# Design and Implementation of a Strong Representation System for Network Policies

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#### Motivation

- Managing networking policies remains hard:
  - One has to fully understand the policy
  - Understanding SDN is difficult for someone not involved in the coding process

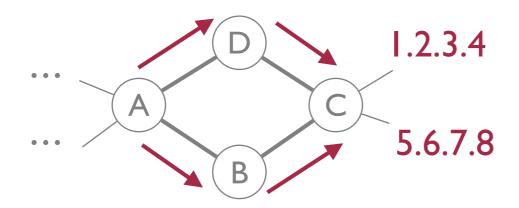
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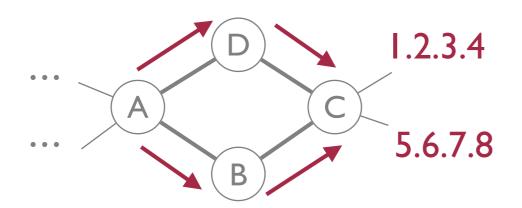
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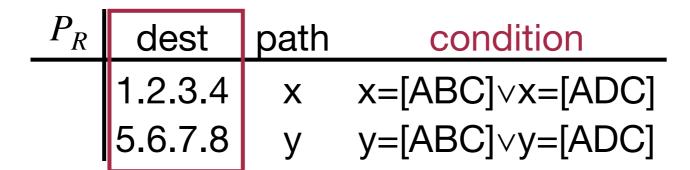
Can we provide a network policy management experience comparable to that on a database?



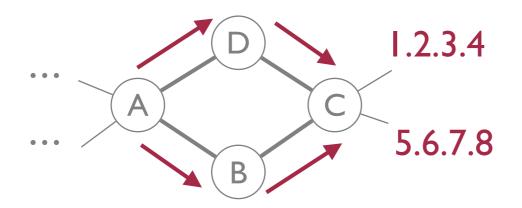
Two prefixes 1.2.3.4 and 5.6.7.8 over two alternative paths [ABC] and [ADC]



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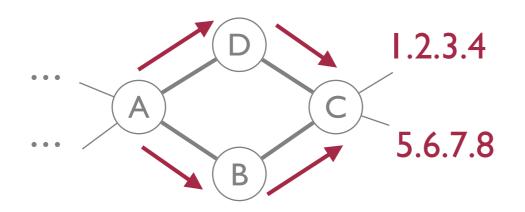
The destination of the path



Two prefixes 1.2.3.4 and 5.6.7.8 over two alternative paths [ABC] and [ADC]

$P_R$	dest	path	condition
	1.2.3.4	Х	x=[ABC]\vx=[ADC] y=[ABC]\vy=[ADC]
	5.6.7.8	У	y=[ABC]vy=[ADC]

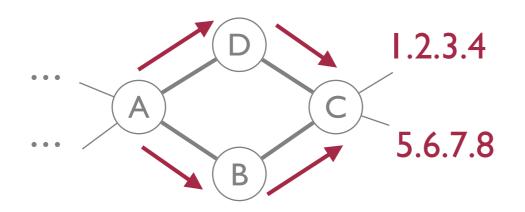
The path to its destination



Two prefixes 1.2.3.4 and 5.6.7.8 over two alternative paths [ABC] and [ADC]

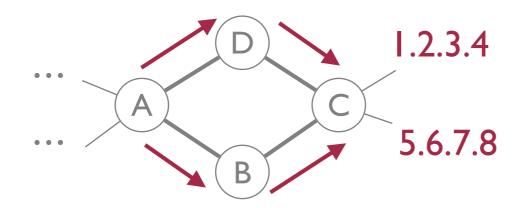
$P_R$	dest	path	condition
	1.2.3.4	X	x=[ABC]vx=[ADC] y=[ABC]vy=[ADC]
	5.6.7.8	У	y=[ABC]vy=[ADC]

The constraints over variables



Two prefixes 1.2.3.4 and 5.6.7.8 over two alternative paths [ABC] and [ADC]

	$P_R$	dest		path		cond	lition	
,		1.2.3.	4	X	X=	[ABC]	/x=[A[	C]
		5.6.7.	8	У	y=	[ABC]	y=[A[	OC]
		I		dest		path	_	
				1.2.3.4	4	[ABC]		
				1.2.3.4	4	[ADC]		
Repre	sent		,	5.6.7.8	3	[ABC]		
				5.6.7.8	3	[ADC]		



Two prefixes 1.2.3.4 and 5.6.7.8 over two alternative paths [ABC] and [ADC]

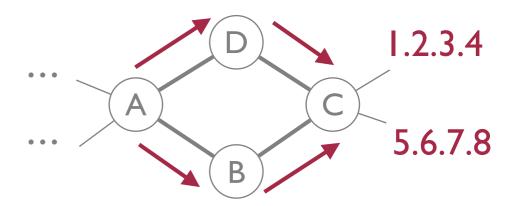
$P_R$	dest	path	condition
	1.2.3.4	X	$x=[ABC] \lor x=[ADC]$
	5.6.7.8	У	y=[ABC]\vy=[ADC]

**Conditional Table** 

Traffic balance policy

Balance traffic to 1.2.3.4 and 5.6.7.8

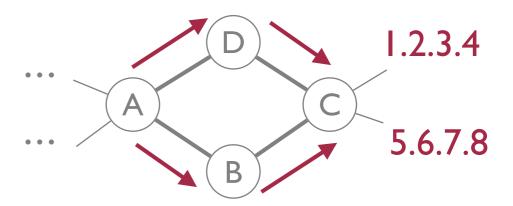
$P_3$		path		
	1.2.3.4 5.6.7.8 1.2.3.4 5.6.7.8	[ABC]	u	u = 1
	5.6.7.8	[ABC]	u	u ≠ 1
	1.2.3.4	[ADC]	V	v = 1
	5.6.7.8	[ADC]	V	v ≠ 1



Traffic balance policy

Balance traffic to 1.2.3.4 and 5.6.7.8

$P_3$	dest	path	flag	
	1.2.3.4	[ABC]	u	u = 1
	5.6.7.8	[ABC]	u	u ≠ 1
	1.2.3.4	[ADC]	V	v = 1
	5.6.7.8	[ADC]	V	v ≠ 1

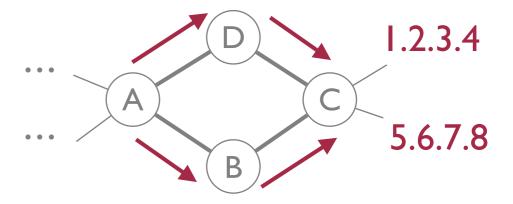


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$P_3$	dest	path	flag	
	1.2.3.4	[ABC]	u	u = 1
	1.2.3.4 5.6.7.8	[ABC]	u	u ≠ 1
	1.2.3.4	[ADC]	V	v = 1
	5.6.7.8	[ADC]	V	V ≠ 1

Cannot assign path [ABC] to 1.2.3.4 and 5.6.7.8 simultaneously.



Traffic balance policy

 $P_3$ 

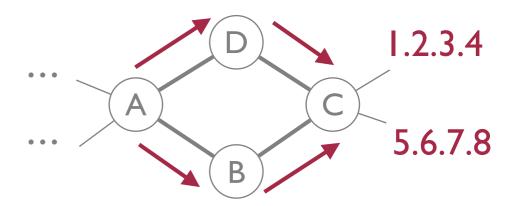
Balance traffic to 1.2.3.4 and 5.6.7.8

dest path flag

1.2.3.4 [ABC] u u = 15.6.7.8 [ABC]  $u u \neq 1$ 1.2.3.4 [ADC] v v = 15.6.7.8 [ADC]  $v \neq 1$ 

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Static route and filter policy

$P_1$	dest	path	
	1.2.3.4	Χ	x=[ABC]
	у	Z	y≠1.2.3.5

Static route and filter policy

Assign a static route [ABC] to destination 1.2.3.4

$$\begin{array}{c|ccccc} P_1 & dest & path \\ \hline & 1.2.3.4 & x & x=[ABC] \\ & y & z & y \neq 1.2.3.5 \end{array}$$

Static route and filter policy

Assign a static route [ABC] to destination 1.2.3.4

Do not allow traffic to destination 1.2.3.5

# Manipulating Network Policies

- Manipulate relational database table:
  - relational algebra: selection, projection, union, join, ...
  - implement SQL query on regular table

# Manipulating Network Policies

- Manipulate relational database table:
  - relational algebra: selection, projection, union, join, ...
  - implement SQL query on regular table
- Our contribution:

Manipulating network policies by simply implementing SQL query on conditional table

- Static route and filter policy
- Traffic balance policy

$P_1$	dest	path		$P_3$		path		
	1.2.3.4	X	L d		1.2.3.4 5.6.7.8 1.2.3.4 5.6.7.8	[ABC]	u	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2.3.4		5.6.7.8	[ABC]	u	u ≠ 1
					1.2.3.4	[ADC]	V	v = 1
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$P_1$	dest	path			$P_3$	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	u	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2	.3.4		5.6.7.8	[ABC]	u	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]		v = 1
satis	sfies $\overline{P}_1$ and	d $P_3$ sin	multaneously?	o <sub>1</sub> join	$P_3$	5.6.7.8	[ADC]	V	v ≠ 1

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$P_1$	dest	path		P	3	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
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Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	V = 1
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$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	X	u	$x=[ABC] \land u=1$
	У	Z	u	y=5.6.7.8∧z=[ABC]∧u≠1
	у	Z	V	$y=5.6.7.8 \land z=[ADC] \land v \neq 1$

- Static route and filter policy
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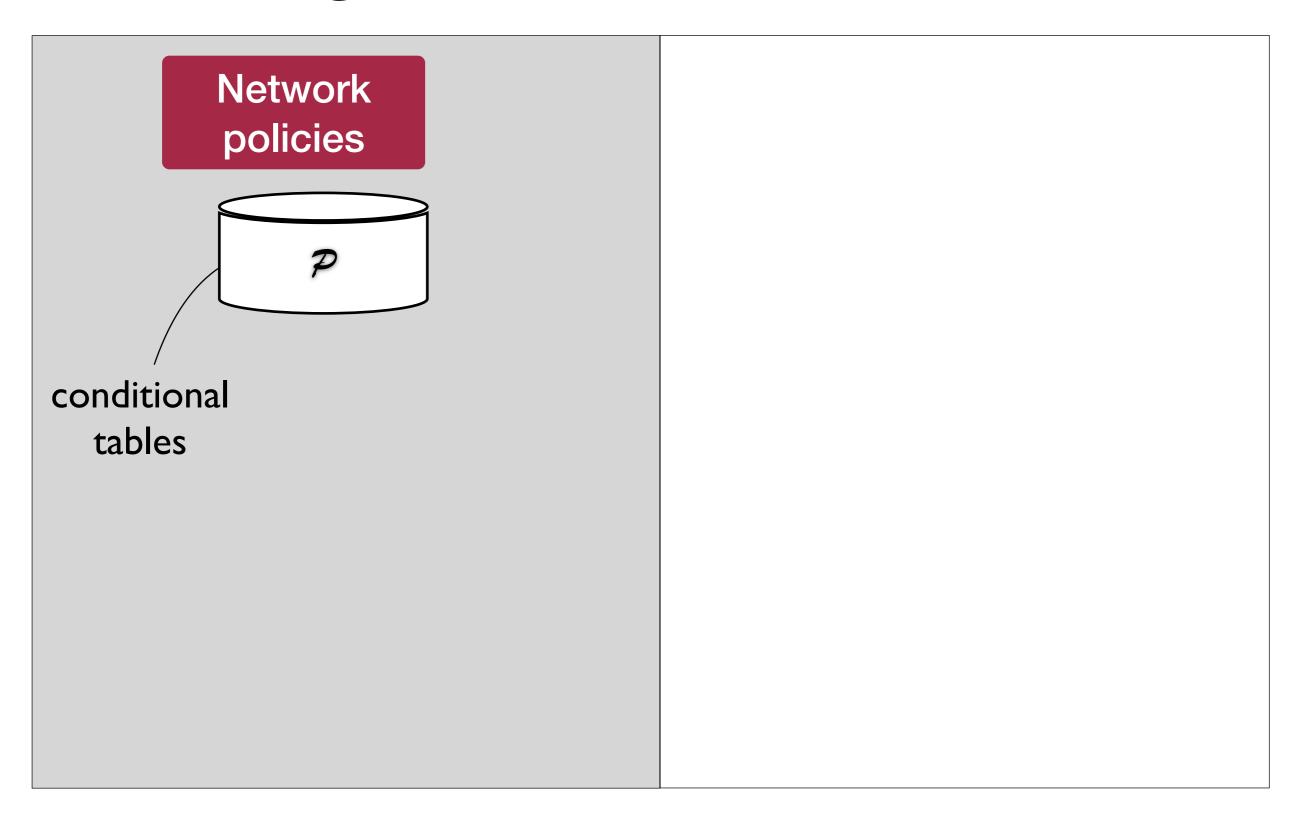
$P_1$	dest	path		P	3	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2.	3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	V = 1
satisfies $P_1$ and $P_3$ simultaneously? $P_1$ join $P_2$						5.6.7.8	[ADC]	V	V ≠ 1

