

Inherent Trade-offs in Language Generation

Anay Mehrotra

Yale University

Alkis Kalavasis



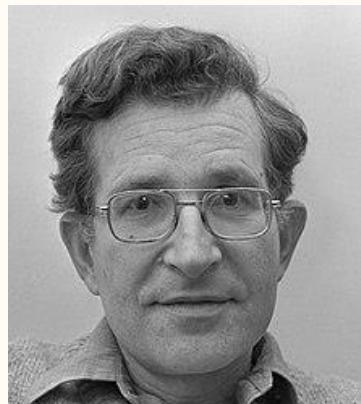
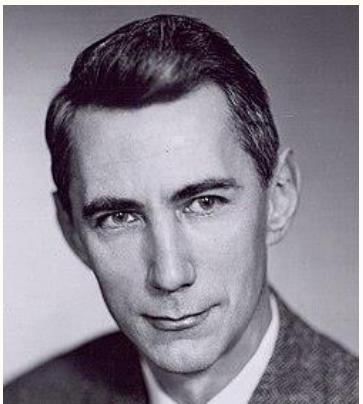
Grigoris Velegkas



May 2025

CS and Language Generation

*Computer scientists have been fascinated by language acquisition
by humans and machines for decades*

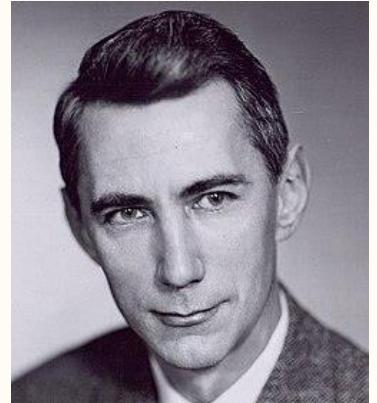


Language Identification in the Limit

E MARK GOLD*

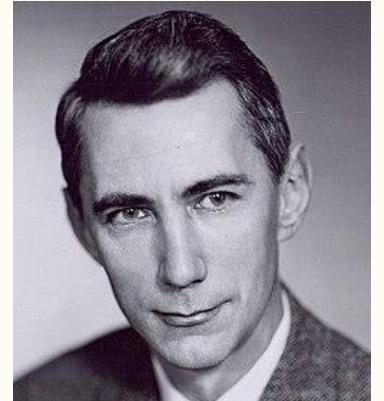
CS and Language Generation

1951 Shannon
*Prediction and
entropy of English*



CS and Language Generation

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- A generation game between Betty and Claude Shannon

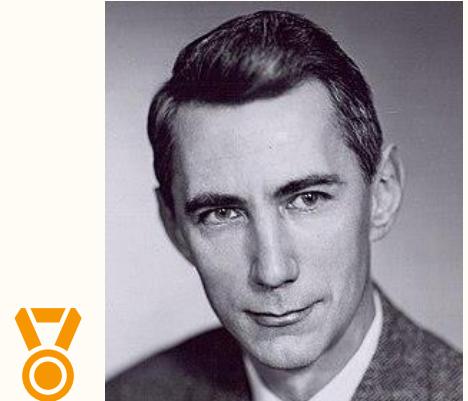
(1) THE ROOM WAS NOT VERY LIGHT A SMALL OBLONG
(2) ----ROO-----NOT-V----I----SM---OBL----

(1) READING LAMP ON THE DESK SHED GLOW ON
(2) REA-----0-----D----SHED-GLO--0--

(1) POLISHED WOOD BUT LESS ON THE SHABBY RED CARPET
(2) P-L-S-----0---BU--L-S--0-----SH-----RE --C-----

CS and Language Generation

1951 Shannon
Prediction and entropy of English



- Introduced n -grams – had tremendous impact in the 1980s!

2-gram model

Rhodesian Army offensive
on average salary increase
it four networks ...

5-gram model

He praised love's ability
to use dialogue to effect
an emotional response...

CS and Language Generation

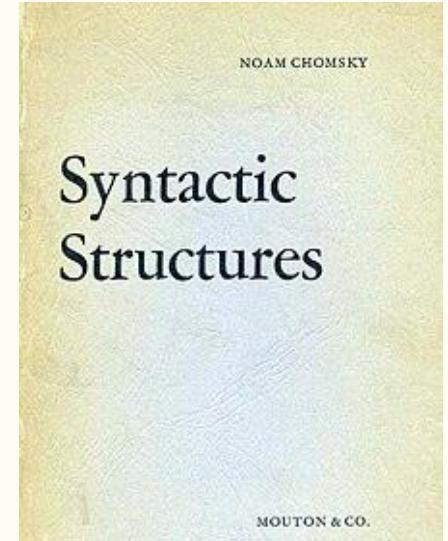
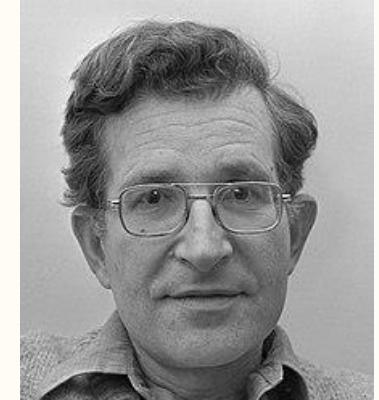
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1957 Chomsky
Syntactic structures & formal grammars



CS and Language Generation

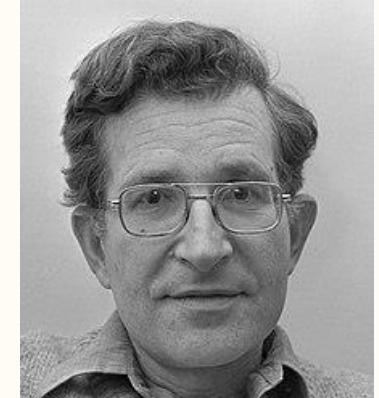
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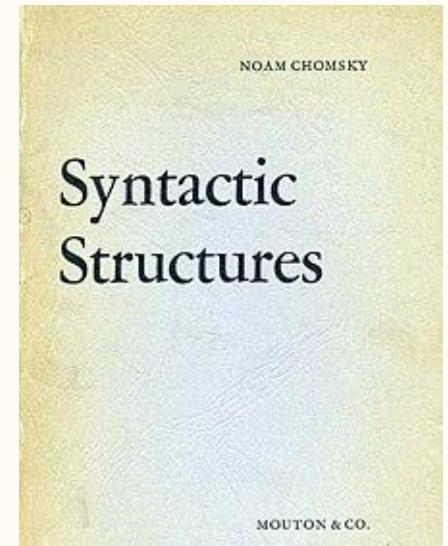
1957 Chomsky

*Syntactic structures
& formal grammars*



- Separated grammar (syntax) and semantics

Colorless green ideas sleep furiously

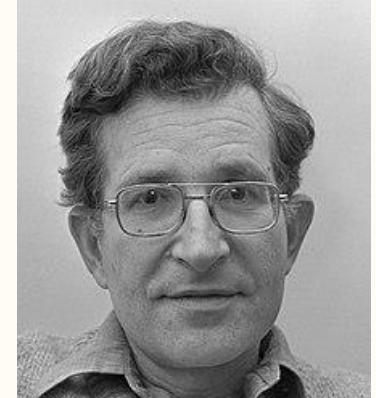


CS and Language Generation

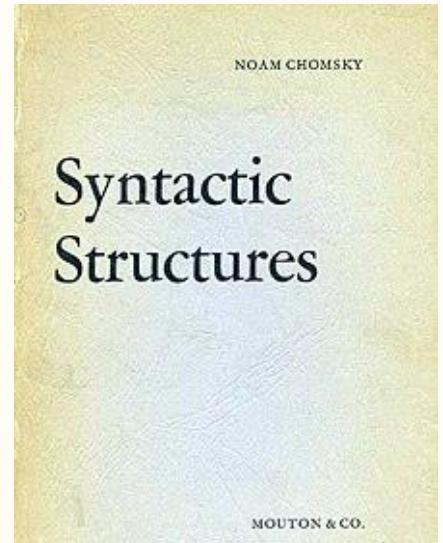
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Syntactic structures & formal grammars



- Separated grammar (syntax) and semantics
 - Colorless green ideas sleep furiously
- Introduced a hierarchy of grammars
- Apart from linguistics also influenced TOC



CS and Language Generation

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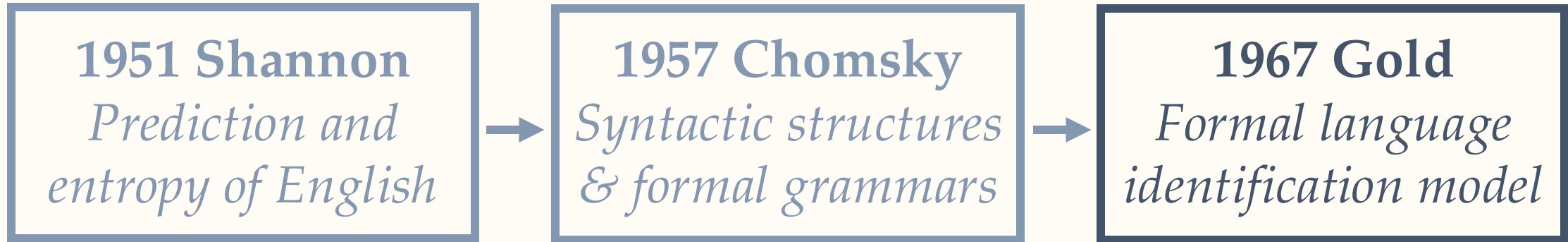
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CS and Language Generation



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1967 Gold
Formal language identification model

I wish to construct a precise model for “able to speak English”...
to investigate theoretically how it can be achieved artificially

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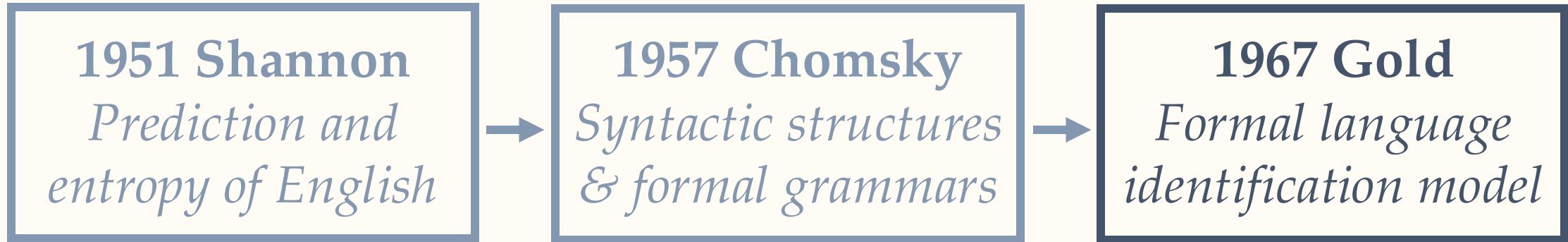
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learning from samples!

