

The *standard* trie

- The *trie*

- Trie:

- The term **trie** comes from (information) **retrieval**.
- Following the etymology, the inventor, **Edward Fredkin**, pronounces it as **"tree"**.
- However, it is often pronounced as **"try"**.

- What is a **trie**:

- The **trie data structure (abstract data type)** is an **specialized (very efficient)** implementation of an **(ordered) index** for **text based keys**

- The *standard* Trie data structure

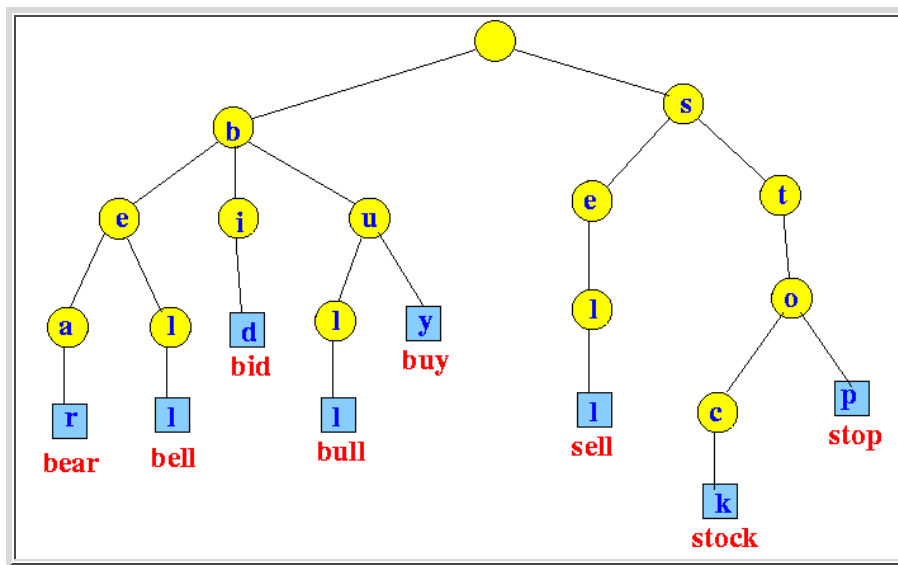
- Definitions:

- **Alphabet** = a set of **characters**
- Let **S** = a set of **s strings** (= **keys**) over an **alphabet Σ**
- A **trie T** that store the **keys** in **S** is a structure where:
 - Each **node** of **T** (except the **root node**) is **labeled** with a **character $c \in \Sigma$**
 - The **root node** has **no label !!!**
 - Each **internal node** of **T** can have $\leq |\Sigma|$ # of keys
 - The **keys** are stored in **alphabetical order** inside an **internal node**
 - The **trie T** has **s external nodes**
 - Each **external node** is associated with **one string** in **S**
(And **index (= location) information** are stored for these **strings**)
 - The **path** from the **root node** to an **external node** yields exactly **one string in S**

- Example:

- **S** = { bear, bell, bid, bull, buy, sell, stock, stop } (**s** = 8)

Trie:



■ Note:

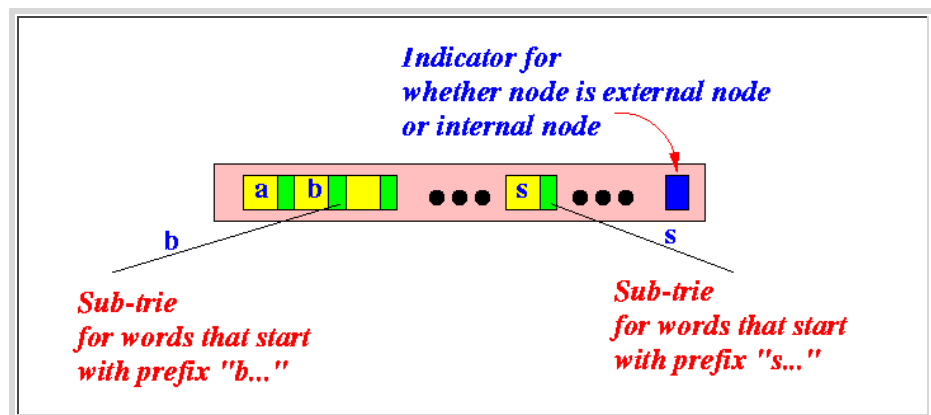
- There are 8 leaf nodes !!

