

Organizers: Moses Charikar, Anay Mehrotra, Charlotte Peale, Chirag Pabbaraju, Grigoris Velegkas

Session 1 (40/20 mins)

- > Language Generation and Learning Theory
- > Stronger Notions of Generation and Comparison to Prediction

Moses

Chirag

Session 2 (20/20/20 mins)

- ➤ Validity—Breadth Trade-Off (Part I)
- ➤ Validity—Breadth Trade-Off (Part II)
- Diverse and Robust Generation

Anay Grigoris Charlotte

Schedule



Tutorial on Language Generation



Visit: LanguageGeneration.github.io

Organizers:

Moses Charikar Stanford



Anay Mehrotra Yale University



Chirag Pabbaraju Stanford



Charlotte Peale Stanford



Grigoris Velegkas Yale → Google Research



COLT 2025

LLMs and Language Generation

Computer scientists have been fascinated by language acquisition by humans and machines for decades

LLMs: what is the problem?

Kleinberg, Mullainathan, 2024



LLMs: what is the problem?

Kleinberg, Mullainathan, 2024

From a large collection of text:

- 1. assign probabilities to every sequence of words
- produce strings with high assigned probability



LLMs: simplest problem statement?

Kleinberg, Mullainathan, 2024

Given strings from an unknown language, produce valid unseen strings from the language.



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string: sequence of symbols

language: set of strings e.g. C programs



E MARK GOLD*

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I wish to construct a precise model for "able to speak English"... to investigate theoretically how it can be achieved artificially

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I wish to construct a precise model for "able to speak English" ... to investigate theoretically how it can be achieved artificially

Since we cannot explicitly write down the rules of English... *artificial intelligence... will have to learn... from implicit data...*

Gold 1967, Angluin 1979

Game between adversary and algorithm.

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Adversary thinks of target language K from countable list (e.g. all context-free grammars)

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Adversary enumerates strings one by one

In each step, algorithm guesses index $\mathbf{i}_{\mathbf{i}}$ (goal: $\mathbf{L}_{\mathbf{i}_{\mathbf{i}}} = \mathbf{K}$)

Gold 1967, Angluin 1979

Game between adversary and algorithm.

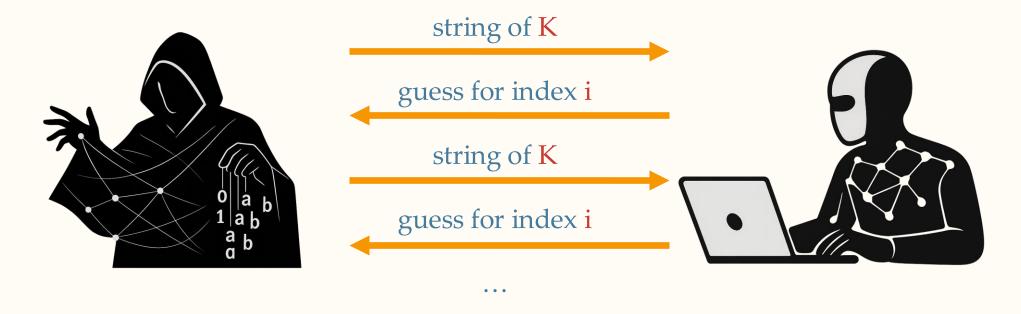
Adversary thinks of target language K from countable list (e.g. all context-free grammars)

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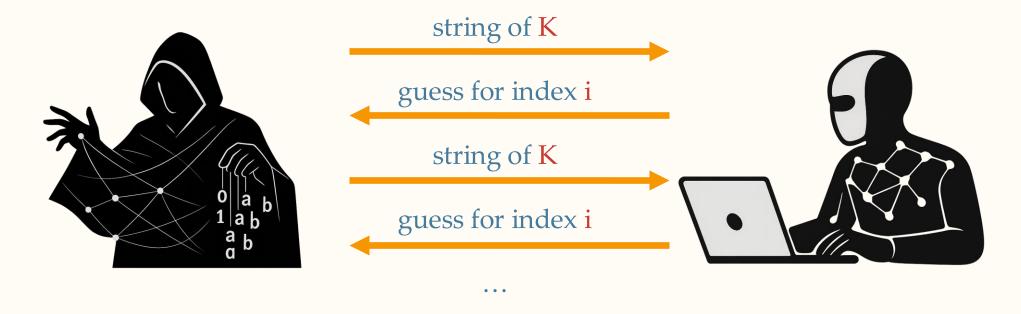
In each step, algorithm guesses index $\mathbf{i}_{\mathbf{i}}$ (goal: $\mathbf{L}_{\mathbf{i}_{\mathbf{i}}} = \mathbf{K}$)

Success: guess correct for every **!** > **!** (We say that algorithm has identified K in the limit)

