TTIC 31230, Fundamentals of Deep Learning

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Information Theory Revisited

Shannon's Source Coding Theorem

Avoiding Differential Entropy

Measuring Cross Entropy of an Exponential Softmax

We typically cannot measure cross-entropy for a graphical model.

Although we can train using pseduo-likelihood, it remains unclear how to measure the resulting cross-entropy loss.

One approach is to construct a compression algorithm.

Let $|z_\Phi(y)|$ be the bit length of the compression of y.

$$\Phi^* = \operatorname*{argmin}_\Phi \; E_{y \sim \mathrm{Pop}} \; -\ln P_\Phi(y)$$

$$\Phi^* = \operatorname*{argmin}_\Phi \; E_{y \sim \mathrm{Pop}} \, |z_\Phi(y)| \quad \text{such that } \forall y \quad y = y_\Phi(z_\Phi(y))$$

Sparse Labeling Compression (TZ)

After training a graphical model Φ on semantic segmentations we can code a segmentation y by a sparse segmentation $z_\Phi(y)$ assigning a label to only a small fraction of the pixels.

We then define the decoding $y_\Phi(z_\Phi(y))$ to be the result of running deterministic local search over the labels of the unspecified pixels to find a locally best-scoring full semantic segmentation.

We can define $z_\Phi(y)$ by some heuristic approximation to

$$z_\Phi(y) = \underset{z:\, y_\Phi(z)=y}{\text{argmin}} \quad |z|$$

where $|z|$ is the number of pixels assigned by z.

Entropy and Compressibility

Let S be a finite set.

Let z be a compression (or coding) function assigning a bit string $z(y)$ to each $y \in S$.

The compression function z is called *prefix-free* if for $y' \neq y$ we have that $z(y')$ is not a prefix of $z(y)$.

Null-terminated byte strings are prefix-free bit strings.

Prefix-Free Codes as Probabilities

A prefix-free code defines a binary branching tree — branch on the first code bit, then the second, and so on.

For a prefix-free code, only the leaves of this tree can be labeled with the elements of S.

The code defines a probability distribution on S by randomly selecting branches.

We have $P_z(y) = 2^{-|z(y)|}$. → Problems

The Source Coding (compression) Theorem

(1) There exists a prefix-free code z such that

$$|z(y)| <= (-\log_2 \mathrm{Pop}(y)) + 1$$

and hence

$$E_{y \sim \mathrm{Pop}} |z(y)| \leq H_2(\mathrm{Pop}) + 1$$

(2) For any prefix-free code z

$$E_{y \sim \mathrm{Pop}} \, |z(y)| \geq H_2(\mathrm{Pop})$$

Code Construction

We construct a code by iterating over $y \in S$ in order of decreasing probability (most likely first).

For each y select a code word $z(y)$ (a tree leaf) with length (depth)

$$|z(y)| = \lceil -\log_2 \mathrm{Pop}(y) \rceil$$

and where $z(y)$ is not an extension of (under) any previously selected code word.

Code Existence Proof

At any point before coding all elements of S we have

$$\sum_{y \in \mathrm{Defined}} 2^{-|z(y)|} \leq \sum_{y \in \mathrm{Defined}} \mathrm{Pop}(y) < 1$$

Therefore there exists an infinite descent into the tree that misses all previous code words.

Hence there exists a code word $z(x)$ not under any previous code word with $|z(x)| = \lceil - \log_2 \mathrm{Pop}(y) \rceil$.

Furthermore $z(x)$ is at least as long as all previous code words and hence $z(x)$ is not a prefix of any previously selected code word.

No Better Code Exists

Let z be an arbirtary coding.

$$
\begin{aligned}
E_y\ |z(y)| &= E_y\ -\log_2 P_z(y) \\
&= H_2(\mathrm{Pop}, P_z) \\
&= H_2(\mathrm{Pop}) + KL_2(\mathrm{Pop}, P_z) \\
&\geq H_2(\mathrm{Pop})
\end{aligned}
$$

Huffman Coding

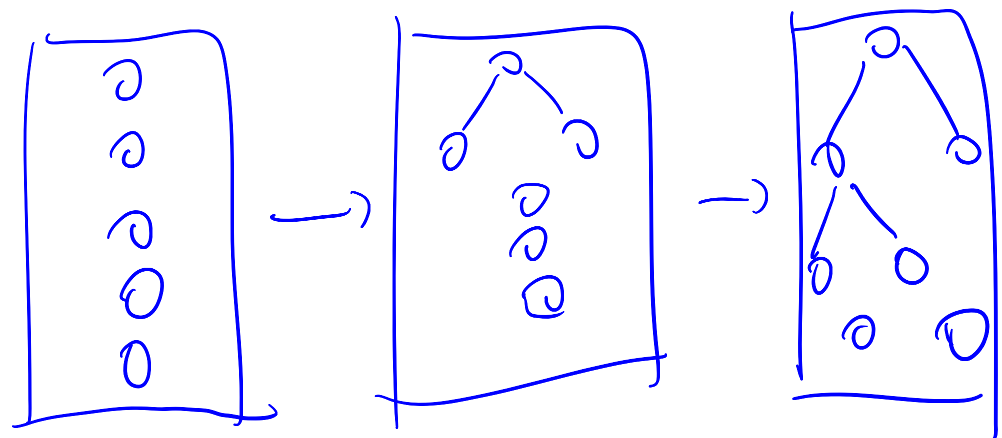
Maintain a list of trees $T_1, \ldots, T_N$.

