. We write $c(G)$ for the number of components of G . Since $(2 - \alpha_k) \geq 2(1 - \alpha_k)$, we have $\sigma_k(G) \geq (k - 3)\alpha_k |\mathcal{H}_k(G)| + 2(1 - \alpha_k)c(\mathcal{L}(G))$.

Corollary 2.14. *Let G be a graph with $\delta := \delta(G) \geq 7$ and $K_{\delta+1} \not\subseteq G$ such that $\mathcal{H}_{\delta+1}(G)$ is edgeless. If G does not have a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$, then $\sigma_{\delta+1}(G) \geq (\delta - 2)\alpha_{\delta+1} |\mathcal{H}_{\delta+1}(G)| + 2(1 - \alpha_{\delta+1})c(\mathcal{L}(G))$.*

3 The bounds for critical graphs

3.1 List critical graphs

Theorem 3.1. *For $k \geq 6$ and $G \neq K_k$ an L -critical graph where L is a $(k - 1)$ -assignment on G , we have $2\|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.*

Proof. Since G is L -critical, we have $\delta(G) \geq k - 1 \geq 5$. If $\delta(G) \geq k$, then $2\|G\| \geq k|G| \geq g_k(|G|, k)$ and we are done. Hence we may assume that $\delta(G) = k - 1$. Since $G \neq K_k$ and G is L -critical, we have $K_{\delta(G)+1} \not\subseteq G$.

Suppose G has a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$. Let π be a coloring of $G - H$ from L and let L' be the list assignment on H defined by $L'(v) := L(v) - \pi(N(v) \cap V(G - H))$ for $v \in V(H)$. Then $|L'(v)| \geq |L(v)| - (d_G(v) - d_H(v)) = k - 1 + d_H(v) - d_G(v)$. By Lemma 1.1, H is f_H -choosable and hence H is L' -colorable. Therefore G is L -colorable, a contradiction.

Now applying Lemma 2.13 proves the theorem. \square

3.2 Online list critical graphs

We will need the following lemma from [21] allowing us to patch together online list colorability of parts into online list colorability of the whole.

Lemma 3.2. *Let G be a graph and $f: V(G) \rightarrow \mathbb{N}$. If H is an induced subgraph of G such that $G - H$ is online $f|_{V(G-H)}$ -choosable and H is online f_H -choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$, then G is online f -choosable.*

Note that applying Lemma 3.2 where H has a single vertex shows that $\delta(G) \geq k - 1$ for an online f -critical graph G where $f(v) := k - 1$ for all $v \in V(G)$.

Theorem 3.3. *For $k \geq 6$ and $G \neq K_k$ an online f -critical graph where $f(v) := k - 1$ for all $v \in V(G)$, we have $2 \|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.*

Proof. Since G is online f -critical, we have $\delta(G) \geq k - 1 \geq 5$. If $\delta(G) \geq k$, then $2 \|G\| \geq k |G| \geq g_k(|G|, k)$ and we are done. Hence we may assume that $\delta(G) = k - 1$. Since $G \neq K_k$ and G is f -critical, we have $K_{\delta(G)+1} \not\subseteq G$.

Suppose G has a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$. Let π be a coloring of $G - H$ from L and let L' be the list assignment on H defined by $L'(v) := L(v) - \pi(N(v) \cap V(G - H))$ for $v \in V(H)$. Then $|L'(v)| \geq |L(v)| - (d_G(v) - d_H(v)) = k - 1 + d_H(v) - d_G(v)$. By Lemma 1.2, H is online f_H -choosable. But then applying Lemma 3.2 shows that G is online f -choosable, a contradiction.

Now applying Lemma 2.13 proves the theorem. \square

3.3 Alon-Tarsi critical graphs

We need a lemma that serves the same purpose as Lemma 3.2 for orientations.