Gold 1967, Angluin 1979

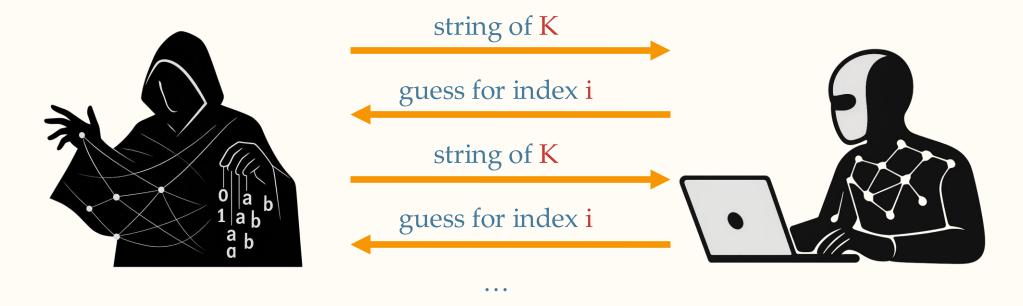


Gold 1967, Angluin 1979



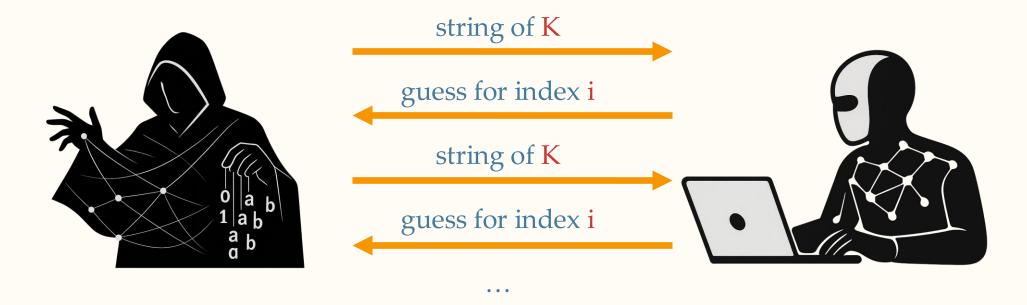
Algorithm never sees string not in K

Gold 1967, Angluin 1979



Algorithm never sees string not in K Never told whether guess is correct

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Algorithm never sees string not in K

Never told whether guess is correct

Cannot ask if string is in K

Theorem [Gold 1967]

Language identification in the limit is impossible even for simple collections such as all regular languages

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Language Generation Kleinberg, Mullainathan, 2024

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5₁ ∈ **K** : strings enumerated up to time t

Kleinberg, Mullainathan, 2024

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 $S_t \in K$: strings enumerated up to time t

In each step, algorithm guesses string $\mathbf{a}_{\mathbf{i}}$ (goal: $\mathbf{a}_{\mathbf{i}} \in \mathbf{K} \setminus \mathbf{S}_{\mathbf{i}}$)

Kleinberg, Mullainathan, 2024

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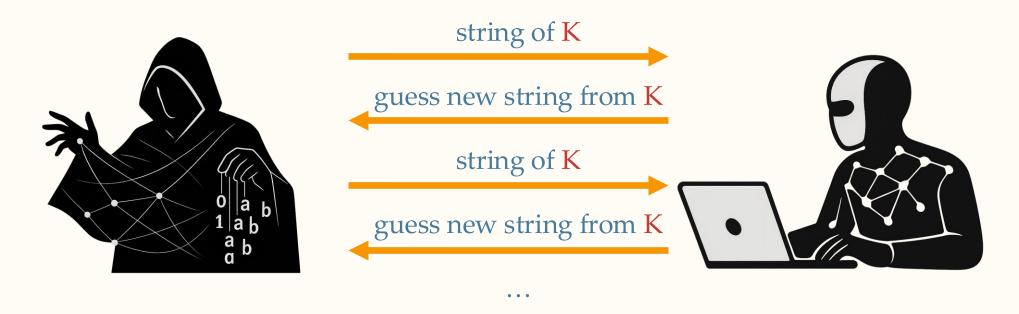
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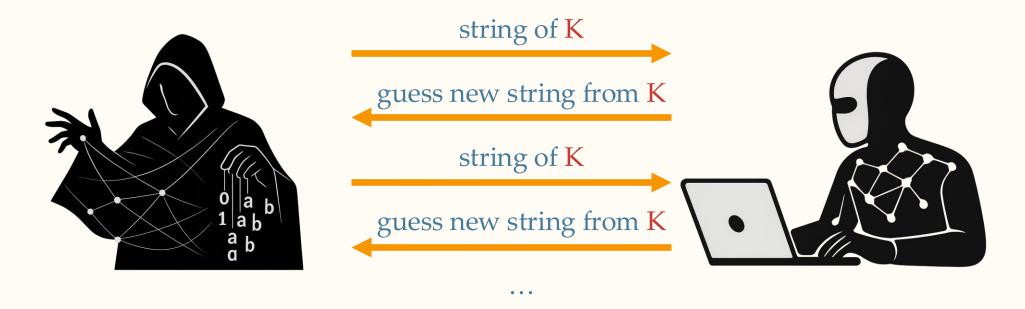
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Kleinberg, Mullainathan, 2024

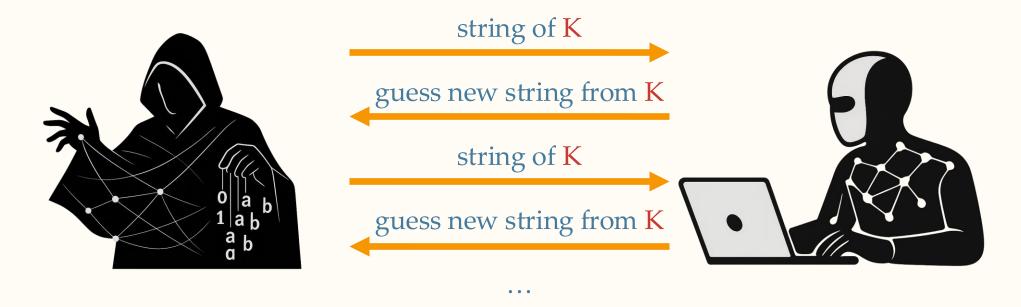


Kleinberg, Mullainathan, 2024



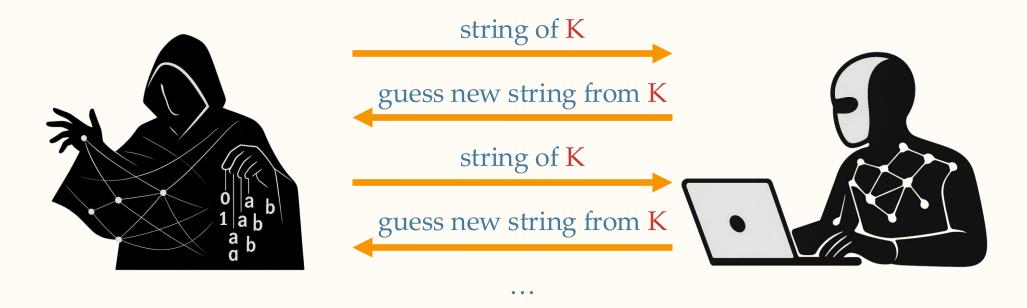
Algorithm never sees negative examples

Kleinberg, Mullainathan, 2024



Algorithm never sees negative examples
No feedback

Kleinberg, Mullainathan, 2024



Algorithm never sees negative examples

No feedback

Assume all languages infinite

Theorem [Kleinberg, Mullainathan 2024]