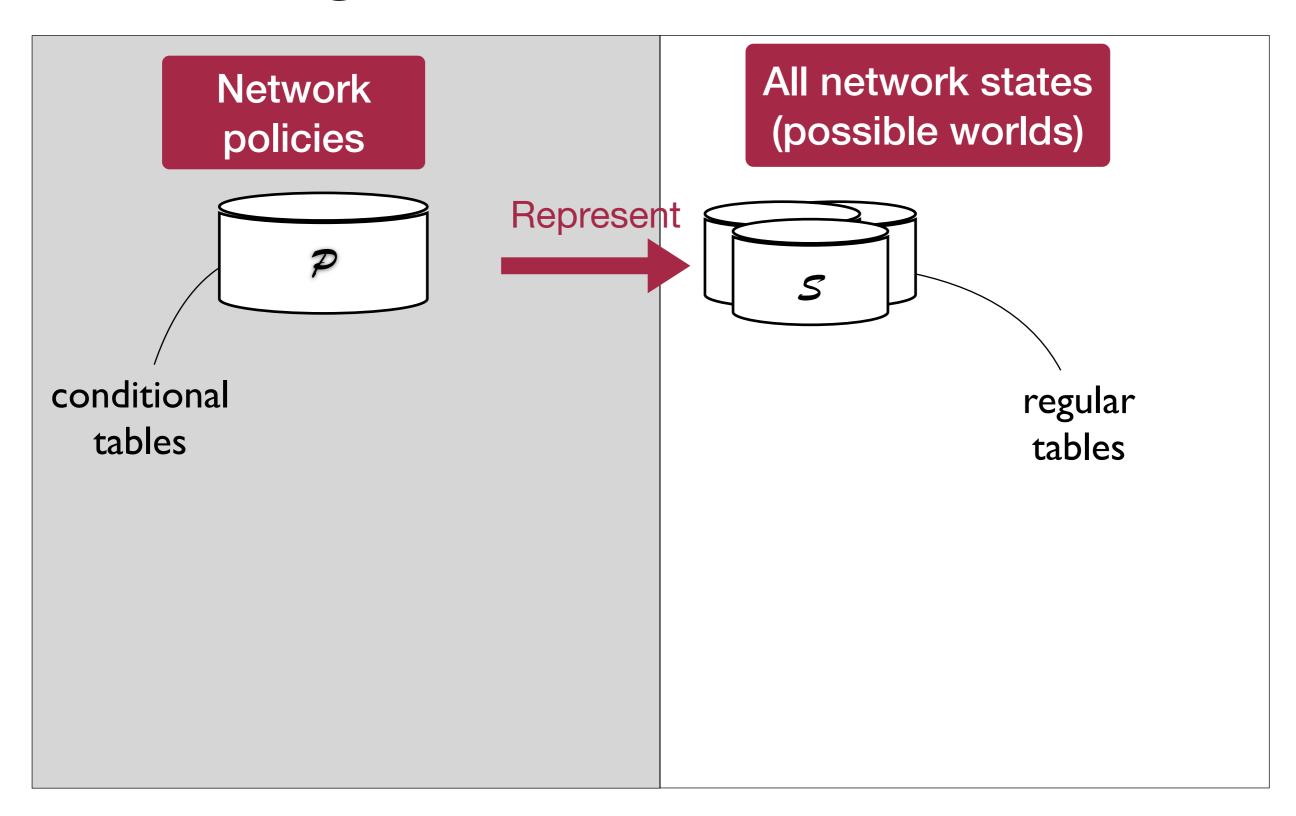
$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4			$x=[ABC] \land u=1$
	У	Z	U	$y=5.6.7.8 \land z=[ABC] \land u \neq 1$
	у	Z	V	$y=5.6.7.8 \land z=[ADC] \land v \neq 1$
		I	dest	path
Valid forwarding			1.2.3.4	4 [ABC]
state			5.6.7.8	B [ADC] <sub>7</sub>

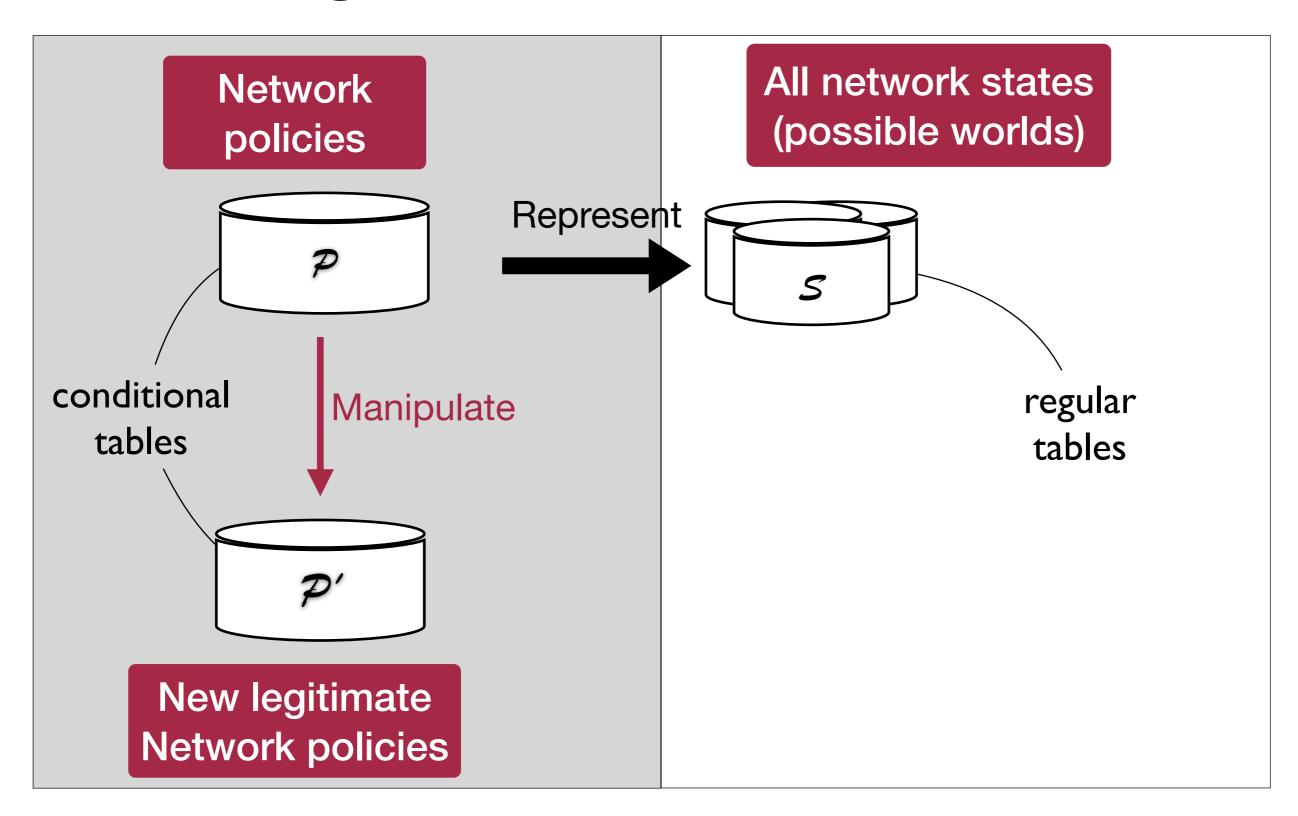
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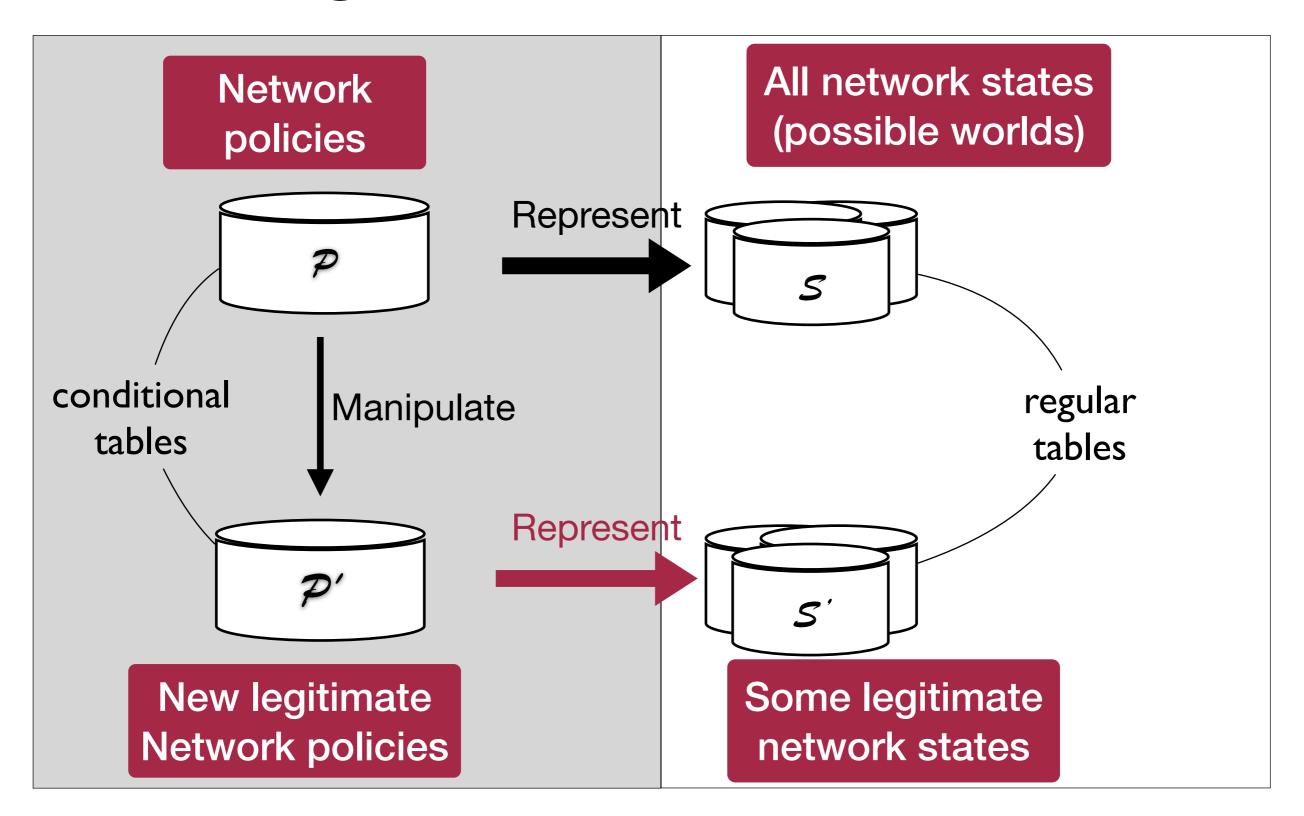
$P_1$	dest	path		$P_3$		dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	u	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2.3.4	1		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	v = 1
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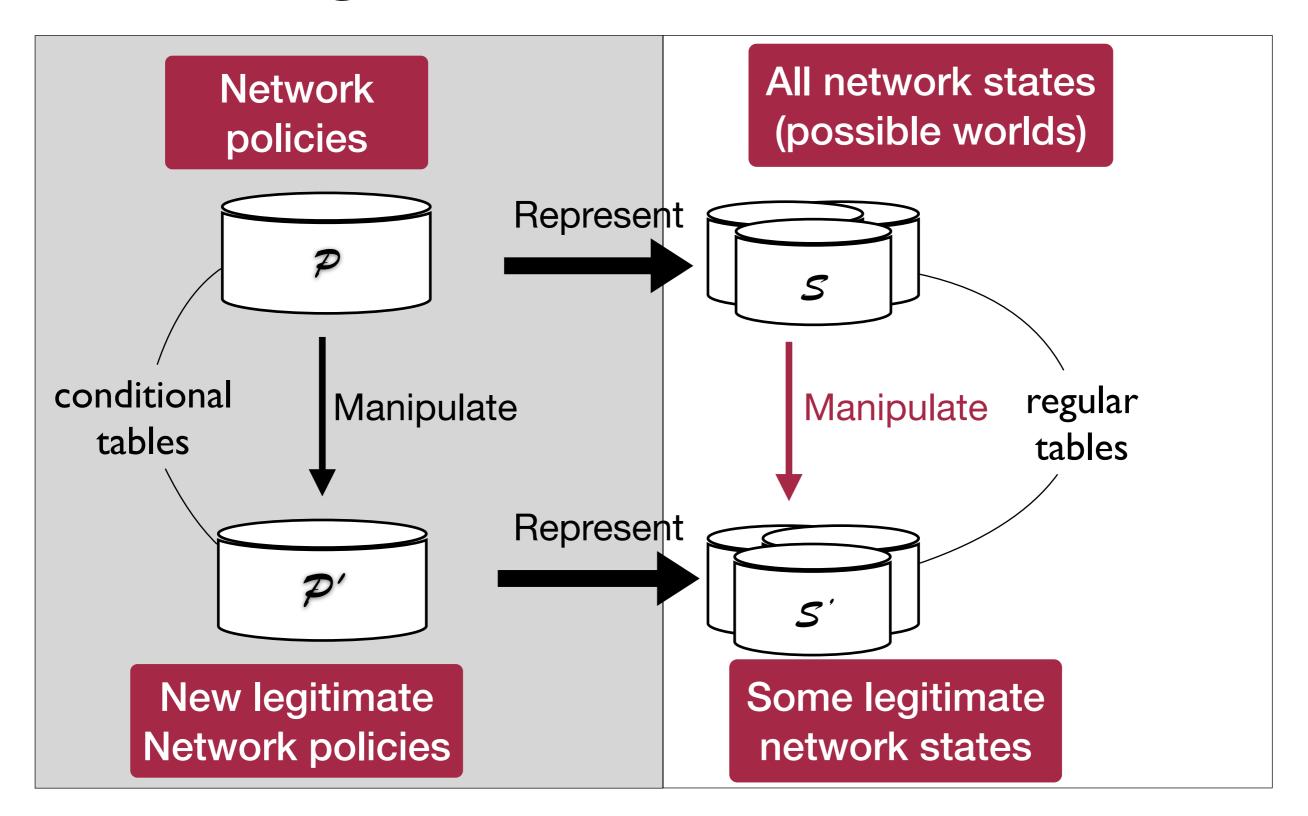
$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	X	u	$x=[ABC] \land u=1$
	•			y=5.6.7.8∧z=[ABC]∧u≠1
	У	Z	V	$y=5.6.7.8 \land z=[ADC] \land v \neq 1$
		I	dest	t path_
Vali	Valid forwarding			.8 [ABC]
	state			7