• How to implement a trie

- o There are many ways to implement a trie

■ Use an array of references

- Each array element represent one letter of the alphabet
- Each array reference will point to a sub-trie that corresponds to strings that starts with the corresponding letter



■ Use a binary tree

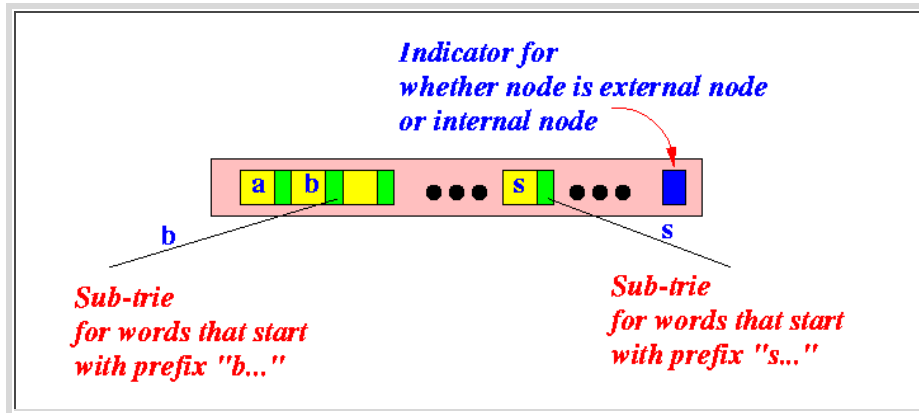
- This implementation is called a bitwise trie

- Implementing a *trie* using an *array* of references

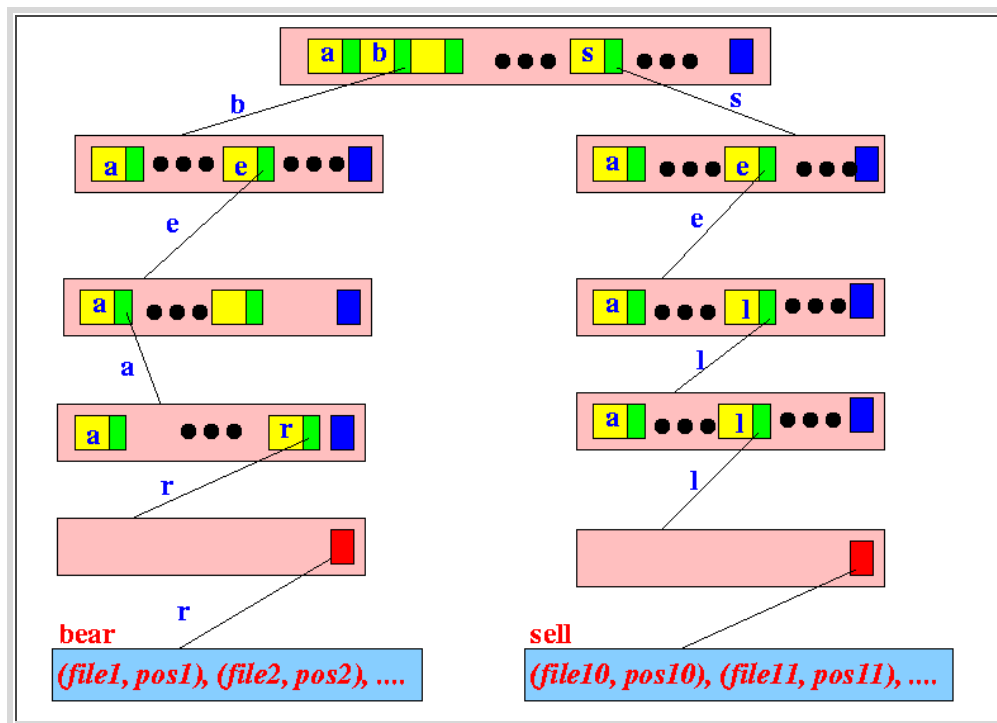
- Array implementation:

- If the **trie** is **dense** and/or **alphabet** Σ is **small** (e.g., $\Sigma = \{a, b, c, \dots, z\}$), you can use an **array** to store the **labels**