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- Laid the groundwork for the celebrated PAC framework [Valiant, 1984] (Turing Award, 2010) 🏅

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Formal language identification model

- Laid the groundwork for the celebrated PAC framework [Valiant, 1984] (Turing Award, 2010) 🏅
- Contains many ideas developed much later in learning theory
 - Learning from samples,
 - Hypothesis class,
 - Two-player online games, and even active learning!

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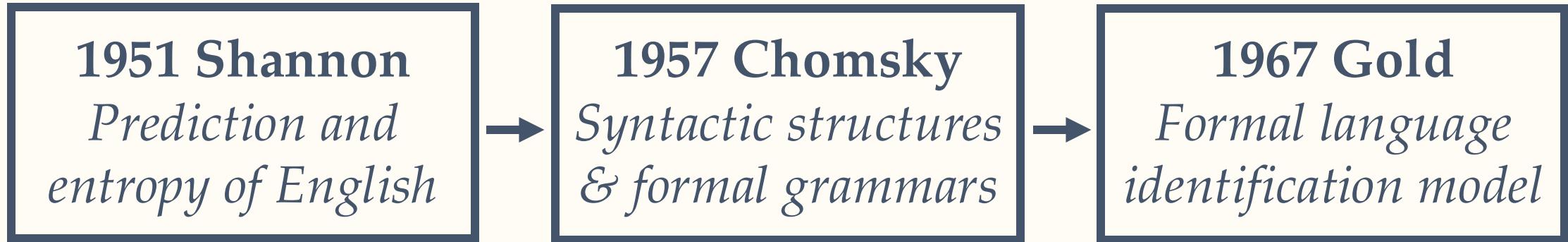


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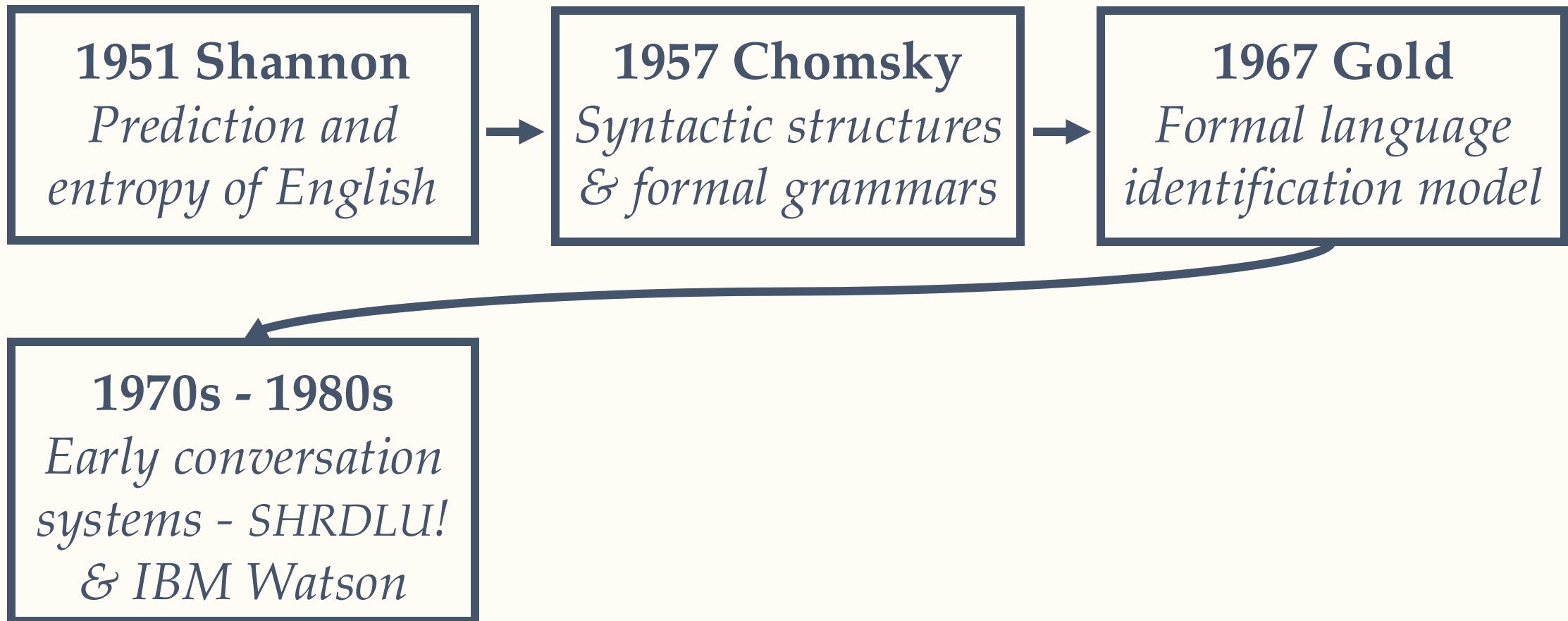
*Formal language
identification model*

- Also had a significant impact in *linguists*
 - Do inductive biases of humans help them learn to speak?
 - Do children need interaction to learn to speak?
 - ...