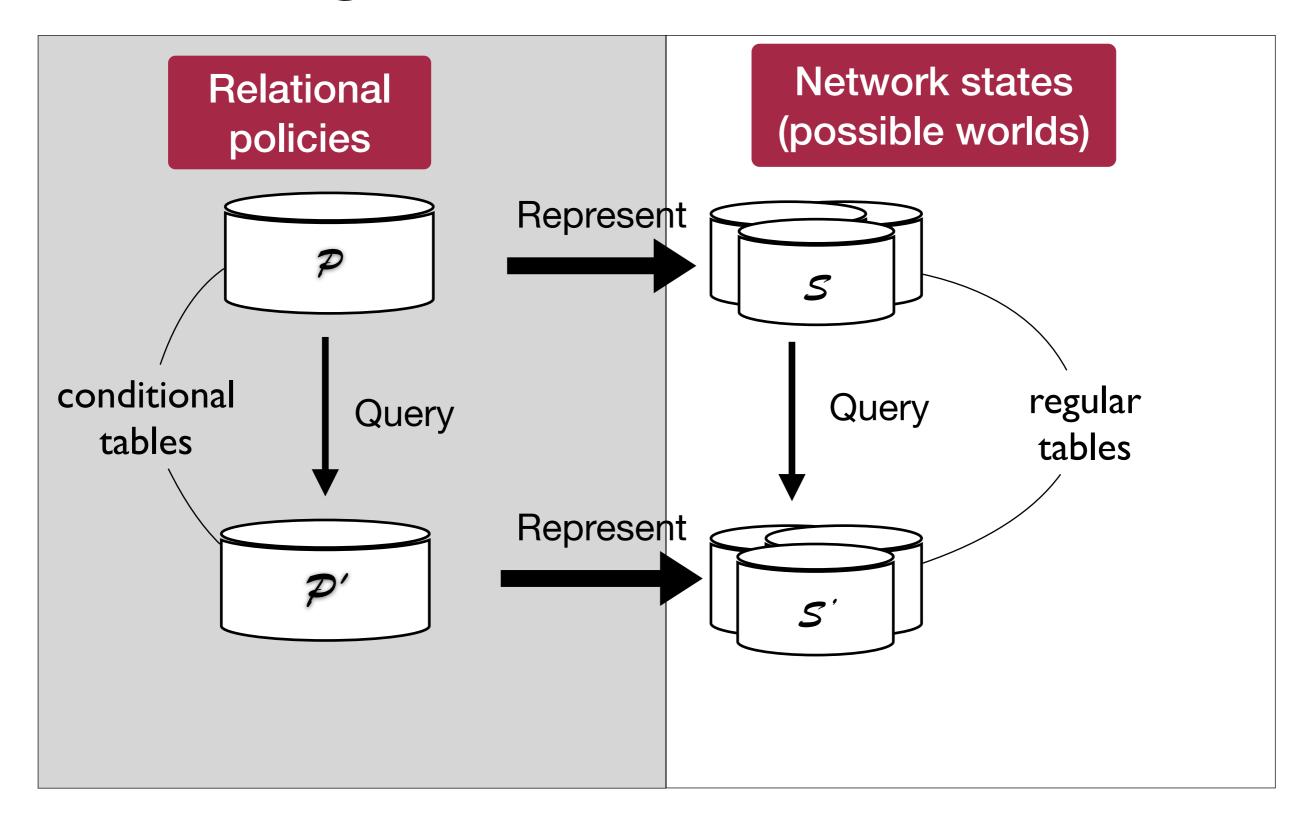












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Using user-defined functions and aggregates

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#### Using user-defined functions and aggregates

$$P_2$$
 dest path s(path)  $x$   $y$   $z$   $I(y) \le z$ 

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 dest path s(path)  $x$  y  $z$   $x$   $y \neq z$ 

Function s: return the lowest hop counts among all alternative paths

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#### Using user-defined functions and aggregates

$$P_2$$
 dest path s(path)  $x$   $y$   $z$   $I(y) \le z$ 

Function s: return the lowest hop counts among all alternative paths

Function I: return hops of the path

# **Support Network Addressing**

- Network addressing is a basic but critical feature to network area
- Allow sets in conditional table
  - i.e. accommodate variables and conditions

# **Support Network Addressing**

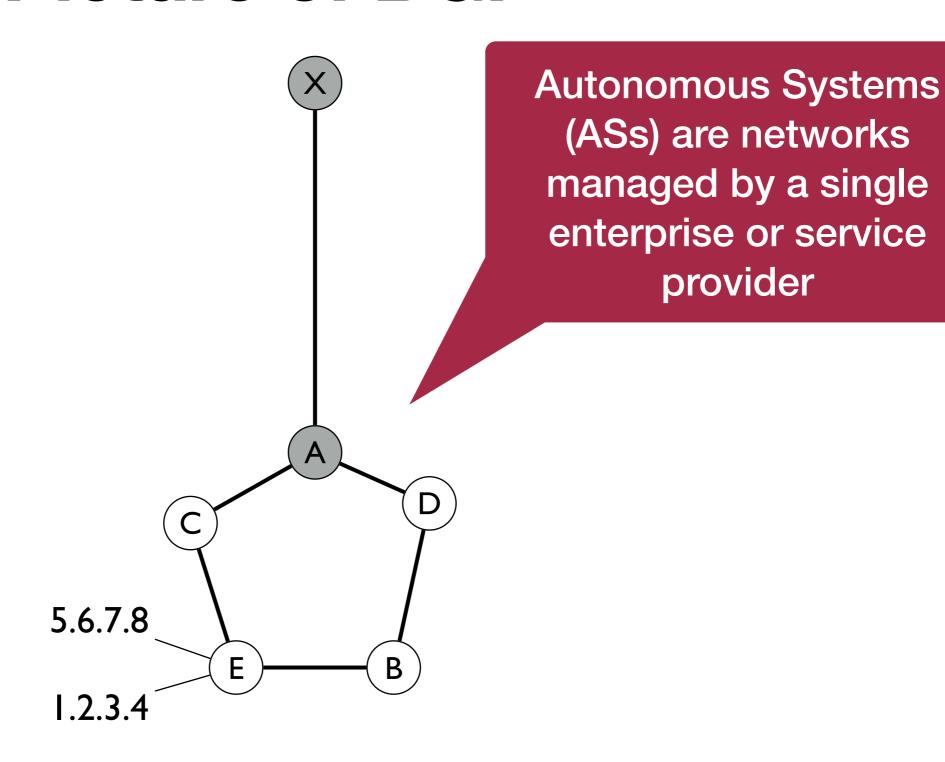
- Network addressing is a basic but critical feature to network area
- Allow sets in conditional table
  - i.e. accommodate variables and conditions
- Two methods:
  - Naive support for sets
  - Leverage SMT(Satisfiability Modulo Theories) solver for sets