- Node structure:



Example:



- Implementing a *trie* using a *binary tree*

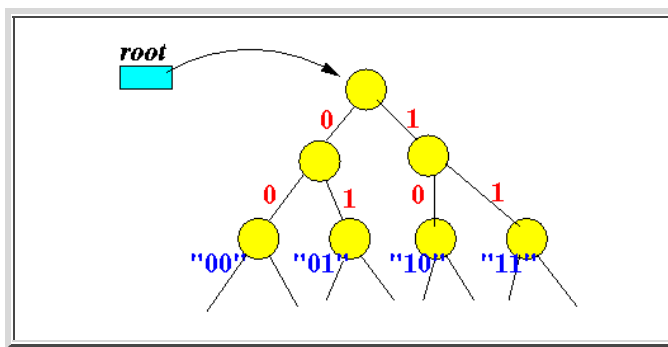
- The **binary tree** implementation:

- The **binary tree** implementation of a **trie** is known as a **bitwise trie**
- The **implementation** uses the **alphabet**: $\Sigma = \{0, 1\}$
 - In other words, the implementation stores a **sequences of bits**
- The **keys** are **read** as a **sequence of bits**

Example:

```
bear = 'b' 'e' 'a' 'r'
      = 01100010 01100101 01100001 01110010
```

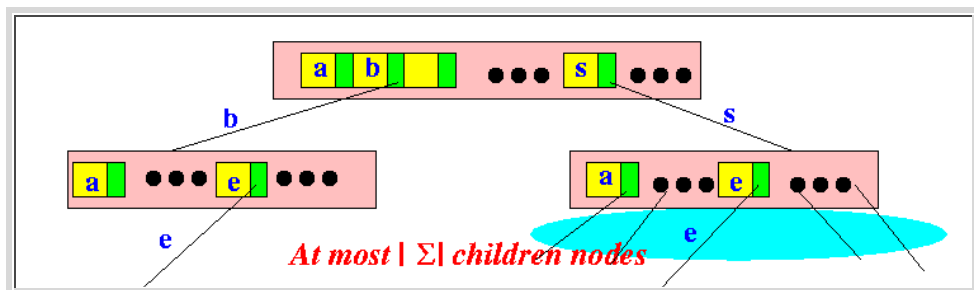
- A **bitwise trie** is a **binary tree**:



- **Structural properties of the standard trie**

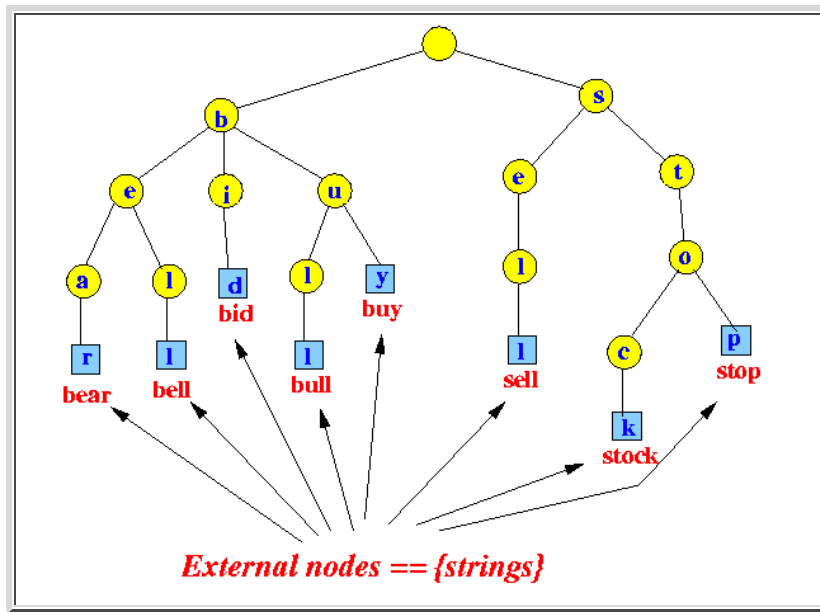
- **Properties of the standard trie:**

- Every **internal node** has $\leq |\Sigma|$ **children**



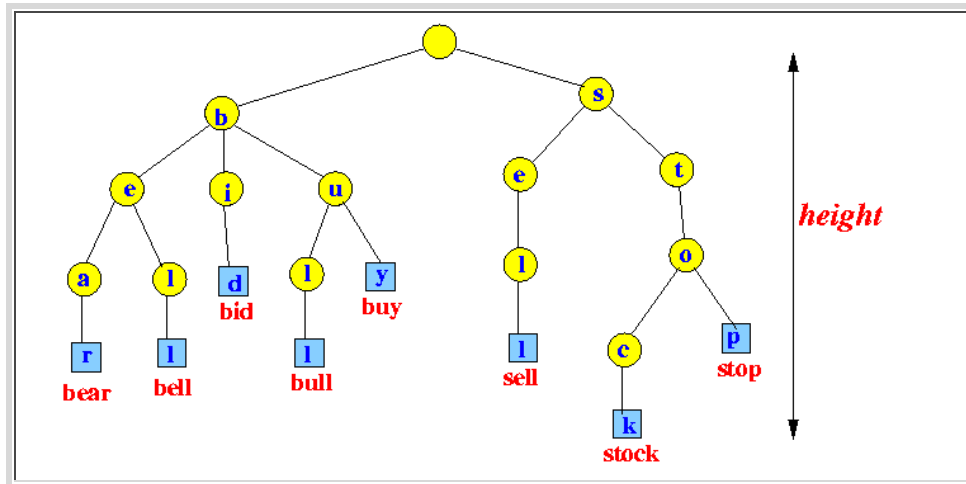
- This follows from the **way** that the **trie** is **constructed**

- The **trie T** on the **set S** with **s strings (keys)** has **exactly s external nodes**



- This is because a **path** from the **root node** to **one external node** corresponds to **1 key**

- The **height** of the **trie T** on the **set S** = the **length of the longest string** $\in S$



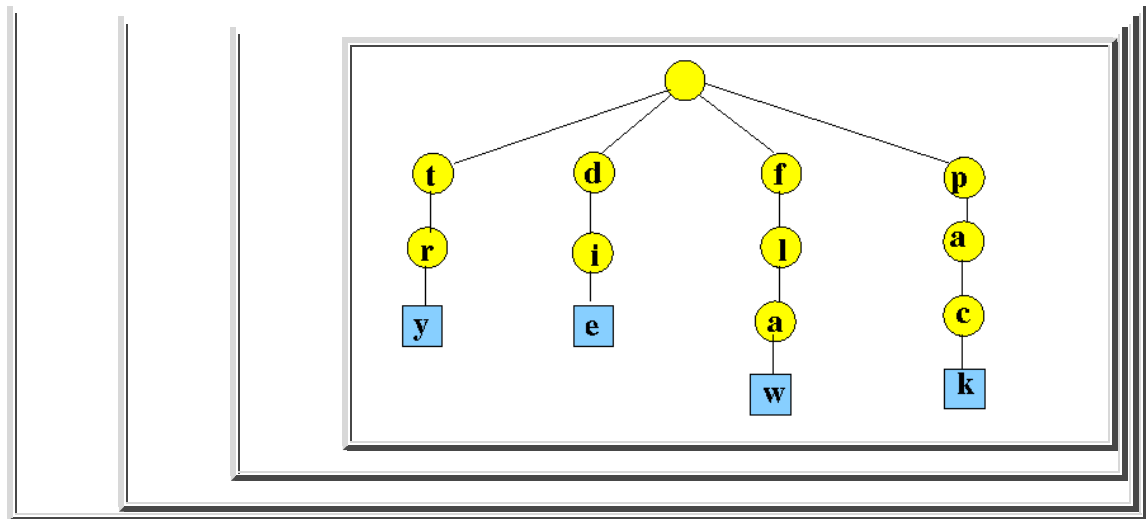
- This is because a **path** from the **root node** to **one external node** corresponds to **1 key**
- The **longest path** = **longest key** $\in S$

- The **number of nodes** in the **trie T** on the **set S** = $O(n)$, where **n** = **# characters in the strings** $\in S$

- In the **worst case**, **every character** in the **keys** are **different**

E.g.:

- **S** = { try, die, flaw, pack }



- **Inserting into a *standard* trie**

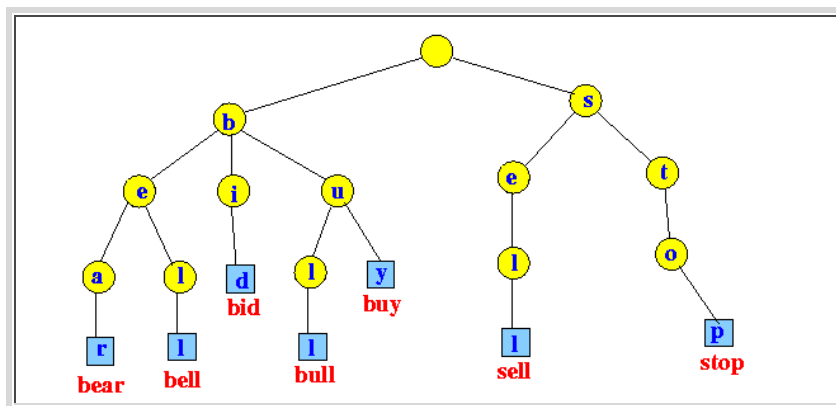
- **High level description:**

```

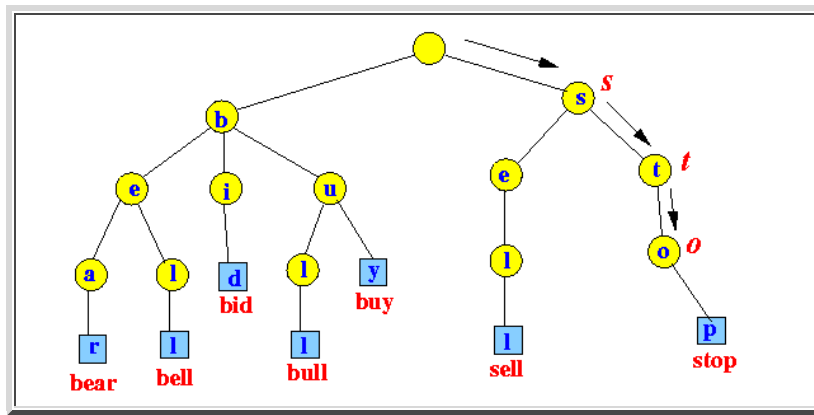
p = root;  // Start at the root
for ( each character c in the keyword ) do
{
  if ( c is already stored of the sub-trie )
    traverse to that subtrie
  else
    create a new node for c
}
insert value into leaf node
    
```

Example:

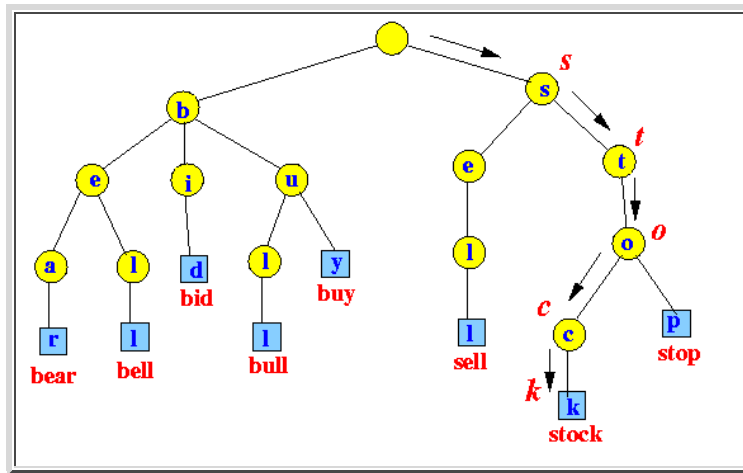
- Insert **stock** in this **trie**:



- First we traverse the **prefix** that is already **stored** in the trie: "**sto**":



- Then we create **nodes** for letters are **not** in the trie: "ck":



o Psuedo code:

```

Value put( Key k, Value v )
{
    int i;
    TrieNode p;

    p = root;

    for ( i = 0; i < k.length(); i++ )
    {
        nextChar = k[i];

        if ( nextChar ∈ p.char[] )
        {
            nextLink = link associated with the char entry;
            p = nextLink;    // Traverse
        }
        else
        {
            p = new TrieNode();    // Create new node

            insert "nextChar" into p;
        }
    }

    /* -----
    When the while loop ends, we have found or created
    */
}

```

```

the external node for the key "k"
----- */
insert v into node p;
}

```

- Advantages of Tries over an ordinary map

- Advantages:

- Looking up data in a **trie** is in the **worst case** = $O(m)$, where **m** = length of the key

- Look up in a **map** is $O(\lg(n))$ where **n** = # entries !!!