Lemma 3.4. *Let G be a graph and $f: V(G) \rightarrow \mathbb{N}$. If H is an induced subgraph of G such that $G - H$ is $f|_{V(G-H)}$ -AT and H is f_H -AT where $f_H(v) := f(v) + d_H(v) - d_G(v)$, then G is f -AT.*

Proof. Take an orientation of $G - H$ demonstrating that it is $f|_{V(G-H)}$ -AT and an orientation of H demonstrating that it is f_H -AT. Now orient all the edges between H and $G - H$ into $G - H$. Call the resulting oriented graph D . Then D satisfies the out degree requirements of being f -AT since the out degree of the vertices in $G - H$ haven't changed and the out degree of each $v \in V(H)$ has increased by $d_G(v) - d_H(v)$. Since no directed cycle in D has vertices in both H and $D - H$, the eulerian subgraphs of D are just all pairings of eulerian subgraphs of H and $D - H$. Therefore $EE(D) - EO(D) = EE(H)EE(D - H) + EO(H)EO(D - H) - (EE(H)EO(D - H) + EO(H)EE(D - H)) = (EE(H) - EO(H))(EE(D - H) - EO(D - H)) \neq 0$. Hence G is f -AT. \square

Note that applying Lemma 3.4 where H has a single vertex shows that $\delta(G) \geq k - 1$ for a k -AT-critical graph G .

Theorem 3.5. *For $k \geq 6$ and $G \neq K_k$ a k -AT-critical graph, we have $2 \|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.*

Proof. Since G is k -AT-critical, we have $\delta(G) \geq k - 1 \geq 5$. If $\delta(G) \geq k$, then $2 \|G\| \geq k |G| \geq g_k(|G|, k)$ and we are done. Hence we may assume that $\delta(G) = k - 1$. Since $G \neq K_k$ and G is k -AT-critical, we have $K_{\delta(G)+1} \not\subseteq G$.

Suppose G has a nonempty induced subgraph H which is f_H -AT where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$. Since $G - H$ is $(k - 1)$ -AT, applying Lemma 3.4 shows that G is $(k - 1)$ -AT, a contradiction.

Now applying Lemma 2.13 proves the theorem. \square

4 Good orientations from big cuts

Here we strengthen the lemmas on list colorings and orientations of Kostochka and Yancey [12]. The proofs are identical except we replace their specific method of getting an orientation by the following general lemma. This lemma can be derived from Hall's theorem via a similar vertex splitting method to that used by Kostochka and Yancey. It also follows by taking an arbitrary orientation and repeatedly reversing paths if doing so gets a gain (really, this is just the proof of the max-flow min-cut theorem). We give the proof for completeness.

Lemma 4.1. *Let G be a graph, $S \subseteq V(G)$ and $g: S \rightarrow \mathbb{N}$. Then G has an orientation such that $d^-(v) \geq g(v)$ for all $v \in S$ iff for every $H \trianglelefteq G[S]$ we have*

$$\|H\| + \|H, G - H\| \geq \sum_{v \in V(H)} g(v).$$

Proof. First, suppose G has such an orientation and let $H \trianglelefteq G[S]$. Then $\|H\| + \|H, G - H\| = \sum_{v \in V(H)} d_H^-(v) + \|H, G - H\| \geq \sum_{v \in V(H)} d_G^-(v) \geq \sum_{v \in V(H)} g(v)$.

For the other direction, pick an orientation of G minimizing

$$\Theta := \sum_{v \in S} \max \{0, g(v) - d^-(v)\}.$$

Note that the orientation has the desired property iff $\Theta = 0$. Suppose $\Theta > 0$. Then we may choose $x \in S$ with $d^-(x) < g(x)$. Put

$$X := \{v \in V(G) \mid \text{there exists an } xv\text{-directed-path}\}.$$

For any $v \in X$ we must have $v \in S$ and $d^-(v) \leq g(v)$ for otherwise reversing all the edges on an xv -path violates minimality of Θ . By definition, all edges between X and $G - X$ are directed into X and hence with $H := G[X]$ we have

$$\|H\| + \|H, G - H\| = \sum_{v \in V(H)} d^-(v) = d^-(x) + \sum_{v \in V(H) - x} d^-(v) < \sum_{v \in V(H)} g(v).$$

This contradicts our supposition and thus completes the proof. \square

A *kernel* in a digraph D is an independent set $I \subseteq V(D)$ such that each vertex in $V(D) - I$ has an edge into I . A digraph in which every induced subdigraph has a kernel is called *kernel-perfect*. Schauz [21] shows that the well-known Kernel Lemma extends to online choosability.

