LAPO SANTI

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1. Introduction

When faced with a multitude of alternatives, individuals often strive to organize them into coherent blocks or groups to better understand the decision landscape. Furthermore, they aim to establish a meaningful order within these blocks, enabling them to prioritize alternatives based on preference. This process involves two fundamental tasks: block clustering and order-based ranking. While clustering involves categorizing alternatives into distinct blocks, ranking focuses on arranging these blocks in a specific order. These tasks are typically accomplished through human judgments, often in the form of pairwise comparisons.

In block clustering, the aim is to determine the inherent similarities among alternatives and group them accordingly. By comparing pairs of alternatives, individuals can identify common characteristics, shared attributes, or comparable features that contribute to their clustering. This process helps unveil the underlying structure of the alternatives, allowing decision-makers to comprehend the relationships and associations between them. Several techniques, such as hierarchical clustering and k-means clustering, have been employed to address this task effectively.

Conversely, in order-based ranking, the primary objective is to establish a preference-based order among the identified blocks of alternatives. By comparing pairs of blocks, individuals can discern the relative favorability of one block over another. These pairwise comparisons generate a ranking list that encapsulates the perceived preference or priority of each block. Various methodologies, including the Bradley-Terry model and pairwise comparison matrices, have been utilized to derive meaningful rankings from the collected preferences.

In this article, we propose a novel approach, termed Block Clustering and Order-based Ranking (BCOR), which unifies the tasks of block clustering and order-based ranking into a cohesive framework. The BCOR model introduces a dynamic parameter that governs the granularity of block clustering, allowing decision-makers to explore a spectrum of clustering options. By iteratively adjusting this parameter, the model can encompass a wide range of decision-making scenarios, from finely differentiated blocks to coarser groupings.

A key insight of the BCOR model lies in its ability to relate the number of blocks to the underlying ranking structure. As the number of blocks converges to the total number of alternatives, the model effectively transitions into a traditional ranking approach, providing a complete ordering of the alternatives. Conversely, by intentionally reducing the number of blocks, decision-makers are presented with distinct groups of choices, each requiring preference considerations within its own subset. This approach offers a nuanced perspective on decision-making, allowing individuals to differentiate between highly favored groups and those that are comparatively less preferred.

The BCOR model provides a flexible and adaptive solution for organizing and prioritizing alternatives in various decision-making contexts. Its application extends beyond conventional clustering and ranking tasks, empowering decision-makers to explore the continuum between comprehensive rankings and granular groupings. Additionally, the model can be

tailored to incorporate different types of pairwise comparisons, enabling its utilization in diverse domains and decision scenarios.

To evaluate the effectiveness of the BCOR model, we conducted experiments using real-world datasets encompassing a wide range of decision contexts. The results demonstrate the model's ability to generate meaningful block clusters and order-based rankings, outperforming traditional approaches that solely focus on clustering or ranking tasks.

The contributions of this work can be summarized as follows:

We introduce the novel problem of block clustering and order-based ranking, bridging the gap between these two fundamental decision tasks. We propose the BCOR model, which provides a unified framework to accommodate various levels of granularity in decision-making, from complete rankings to distinct preference-based groups. We showcase the versatility of the BCOR model through experiments on real-world datasets, highlighting its superior performance compared to existing methods. By integrating block clustering and order-based ranking, the BCOR model offers decision-makers a comprehensive tool to navigate complex decision landscapes

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2. The Simple Model

The following Bayesian model is used to describe and analyze pairwise data, with the specific aim to identify clusters of points with similar connectivity patterns. Consider a scenario where entities engage in pairwise comparisons, and the outcomes are asymmetric – meaning that the probability of entity i winning over entity j is not necessarily the same as the probability of j winning over i.

The model uses a Poisson distribution to model the number of blocks, a Dirichlet-multinomial distribution to model the distribution of nodes' assignment across blocks, and a binomial distribution to model the distribution of edges within blocks. Additionally, the model includes a POMM process to model the probability of edge formation between nodes within blocks.

The goal of the model is to estimate the number of blocks, the distribution of nodes across blocks, and the probability of edge formation between nodes within blocks, given observed network data. The Bayesian approach allows for uncertainty in these estimates and provides a framework for incorporating prior knowledge and updating beliefs as new data becomes available.

2.1. Simple Model specification. This is a model for pairwise count data. We explicitly model the results of the interactions between two individuals i and j. Given N observations, the likelihood is

(1)
$$p(y|z, P, K) = \prod_{i=2}^{N-1} \prod_{j=i}^{N} p(y_{ij}|z, P, K)$$

(2)
$$= \prod_{i=2}^{N-1} \prod_{j=i}^{N} \binom{n_{ij}}{y_{ij}} p_{z_i, z_j}^{y_{ij}} (1 - p_{z_i, z_j})^{n_{ij} - y_{ij}}$$

where n_{ij} denotes the total number of interactions between the two individuals i and j and j and j is the number of successes of the individual i in interacting with j. The probability of success is given by p_{z_i,z_j} which consists of two parameters. The $K \times K$ matrix P and the $N \times 1$ vector z.