#### Application in inter-domain routing

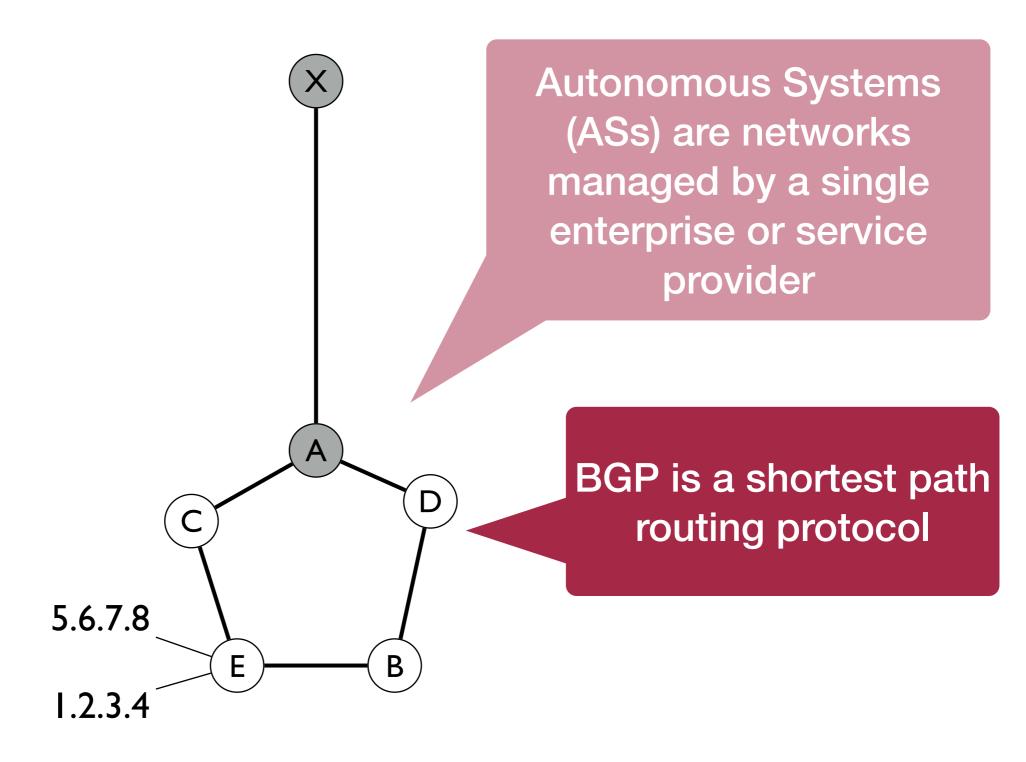
- We show the possibility of coordinating knowledge-driven policies from a single central point
- What about in realistic network scenarios like inter-domain routing?

We show an application - **Policy Exchange** 

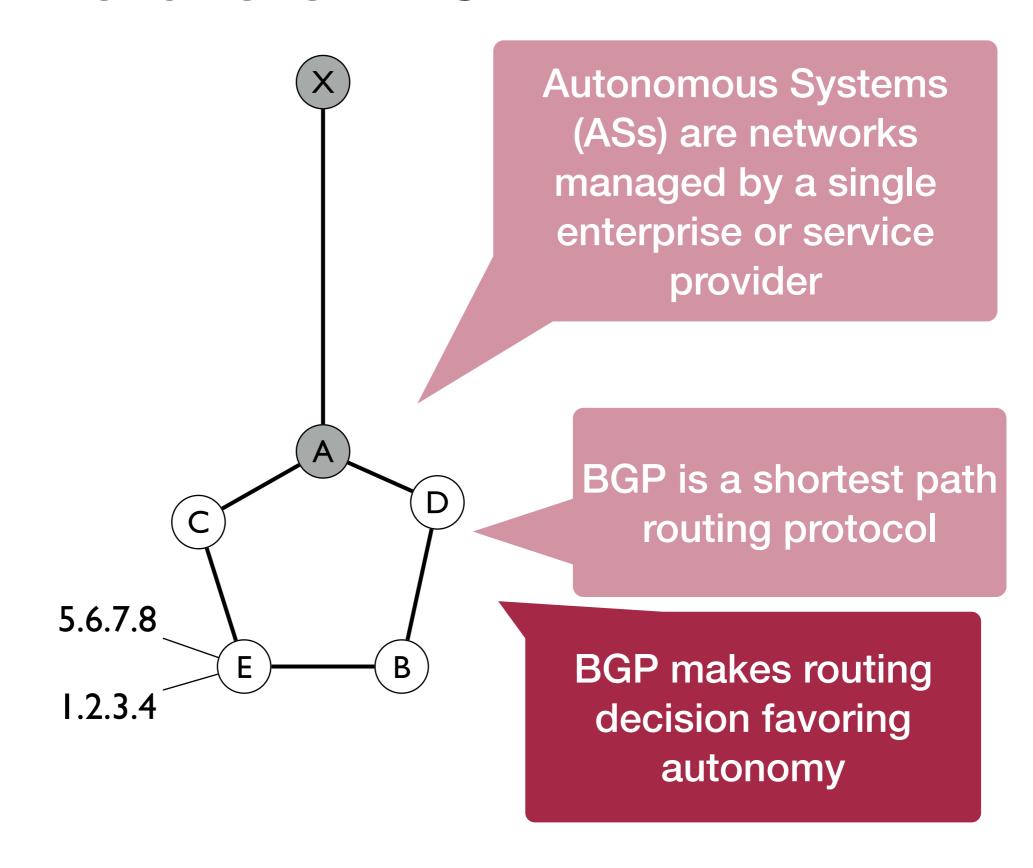
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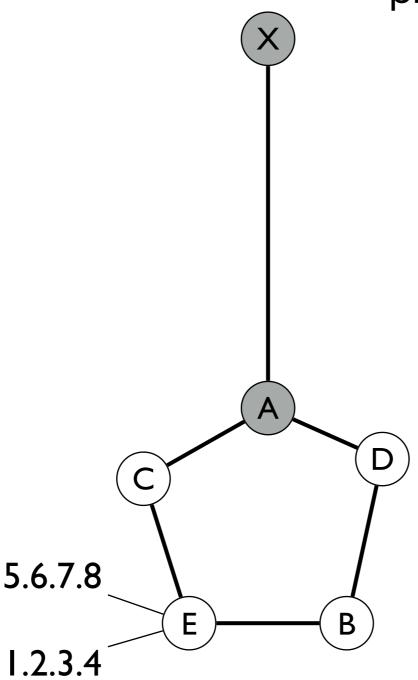


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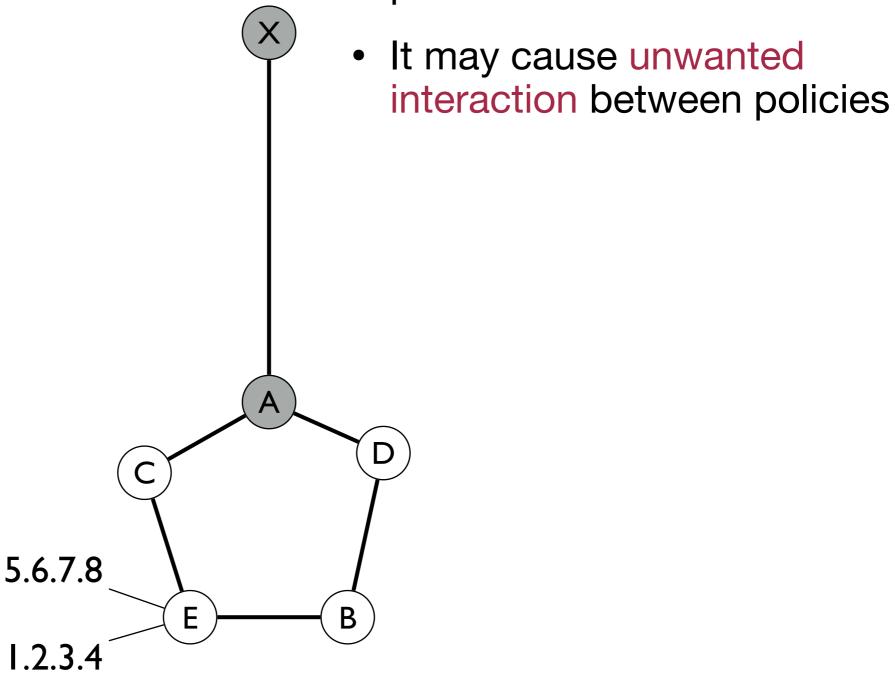
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X: requires routes no more than 3 hops to 1.2.3.4