So a **trie** has performance levels that is similar to a **hash table** !!!

- Unlike a **hash table**, a **trie** can provide an **alphabetical ordering** of the **entries by key**
(I.e., A **trie** implements an **ordered map** while a **hash table** cannot !)

- Handling keys that are prefixes of another key

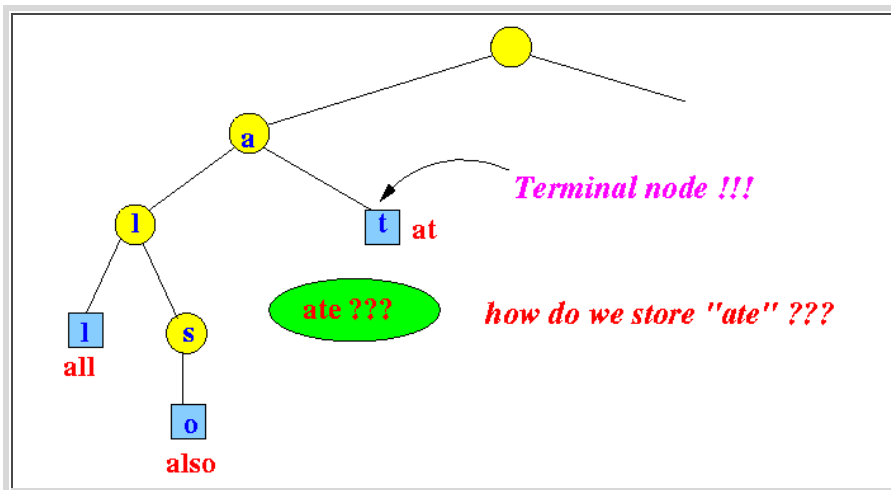
- The **standard trie** has the **property** that:

- Only the **external nodes** can store **information**

(The **path** formed by the **internal nodes** represents the **key**)

- When a **key (string)** is a **prefix** of **another key**, the **path** of the **first key** would **end** in an **internal node**

Example: **at** and **ate**



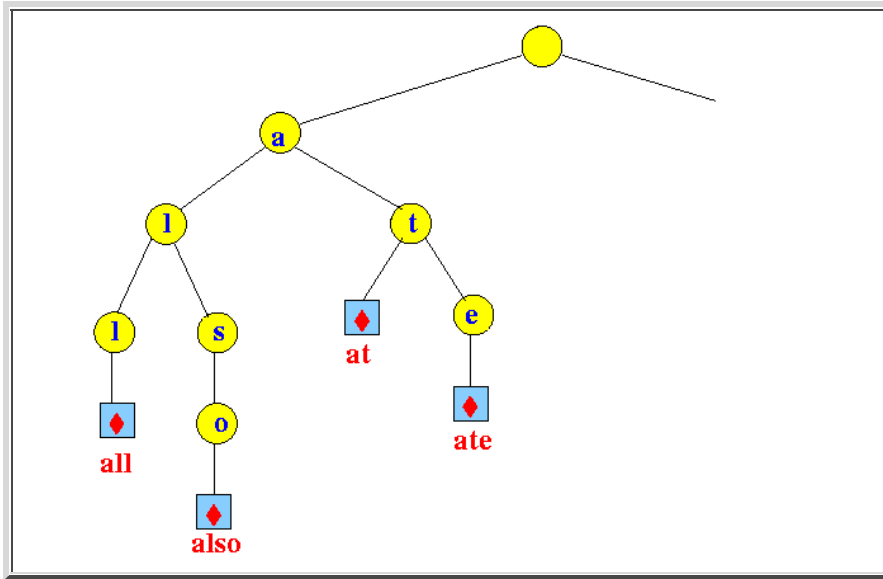
- Solution:

- Add a **special termination symbol** \diamond to the **alphabet** Σ

- The **termination symbol** \diamond has the **lower value** in the **alphabet**
I.e.,: **termination symbol** \diamond **preceeds** every character in the **alphabet** Σ !!