Language generation in the limit is possible for any countable collection of languages

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Generation vs Identification

After seeing many C programs

(generation) output valid C programs

(identification) output valid grammar for C

Theorem [Kleinberg, Mullainathan 2024]

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Algorithm only needs to generate from infinite subset of K

Language Generation in the Limit

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validity (only generate valid strings) vs breadth (large subset of K) tradeoff:

Language Generation in the Limit

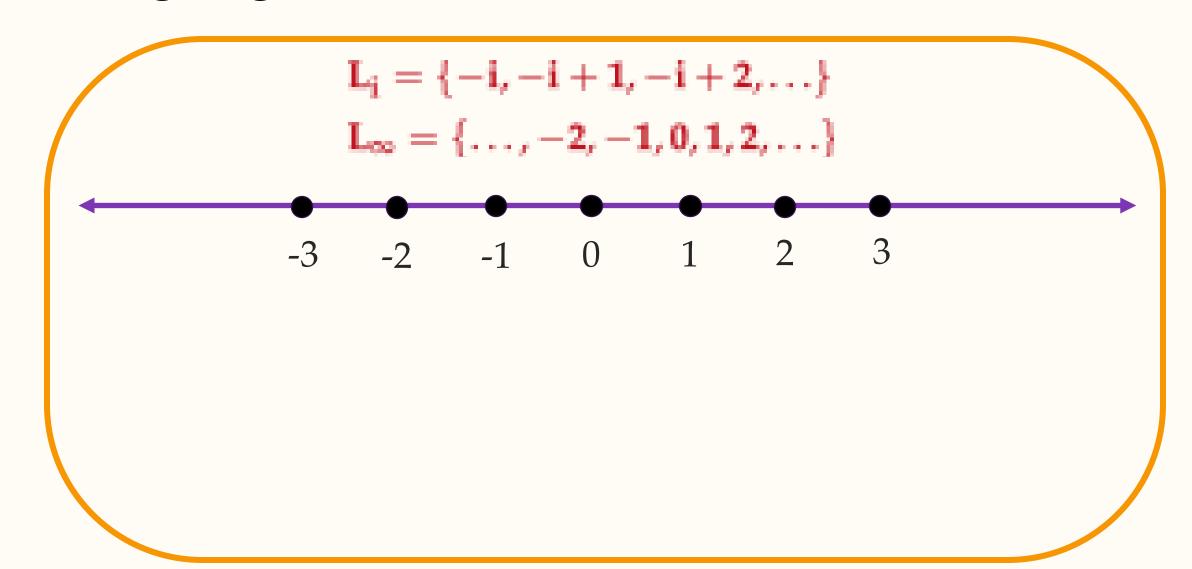
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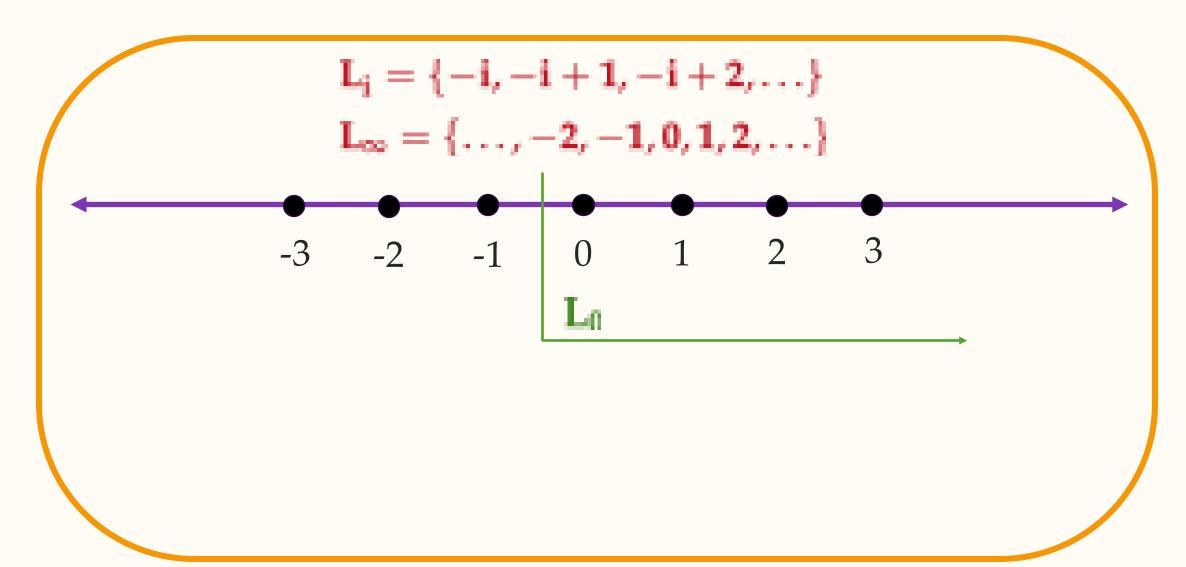
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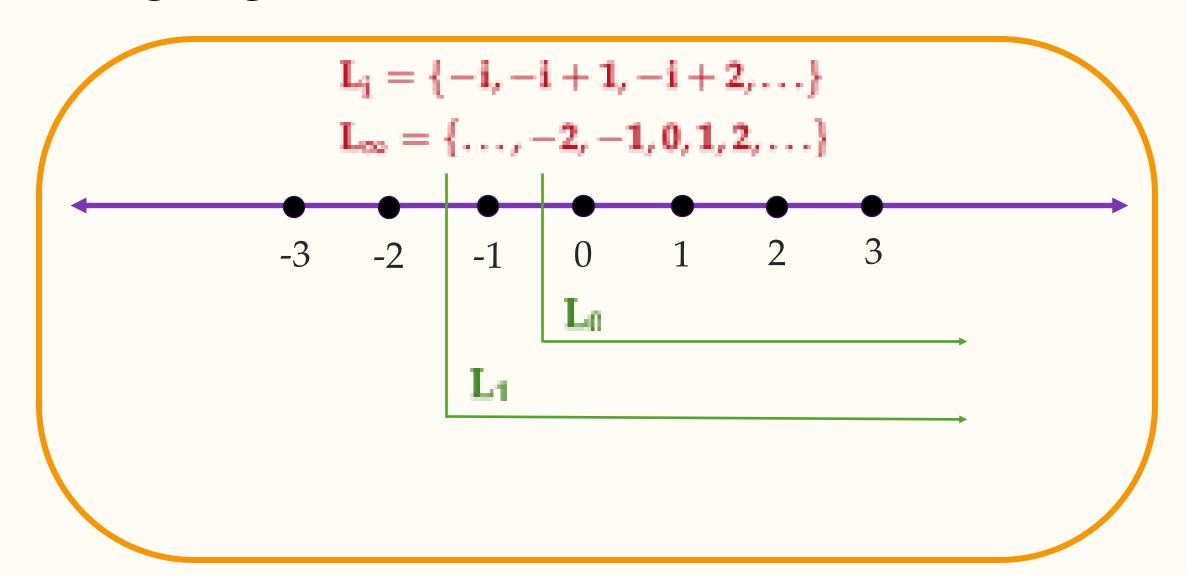
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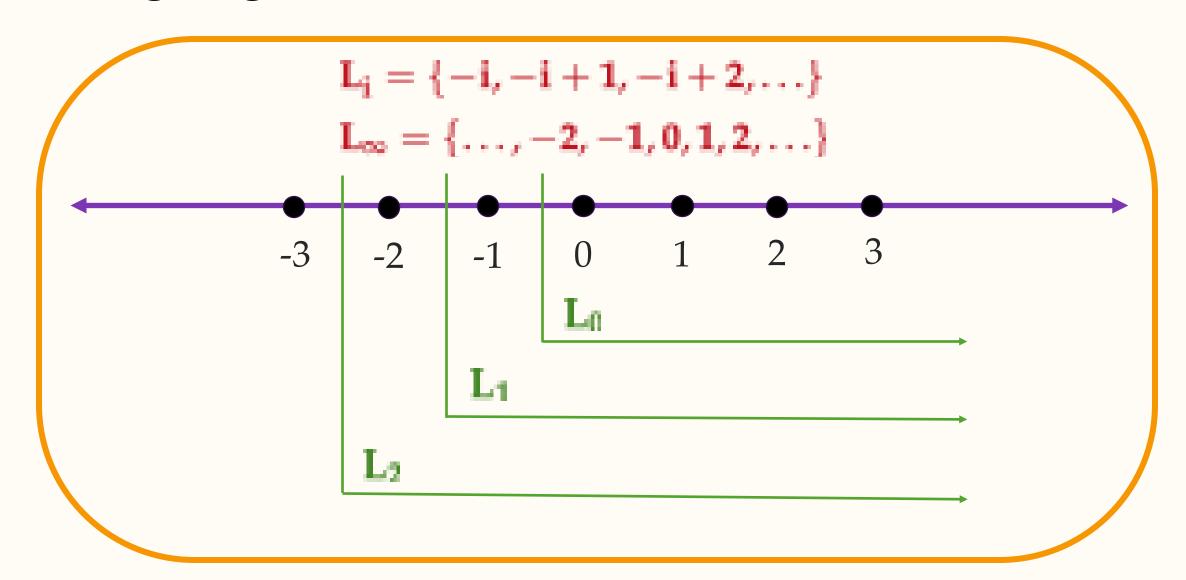
validity (only generate valid strings) vs breadth (large subset of K) tradeoff:

hallucination vs mode-collapse









Closure operation

(5): intersection of all languages containing S

Output an element of $(S_1) \setminus S_1$?

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Problem: $(S_1) \setminus S_1$ can be empty for all t!

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Problem: $(\mathbf{5}_{\mathbf{1}}) \setminus \mathbf{5}_{\mathbf{1}}$ can be empty for all t!

e.g. modify previous example

$$L_{i,V}=V\cup\{-i,-i+1,-i+2,\ldots\}$$

where V ranges over all finite sets







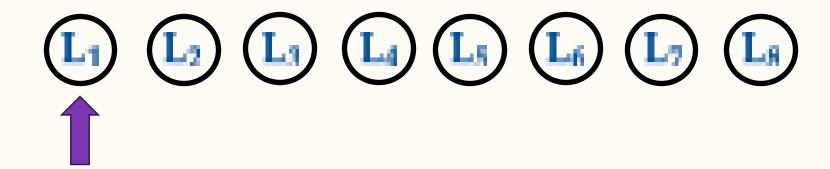


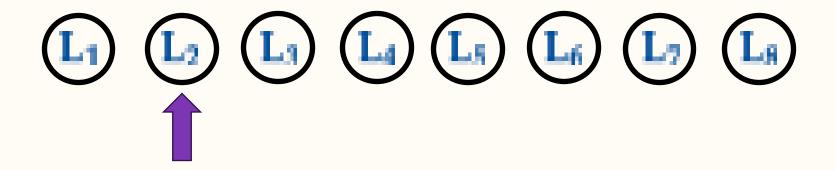


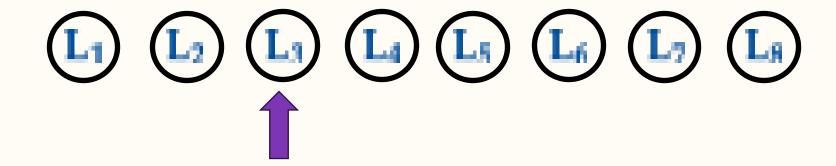


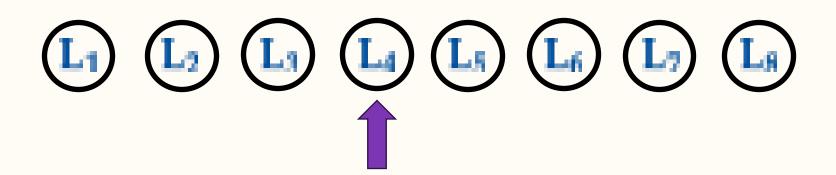


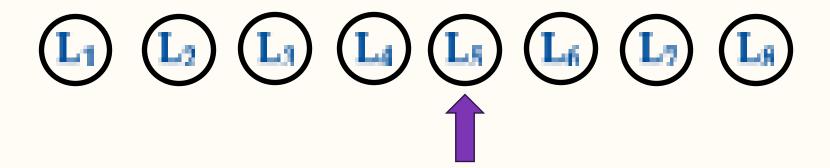


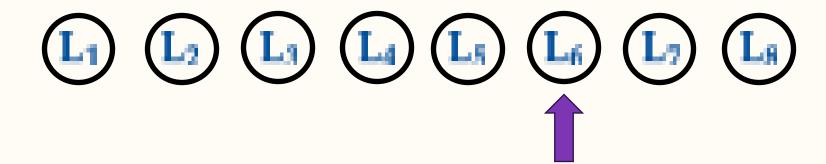












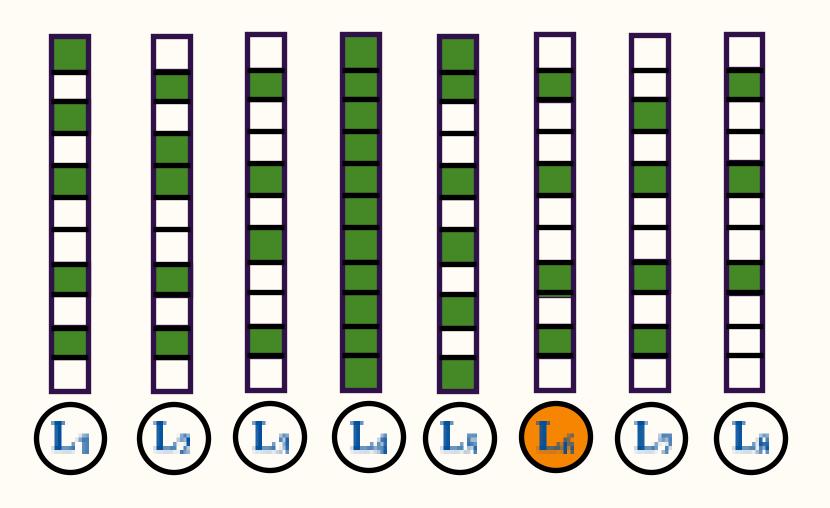


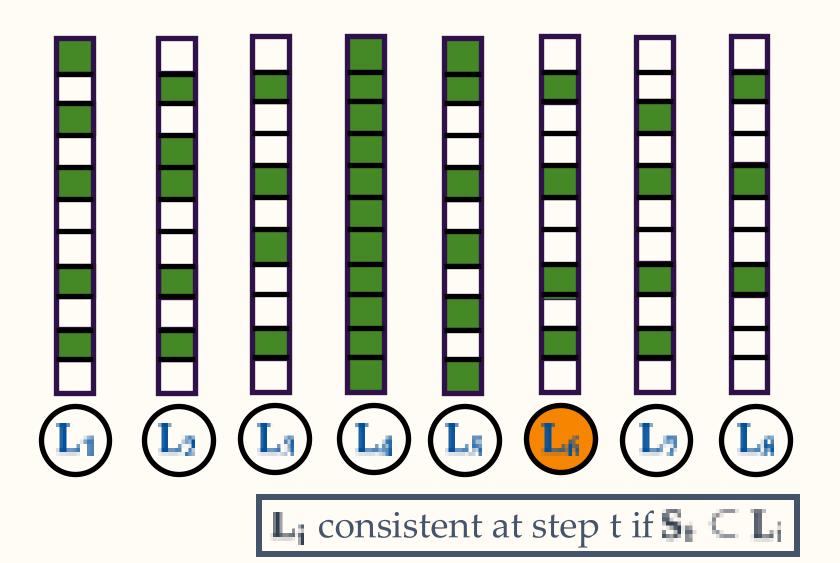
(wrong) Proof:

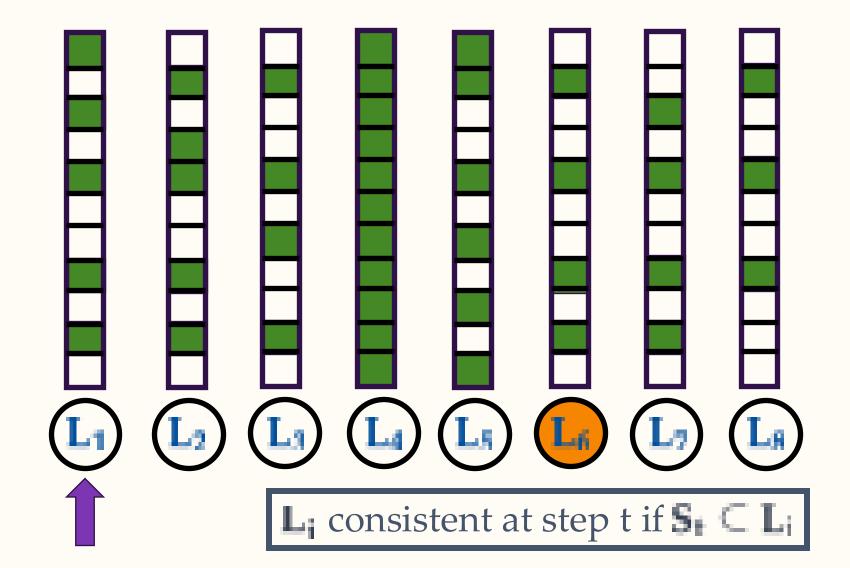
If $\mathbf{L}_{i} \neq \mathbf{K}$ enumeration of K will eventually reveal this