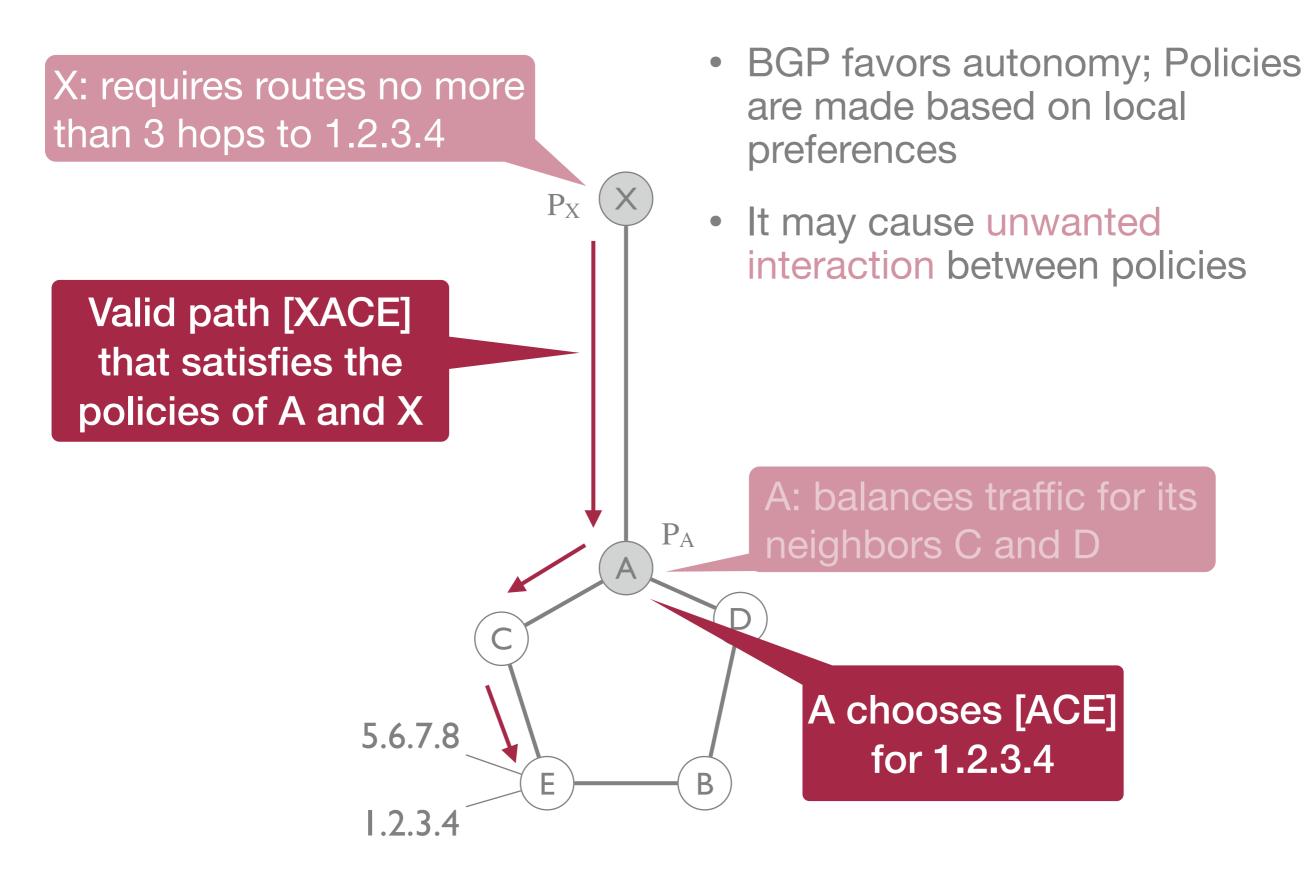
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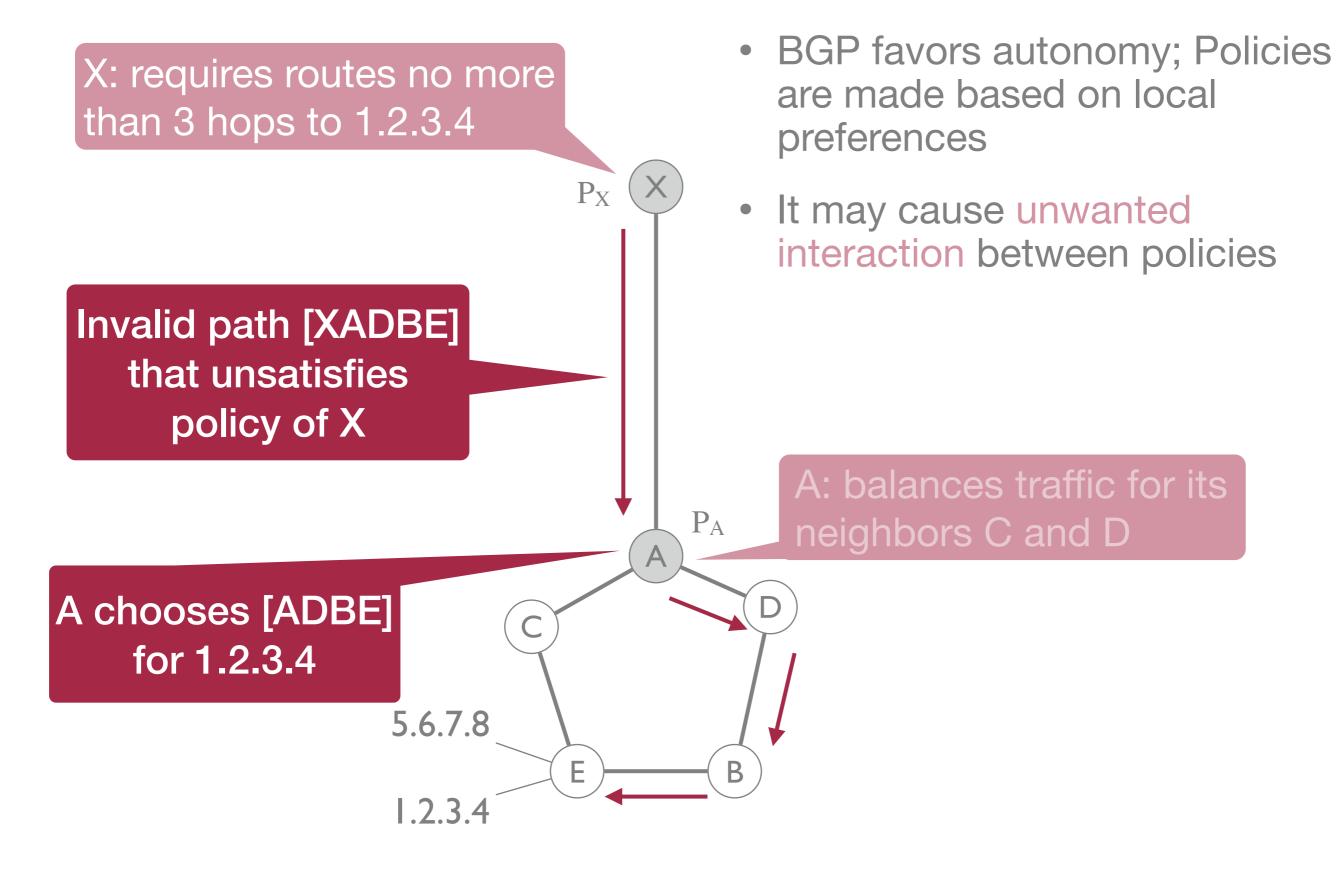
 It may cause unwanted interaction between policies

5.6.7.8 E B

 $P_{X}$ 

A: balances traffic for its neighbors C and D





X: requires routes no more than 3 hops to 1.2.3.4

Require a

proper decision

5.6.7.8

1.2.3.4

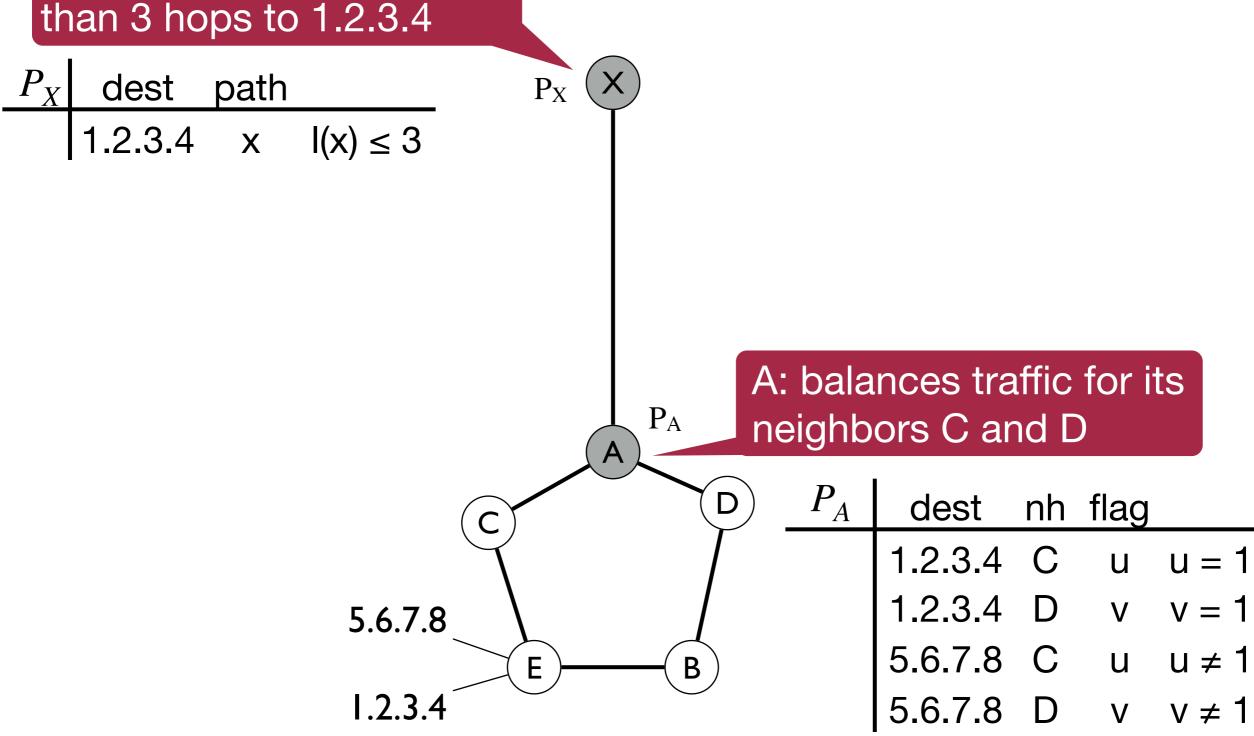
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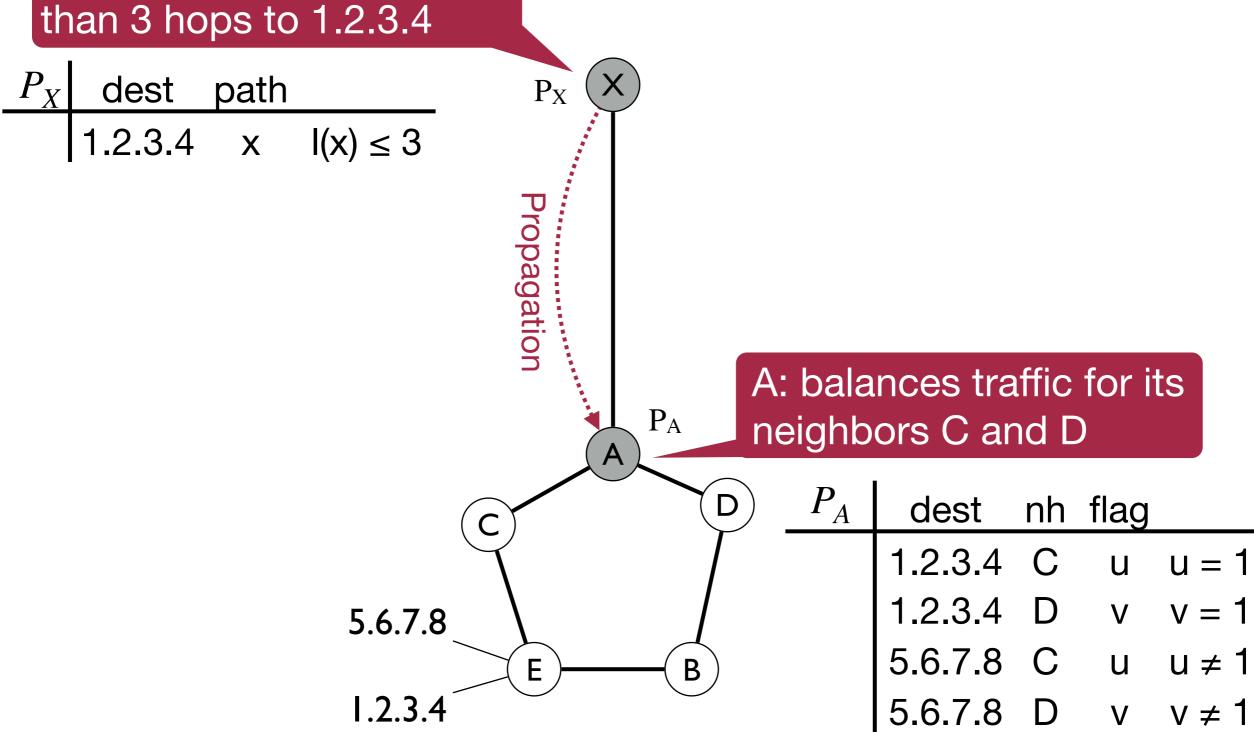
Approach: policy exchange