Kernel Lemma. *If D is a kernel-perfect digraph and $f: V(D) \rightarrow \mathbb{N}$ such that $f(v) \geq d^+(v) + 1$ for each $v \in V(D)$, then G is online f -choosable.*

All bipartite graphs are kernel-perfect, the following lemma from [12] generalizes this fact.

Lemma 4.2. *Let B be an independent set in a graph G and put $A := V(G) - B$. Any digraph created from G by replacing each edge in $G[A]$ by a pair of opposite arcs and orienting the edges between A and B arbitrarily is kernel-perfect.*

Now we have everything needed to prove the strengthening. For independent $A \subseteq V(G)$, we write G_A for the bipartite subgraph $G - E(G - A)$ of G , so just the edges between A and $G - A$ remain.

Lemma 4.3. *Let G be a graph and $f: V(G) \rightarrow \mathbb{N}$ with $f(v) \leq d_G(v) + 1$ for all $v \in V(G)$. If there is independent $A \subseteq V(G)$ such that for each induced subgraph Q of G_A , we have*

$$\|Q\| + \|Q, G_A - Q\| \geq \sum_{v \in V(Q)} d_G(v) + 1 - f(v).$$

then G is online f -choosable.

Proof. Applying Lemma 4.1 on G_A with $g(v) := d_G(v) + 1 - f(v)$ for all $v \in V(G_A)$ gives an orientation of G_A where $d^-(v) \geq d_G(v) + 1 - f(v)$ for each $v \in V(G_A)$. Make an orientation D of G by using this orientation of G_A for the edges between A and $V(G - A)$ and replacing each edge in $G - A$ by a pair of opposite arcs. For $v \in V(D)$ we have $d^+(v) \leq d_{G-A}(v) + d_{G_A}(v) - (d_G(v) + 1 - f(v)) = f(v) - 1$ and hence $f(v) \geq d^+(v) + 1$. By Lemma 4.2, D is kernel-perfect, so the Kernel Lemma shows that G is online f -choosable. \square

Lemma 4.4. *Let G be a nonempty graph and $f: V(G) \rightarrow \mathbb{N}$ with $f(v) \leq d_G(v) + 1$ for all $v \in V(G)$. If there is independent $A \subseteq V(G)$ such that*

$$\|A, G - A\| \geq \sum_{v \in V(G)} d_G(v) + 1 - f(v),$$

then G has a nonempty induced subgraph H that is online f_H -choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$ for $v \in V(H)$.

Proof. Let $A \subseteq V(G)$ be an independent set with

$$\|A, G - A\| \geq \sum_{v \in V(G)} (d_G(v) + 1 - f(v)).$$

Choose a nonempty induced subgraph H of G with $\|H_A\| \geq \sum_{v \in V(H)} (d_H(v) + 1 - f_H(v))$ minimizing $|H|$ (we can make this choice since G is a such a subgraph). Suppose H is not online f_H -choosable. Then, by Lemma 4.3, we have an induced subgraph Q of H_A with $\|Q\| + \|Q, H_A - Q\| < \sum_{v \in V(Q)} (d_H(v) + 1 - f_H(v))$. Now $Q \neq H$ by our assumption on $\|H_A\|$, hence $Z := H - Q$ is a nonempty induced subgraph of G with

$$\begin{aligned} \|Z_A\| &= \|H_A\| - \|Q\| - \|Q, H_A - Q\| \\ &> \sum_{v \in V(H)} (d_H(v) + 1 - f_H(v)) - \sum_{v \in V(Q)} (d_H(v) + 1 - f_H(v)) \\ &= \sum_{v \in V(Z)} (d_Z(v) + 1 - f_Z(v)), \end{aligned}$$

contradicting the minimality of $|H|$. \square

As a special case we get the following lower bound on the number of edges. Recall that $\text{mic}(G)$ is the maximum of $\sum_{v \in I} d_G(v)$ over all independent sets I of G .

Theorem 4.5. *For any graph G we have either:*

1. *G has a nonempty induced subgraph H which is online f_H -choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or*
2. $2 \|G\| \geq (\delta(G) - 1) |G| + \text{mic}(G) + 1$.

We now give a few applications of these ideas.