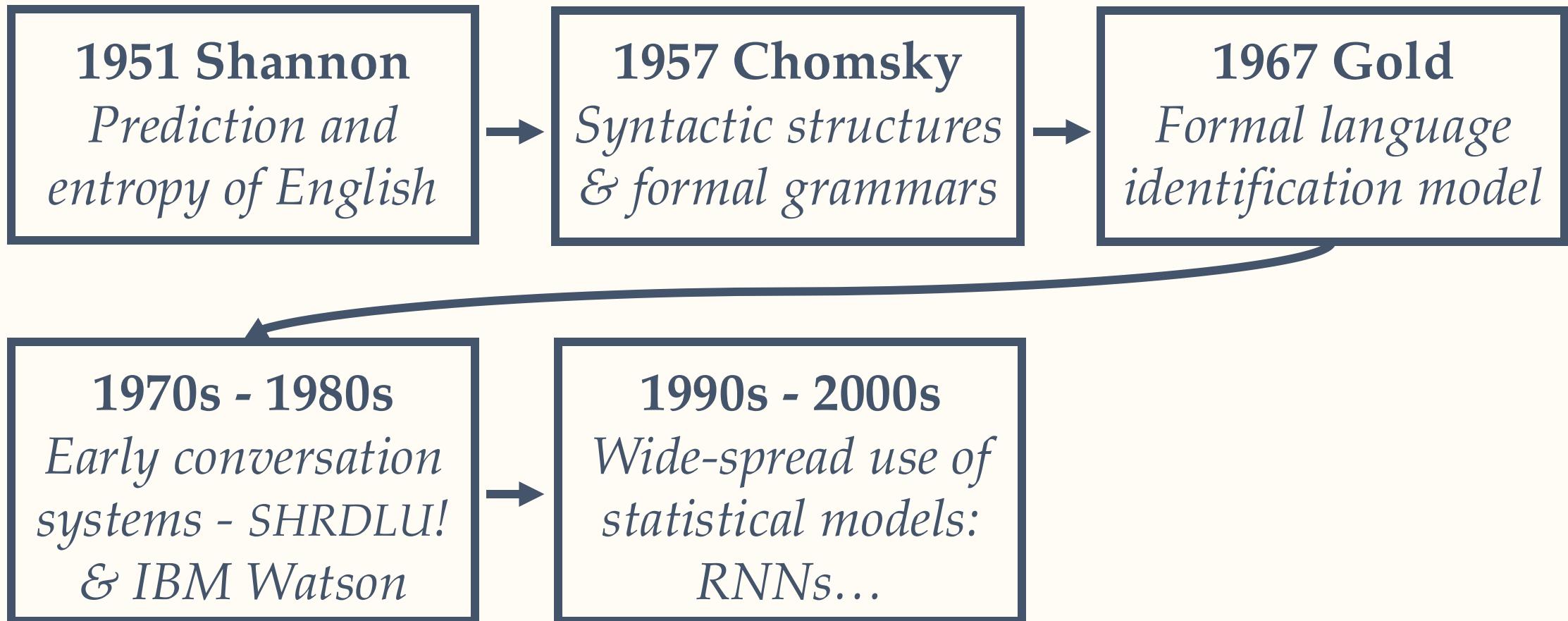
CS and Language Generation



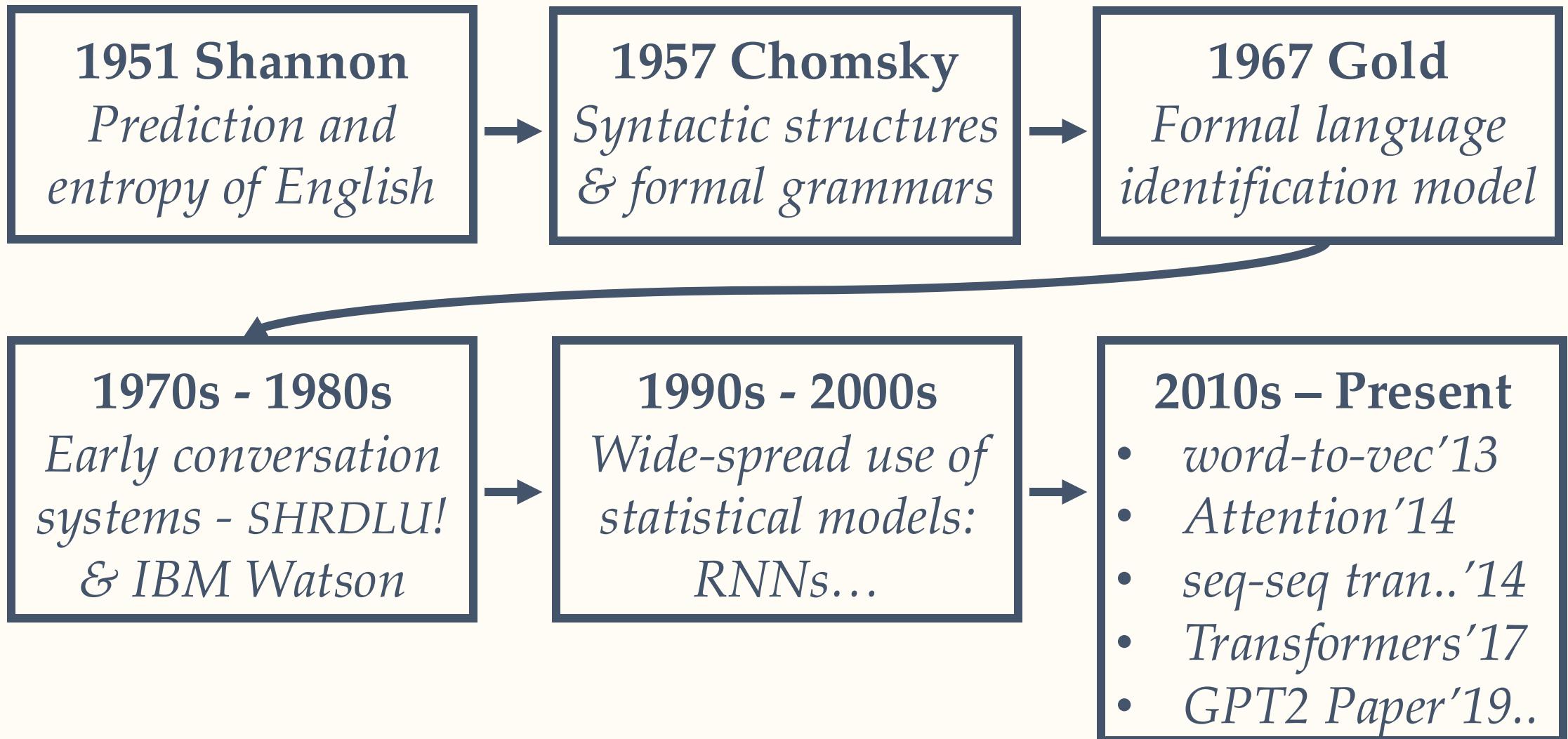
CS and Language Generation



CS and Language Generation



CS and Language Generation



Modern Language Generators – LLMs

Modern Language Generators – LLMs

I am giving a talk about language generation. Can you write something brief (it has to go on a slide) and creative to demonstrate what today's language models are capable of?

Reasoned about language model capabilities for 5 seconds >

Language Generation Magic

From code to poetry, AI weaves words into art. For example:

"In circuits' hum, a dream takes flight—
A digital muse sparks worlds of light.
Infinite stories in a single line,
Where thought and code beautifully entwine."

This is today's language model: creativity, computation, and endless possibilities.

Modern Language Generators – LLMs



Source: UBS and Yahoo Finance

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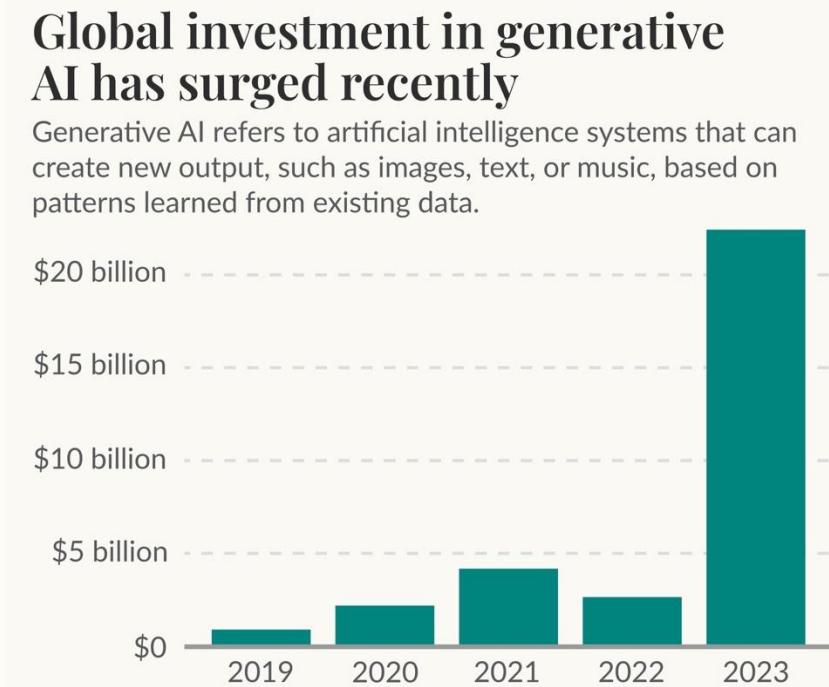
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Modern Language Generators – LLMs



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Modern Language Generators – LLMs



Source: Quid via AI Index (2024) and US Bureau of Labor Statistics (2024)

Hallucinations in LLMs

Give a brief definition of hallucinations in language models.

Reasoned about hallucinations in language models for a few seconds >

Hallucinations in language models refer to instances when the model generates text that appears plausible but is actually fabricated, inaccurate, or not supported by the input or training data.

Hallucinations in LLMs

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Early days ... “simple” errors

How many Rs in the word
strawberry?



There are **two** Rs in the word "strawberry."

Hallucinations in LLMs

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Subsequently ... hallucinations in products

Google still recommends glue for your pizza
/ It's almost like AI answers aren't fully baked!

by [Elizabeth Lopatto](#)

Jun 11, 2024, 6:24 PM EDT

Hallucinations in LLMs

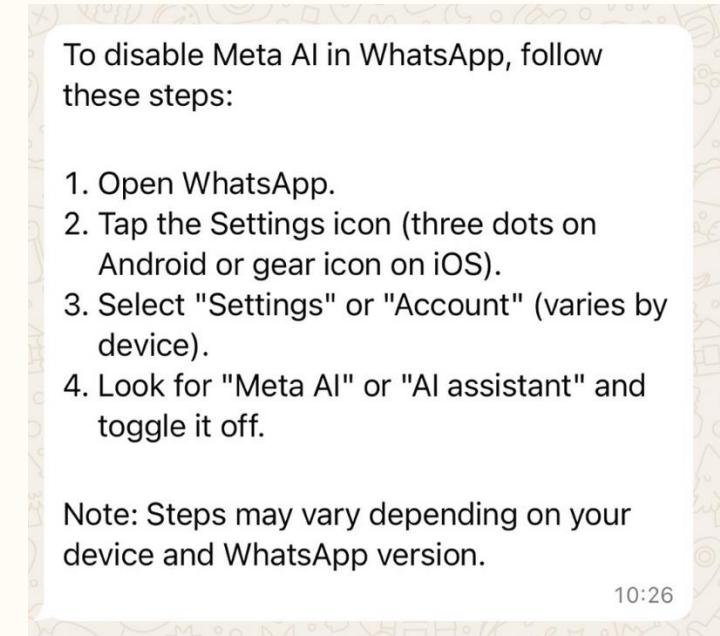
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Source: Twitter / X

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Today, hallucinations are more rare on “common tasks”

- Due to a variety of techniques:
 - chain-of-thought,
 - auxiliary tools (e.g., web search), ...
- Models still hallucinate and make errors on more complex tasks (e.g., proofs, real world tasks, ...)

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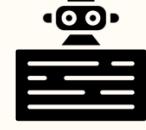
- Due to a variety of techniques:
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- Models still hallucinate and make errors on more complex tasks (e.g., proofs, real world tasks, ...)

Question. Can hallucinations be avoided by better (but “*similar*”) models/training or is fundamental change needed?