A: balances traffic for its neighbors C and D

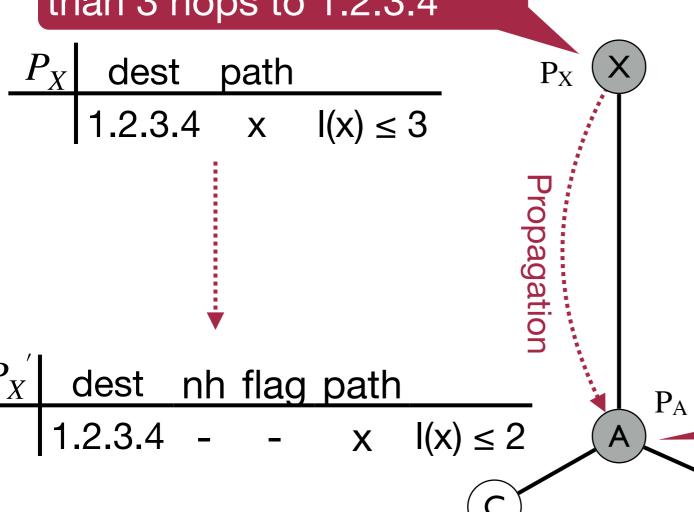




X: requires routes no more than 3 hops to 1.2.3.4





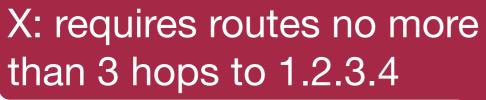


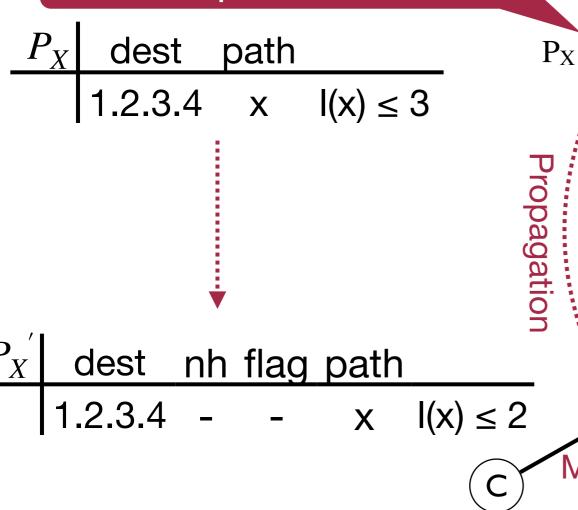
5.6.7.8

1.2.3.4

A: balances traffic for its neighbors C and D

D	$P_A$	dest	nh	flag	
		1.2.3.4 1.2.3.4 5.6.7.8 5.6.7.8	С	u	u = 1
		1.2.3.4	D	V	v = 1
(B)		5.6.7.8	C	u	u ≠ 1
		5.6.7.8	D	V	v ≠ 1

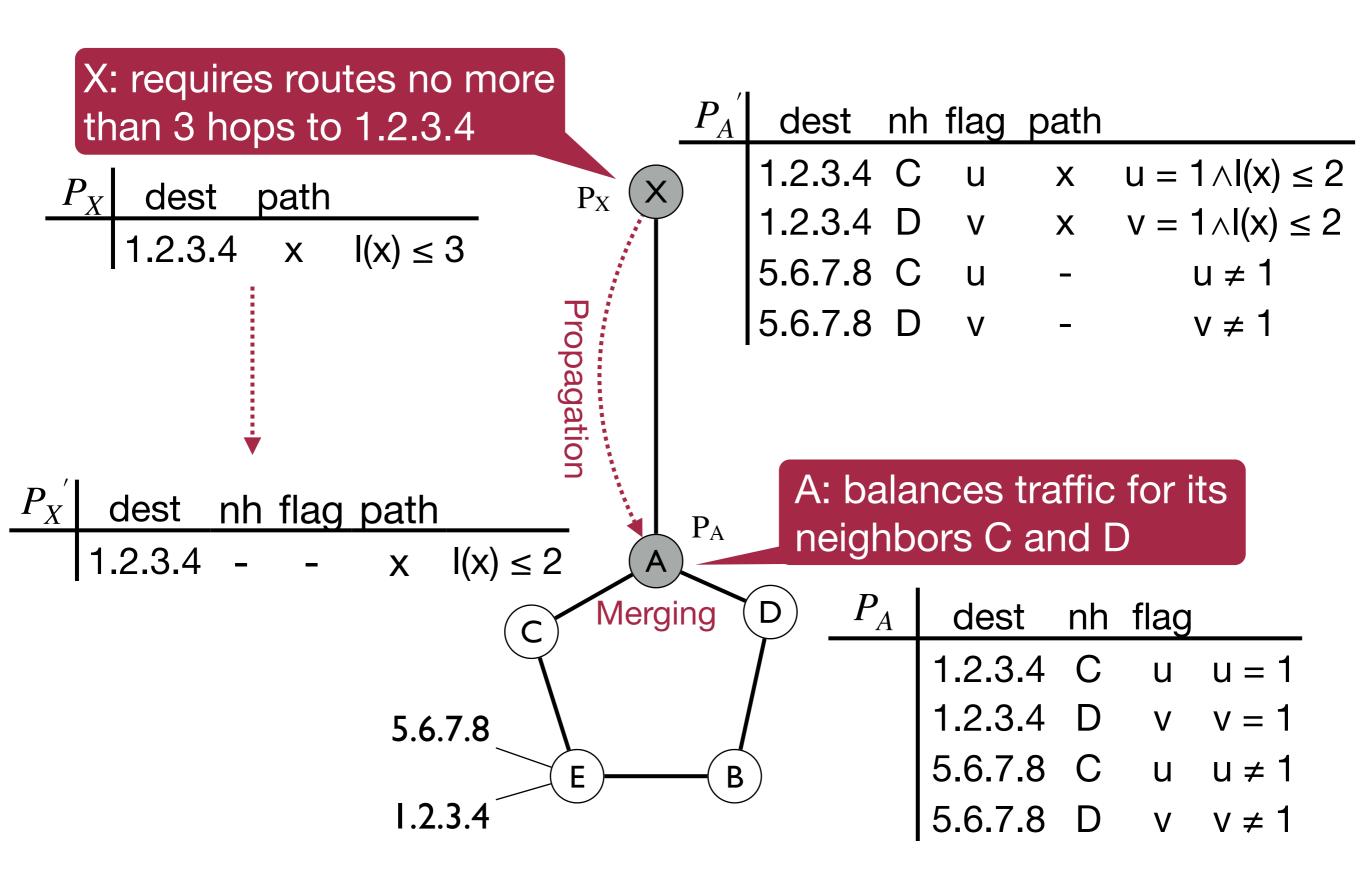


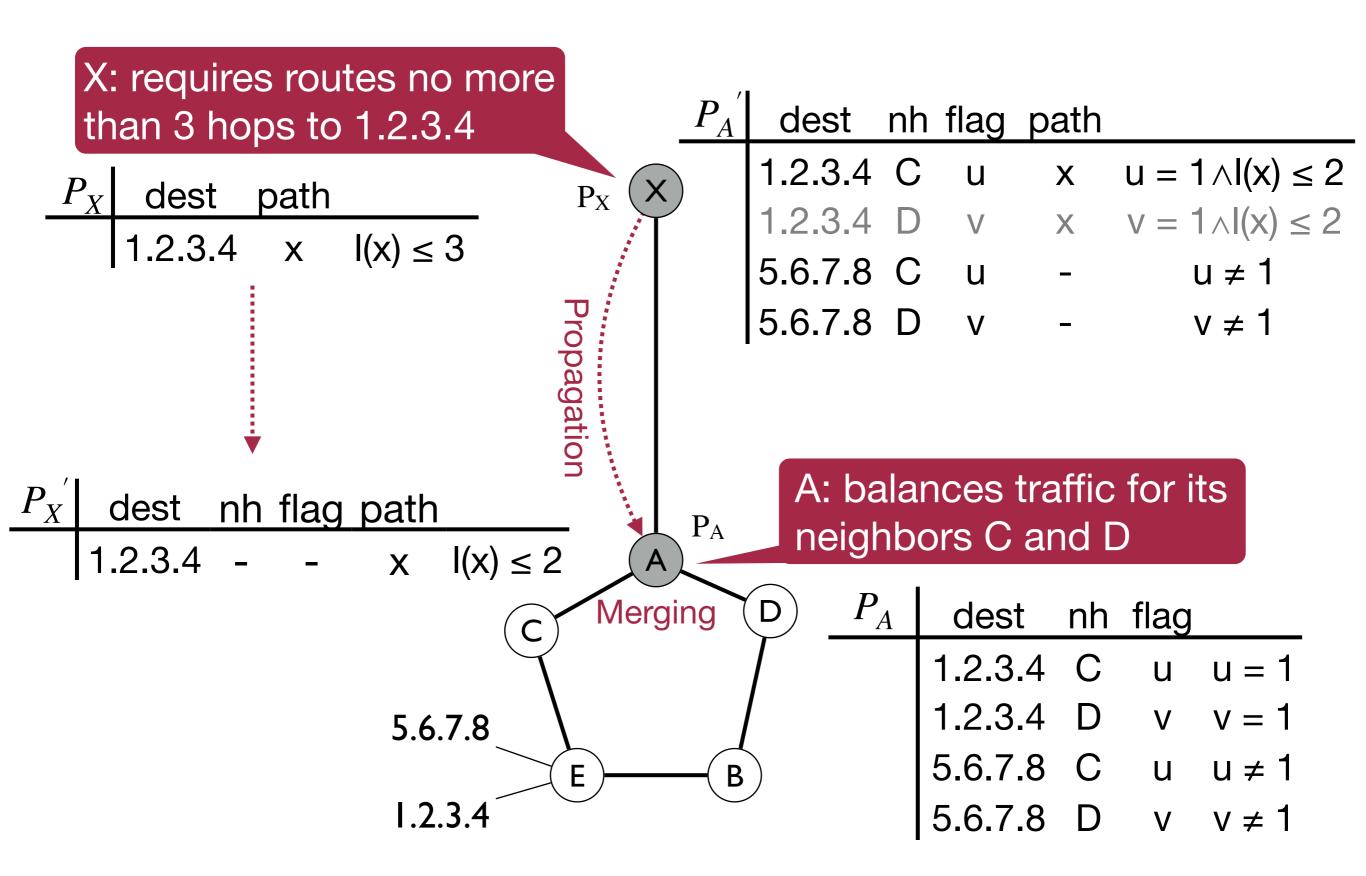


A: balances traffic for its neighbors C and D

Merging D	$P_A$	dest	nh	flag	
		1.2.3.4 1.2.3.4	C	u	u = 1
5.6.7.8		1.2.3.4	D	V	v = 1
(E)		5.6.7.8	C	u	u ≠ 1
1.2.3.4		5.6.7.8	D	V	v ≠ 1

 $P_{A}$ 





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- Applications:
  - Policy Exchange
  - BGP simulation

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Mimicking the behavior of the BGP speaker as it receives announcement from their neighbors

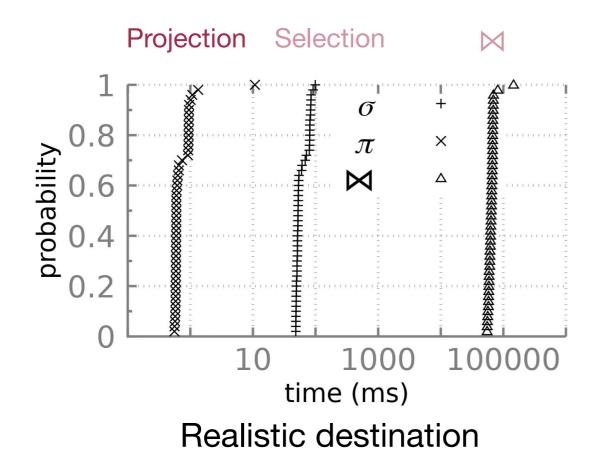
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- Applications:
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#### **Benchmark**

 Using MRT format RouteView BGP data - RIBs and UPDATEs from <u>route-view2.oregon-ix.net</u> to generate synthetic policies and realistic topologies