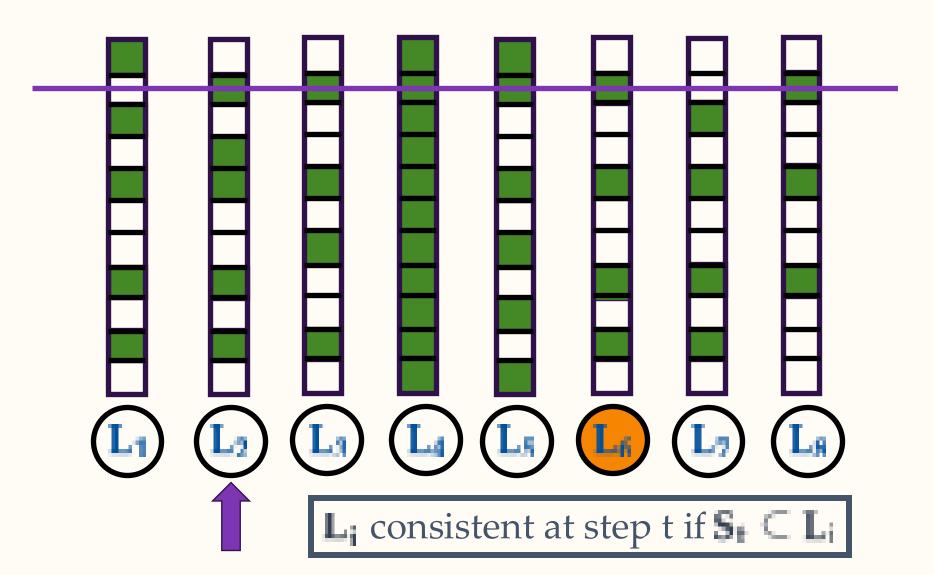
Since $K = L_1$ for some z, we will eventually get to $K = L_1$

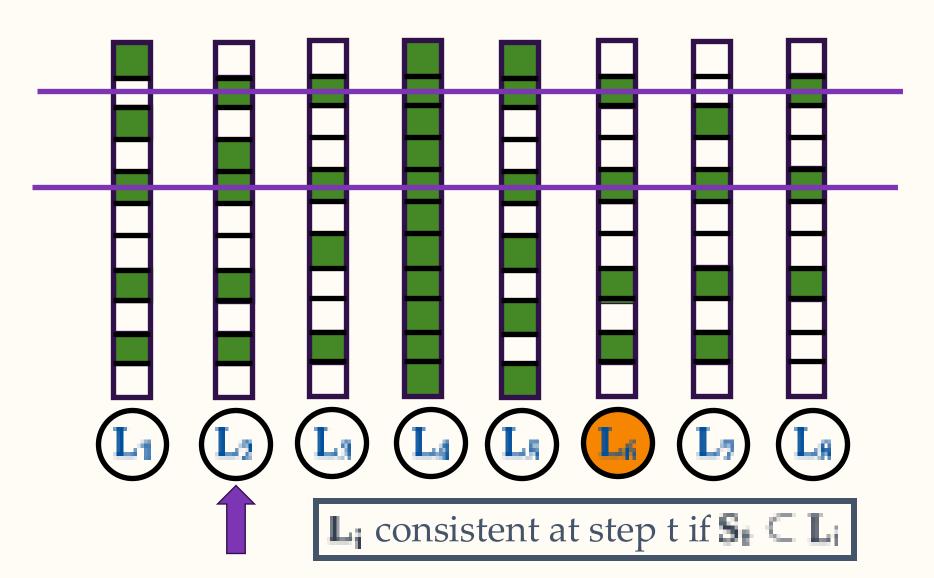
We will never move beyond K

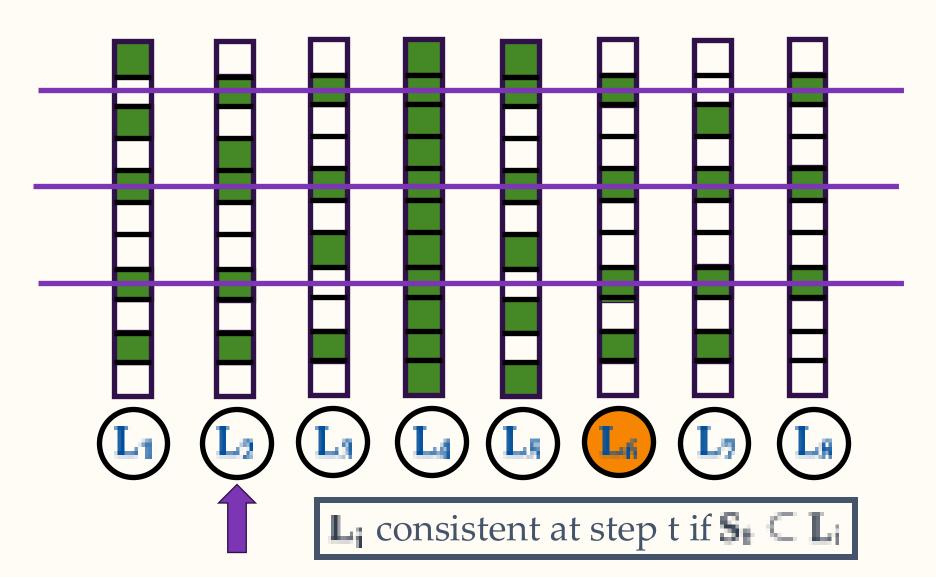












Minimal Consistent Languages?