Outline of the Talk

1. Motivation: CS and Language Generation
2. Model
3. Overview Our Definitions and Results
 - a. Characterizations I (Generation with Breadth)
 - b. Characterizations II (*Stable* Generation with Breadth)
 - c. Beyond Characterizations
 - d. Learning Curves
4. Overview of *Some* Proofs

A Model of Language Generation



We introduce the model of language generation by Kleinberg and Mullainathan (2024), which builds on [Gold'67] and [Angluin'79]

Language Generation in the Limit

Jon Kleinberg

Departments of Computer Science
and Information Sciene
Cornell University
Ithaca NY

Sendhil Mullainathan

Booth School of Business
University of Chicago
Chicago IL

The Core Task in Language Generation

Language Generation [Kleinberg-Mullainathan, 2024]

Given text from unknown language, produce *valid* and *unseen* text

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Simplifications

- Do not need to learn *all* of the language, okay to learn a subset
- It is a prompt-less model – can be extended

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Model of Language Identification

Language Identification in the Limit [Gold, 1967]

- Domain \mathcal{X} , e.g., $\{a\text{-}z, A\text{-}Z\}^*$ or \mathbb{N}
- Collection of languages $\mathcal{L} = \{L_1, L_2, \dots\}$

Model of Language Identification

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Game between adversary  and learner 

1. Adversary picks target $K = L_{i^*}$

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Adversary has to present a complete enumeration

Example: $K = \mathbb{N}$,

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Example: $K = \mathbb{N}, 2, 4, 6, \dots,$



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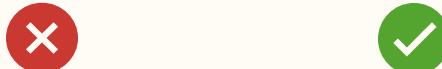
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Example: $K = \mathbb{N}, \quad 2, 4, 6, \dots, \quad 1, 2, 3, \dots \quad \text{and} \quad 2, 4, 6, \dots, 1, 2, 3, \dots$



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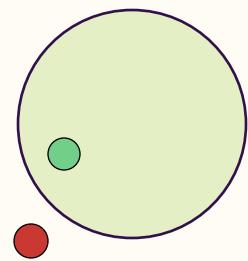
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Learners access:

Is $w \in L_i?$

Membership Query



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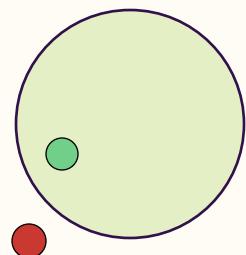
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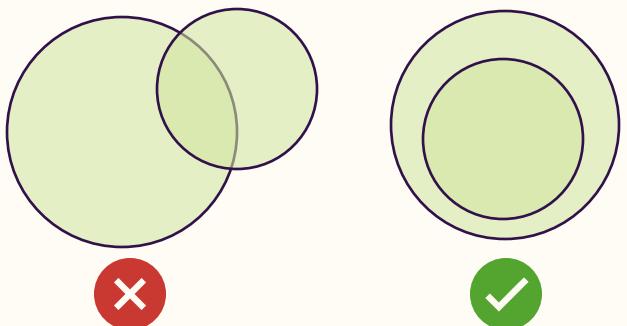
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Membership Query 

Is $L_i \subseteq L_j?$

Subset Query



Model of Language Generation

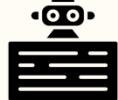
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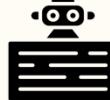
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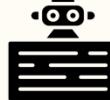
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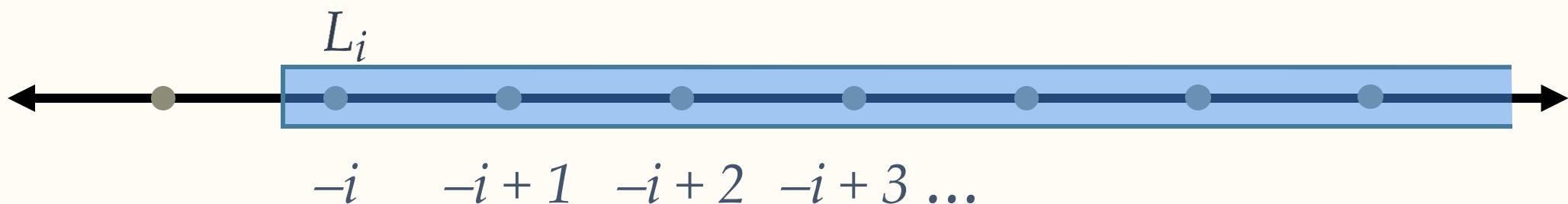
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3. Generator wins if guesses are eventually in K : $K \ni g_t, g_{t+1}, \dots$ after some finite time $t < \infty$

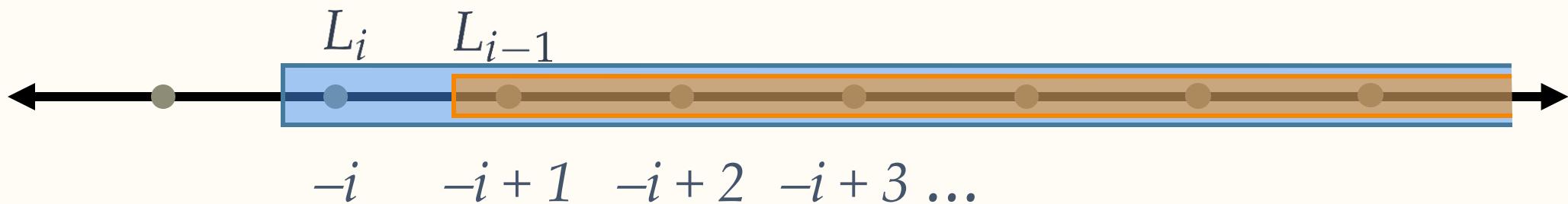
Example [Kleinberg-Mullainathan' 24] [Charikar-Pabbaraju'24]

$\mathcal{L} = \{\mathbb{Z}, L_1, L_2, \dots\}$ where $L_i = \{-i, -i + 1, -i + 2, \dots\}$.



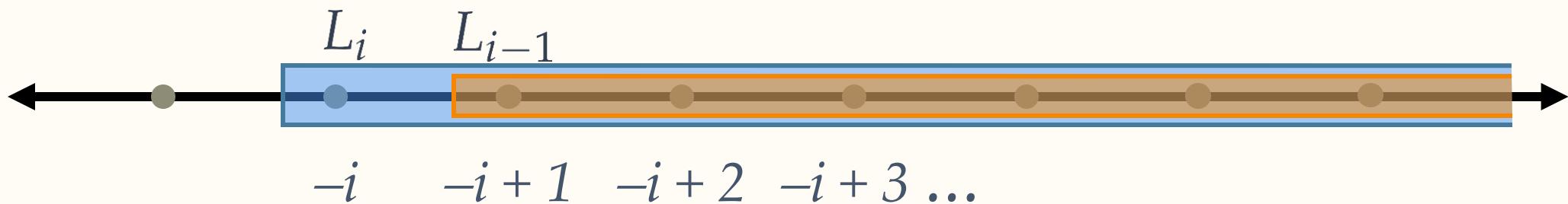
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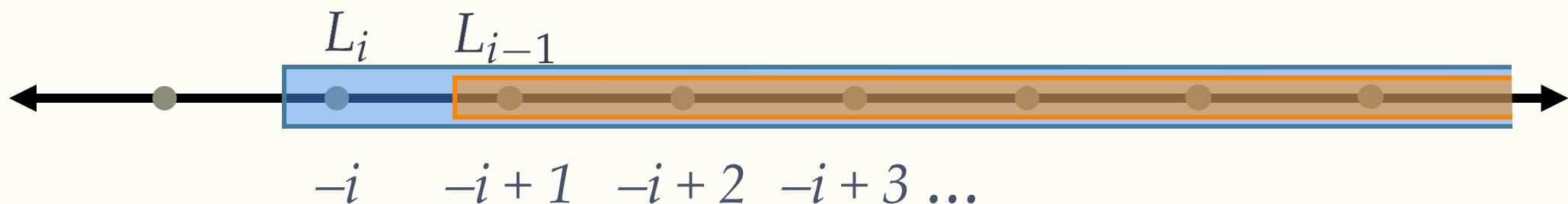
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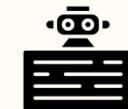
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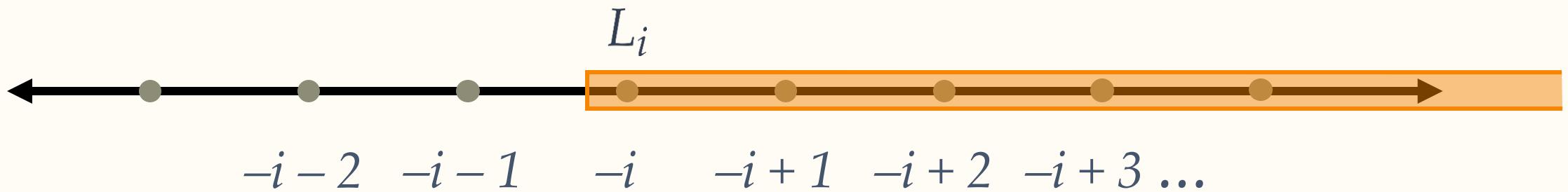
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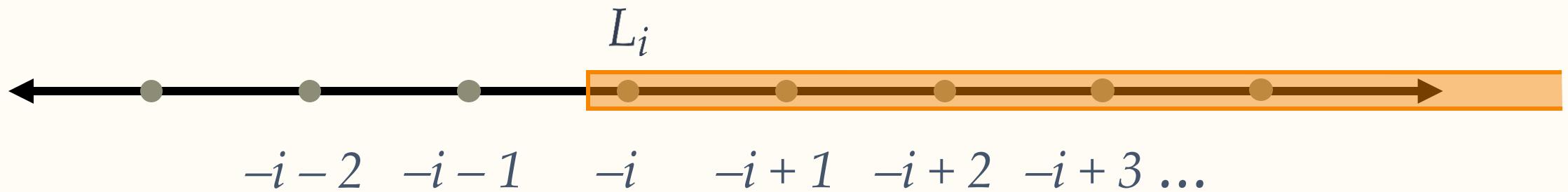
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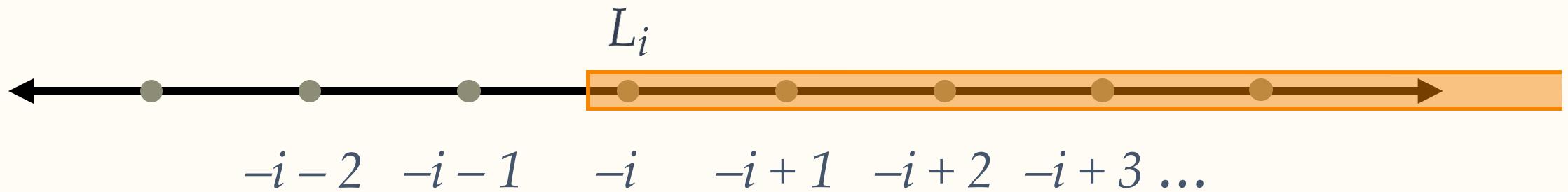
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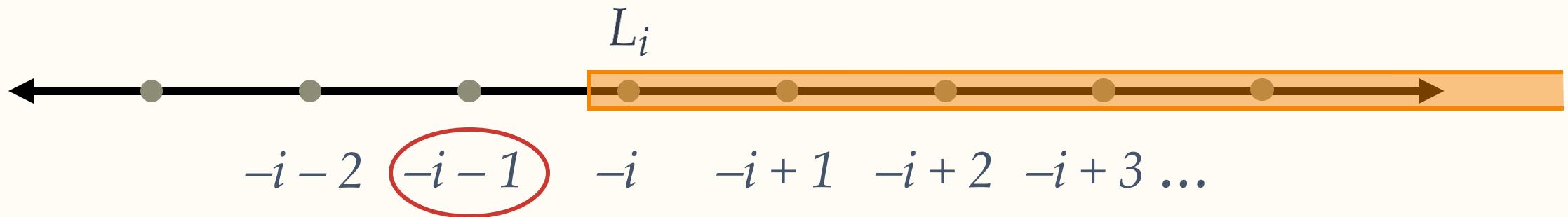