4.1 Classification of (online) degree-choosable graphs

Using Theorem 4.5, it is easy to prove Brooks' theorem for online list coloring given Brooks' theorem for ordinary coloring; to wit: a minimum counterexample G is $\Delta(G)$ -regular, applying Brooks' theorem gives $\alpha(G) \geq \frac{|G|}{\Delta(G)}$ and hence $\text{mic}(G) \geq |G|$ contradicting Theorem 4.5.

In fact, from Lemma 4.4 we can derive the classification of (online) d_0 -choosable graphs. The following generalizes this classification.

Theorem 4.6. *For a connected graph G , we have $\text{mic}(G) \geq |G| - 1$ with equality only when G is a Gallai tree.*

Proof. Suppose not and choose a counterexample G minimizing $|G|$. Plainly, we must have $|G| \geq 2$. First, suppose $\text{mic}(G) < |G| - 1$. Let v be a noncutvertex in G . By minimality of $|G|$, we have $\text{mic}(G - v) \geq |G| - 2$. Let A be an independent set in $G - v$ with $\|A, G - v - A\| = \text{mic}(G - v)$. Put $A' := A \cup \{v\}$ if $\|v, A\| = 0$ and $A' := A$ otherwise. Then A' is independent in G and $\|A', G - A'\| \geq |G| - 1$, a contradiction.

Now suppose G is not a Gallai tree and $\text{mic}(G) = |G| - 1$. As above we know that for any noncutvertex v we have $\text{mic}(G) \geq \text{mic}(G - v) + 1$; it follows, by minimality of G , that $G - v$ is a Gallai tree. So, if G had more than one block, we could remove a vertex from each of two endblocks to show that every block of G is either complete or an odd cycle contradicting our assumption. Hence G is 2-connected and in particular $\delta(G) \geq 2$. Pick $v \in V(G)$. Suppose $G - v$ is 2-connected. Then $G - v$ is complete or an odd cycle. If $G - v$ is complete then taking v together with a nonneighbor of v shows $\text{mic}(G) \geq |G|$. Suppose $G - v$ is an odd cycle. If $d(v) = 2$, then pick a maximum independent set avoiding $N(v)$ together with v to show $\text{mic}(G) \geq |G|$. Otherwise we may pick maximum independent set containing 2 elements of $N(v)$ again showing $\text{mic}(G) \geq |G|$. Hence $G - v$ has a cutvertex for every $v \in V(G)$. Therefore v is adjacent to every noncutvertex in every endblock of $G - v$ and hence $G - v$ must have exactly two endblocks. Since internal noncutvertices could be removed leaving a 2-connected graph, we conclude that $G - v$ must be a path. Since this was for any v , G must be an even cycle. But then $\text{mic}(G) \geq |G|$, a contradiction.

It remains to show that $\text{mic}(G) = |G| - 1$ when G is a Gallai tree. This can be proved directly, but we won't do so since it follows immediately from the construction of bad d_0 -assignments on Gallai trees and applying Lemma 4.4 as in the proof of the following Corollary. \square

Corollary 4.7. *A connected graph is online d_0 -choosable if it is not a Gallai tree.*

Proof. Let G be a connected graph that is not a Gallai tree. By Theorem 4.6, we have $\text{mic}(G) \geq |G|$. Applying Lemma 4.4 with $f(v) := d_G(v)$ for all $v \in V(G)$ gives a nonempty induced subgraph H of G that is online f_H -choosable where $f_H(v) := f(v) + d_H(v) - d_G(v) = d_H(v)$ for $v \in V(H)$. Hence G is online d_0 -choosable by Lemma 3.2. \square

4.2 Online choosability of triangle-free graphs

We write $\lg(x)$ for the base 2 logarithm of x . We can get a reasonably good lower bound on $\text{mic}(G)$ for triangle-free graphs using a simple probabilistic technique of Shearer and its modification by Alon (see [2]).

Lemma 4.8. *If G is a triangle-free graph, then $\text{mic}(G) \geq \frac{1}{4} \sum_{v \in V(G)} \lg(d(v))$.*

Proof. Let W be a random independent set in G chosen uniformly from all independent sets in G . For each $v \in V(G)$ put $X_v := d(v) |\{v\} \cap W| + |N(v) \cap W|$. We claim that $E(X_v) \geq \frac{1}{2} \lg(d(v))$. This implies the lemma since by linearity of expectation $2 \text{mic}(G) \geq E\left(\sum_{v \in V(G)} X_v\right) \geq \frac{1}{2} \sum_{v \in V(G)} \lg(d(v))$.

