

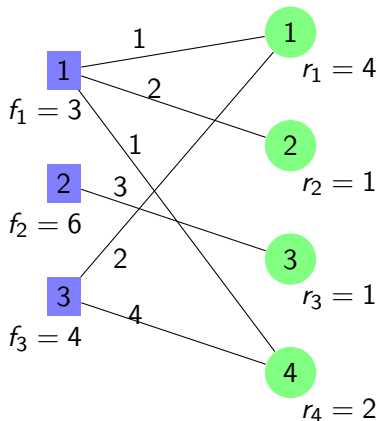
Approximation Algorithms for the Fault-Tolerant Facility Placement Problem

Li Yan and Marek Chrobak

Computer Science
University of California Riverside

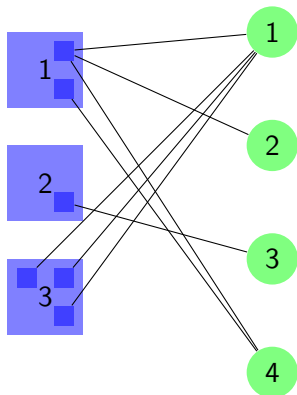
- ① Problem Definition
- ② Related Work
 - The Uncapacitated Facility Location problem (UFL)
 - The Fault-tolerant Facility Location problem (FTFL)
 - The Fault-tolerant Facility Placement problem (FTFP)
- ③ General Approach
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The Fault-tolerant Facility Placement Problem (FTFP)



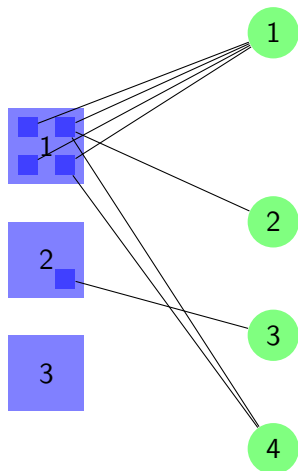
Instance

A Feasible Integral Solution



Cost is $2f_1 + f_2 + 3f_3 + d_{11} + d_{12} + 2d_{14} + d_{23} + 3d_{31} = 38$.

An Optimal Integral Solution



Cost is $4f_1 + 1f_2 + 0f_3 + 4d_{11} + d_{12} + 2d_{14} + d_{23} = 29$.

The Fault-Tolerant Facility Placement Problem (FTFP)

Given

- \mathbb{F} , a set of sites can have facilities built,
- \mathbb{C} , a set of clients with demands,
- r_j , demand for client j ,
- f_i , cost to build one facility at site i ,
- d_{ij} , cost to connect one demand from client j to facility at site i . Distances form a metric.

Find

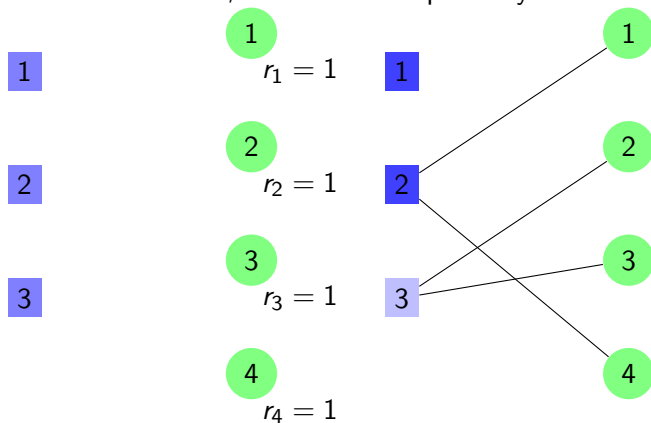
- the number of facilities to build at each site,
- the number of connections between site i and client j .

Goal: Minimize the total cost of opening facilities and connecting clients.

- The Uncapacitated Facility Location problem (UFL), all $r_j = 1$.
- The Fault-tolerant Facility Location problem (FTFL), each site can have at most one facility.