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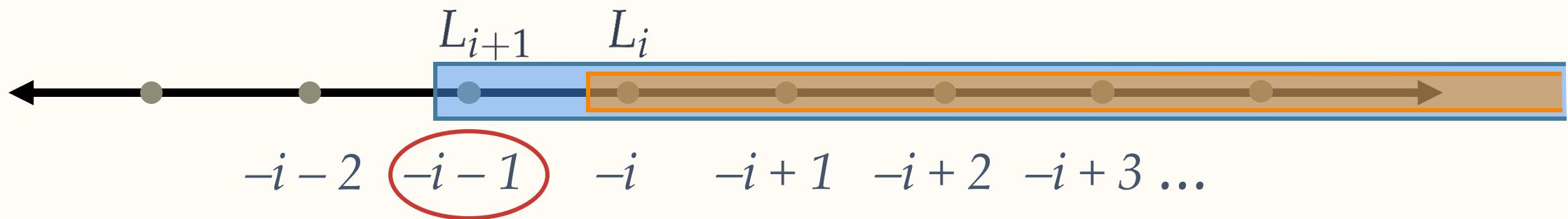


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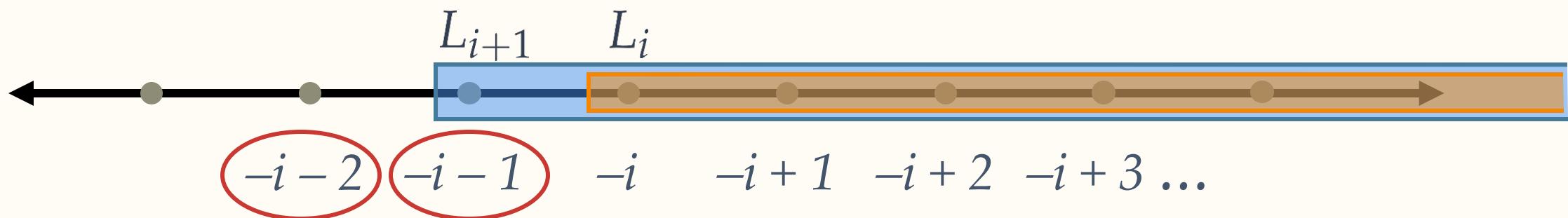


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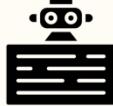
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Model of Language Generation

Language Generation [KM'24]. **Generator** wins if guesses are eventually in K : $K \ni g_t, g_{t+1}, \dots$ after some finite time $t < \infty$

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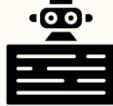
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Details abstracted away: computation, next-token-prediction, ...

These are important ...[Bhattamishra, Ahuja, and Goyal'20] [Sanford, Hsu, Telgarsky'23] [Peng, Narayanan, and Papadimitriou'24] [Chen, Peng, and Wu'24]...

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Question. *Even in an idealized model,* can hallucinations be avoided with better models/training or is fundamental change needed?

Is Language Generation Feasible?

Informal Theorem [Gold'67, Angluin'79, '80] Almost all interesting language collections \mathcal{L} are *not* identifiable

Even regular languages are non-identifiable... simpler than English

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Informal Theorem [Kleinberg-Mullainathan'24] All (countable) language collections \mathcal{L} are generatable

Outline of the Talk

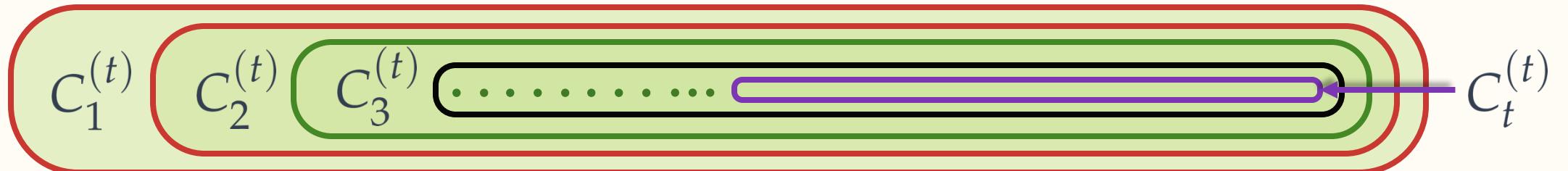
1. Motivation: CS and Language Generation
2. Model
3. Overview Our Definitions and Results
 - a. Characterizations I (Generation with Breadth)
 - b. Characterizations II (*Stable* Generation with Breadth)
 - c. Beyond Characterizations
 - d. Learning Curves
4. Overview of *Some* Proofs

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[KM'24]**'s Generator:** Construct (dynamic) list of *critical languages*

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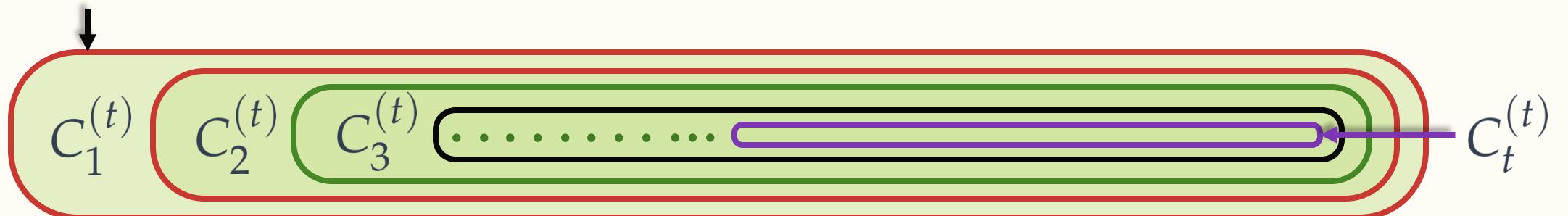


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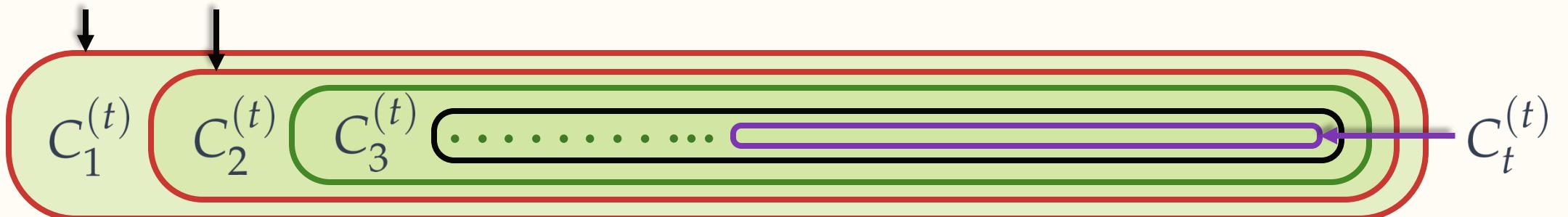


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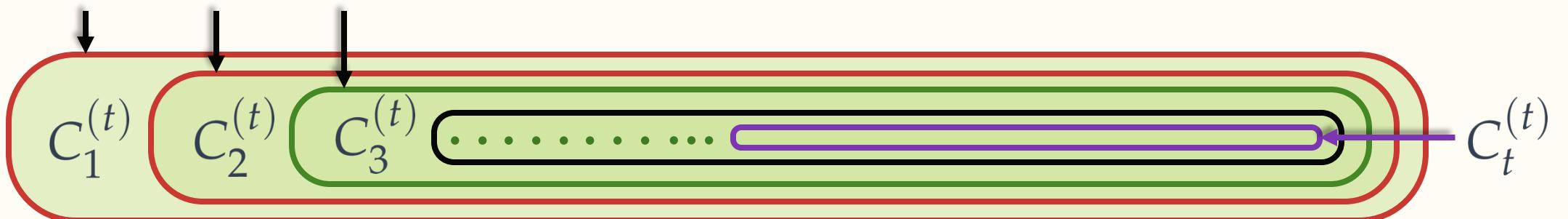