Consistent language L_i such that there is no consistent L_i with $L_i \subseteq L_i$

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Problem: There may not be a minimal consistent language!

Critical Languages

Definition: L_n is critical at step t if

- 1. L_n is consistent at step t
- 2. If L_i is consistent, i < n, then $L_n \subseteq L_i$

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Key facts:

1. critical language exists: lowest-indexed consistent language

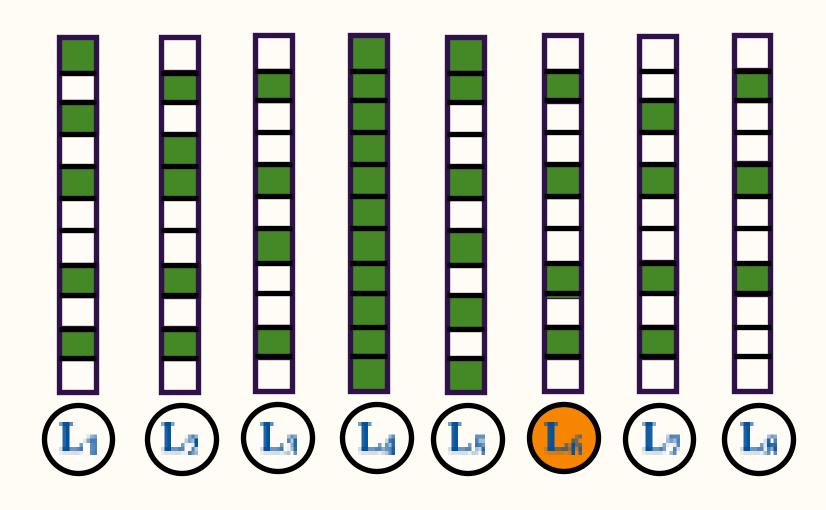
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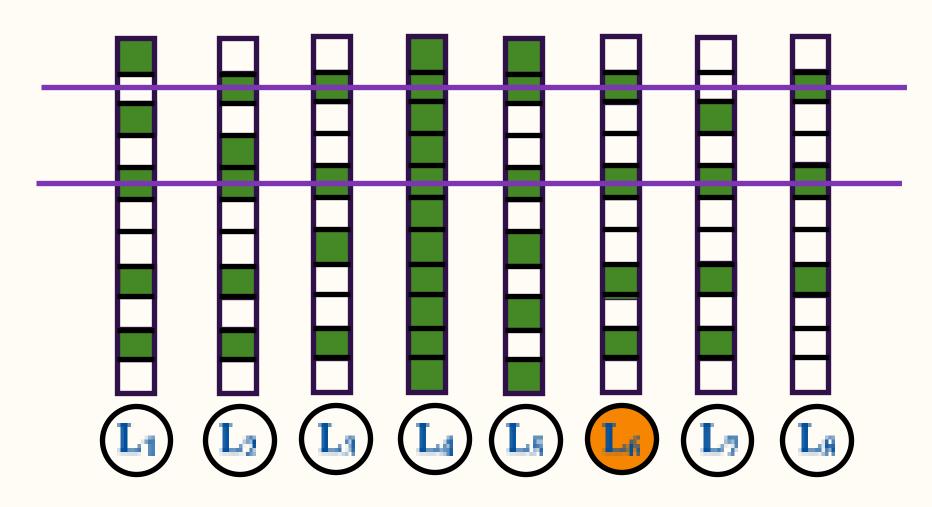
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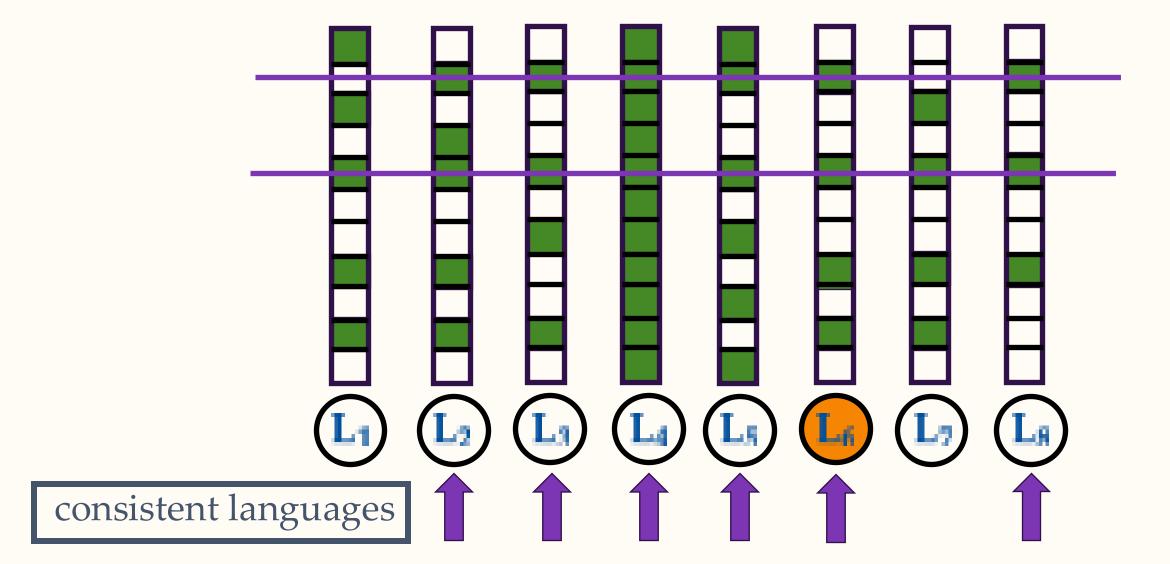
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- 2. If L_i is consistent, i < n, then $L_n \subseteq L_i$