The vector z takes values over the discrete and finite set $\{1, \ldots, K\}$, and it is an indicator variable such that if $z_i = k$ individual i belongs to block k.

The matrix P contains the probabilities of success for individuals belonging to each possible blocks combination. For this reason P is $K \times K$. Therefore, the parameter p_{z_i,z_j} consists in the probability of success in an interaction between one individual belonging to block z_i and another of block z_j .

2.2. **Prior Specification.** This model has three parameters, and we put a prior on each of them.

Starting with P, we assume that its entries, namely $p_{k,k'}$, are independent and identically Beta(a,b) distributed random variable. By setting a=b=1 they collapse to a uniform

distribution.

(3)
$$p_{k,k'} \sim Beta(1,1) \text{ for } k,k'=1,\ldots,K$$

Second, we assume that the z_i s are independent and identically drawn from a multinomial distribution with one trial and probability vector $(\theta_1, \ldots, \theta_K)$. We can write then:

(4)
$$z_i | \boldsymbol{\theta} \sim \text{Multinomial}(1, \boldsymbol{\theta}) \text{ for } i = 1, \dots, N$$

To have more flexibility in the blocks sizes, we put an hyper-prior on the $\theta_1, \ldots, \theta_K$, assuming that they are drawn from a Dirichlet distribution with parameter the $K \times 1$ vector γ .

By marginalizing out θ , following the common practice in the literature, we can express the marginal distribution of z as:

(5)
$$p(\mathbf{z}|\boldsymbol{\gamma}) = \frac{\Gamma(\sum_{k=1}^{K} \gamma_k)}{\prod_{k=1}^{K} \Gamma(\gamma_k)} \frac{\prod_{k=1}^{K} \Gamma(n_k + \gamma_k)}{\Gamma(\sum_{k=1}^{K} (n_k + \gamma_k))}$$

where n_k is the number of players assigned to block k.

Finally, we assume that the number of clusters K follow a Poisson distribution Poisson($\lambda = 1$), subject to the condition K > 0.

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3. The POMM Model

In this section, we introduce a novel pairwise model that combines the notions of ordering and community structure to capture the complex interactions observed in various systems. Just as in network analysis, where directed hierarchies and community patterns play crucial roles in understanding interactions among entities (see Peixoto), we seek to develop a framework that simultaneously incorporates these two fundamental aspects in the context of pairwise data.

Drawing inspiration from the directed stochastic block model (SBM) used in network analysis, we develop a modified version tailored specifically for pairwise data. The model allows us to identify directed hierarchies among the entities while simultaneously detecting community structures that influence their interactions. By incorporating these two elements into a unified framework, our model offers a more comprehensive understanding of the underlying dynamics governing pairwise comparisons.

Several models of stochastic transitivity exist and have been used to describe the probabilities involved in experiments of paired comparisons, specifically in scenarios where transitivity is expected, however, empirical observations of the binary relation is probabilistic. This allows for a weaker player to win with a positive probability against a stronger opponent.

Tightly matched players might have a higher chance of observing this inversion while players with large differences in their skills might only see these inversions happen seldom. Stochastic transitivity models formalize such relations between the probabilities (e.g. of an outcome of a match) and the underlying transitive relation (e.g. the skills of the players).

Tightly matched players might have a higher chance of observing this inversion while players with large differences in their skills might only see these inversions happen seldom. Stochastic transitivity models formalize such relations between the probabilities (e.g. of an outcome of a match) and the underlying transitive relation (e.g. the skills of the players).

A binary relation \succeq on a set \mathcal{A} is called *transitive*, in the standard *non-stochastic* sense, if $a \succeq b$ and $b \succeq c$ implies $a \succeq c$ for all members a, b, c of \mathcal{A} .

Stochastic versions of transitivity include:

- (1) Weak Stochastic Transitivity (WST): $\mathbb{P}(a \succsim b) \ge \frac{1}{2}$ and $\mathbb{P}(b \succsim c) \ge \frac{1}{2}$ implies $\mathbb{P}(a \succsim c) \ge \frac{1}{2}$, for all $a, b, c \in \mathcal{A}$;
- (2) Strong Stochastic Transitivity (SST): $\mathbb{P}(a \succsim b) \geq \frac{1}{2}$ and $\mathbb{P}(b \succsim c) \geq \frac{1}{2}$ implies $\mathbb{P}(a \succsim c) \geq \max{\{\mathbb{P}(a \succsim b), \mathbb{P}(b \succsim c)\}}$, for all $a, b, c \in \mathcal{A}$;
- (3) Linear Stochastic Transitivity (LST): $\mathbb{P}(a \succeq b) = F(\mu(a) \mu(b))$, for all $a, b \in \mathcal{A}$, where $F : \mathbb{R} \to [0, 1]$ is some *increasing* and symmetric function (called a *comparison function*), and $\mu : \mathcal{A} \to \mathbb{R}$ is some mapping from the set \mathcal{A} of alternatives to the real line (called a *merit function*).