# Performance - Relational Operators

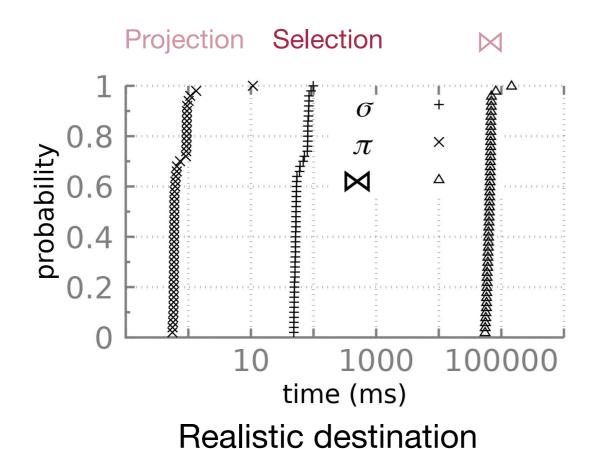


 $\bowtie: P_1 \bowtie P_3$ 

 $P_1$ : static-route & filter policy

 $P_3$ : traffic balance policy

#### Performance - Relational Operators

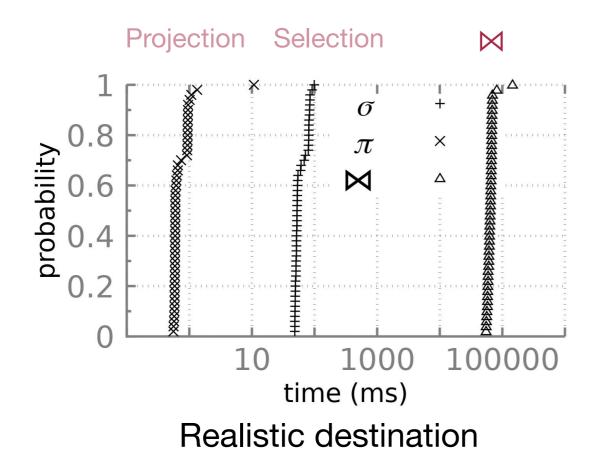


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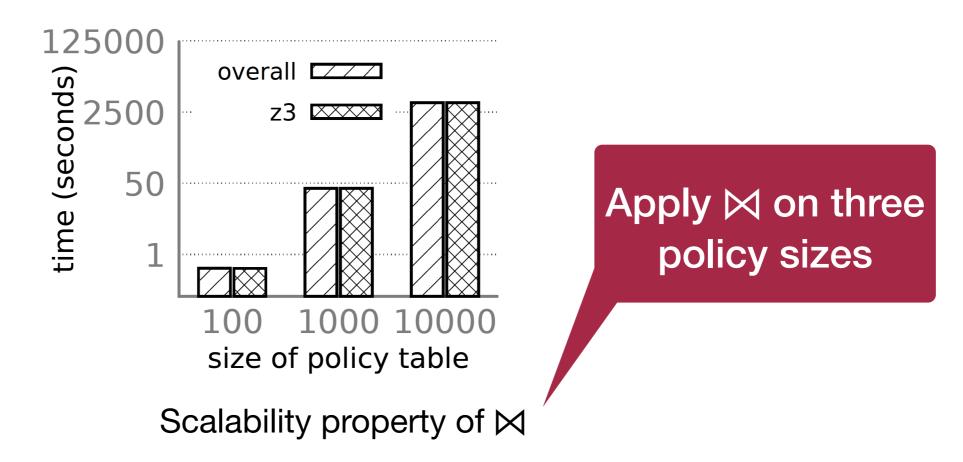


 $\bowtie: P_1 \bowtie P_3$ 

 $P_1$ : static-route & filter policy

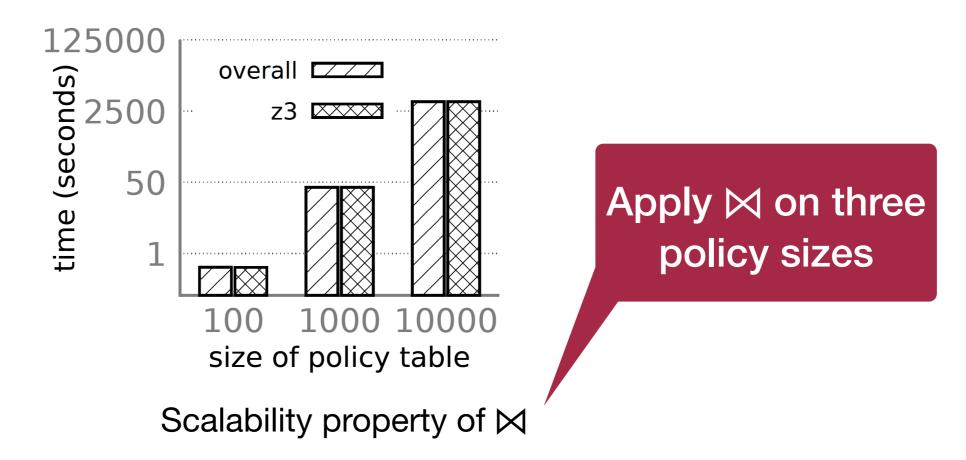
 $P_3$ : traffic balance policy

#### Performance - Scalability



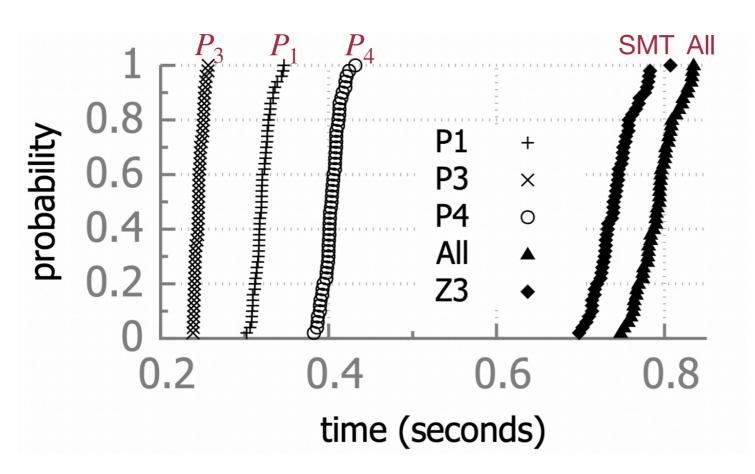
- SMT(Z3) takes 96% of the running time
- Our implementation can handle 10,000 policies in ≤ 70 minutes

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#### Performance - More Policy Examples



 $P_1$ : static route policy

 $P_3$ : traffic balance policy

 $P_4$ : filter policy

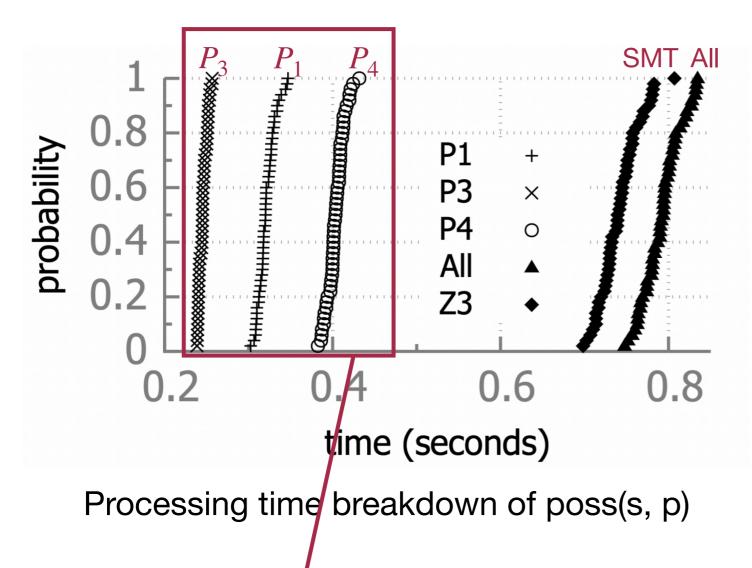
All: total time

SMT: reasoning time

Processing time breakdown of poss(s, p)

Primitive that validates a relational policy (p) on a given network state (s)

#### Performance - More Policy Examples



 $P_1$ : static route policy

 $P_3$ : traffic balance policy

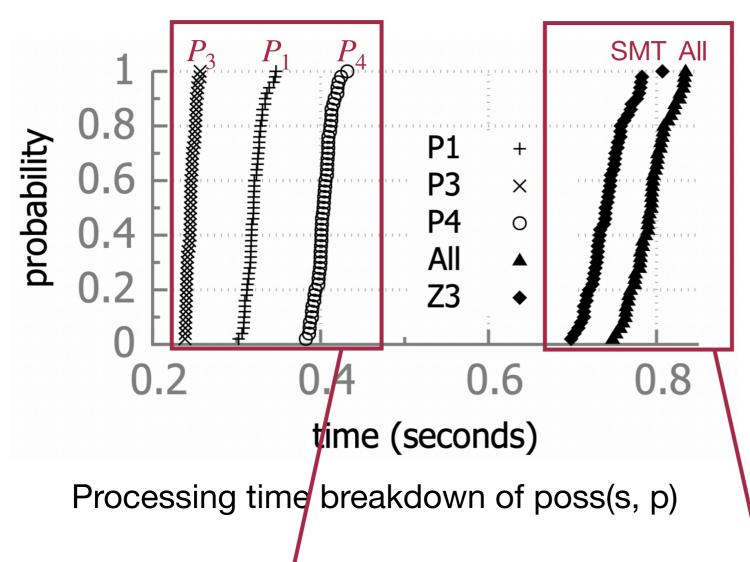
 $P_4$ : filter policy

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- Running time: P3 < P1 < P4</li>

#### Performance - More Policy Examples



 $P_1$ : static route policy

 $P_3$ : traffic balance policy

 $P_4$ : filter policy

All: total time

SMT: reasoning time

- Primitive that validates a relational policy (p) on a given network state (s)
- Running time: P3 < P1 < P4</li>

SMT dominates the source of delay

#### **Related Work**

- Inter-domain Routing Protocols and Architectures
  - D-BGP and Trotsky: partial deployments of protocols, requires their co-existence on the global internet

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  - D-BGP and Trotsky: partial deployments of protocols, requires their co-existence on the global internet
- Declarative Networking:
  - We share database usage in networking
  - We introduce and implement a novel use of conditional tables

# Thank you

http://ravel-net.org/

Support full set of relational operators

Union: ∪

Selection:  $\sigma$ 

Difference: -

Projection:  $\pi$ 

Rename:  $\rho$ 

Support full set of relational operators

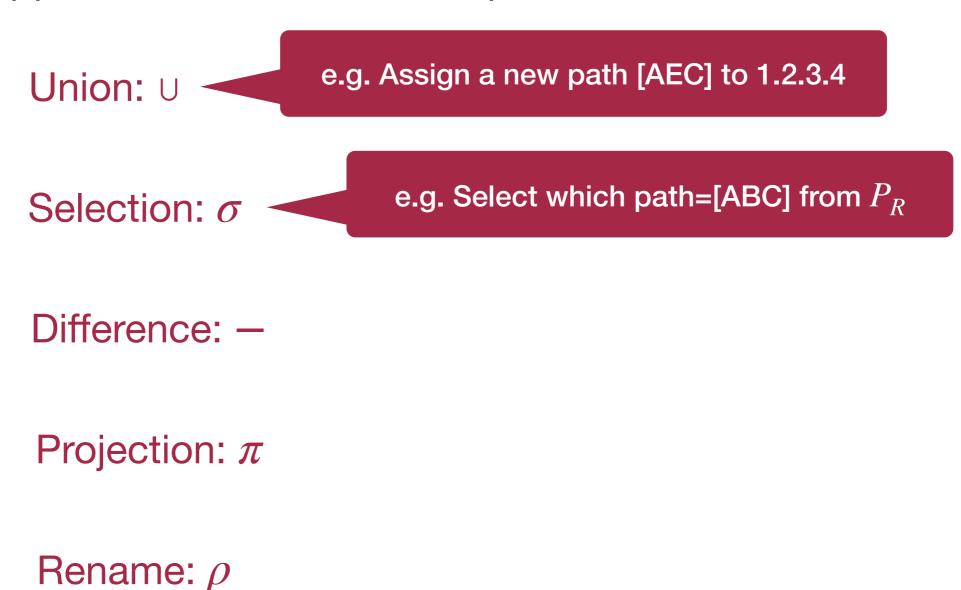
Union: U e.g. Assign a new path [AEC] to 1.2.3.4