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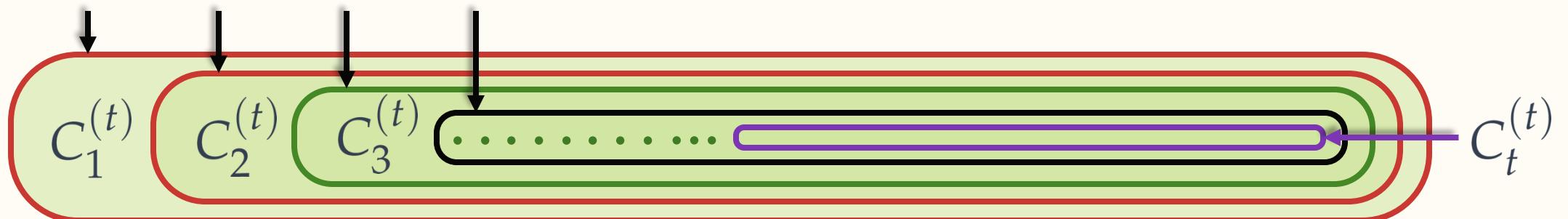


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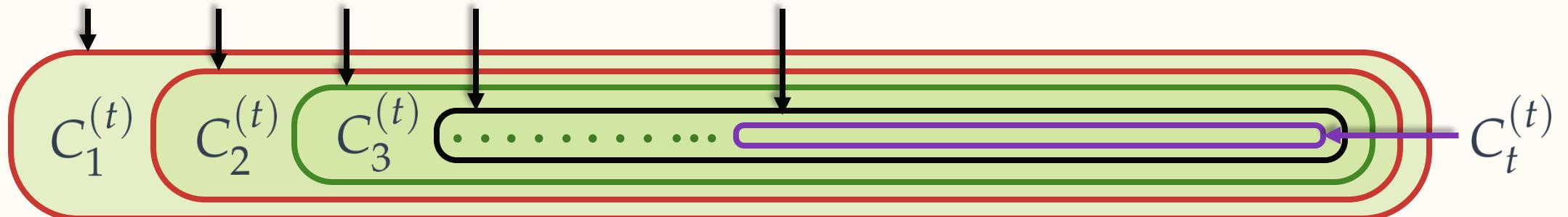


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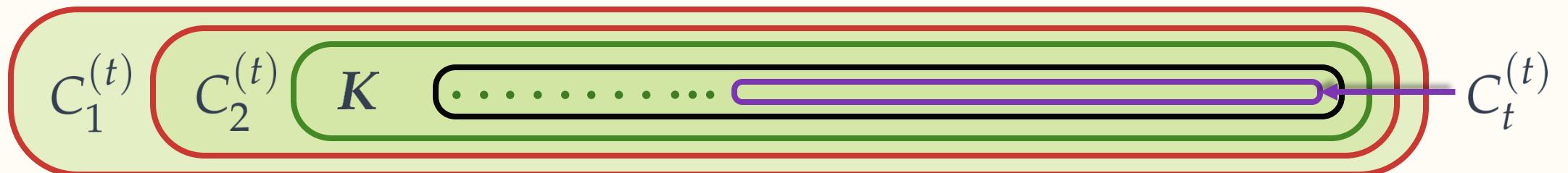
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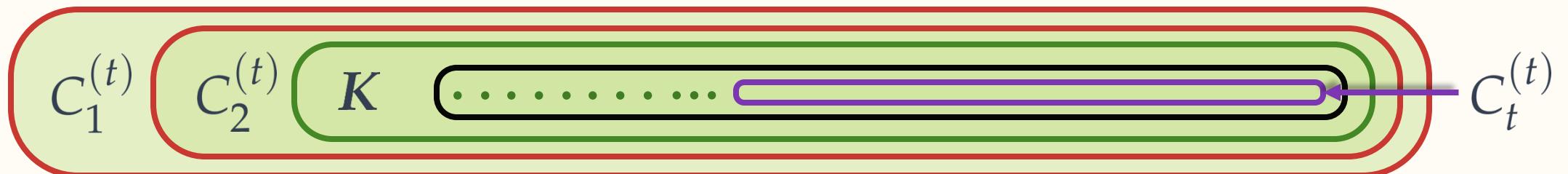
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Adaptation of our main question [KM24]: Can a generator avoid hallucinations while maintaining some notion of “breadth”?

Language Generation

Characterizations I

- *We introduce several notions of breadth*
- *We show that achieving breadth + no hallucinations is impossible for most language collections*

Notions of Coverage or Breadth

Let generator G be a mapping from training data to subsets of the domain, *i.e.*, $G(S)$ is output-set of G trained on S

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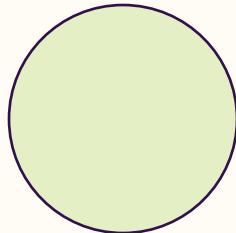
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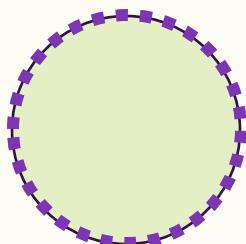


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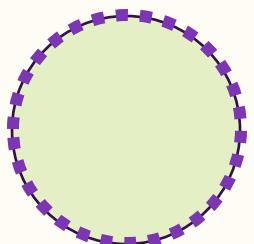


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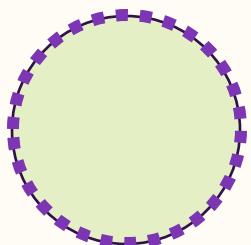
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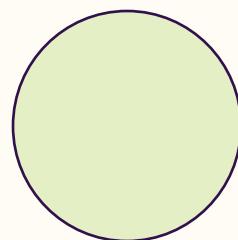
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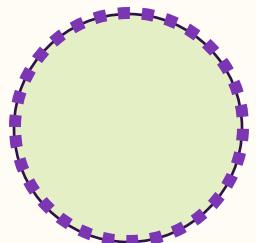


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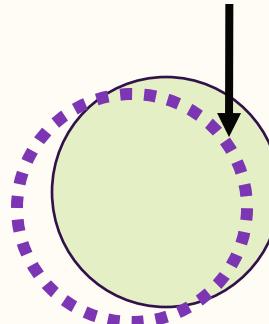
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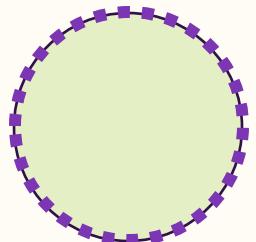


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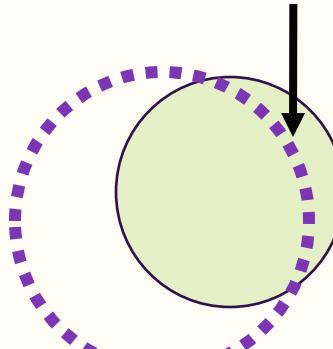
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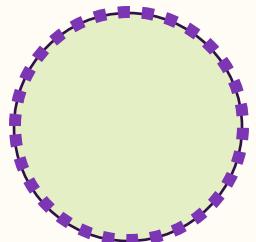


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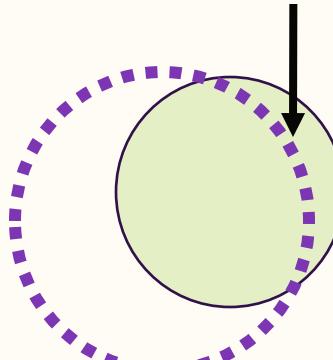
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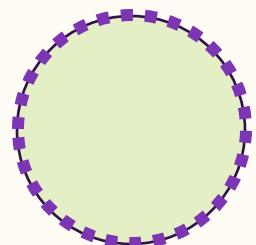
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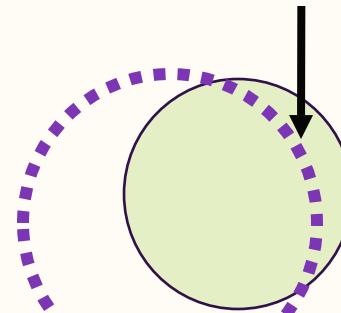
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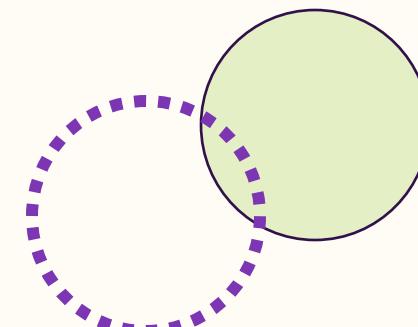
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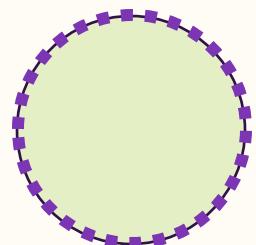


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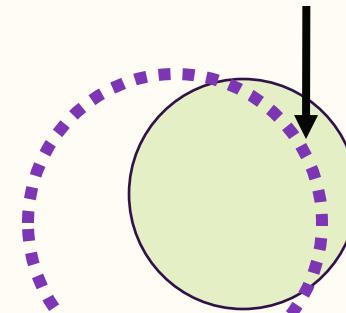
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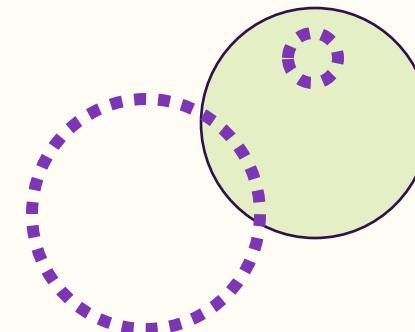
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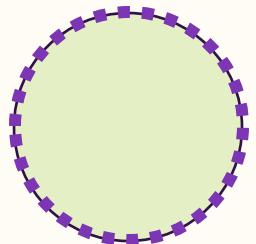