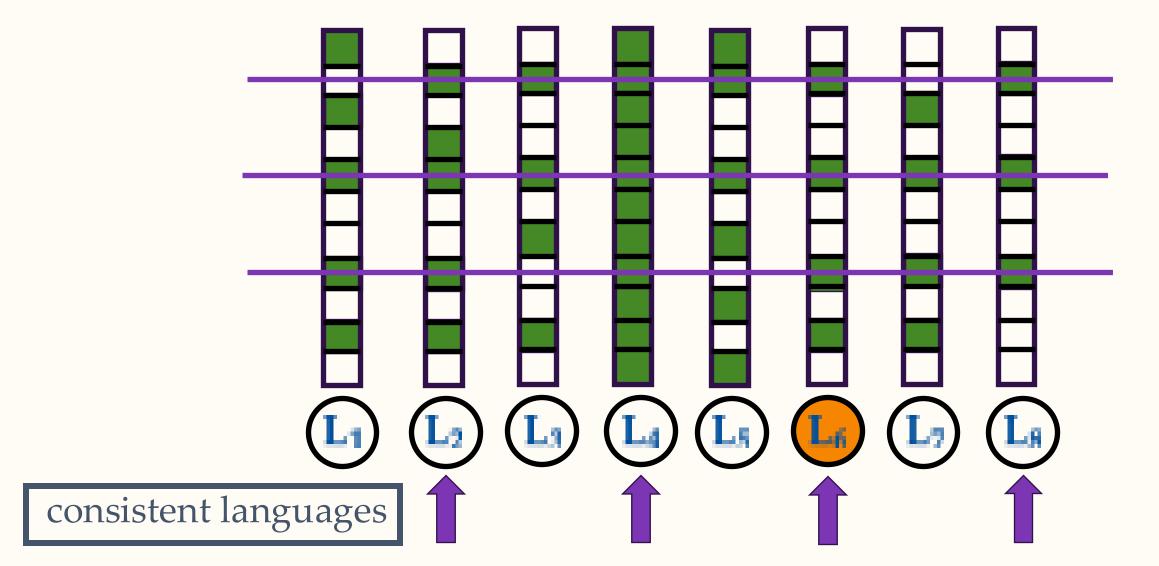
Key facts:

- 1. critical language exists: lowest-indexed consistent language
- 2. Target language will eventually become critical









Algorithm for Generation in the Limit

At step t, only consider L₁....L₁

 L_{n_i} : critical language with highest index $n_i \le 1$

Generate a string from L_{ni} \ S₁

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At step t, only consider L_1, \ldots, L_1

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Proof sketch:

For large enough t, target language $L_{\mathbf{r}}$ is critical and in L_{1}, \ldots, L_{l}

Algorithm for Generation in the Limit

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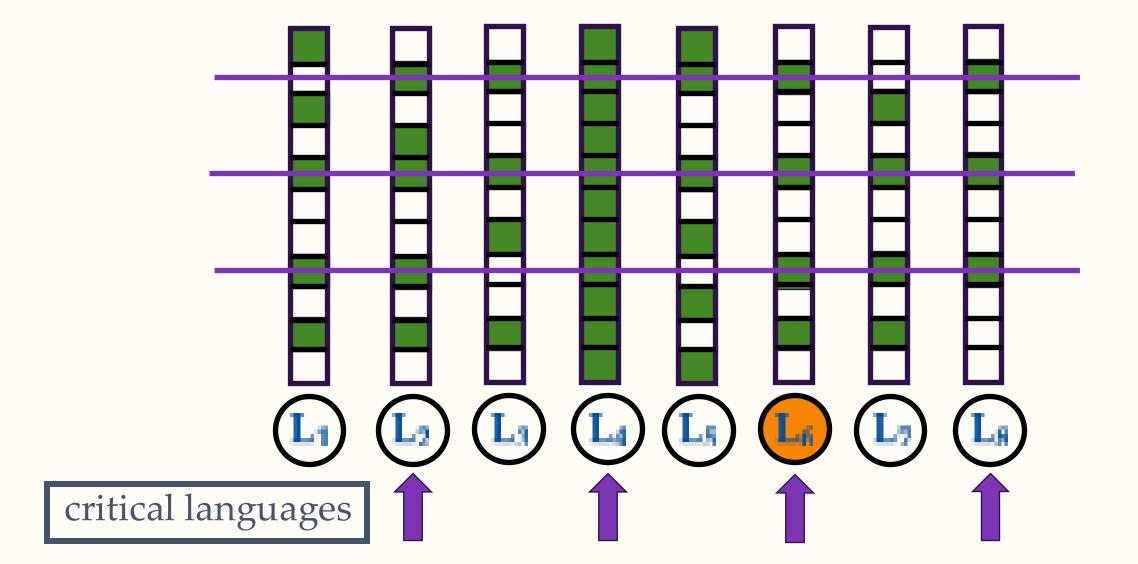
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Proof sketch:

For large enough t, target language $L_{\mathbf{I}}$ is critical and in L_{1}, \ldots, L_{l}

 $L_{n_i} \subseteq L_z$, so any string from $L_{n_i} \setminus S_i$ also belongs to $L_z \setminus S_i$

Generation in the Limit correctness



Language Generation: Tradeoffs & Extensions

Mode-collapse in [Kleinberg, Mullainathan 2024] algorithm. validity vs breadth tradeoff: Anay and Grigoris

Stronger requirements (uniform/non-uniform generation): Chirag

Diversity constraints and noisy data: Charlotte