To prove the claim, let H be the subgraph of G induced on $V(G) - (N(v) \cup \{v\})$, fix an independent set S in H and let X be the set of all nonneighbors of S in $N(v)$. Put $x := |X|$. It will suffice to bound the conditional expectation for each possible S as follows:

$$E(X_v \mid W \cap V(H) = S) \geq \frac{\lg(d(v))}{2}.$$

For each S , there are exactly $2^x + 1$ possibilities for W and we see that the conditional expectation is exactly $\frac{d(v) + x2^{x-1}}{2^x + 1}$. Suppose this is less than $\frac{\lg(d(v))}{2}$ for some x . Then $2^x \left(\frac{\lg(d(v))}{2} - \frac{x}{2} \right) > d(v) - \frac{\lg(d(v))}{2}$. Put $t := \lg(d(v)) - x$. We have $\frac{td(v)}{2^{t+1}} = \frac{d(v)}{2^t} \left(\frac{\lg(d(v))}{2} - \frac{\lg(d(v)) - t}{2} \right) > d(v) - \frac{\lg(d(v))}{2} > \frac{d(v)}{2}$ and hence $\frac{t}{2^t} > 1$, a contradiction. \square

Theorem 4.9. *Let G be a triangle-free graph and define $f: V(G) \rightarrow \mathbb{N}$ by $f(v) := d_G(v) + 1 - \lfloor \frac{1}{4} \lg(d_G(v)) \rfloor$. Then G has a nonempty induced subgraph H that is online f_H -choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$ for $v \in V(H)$.*

Proof. Immediate upon applying Lemma 4.4 to G since

$$\sum_{v \in V(G)} d_G(v) + 1 - f(v) = \sum_{v \in V(G)} \left\lfloor \frac{1}{4} \lg(d_G(v)) \right\rfloor \leq \text{mic}(G).$$

\square

Corollary 4.10. *If G is a triangle-free graph with $\Delta(G) \leq t$ for some $t \in \mathbb{N}$, then G is online $(t + 1 - \lfloor \frac{1}{4} \lg(t) \rfloor)$ -choosable.*

Proof. Suppose not and choose a counterexample G and $t \in \mathbb{N}$ so as to minimize $|G|$. Put $f(v) := d_G(v) + 1 - \lfloor \frac{1}{4} \lg(d_G(v)) \rfloor$. By Theorem 4.9, G has a nonempty induced subgraph H that is online f_H -choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$ for $v \in V(H)$. Since $t + 1 - \lfloor \frac{1}{4} \lg(t) \rfloor \geq d_G(v) + 1 - \lfloor \frac{1}{4} \lg(d_G(v)) \rfloor$ for all $v \in V(G)$, we have that H is $g(v)$ -choosable where $g(v) := t + 1 - \lfloor \frac{1}{4} \lg(t) \rfloor + d_H(v) - d_G(v)$. Now applying minimality of $|G|$ and Lemma 3.2 gives a contradiction. \square

The best known bounds for the chromatic number of triangle-free graphs are Kostochka's upper bound of $\frac{2}{3}\Delta + 2$ in [13] (see [19] for a proof in English) for small Δ and Johansson's upper bound of $\frac{9\Delta}{\ln(\Delta)}$ for large Δ . Johansson's proof also works for list coloring, but not for online list coloring. To the best of our knowledge Corollary 4.10 is the best known upper bound for online list colorings of triangle-free graphs. Additionally, Corollary 4.10 improves on Johansson's bound for list coloring for $\Delta \leq 8000$. The bound can surely be improved by a more complicated computation of $\text{mic}(G)$, but not beyond around $\Delta + 1 - \lfloor 2 \ln(\Delta) \rfloor$ via this method as can be seen by examples of triangle-free graphs with independence number near $\frac{2 \ln(\Delta)}{\Delta} n$.

4.3 Ore Brooks for maximum degree four

Kostochka and Yancey's bound [12] shows that if G is 4-critical, then $\|G\| \geq \left\lceil \frac{5|G|-2}{3} \right\rceil$. If we try to analyze 4-critical graphs with edgeless high vertex subgraphs by putting this lower bound on the number of edges together with the results on orientations and list coloring obtained in [12], the bounds miss each other. Using the improved bound from Lemma 5.1 we get an exact bound on the number of edges in such a graph.