The Uncapacitated Facility Location Problem (UFL)

All demands are 1, each site can open only one facility.

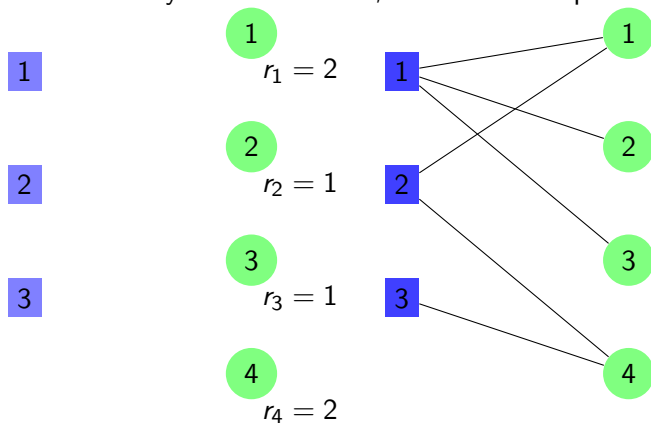


Instance

Solution

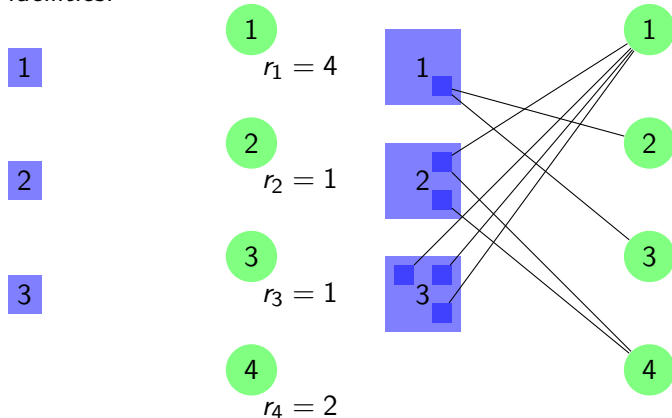
The Fault-tolerant Facility Location Problem (FTFL)

Demands may be more than 1, each site can open only one facility.



The Fault-tolerant Facility Placement Problem (FTFP)

Demands may be more than 1, each site can open multiple facilities.



Instance

Solution

- For the UFL problem,
 - First $O(1)$ -approximation result due to Shmoys, Tardos and Aardal.
 - Best known approximation ratio 1.488 due to Li, combining Byrka's LP-rounding algorithm and Jain, Mahdian and Saberi's greedy algorithm.
- For the FTFP problem, best known approximation ratio 1.7245 due to Byrka, Srinivasan and Swamy, using LP-rounding.

For the FTFP problem, we show

- An LP-rounding algorithm with approximation ratio 1.575.
- A reduction from FTFP to FTFL.
- Our approximation ratio for FTFP matches the best known LP-based approximation ratio for UFL.

General Approach

- Generalize the LP-rounding algorithms to the FTFP problem with fault-tolerant requirement.
- Main techniques:
 - Demand Reduction.
 - Adaptive Partition.

The LP Formulation for FTFP

- y_i represent the number of facilities built at site i .
- x_{ij} represent the number of connections from client j to facilities at site i .

$$\begin{aligned} \text{minimize} \quad & \sum_{i \in \mathbb{F}} f_i y_i + \sum_{i \in \mathbb{F}, j \in \mathbb{C}} d_{ij} x_{ij} & (1) \\ \text{subject to} \quad & y_i - x_{ij} \geq 0 & \forall i \in \mathbb{F}, j \in \mathbb{C} \\ & \sum_{i \in \mathbb{F}} x_{ij} \geq r_j & \forall j \in \mathbb{C} \\ & x_{ij} \geq 0, y_i \geq 0 & \forall i \in \mathbb{F}, j \in \mathbb{C} \end{aligned}$$

$$\begin{aligned} \text{maximize} \quad & \sum_{j \in \mathbb{C}} r_j \alpha_j & (2) \\ \text{subject to} \quad & \sum_{j \in \mathbb{C}} \beta_{ij} \leq f_i & \forall i \in \mathbb{F} \\ & \alpha_j - \beta_{ij} \leq d_{ij} & \forall i \in \mathbb{F}, j \in \mathbb{C} \\ & \alpha_j \geq 0, \beta_{ij} \geq 0 & \forall i \in \mathbb{F}, j \in \mathbb{C} \end{aligned}$$