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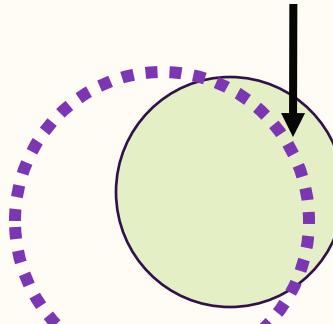
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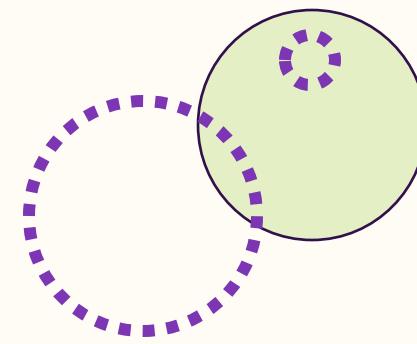
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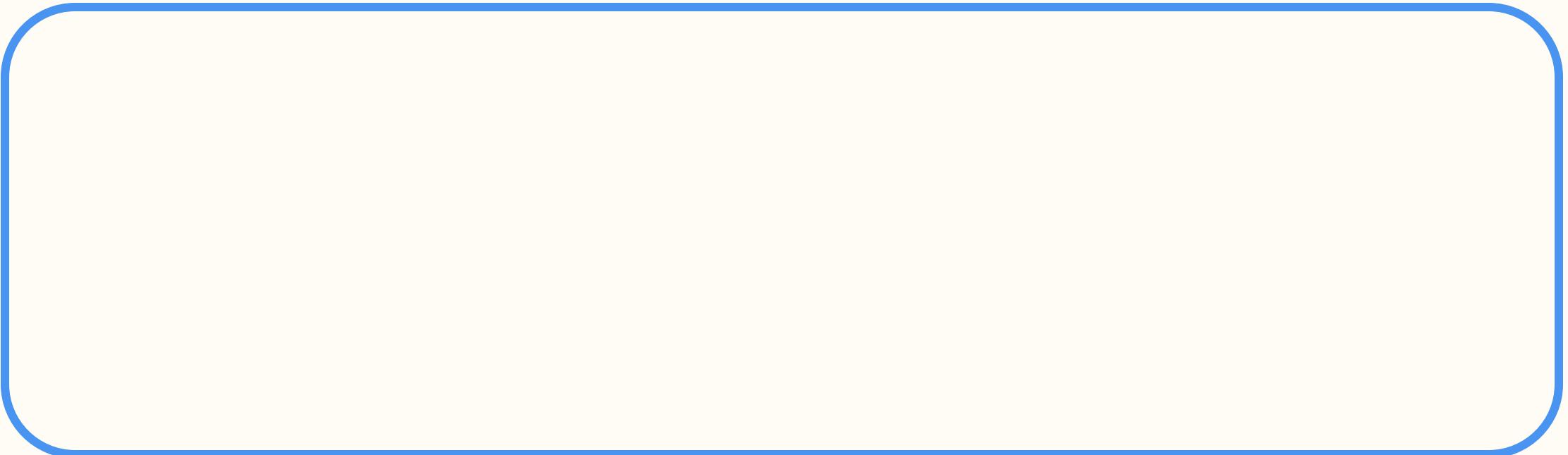
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Consider $K = \mathbb{N}$, $G(S) = \{i, i + 1, \dots\}$ and $G(S) = \{2, 4, 6, \dots\}$

Results for Breadth with *no* Hallucination



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Main Takeaway [Kalavasis, M., Velegkas'24 and '25]. For most interesting language collections, LLMs cannot avoid hallucination while achieving any of these notions of breadth

Technical Vignette: Properties of Breadth

Definitions. A relation P satisfies:

- ▷ Uniqueness if $(\mathcal{G}, L), (\mathcal{G}, L') \in P$, then $L = L'$
- ▷ Finite non-uniqueness if $(\mathcal{G}, L), (\mathcal{G}, L') \in P$, then $|L \Delta L'| < \infty$

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Fact. Exact breadth satisfies uniqueness

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Fact. Exact breadth satisfies uniqueness

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Theorem. [Kalavasis, M., Velekcas, 24] Consider collection \mathcal{L} .

- ▷ If P is unique and \mathcal{L} violates Angluin's condition, or
- ▷ If P is finite-non-unqiue and \mathcal{L} violates weak Angluin's condition,

Then, no generator can satisfy P in the limit

Results for Breadth with *some* Hallucination

	No Hallucinations $ G(S_t) \setminus K = 0$	Finite Hallucinations $ G(S_t) \setminus K < \infty$	Infinite Hallucinations $ G(S_t) \setminus K = \infty$
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Progress. [Kleinberg and Wei, 2025] and [Peale, Raman, Reingold, 2025]

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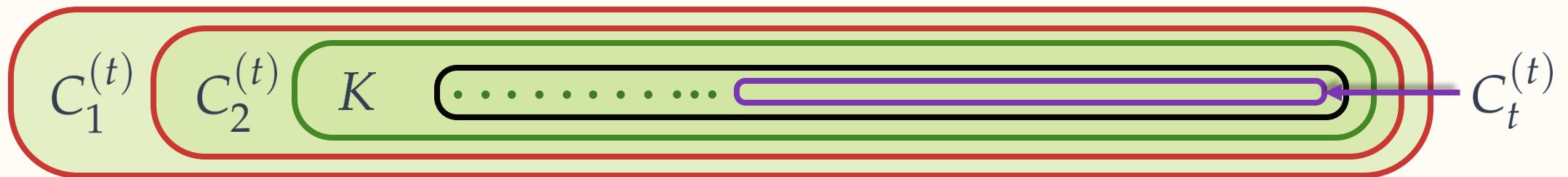
Language Generation

Characterizations II

- *We introduce stable generation*
- *Related to whether the learner can recognize that it has learnt?*
- *Requiring stability makes language generation much harder*

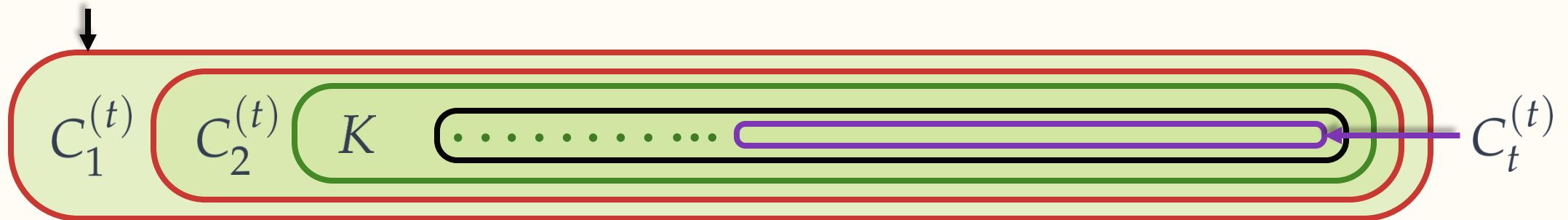
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[KM'24]'s and our generator change output set infinitely often