Lemma 4.11. *For a critical graph G with $\Delta(G) \leq \chi(G) = 4$ such that $\mathcal{H}(G)$ is edgeless we have $\|G\| = \left\lceil \frac{5|G|-2}{3} \right\rceil$ and $|G|$ is not a multiple of 3.*

Proof. Since G is 4-critical, applying Lemma 5.1 gives $2\|G\| < (3 + \frac{1}{3})|G| = \frac{10}{3}|G|$. By Kostochka and Yancey's bound we have $\left\lceil \frac{5|G|-2}{3} \right\rceil \leq \|G\| < \frac{5}{3}|G|$. Hence $\|G\| = \left\lceil \frac{5|G|-2}{3} \right\rceil$ and $|G|$ is not a multiple of 3. \square

It is easy to see that contracting a diamond in a critical graph G with $\Delta(G) \leq \chi(G) = 4$ such that $\mathcal{H}(G)$ is edgeless gives another such graph. The following characterization of these graphs is natural. We have recently learned that Postle has proven this using an extension of the potential method of Kostochka and Yancey.

Conjecture 4.12. *Every critical graph G with $\Delta(G) \leq \chi(G) = 4$ such that $\mathcal{H}(G)$ is edgeless, except K_4 , has an induced diamond. In particular, any such G can be reduced to K_4 by a sequence of diamond contractions.*

5 Ore Brooks for online list coloring

For a graph G , let $\mathcal{H}(G)$ be the subgraph of G induced on the vertices of degree greater than $\delta(G)$ and $\mathcal{L}(G)$ the subgraph of G induced on the vertices of degree $\delta(G)$.

Lemma 5.1. *If G is a graph such that $\mathcal{H}(G)$ is edgeless and $\Delta(G) = \delta(G) + 1$, then either:*

1. *G has a nonempty induced subgraph H which is online f_H -choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or*
2. $2\|G\| < \left(\delta(G) + \frac{1}{\delta(G)} \right) |G|$.

Proof. Put $\delta := \delta(G)$ and suppose $2 \|G\| \geq (\delta + \frac{1}{\delta}) |G|$. Then $|\mathcal{H}(G)| + \delta |G| = 2 \|G\| \geq (\delta + \frac{1}{\delta}) |G|$ and hence $|G| \leq \delta |\mathcal{H}(G)|$. Therefore $\|\mathcal{L}(G), \mathcal{H}(G)\| \geq (\delta + 1) |\mathcal{H}(G)| \geq |\mathcal{H}(G)| + |G|$. Plugging into (2) of Lemma 4.5 gives $2 \|G\| \geq 2 \|G\| + 1$, which is impossible; hence we get the desired induced subgraph H and (1) holds. \square

To break up our computations we reformulate Lemma 5.1 as an upper bound on σ .

Lemma 5.2. *If G is a graph such that $\mathcal{H}(G)$ is edgeless and $\Delta(G) = \delta(G) + 1$, then either:*

1. *G has a nonempty induced subgraph H which is online f_H -choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or*
2. $\sigma_{\delta(G)+1}(G) < \left(4 - \frac{2}{\delta(G)}\right) |\mathcal{H}(G)|$.

Proof. Put $\delta := \delta(G)$. Suppose (1) does not hold. By Lemma 5.1, we have $2 \|G\| < (\delta + \frac{1}{\delta}) |G|$. We have $\delta |G| + |\mathcal{H}(G)| = 2 \|G\| < (\delta + \frac{1}{\delta}) |G|$ giving $|\mathcal{H}(G)| < \frac{|G|}{\delta}$ and hence $|\mathcal{L}(G)| > (\delta - 1) |\mathcal{H}(G)|$. The lemma follows from the following computation:

$$\begin{aligned}
\sigma_{\delta+1}(G) &= 2\delta |\mathcal{L}(G)| - 2 \|\mathcal{L}(G)\| - \left(\delta + 1 - \frac{2}{\delta}\right) |\mathcal{L}(G)| \\
&= 2 \|G\| - \left(\delta + 1 - \frac{2}{\delta}\right) |\mathcal{L}(G)| \\
&< \left(\delta + \frac{1}{\delta}\right) |G| - \left(\delta + 1 - \frac{2}{\delta}\right) |\mathcal{L}(G)| \\
&= -\left(1 - \frac{3}{\delta}\right) |\mathcal{L}(G)| + \left(\delta + \frac{1}{\delta}\right) |\mathcal{H}(G)| \\
&< -\left(1 - \frac{3}{\delta}\right) (\delta - 1) |\mathcal{H}(G)| + \left(\delta + \frac{1}{\delta}\right) |\mathcal{H}(G)| \\
&= \left(4 - \frac{2}{\delta}\right) |\mathcal{H}(G)|.
\end{aligned}$$

\square

By combining Lemma 5.2 with Corollary 2.14 we can prove the Ore version of Brooks' theorem for online list coloring for $\Delta \geq 11$. With a bit more work we will improve this to $\Delta \geq 10$. First, we can actually get a bit more out of Lemma 4.5 by considering independent sets of low vertices that have no high neighbors. Such sets can be added to $V(\mathcal{H}(G))$ to get a cut with more edges. To apply this idea we need the following counting lemma. For a graph G and $t \in \mathbb{N}$, let $\beta_t(G)$ be the size of a largest independent set of degree t vertices in G ; that is, $\beta_t(G) := \alpha(G[x | d(x) = t])$.

Lemma 5.3. *Fix $k \geq 6$. Let G be a Gallai forest with maximum degree at most $k - 1$. We have the following inequality:*

$$(k - 1)\beta_{k-1}(G) + \sum_{v \in V(G)} k - 1 - d(v) \geq \frac{2(k - 3)}{k - 2} |G| - \frac{(k - 1)(k - 4)}{k - 2} c(G).$$

Proof. It will suffice to prove that for any Gallai tree with maximum degree at most $k - 1$ we have:

$$(k - 1)\beta_{k-1}(T) + \sum_{v \in V(T)} k - 1 - d(v) \geq \frac{2(k - 3)}{k - 2} |T| - \frac{(k - 1)(k - 4)}{k - 2}.$$

Suppose not and choose a counterexample T minimizing $|T|$. First, if T has only one block it is easy to see that the inequality is satisfied. Let B be an endblock of T and say x is the cutvertex in B . Suppose $\chi(B) \leq k - 3$. Put $T' := T - (B - x)$. By minimality of $|T|$, T' satisfies the inequality. When we add $B - x$ back in, the left side increases by $(k - \chi(B))(|B| - 1) - (|B| - 1) \geq 2(|B| - 1)$. But the right side increases by only $\frac{2(k-3)}{k-2}(|B| - 1)$ and hence T is not a counterexample, a contradiction.

Hence B is either K_{k-2} or K_{k-1} . Consider $T' := T - B$. Suppose $d_T(x) = k - 1$. Note that none of x 's neighbors in T' have degree $k - 1$ in T' and thus are in no maximum independent set of degree $k - 1$ vertices in T' . Therefore, we can add x to any such independent set, giving $\beta_{k-1}(T) > \beta_{k-1}(T')$. Hence, after applying minimality to T' , we see that adding back B increases the left side by $k - 1 + (k - 2) - 1$ if B is K_{k-1} and by $k - 1 + 2(k - 3) - 2$ if B is K_{k-2} . Since the right side increases by only $\frac{2(k-3)}{k-2}|B|$ in both cases, T satisfies the inequality, a contradiction.

Therefore, it must be that B is K_{k-2} and $d_T(x) = k - 2$. Now when we add B back, the left side increases by $2(k - 3) + 1 - 1$ and the right side increases by only $2(k - 3)$ and again T satisfies the inequality, a contradiction. \square

Lemma 5.4. *If G is a graph such that $\mathcal{H}(G)$ is edgeless and $\Delta(G) = \delta(G) + 1 \geq 7$, then either:*

1. G has a nonempty induced subgraph H which is online f_H -choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or
2. $|\mathcal{H}(G)| < \frac{\delta(G)(\delta(G)-3)}{(\delta(G)-1)(\delta(G)-5)} c(\mathcal{L}(G))$.