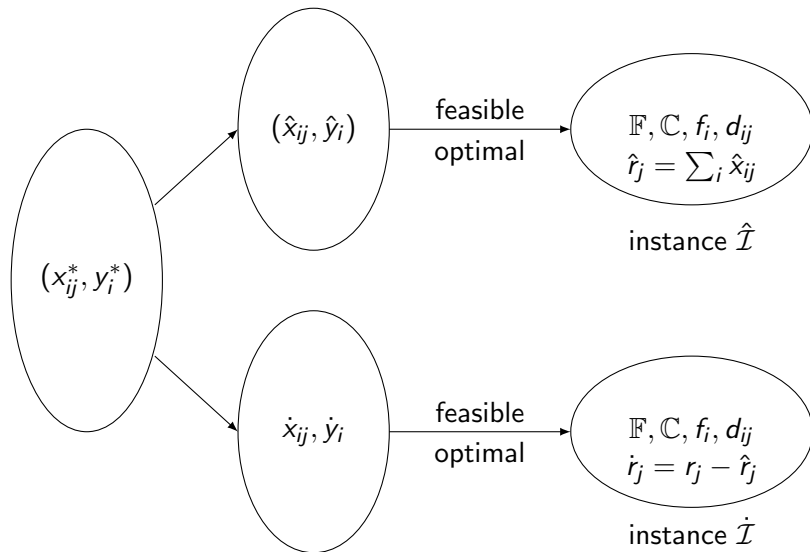
- Demand Reduction.
- Adaptive Partition.

- Reduce a general FTFP instance to a restricted FTFP instance with $r_j \leq |\mathbb{F}|$ for all clients j .
- Solving LP to obtain $(\mathbf{x}^*, \mathbf{y}^*)$.
- Round down $(\mathbf{x}^*, \mathbf{y}^*)$ to obtain integral part $(\hat{\mathbf{x}}, \hat{\mathbf{y}})$. Define $\hat{r}_j = \sum_i \hat{x}_{ij}$.
- The rest form fractional part $(\dot{\mathbf{x}}, \dot{\mathbf{y}})$. Define $\dot{r}_j = r_j - \hat{r}_j$.
- Both parts are feasible and optimal for their respective FTFP instances $\hat{\mathcal{I}}$ and $\dot{\mathcal{I}}$.

Claim

$\dot{r}_j \leq |\mathbb{F}|$ for all clients j in $\dot{\mathcal{I}}$.

Diagram for Demand Reduction



Theorem

Given any ρ -approximation algorithm \mathcal{A} for the restricted FTFP problem with $r_j \leq |\mathbb{F}|$, if ρ is an upper bound on comparing algorithm's cost and the optimal fractional solution's cost, then we have a ρ -approximation algorithm for the general FTFP problem.

Corollary

Using 1.7245-approximation algorithm for FTFL, can have a 1.7245-approximation algorithm for FTFP.

