

# Weekly Assignment 2: Breadth-First Search

September 2025

1. Apply the BFS algorithm to the directed graph displayed in Figure 1 using the node with label 2 as source. Specify:

- The content of the queue, after initialization and after each iteration of the loop,
- For any discovered node, its predecessor and the iteration when it was discovered

You may assume that the adjacent vertices of a node are visited in the order of their sequence number.

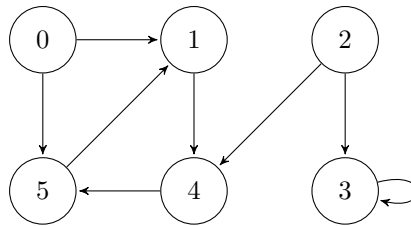


Figure 1: Directed graph

2. Suppose we have a string  $s$  containing lowercase English letters. Design an algorithm, using a stack, that removes all adjacent and equal letters from  $s$ . The final string after removing all adjacent and equal letters should be returned. For example, input  $s = "aacacbbcab"$  leads to output  $"cb"$ . You can use the following functions:
  - `push( $e$ )` takes an element  $e$  and adds it to the end of the stack.
  - `pop` return and removes the last element from the stack.
  - `peek` returns the last element on the stack without removing it from the stack.
  - `toStr` returns the stack in String format (the first element in the stack is also the first element in the string).
3. Consider a directed graph  $G$  and a vertex  $s \in V$ . Let  $C$  be the set of cycles in  $G$  that visit vertex  $s$ . Give an efficient algorithm that takes a graph  $G$  and a vertex  $v$  in  $G$  and prints a cycle from  $C$  in case  $C$  is nonempty, and `false` otherwise. You can use the following functions:
  - `New(queue)` returning an empty queue.
  - `Enqueue( $Q$ ,  $element$ )` adding  $element$  to queue  $Q$ .
  - `Dequeue( $Q$ )` removing and returning  $element$  from  $Q$ .
  - `Print( $w$ )` to print the vertex id

Explain why your algorithm is correct.

4. Given a graph  $G = (V, E)$ , a 2-coloring is a function  $c: V \rightarrow \{blue, red\}$  assigning colors to vertices such that  $(u, v) \in E$  implies  $c(u) \neq c(v)$ , i.e., adjacent vertices have different colors. Let  $G$  be an undirected graph. An odd-length cycle is a sequence  $v_1, v_2, \dots, v_k$  of vertices of  $G$  such that  $(v_i, v_{i+1}) \in E$ , for  $i = 1, \dots, k-1$ ,  $(v_k, v_1) \in E$  and  $k$  is odd. See Figure 2 for an example of a graph with a valid 2-coloring and a graph with an odd-length cycle.

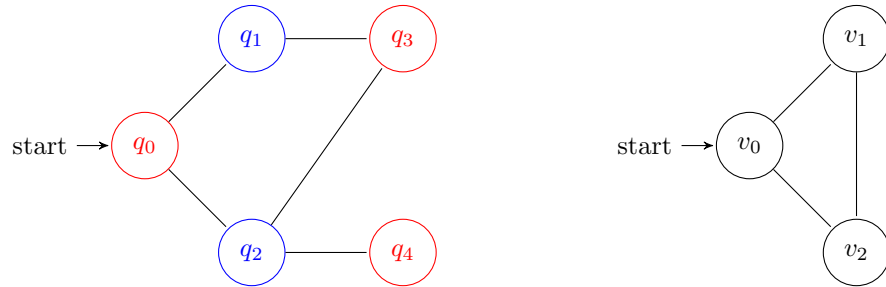


Figure 2: A graph with a 2-coloring and a graph with odd-length cycle  $v_0, v_1, v_2$ .

- a. Show that, if  $G$  has an odd-length cycle, then there is no 2-coloring.
  - b. Assume  $G$  has no odd-length cycles. Use BFS algorithm to find an appropriate 2-coloring for  $G$  and prove that it is correct. Hints:
    - i. Assume that your algorithm does not produce a valid coloring, i.e., there is a pair of adjacent vertices with the same color, and prove the claim by contradiction;
    - ii. Prove and use the following fact: two vertices with the same color must both have even or odd shortest distance from the BFS source.
5. Give an example of a directed graph  $G = (V, E)$ , a source vertex  $s \in V$ , and a set of tree edges  $E_\pi \subseteq E$  such that for each vertex  $v \in V$ , the unique simple path in the graph  $(V, E_\pi)$  from  $s$  to  $v$  is a shortest path in  $G$ , yet the set of edges  $E_\pi$  cannot be produced by running BFS on  $G$ , no matter how the vertices are ordered in each adjacency list.