Inititally each tree is just one root node labeled with an element of S.

Each tree T_i has a weight equal to the sum of the probabilities of the nodes on the leaves of that tree.

Repeatedly merge the two trees of lowest weight into a single tree until all trees are merged.

See algo hw 5



Optimality of Huffman Coding

Theorem: The Huffman code T for Pop is optimal — for any other tree T' we have $H(\text{Pop}, T) \leq H(\text{Pop}, T')$.

Proof: The algorithm maintains the invariant that there exists an optimal tree including all the subtrees on the list.

To prove that a merge operation maintains this invariant we consider any tree containing the given subtrees.

Consider the two subtrees T_i and T_j of minimal weight. Without loss of generality we can assume that T_i is at least as deep as T_j.

Lowering T_j to be the sibling of T_i while raising the old sibling of T_i to T_j's original position brings T_i and T_j together and can only improve the average depth.

$$E_{y \sim \text{Pop}} |c(y)| \leq E_{y \sim \text{Pop}} |c'(y)|$$

Avoiding Differential Entropy

Consider a continuous density $p(x)$. For example

density $\to \infty$

$$p(x) = \frac{1}{\sqrt{2\pi}\ \sigma}\ e^{\frac{-x^2}{2\sigma^2}}$$

$\to$ cross entropy $\to -\infty$

Differential entropy is often defined as

$$H(p) \doteq \int \left(\ln \frac{1}{p(x)}\right) p(x) dx$$

entropy ↓ when gaussian narrow

Differential Entropy Depends on the Choice of Units

$$
\begin{aligned}
H(\mathcal{N}(0,\sigma)) &= +\int \left(\ln(\sqrt{2\pi}\sigma) + \frac{x^2}{2\sigma^2}\right) p(x)dx \\
&= \ln\sigma + \ln\sqrt{2\pi} + \frac{1}{2}
\end{aligned}
$$

$\lim_{\sigma \to 0} = -\infty$

But the numerical value of σ depends on the choice of units.

A distributions on lengths will have a different entropy when measuring in inches than when measuring in feet.

Also, for σ small we get $H(\mathcal{N}(0,\sigma)) < 0$

$\Phi^* = \operatorname{argmin}_{\Phi} E_{y \sim \text{train}} - \ln P_\Phi(y)$

Cannot prevent from this go to infinity

More Problems with Differential Entropy

There are also other problems with continuous entropy and cross-entropy.

- Differential entropy violates the source coding theorem — it takes an infinite number of bits to code a real number.

- Differential entropy violates the data processing inequality that $H(f(x)) \leq H(x)$. For a continuous random variable x under finite continuous entropy we can have $H(f(x)) > H(x)$.

For these reasons it seems advisable to avoid differential entropy and differential cross entropy.

Differential entropy of continuous pdf is actually infinity (Shannon coding theorem) but

Differential KL-divergence is Independent of Units

$$KL(p, q) = \int \left(\ln \frac{p(x)}{q(x)} \right) p(x) dx$$ → meaningful

If x has units of length then $p(x)$ has units of probability mass per length.

In this case $p(x)/q(x)$ is dimensionless.

Avoiding Differential Entropy

To avoid differential entropy we can use a rate-distortion objective.

$$\Phi^* = \operatorname*{argmin}_{\Phi} \; E_{y\sim\mathrm{Pop}} \; -\ln P_\Phi(y) \quad y \text{ discrete}$$

$$\Phi^* = \operatorname*{argmin}_{\Phi} \; E_{y\sim\mathrm{Pop}} \left(\begin{array}{c} -\ln P_\Phi(z_\Phi(y)) \\ +\lambda \mathrm{Dist}(y, y_\Phi(z_\Phi(y))) \end{array}\right) \left\{\begin{array}{l} y \text{ continuous} \\ z \text{ discrete} \end{array}\right.$$

Lossy Compression

Lossy compression combines compression for measuring cross-entropy with distortion for avoiding differential entropy.

$$\Phi^* = \operatorname*{argmin}_\Phi \; E_{y \sim \mathrm{Pop}} \; -\ln P_\Phi(y) \;\; y \text{ discrete}$$

$$\Phi^* = \operatorname*{argmin}_\Phi \; E_{y \sim \mathrm{Pop}} \left(\begin{array}{c} |\tilde{z}_\Phi(y)| \\ +\lambda \mathrm{Dist}(y, y_\Phi(\tilde{z}_\Phi(y))) \end{array} \right) \left\{ \begin{array}{l} y \text{ continuous} \\ \tilde{z} \text{ discrete} \end{array} \right.$$

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