Adaptive Partition

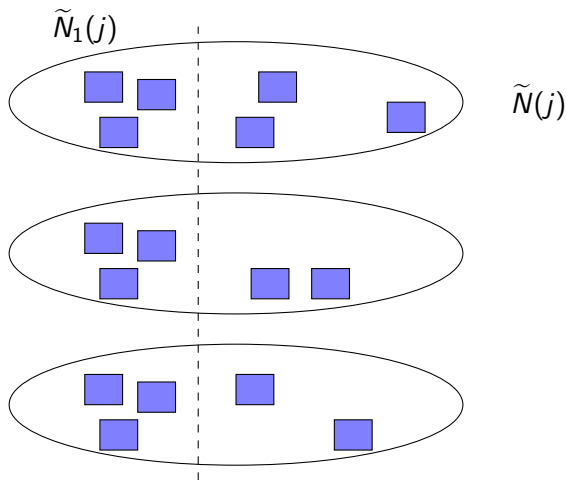
- Begin with a fractional complete solution (\mathbf{x}, \mathbf{y}) .
- In the partitioned solution,
 - Each site i has facilities μ .
 - Each client j has r_j demand points ν .
 - Each facility μ has fractional opening \bar{y}_μ .
 - Each demand point connects to each facility with value $\bar{x}_{\mu\nu}$.
- The partitioned solution $(\bar{\mathbf{x}}, \bar{\mathbf{y}})$ satisfies a number of properties.
 - y_i^* distributed among facilities at site i ,
 - x_{ij}^* distributed among sibling demands of client j ,
 - $\bar{x}_{\mu\nu} = \bar{y}_\mu$ or 0 (completeness),
 - Each demand ν is assigned to a primary demand κ with a low cost.

Two phases

- Phase 1, the partitioning phase, define demands and allocate facilities.
- Phase 2, the augmenting phase, allocate additional facilities to make total connection value unit.

Phase 1, Step 1

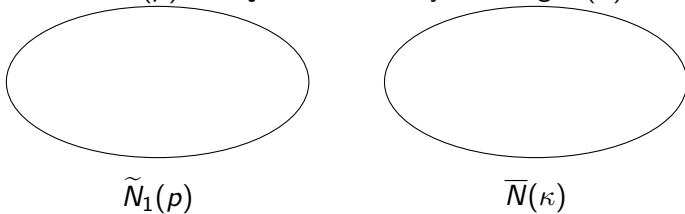
For each client j with residual demand $\bar{r}_j > 0$, arrange neighboring facilities from near to far. The nearest few with total connection value 1 defines $\bar{N}_1(j)$.



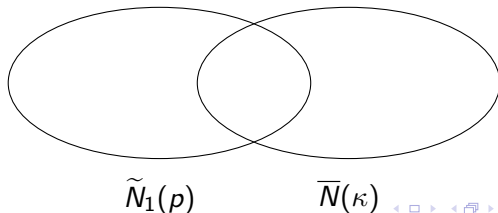
Phase 1, Step 2

Select client p such that the sum of the average distance to $\tilde{N}_1(p)$ and α_p^* is minimized. Now we have two cases to proceed.

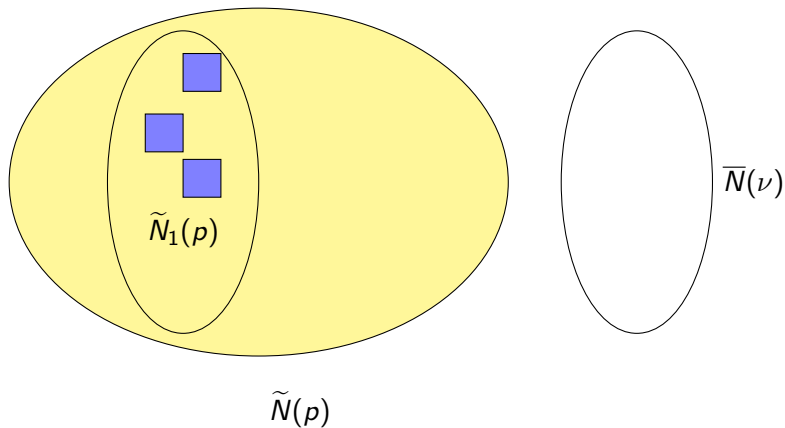
- Case 1: $\tilde{N}_1(p)$ is disjoint from every existing $\overline{N}(\kappa)$.



- Case 2: $\tilde{N}_1(p)$ overlaps with some $\overline{N}(\kappa)$.