Proof. Put $\delta := \delta(G)$ and suppose (1) does not hold. By Lemma 4.5 we have $\text{mic}(G) < |\mathcal{L}(G)| + 2|\mathcal{H}(G)| < |\mathcal{L}(G)| + \frac{2}{\delta-1}|\mathcal{L}(G)| = \frac{\delta+1}{\delta-1}|\mathcal{L}(G)|$. But applying Lemma 5.3 to $\mathcal{L}(G)$ gives $\text{mic}(G) \geq 2\frac{\delta-2}{\delta-1}|\mathcal{L}(G)| - \frac{\delta(\delta-3)}{\delta-1}c(\mathcal{L}(G))$. Combining these inequalities and $|\mathcal{L}(G)| > (\delta - 1)|\mathcal{H}(G)|$ proves the lemma. \square

Lemma 5.5. *Every graph G with $\delta(G) + 1 = \Delta(G) \geq 10$ such that $\mathcal{H}(G)$ is edgeless either contains $K_{\Delta(G)}$ or has a nonempty induced subgraph H that is online f_H -choosable where $f_H(v) := \Delta(G) - 1 + d_H(v) - d_G(v)$ for all $v \in V(H)$.*

Proof. Suppose not and let G be a counterexample. Put $\delta := \delta(G)$. By Corollary 2.14 we have $\sigma_{\delta+1}(G) \geq (\delta - 2)\alpha_{\delta+1}|\mathcal{H}(G)| + 2(1 - \alpha_{\delta+1})c(\mathcal{L}(G))$. By Lemma 5.4 we have $c(\mathcal{L}(G)) > \frac{(\delta-1)(\delta-5)}{\delta(\delta-3)}|\mathcal{H}(G)|$ and hence. Also, by Lemma 5.2, we have $\sigma_{\delta+1}(G) < (4 - \frac{2}{\delta})|\mathcal{H}(G)|$. Putting these together, we get $4 - \frac{2}{\delta} > \alpha_{\delta+1}(\delta - 2) + 2(1 - \alpha_{\delta+1})\frac{(\delta-1)(\delta-5)}{\delta(\delta-3)}$. But then $\delta \leq 8$, a contradiction. \square

Theorem 5.6. *If G is a graph with $\Delta(G) \geq 10$ not containing $K_{\Delta(G)}$ such that $\mathcal{H}(G)$ is edgeless, then G is online $(\Delta(G) - 1)$ -choosable.*

Proof. Suppose not and choose a counterexample G minimizing $|G|$. Then G is online f -critical where $f(v) := \Delta(G) - 1$ for all $v \in V(G)$. Hence $\delta(G) \geq \Delta(G) - 1$ and we may apply Lemma 5.5 to get a nonempty induced subgraph H of G that is online f_H -choosable where $f_H(v) := \Delta(G) - 1 + d_H(v) - d_G(v)$ for all $v \in V(H)$. But then applying Lemma 3.2 shows that G is $(\Delta(G) - 1)$ -choosable, a contradiction. \square

Combining Lemma 5.6 with the following version of Brooks' theorem for online list coloring (first proved in [9]) we get Theorem 5.8.

Lemma 5.7. *Every graph with $\Delta \geq 3$ not containing $K_{\Delta+1}$ is online Δ -choosable.*

Theorem 5.8. *Every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is online $\lfloor \frac{\theta}{2} \rfloor$ -choosable.*

Proof. Suppose not and choose a counterexample G minimizing $|G|$. Put $k := \lfloor \frac{\theta(G)}{2} \rfloor$. Then G is online f -critical where $f(v) := k$ for all $v \in V(G)$. Hence $\delta(G) \geq k$ and thus $\Delta(G) \leq k + 1$. If $\Delta(G) = k$, then the theorem follows from Lemma 5.7. Hence we must have $\Delta(G) = k + 1$. Therefore $\mathcal{H}(G)$ is edgeless, $\Delta(G) \geq 10$ and $\omega(G) \leq \Delta(G) - 1$. Applying Theorem 5.6 shows that G is online $(\Delta(G) - 1)$ -choosable, a contradiction. \square

The same result for list coloring is an immediate consequence.

Theorem 5.9. *Every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is $\lfloor \frac{\theta}{2} \rfloor$ -choosable.*

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