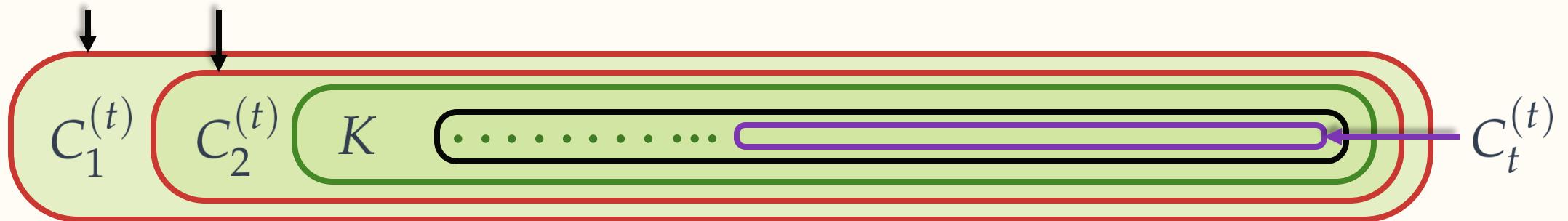
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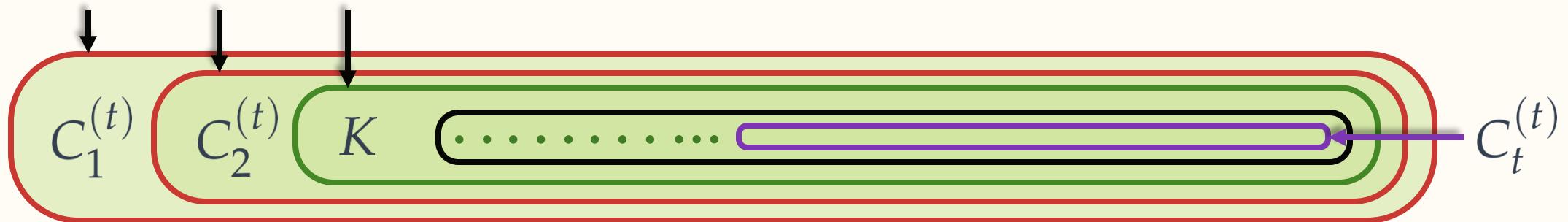
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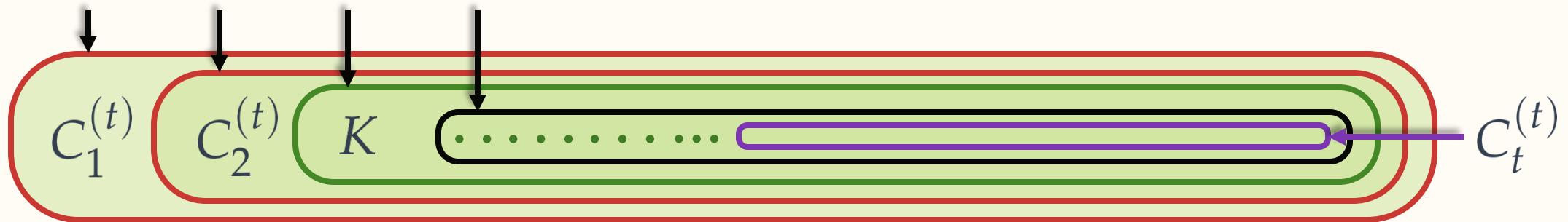
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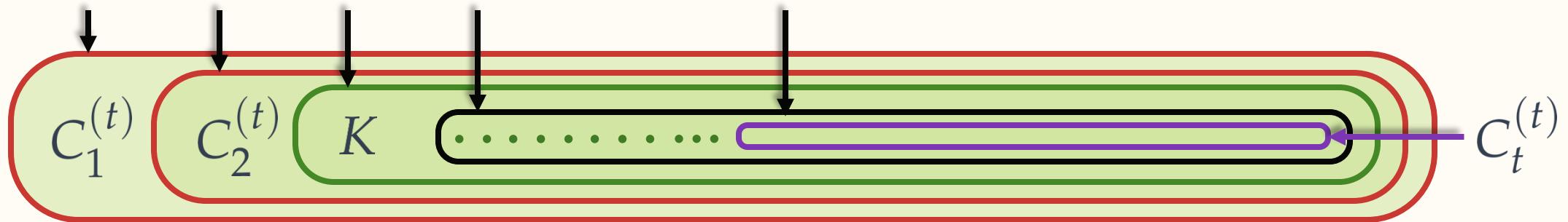
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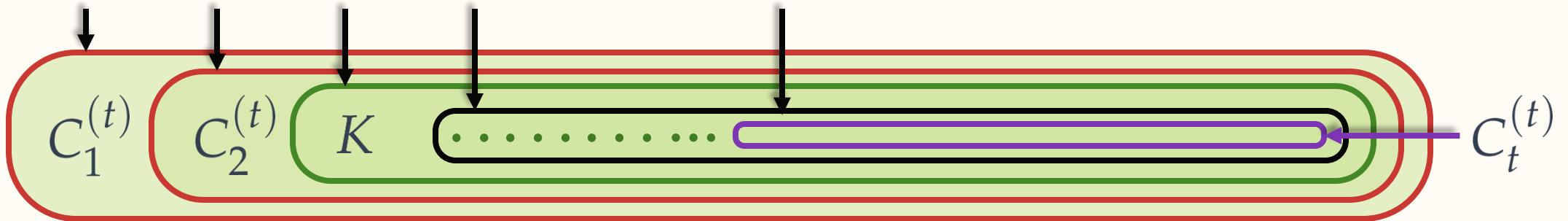
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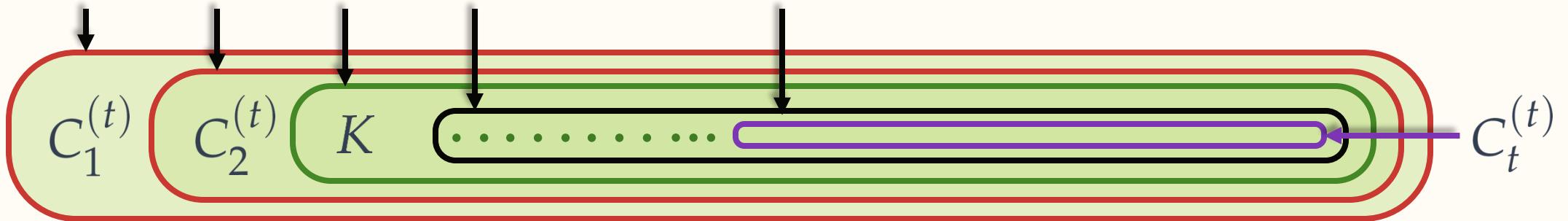


Can generator's stabilize their outputs?

If the generator “knows” it has learnt, then it can stabilize.

Stability in Language Generation

[KM'24]'s and our generator change output set infinitely often



Can generator's stabilize their outputs?

If the generator “knows” it has learnt, then it can stabilize.

Stability. A generator is said to achieve *stability* if for any target K and its enumeration, there is a finite $t < \infty$, after which $G(S_t) = G(S_{t'})$ for all $t' \geq t$

Results with Stability

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Infinite Present Elements $ K \cap G(S_t) = \infty$	Characterization ? (<i>Not all countable collections</i>)	Characterization ?	All Countable Collections

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Language Generation

Beyond Characterizations

Providing the generator with negative examples, enables achieving breadth without hallucinations!

Generation with Negative Examples

Informal Theorem [Gold'67][Kalavasis, M., Velegkas'25] Consider a variation of language generation where adversary enumerates both elements in K and elements outside of K (*negative examples*). Then, *all countable collections \mathcal{L} are generatable with exact breadth.*

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Insights into LLM training

- Perhaps a principled explanation why RLHF is useful
- Does this suggest including negative information in pre-training would be useful?

Generation with Negative Examples

Proxies for negative examples have found to be useful

NEURAL TEXT DEGENERATION WITH UNLIKELIHOOD TRAINING

Sean Welleck^{1,2*}

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Q: *Given high and low-quality data, can one extract negative examples?*

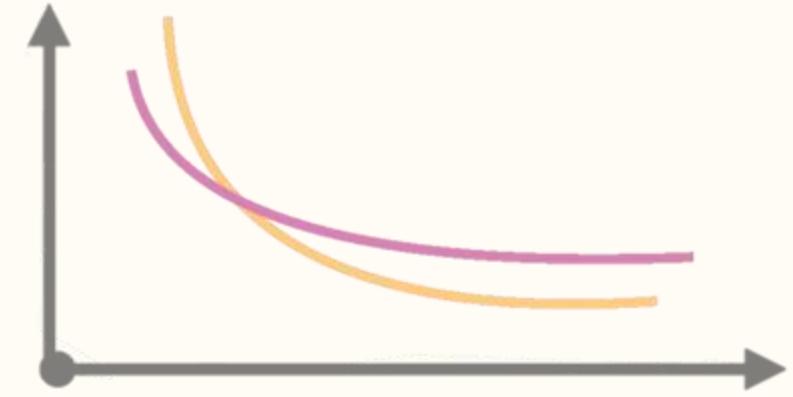
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Learning Curves

We establish universal rates (aka learning curves) for generation with and without breadth



Learning Curves for Generation

How many samples does one need for the generator to generate?

Learning Curves for Generation

How many samples does one need for the generator to generate?

Statistical Model of Language Generation

- Adversary picks a distribution \mathcal{D} supported entirely on $K \in \mathcal{L}$
- Generator gets n i.i.d. samples from \mathcal{D} and outputs $\mathcal{G}(S_n)$
- $\text{Err}_n(\mathcal{G}) := \mathbb{1} \{\mathcal{G}(S_n) \text{ does not satisfy notion of generation}\}$

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Informal Theorem [Kalavasis, M., Velekcas'25] Error either drops exponentially quickly or is arbitrarily slow; where the characterization has tight connections to the online model

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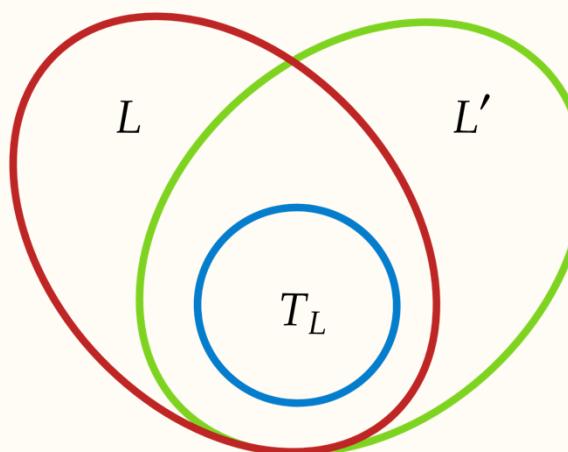


Angluin's Condition

Definition. Language collection \mathcal{L} satisfies Angluin's condition if:

For all $L \in \mathcal{L}$ there is some finite tell-tale subset $T_L \subseteq L$ such that:

For all $L' \neq L$ either $T_L \not\subseteq L'$ or L' is not a proper subset of L



Lower Bound Construction

[Kalavasis, M., Velekcas'24] [Charikar, Pabbaraju'24]

Since \mathcal{L} violates Angluin's condition, there is L^* such that
for all finite subsets $T \subseteq L^*$, there is $L_T \in \mathcal{L}$, $T \subseteq L_T$ and $L_T \subsetneq L^*$

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Either Step2 repeats infinitely often and we enumerate K or we find a language on which \mathcal{G} makes infinitely many mistakes.

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5. Extraction of negative information from available data

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 - ▷ E.g., Clustering, search, algorithmic fairness, robustness...

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Summary

1. TCS can contribute the right *abstractions* for empirical systems
 - ▷ E.g., Clustering, search, algorithmic fairness, robustness...
2. We establish a tension between avoiding hallucinations while achieving breadth for existing language model “frameworks” in a theoretical model by [Kleinberg and Mullainathan’24]
3. *How can theory guide practice?*

Thank you!

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Organized with:



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Thank you!