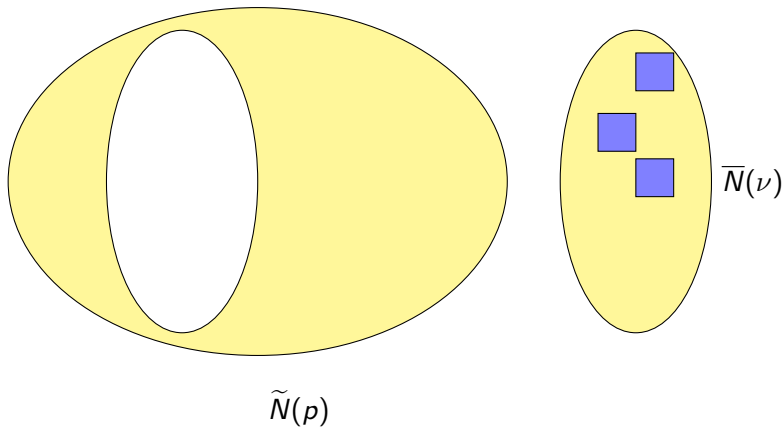


Phase 1, Step 2 (Cont. Case 1)



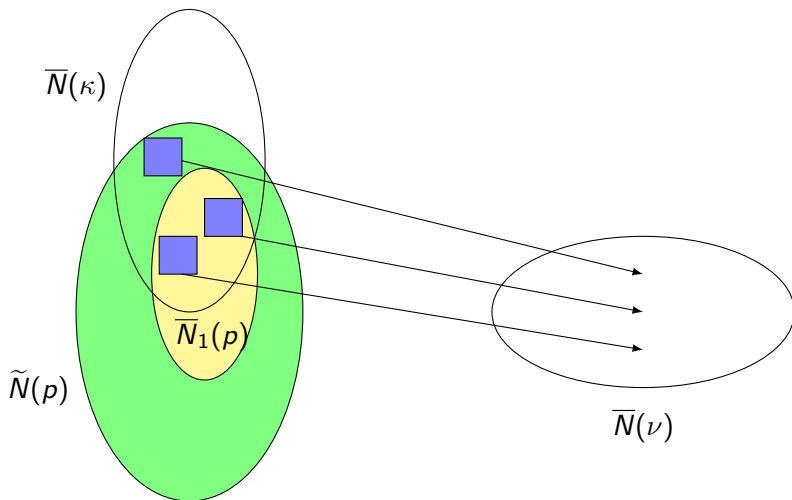
Phase 1, Step 2 (Cont. Case 1)

All facilities in $\tilde{N}_1(p)$ moved to $\overline{N}(\nu)$.



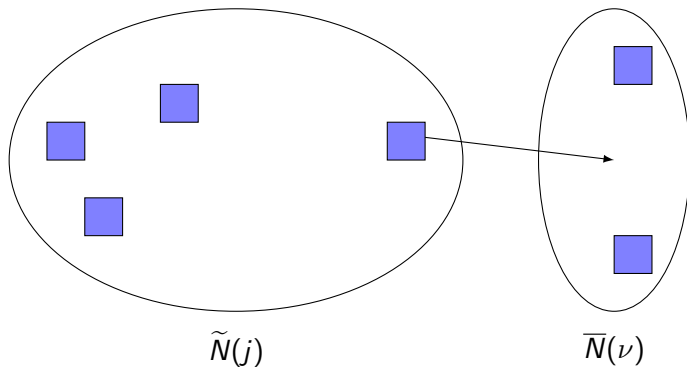
Phase 1, Step 2 (Cont. Case 2)

Move all overlapping facilities in $\tilde{N}(p) \cap \overline{N}(\kappa)$ into $\overline{N}(\nu)$.