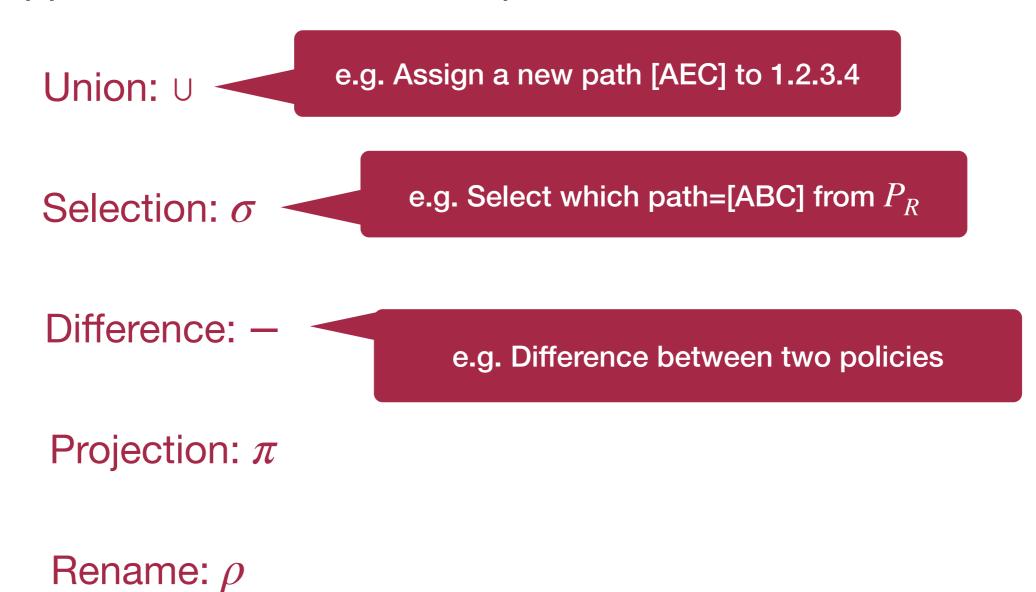
Selection:  $\sigma$ 

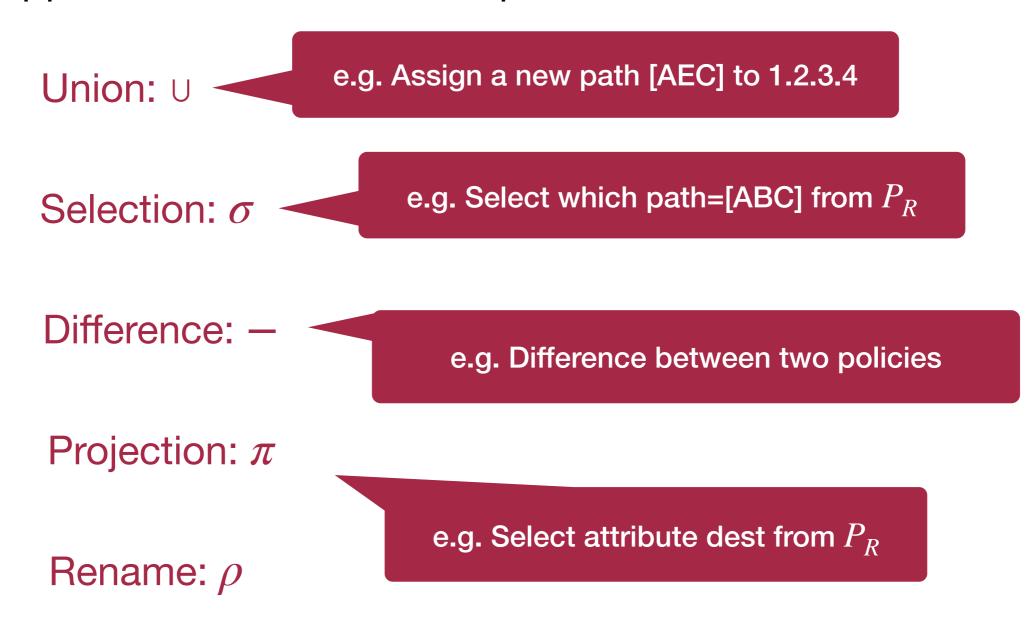
Difference: -

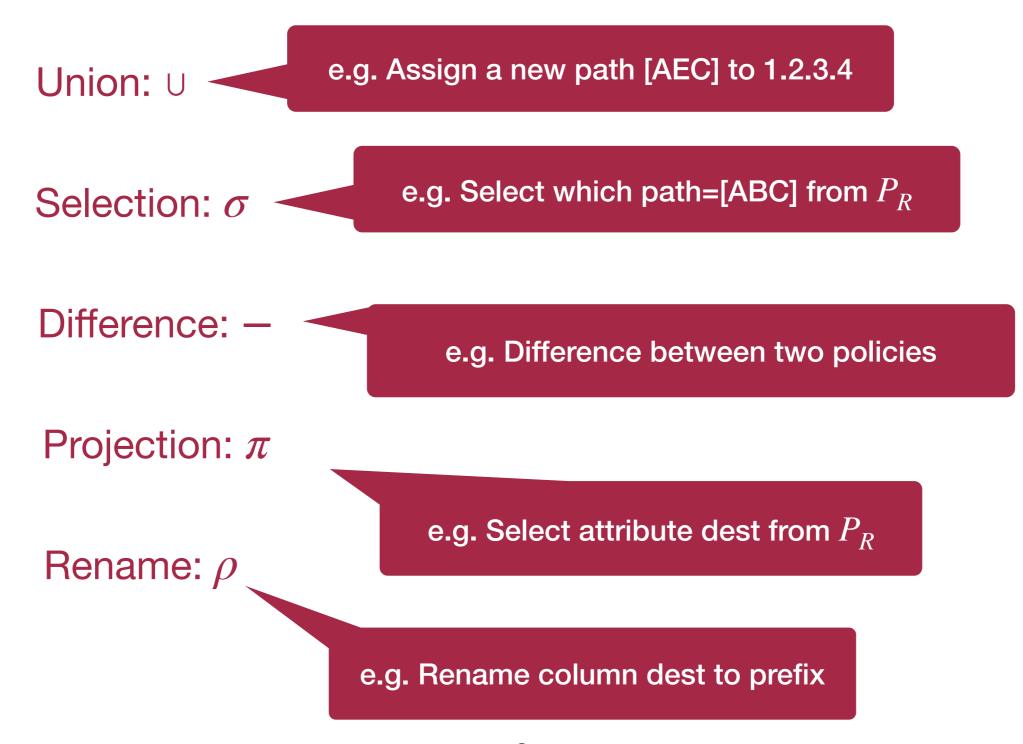
Projection:  $\pi$ 

Rename:  $\rho$ 









- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path			$P_3$	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Z	y≠1.2.3.5∧y≠1	1.2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	v = 1
satis	sfies $\stackrel{\circ}{P}_1$ an	d $P_3$ si	multaneously?	$P_1$ join	$P_3$	5.6.7.8	[ADC]	V	V ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	Χ	u	x=[ABC] ∧ u=1
	1.2.3.4	Χ	V	$x=[ABC] \land x=[ADC] \land v=1$
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ABC] \land u = 1$
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ABC] \land u \neq 1$
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
	у	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$

- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path			$P_3$		dest	path	flag	
	1.2.3.4	X	x=[ABC]				1.2.3.4	[ABC]	U	u = 1
	У	Z	y≠1.2.3.5∧y≠1	.2.3.4	_		5.6.7.8	[ABC]	U	u ≠ 1
Can we generate a new policy that							1.2.3.4	[ADC]	V	v = 1
satisfies $P_1$ and $P_3$ simultaneously?							5.6.7.8	[ADC]	V	V ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
Γ	1.2.3.4	Х	u	x=[ABC] ∧ u=1
L	1.2.3.4	Χ	V	$x=[ABC] \land x=[ADC] \land v=1$
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.4∧y=1.2.3.4∧z=[ABC]∧u=1
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ABC] \land u \neq 1$
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
	у	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$

- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path			$P_3$	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	У	Z	y≠1.2.3.5∧y≠1	.2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a r	new policy that			1.2.3.4	[ADC]	V	v = 1
			multaneously?	$P_1$ join	$P_3$	5.6.7.8	[ADC]	V	v ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	X	u	x=[ABC] ∧ u=1
	1.2.3.4	Χ	V	$x=[ABC] \land x=[ADC] \land v=1$
Γ	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ABC] \land u = 1$
	у	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ABC] \land u \neq 1$
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
L	у	Z	V	y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[ADC]∧v≠1

- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path			$P_3$	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2	2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	V = 1
satis	sfies $\overline{P}_1$ and	$dP_3$ $\sin$	multaneously?	$P_1$ join $I$	3	5.6.7.8	[ADC]	V	V ≠ 1

	_			Contradictory
$P_1 \bowtie P_3$	dest	path	flag	<b>*</b>
	1.2.3.4	X	u	x=[ABC] ∧ u=1
	1.2.3.4	X	V	$x=[ABC] \land x=[ADC] \land v=1$
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y \neq 1.2.3.4 \land z = [ABC] \land u = 1$
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ABC] \land u \neq 1$
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
	у	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$

- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path			$P_3$	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2	2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	V = 1
satis	sfies $\overline{P}_1$ and	$dP_3$ $\sin$	multaneously?	$P_1$ join $I$	3	5.6.7.8	[ADC]	V	V ≠ 1

				Contradictory
$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	X	u	$x=[ABC] \land u=1$
	1.2.3.4	×		$X=[ABC] \land x=[ADC] \land v=1$
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.4∧y≠1.2.3.4∧z=[ABC]∧u=1
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ABC] \land u \neq 1$
	У	Z	<b>── ∀</b>	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
	у	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$

- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path		$P_{i}$	3	dest	path	flag	
	1.2.3.4	X	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2.3	.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	ate a n	ew policy that			1.2.3.4	[ADC]	V	v = 1
satis	sfies $\overline{P}_1$ and	d $P_3$ si	multaneously? $P_1$	oin $P_3$		5.6.7.8	[ADC]	V	V ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	X	u	x=[ABC] ∧ u=1
				$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ABC] \land u \neq 1$
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$

- Static route and filter policy
   Traffic balance policy

$P_1$	dest	path		$P_3$	dest	path	flag	
	1.2.3.4	X	x=[ABC]		1.2.3.4	[ABC]	u	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2.3		5.6.7.8			
Can	we gener	ate a n	ew policy that		1.2.3.4	[ADC]	V	v = 1
satisfies $P_1$ and $P_3$ simultaneously? $P_1$ join $P_2$					5.6.7.8	[ADC]	V	$V \neq 1$

	_				Redur	ndancy
$P_1 \bowtie P_3$	dest	path	flag		1	
	1.2.3.4	X	u	x=[ABC] ∧ u=1		
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.	4∧y=5 <b>.</b>	6.7.8∧z=[ABC]∧u≠1
	У	Z	V	y≠1.2.3.5∧y≠1.2.3.	4∧y=5.	6.7.8∧z=[ADC]∧v≠1