- We **append** the **termination symbol** \diamond to each **keyword** stored in the **trie**

Example:



◦ **Note:**

- We **typically** use the **NUL character** `'\0'` as **termination symbol**

• External material

◦ I found an implementation of **trie** here:

- <http://simplestcodings.blogspot.com/2012/11/trie-implementation-in-c.html>

and modified it to **print** the internal structure...

Caveat:

- The author of this program uses a **non-standard way** to represent nodes.
(I don't want to go into the details of **his** implementation. I do like his code because it can be easily adapted to print out the structure of the **trie**)

◦ **Example Program:** (Demo above code)

Example

- The **trie.c** Prog file: [click here](#)
- The **trie.h** Header file: [click here](#)
- A test program **main.c** Prog file: [click here](#)

How to run the program:

- **Right click** on link(s) and **save** in a scratch directory
- To compile: `gcc -o main main.c trie.c`
- To run: `main`

(Edited) Output:

```
(^*,0) (^b',-1) (^e',-1) (^a',-1) (^r',-1) (^*,36) "bear"
=====
(^*,0) (^b',-1) (^e',-1) (^a',-1) (^r',-1) (^*,36) "bear"

      (^l',-1) (^l',-1) (^*,14) "bell"
=====
(^*,0) (^b',-1) (^e',-1) (^a',-1) (^r',-1) (^*,36) "bear"

      (^l',-1) (^l',-1) (^*,14) "bell"

      (^*,,-1) (^*,4) "be"
=====
(^*,0) (^b',-1) (^e',-1) (^a',-1) (^r',-1) (^*,36) "bear"

      (^l',-1) (^l',-1) (^*,14) "bell"

      (^*,,-1) (^*,4) "be"

      (^u',-1) (^l',-1) (^l',-1) (^*,40) "bull"
=====
(^*,0) (^b',-1) (^e',-1) (^a',-1) (^r',-1) (^*,36) "bear"

      (^l',-1) (^l',-1) (^*,14) "bell"

      (^*,,-1) (^*,4) "be"

      (^u',-1) (^l',-1) (^l',-1) (^*,40) "bull"

      (^y',-1) (^*,40) "buy"
```

• An array implementaion of the standard trie

- I found an implementation of the **standard trie** that uses an **array** as described in the notes above:

■ <http://cristibalas.wordpress.com/2008/05/11/generic-string-trie-implementation/>

- The **node structure** is as follows:

```
typedef struct sTRIE_NODE
{
    char is_final;    // 1 if the node is an external node
                    // ==> void *data will contain data stored
                    // 0 if node is an internal node
                    // ==> use nxt[charCode] to search further

    void *data;       // user data that is stored with this node
                    // Only use this field if is_final == 1

    unsigned char childs; // number of childs, useful for fast deletion
                    // (If childs becomes 0, we delete this node)

    struct sTRIE_NODE *nxt[256]; // child pointers for every ASCII char !
} TRIE_NODE;
```

- o Here is how you **search** the **trie**:

```

/* =====
trie_search_nodes( node, str, pos )

    Search for the string str[pos...]
    in sub-trie rooted at node "node"
===== */
int trie_search_nodes(TRIE_NODE *node, char *str, int pos)
{
    unsigned char ch = (unsigned char)str[pos]; // Use this character to search
                                                // in node "node"

    if( ch == '\0' ) // string is ended
    {
        if( node->is_final == 1 )
        {
            // We reached an external node ==> the string is found !

            crt_search_data = node->data; // Get data assoc. with keyword
            return 1; // Return FOUND indication
        }
        else
        {
            return 0; // Return NOT FOUND indication
        }
    }
    else
    {
        if( node->nxt[ch] != NULL ) // More characters in string
        {
            return trie_search_nodes(node->nxt[ch], str, pos + 1);
            // Search further - using next character position
        }
        else
        {
            return 0; // not found
        }
    }
}

```

- o **Example Program:** (Demo above code)

Example

- The **C** Prog file (implements the **trie**): [click here](#)
- The **header** file: [click here](#)