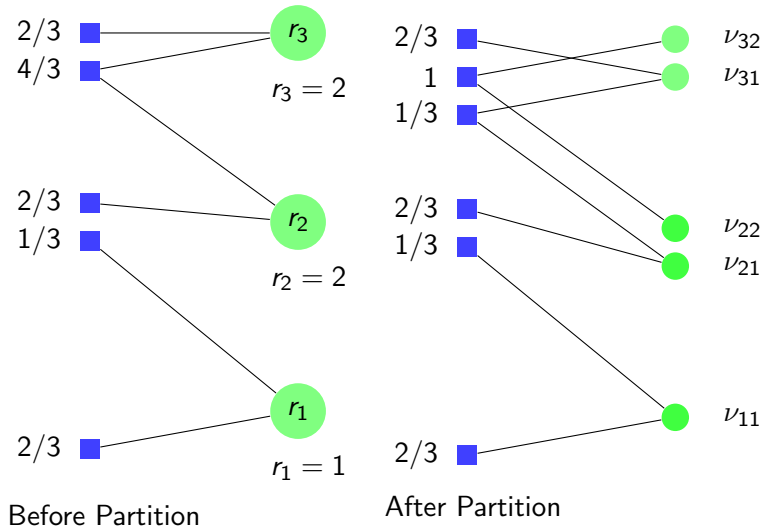


Phase 2

Add facilities from $\tilde{N}(j)$ to $\overline{N}(\nu)$ until total connection value is 1.



An Example of Partition



Properties of Partition

Properties

- Each demand ν assigned to a primary demand κ with overlapping $\overline{N}(\nu)$ and $\overline{N}(\kappa)$.
- For sibling demands ν_1 and ν_2 , $\overline{N}(\nu_1) \cup \overline{N}(\kappa_1)$ is disjoint from $\overline{N}(\nu_2) \cup \overline{N}(\kappa_2)$.
- For a certain cost specified by the approximation algorithm, κ always have a lower cost compared to the assigned ν .

Implication

- The fractional solution can be rounded to a fault-tolerant integral solution.
- The cost of the integral solution can be approximated.

A 3-approximation Algorithm

Given $(\bar{\mathbf{x}}, \bar{\mathbf{y}})$, rounded by

- For each primary κ , choose facility μ in neighborhood with probability \bar{y}_μ .
- For each non-primary ν , connects to $\phi(\kappa)$, the facility chosen in the primary's neighborhood.

The rounded solution satisfies fault-tolerant requirement.

The rounded solution has cost at most $3LP^*$.

- Facility cost is at most F^* .
- For each demand ν , connection cost is at most $\sum_{\mu \in \bar{N}(\nu)} d_{\mu\nu} \bar{x}_{\mu\nu} + 2\alpha_\nu^*$.

A 1.736-approximation Algorithm

- Change in rounding:
 - For facilities μ not in any $\bar{N}(\kappa)$, round independently.
 - each non-primary ν uses nearest neighboring facility if one is open, else use $\phi(\kappa)$.
- The expected connection cost for ν now reduced to $\sum_{\mu \in \bar{N}(\nu)} d_{\mu\nu} \bar{x}_{\mu\nu} + 2/e \cdot \alpha_{\nu}^*$.

Refined Partition for 1.575-approximation

Properties

- $\overline{N}(\nu)$ consists of $\overline{N}_{\text{cls}}(\nu)$ and $\overline{N}_{\text{far}}(\nu)$ and they are disjoint.
- $\overline{N}_{\text{cls}}(\nu)$ overlaps with $\overline{N}_{\text{cls}}(\kappa)$.
- For siblings ν_1, ν_2 , $\overline{N}_{\text{cls}}(\kappa_1) \cup \overline{N}(\nu_1)$ disjoint from $\overline{N}_{\text{cls}}(\kappa_2) \cup \overline{N}(\nu_2)$.
- cost of κ is smaller than cost of ν .

Construction of Partition

- Allocation.
- Augmentation.

1.575-approximation

- Use Byrka's rounding.
- ν uses only facilities in $\overline{N}(\nu) \cup \overline{N}_{\text{cls}}(\kappa)$. Thus no two sibling conflict.
- Cost analysis is similar to Byrka's for UFL.

The End.