In the following sections we will focus on the Strong Stochastic Transitivity version.

3.1. Assuming the SST property for the probabilities in matrix P (SST matrix). Starting from the Simple model, we modify the parameter P, so that it satisfies the SST

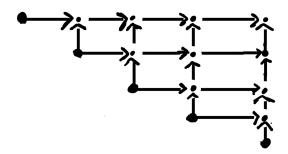


FIGURE 1. Dag representation of the Poset imposed onto the SST matrix

assumption. Recall that $P \atop K \times K$ is a collection of victory probabilities among K blocks or communities as follows:

$$P = \begin{pmatrix} p_{1,1} & p_{1,2} & \dots & p_{1,K} \\ p_{2,1} & p_{2,2} & \dots & p_{2,K} \\ \vdots & \vdots & \vdots & \vdots \\ p_{K,1} & p_{K,2} & \dots & p_{K,K} \end{pmatrix}$$

In order to satisfy the SST property, P entries can be spelled out explicitly as follows:

$$\begin{pmatrix}
0.5 & \leq & p_{1,2} & \leq & \dots & \leq & p_{1,K} \\
& & & & & & & & & & & & \\
1 - p_{1,2} & \leq & 0.5 & \leq & \dots & \leq & p_{2,K} \\
\vdots & \vdots & \vdots & \vdots & \dots & \vdots & \vdots \\
1 - p_{1,K} & \leq & 1 - p_{2,K} & \leq & \dots & \leq & 0.5
\end{pmatrix}$$

We further assume that the P matrix is skew-simmetric, that is, $p_{j,i} = 1 - p_{i,j}$. Furthermore, we set the main diagonal of matrix P to 0.5 for teams and 0 for individual players.

3.2. **Prior on** P. In order to incorporate such structure within a Bayesian framework, we need to impose an appropriate prior and hyper-prior on P so to reproduces such ordering among the probabilities. To this end, we draw inspiration from the the literature on spatial statistics. In particular, we exploit the idea of Partially Order Markov Models (POMMs) to define dependency and structure within the P matrix.

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3.3. **Posets and Dags.** P is a partially ordered set (Poset) (P, \leq) where a Poset is defined as follows.

Definition 1. [Poset]A poset, a partially ordered set, is a set of elements D where a binary relation \prec defined on D itself is such that:

- (6) For any $x \in D, x \prec x$ (reflexivity)
- (7) For any $x, y, z \in D, x \prec y$ and $y \prec z \implies x \prec z$ (transitivity)
- (8) For any $x, y \in D, x \prec y$ and $y \prec x \implies x = y$ (antisymmetry).

Then we call (D, \prec) a partially ordered set, or a poset.

We see that (P, \leq) satisfies the three properties of Definition (1), and therefore, is a poset. Let us now explicit how posets can be represented using directed cyclic graphs; we will see later how this will turn out useful to elicit the dependency structure between P's elements.

Let us denote with (V, E) the vertex and the edge sets of be a directed acyclic graph, respectively. To find the pair (V = P, E), meaning the graph corresponding to the poset P, we define the edge set E so that every pair $\{(p_i, p_i) \in E\} \implies p_i \leq p_j$.

This correspondence is many-to-one. Given a finite poset, one may construct a class of directed acyclic graphs; however, it may be of particular interest the minimal directed acyclic graph since it has the smallest possible directed edge set, which consists in $\{(p_i, p_j) \in E\} \implies p_i \leq p_j \land \{p_k : p_i \leq p_k \leq p_j\} = \emptyset\}$

We represent the minimal directed acyclic graph (P, E) corresponding to P in figure (1) by using the definition of Adjacent Lower Neighborhood that we introduced before.

Definition 2 (Cone). For any $y \in D$, the cone of y is the set

$$cone \ y = \{x \in D : x \prec y; x \neq y\}$$

Definition 3 (Adjl). For any $y \in D$ the adjacent lower neighbourhood of y is the set

$$adjl \ y = \{x \in D : (x, y) \text{ is a directed edge } in(D, F)\}$$

Definition 4 (Cover of a Subset). The cover of a subset B is a set of all elements x in D such that x is adjacent to an element in B and x is not in B. Formally, the cover of B is defined as follows:

$$covr\ B = x \in D : adjl\ x \subset B \ and\ x \notin B$$

where $adjl\ x$ is the set of all adjacent elements of x in D.

Intuitively, the cover of a subset B represents all the elements in D that are outside of B but are adjacent to at least one element in B. In other words, the cover of B captures the neighborhood of B in D.

Definition 5 (Level Sets). The level sets of a poset D are a sequence of nonempty cover sets defined recursively as follows:

$$L^0 = D_{min}; \quad L^i = covr\left(\bigcup_{k=0}^{i-1} L^k\right)$$

where D_{min} is the set of all minimal elements in D.

The first level set L^0 is simply the set of all minimal elements in D. The subsequent level sets are defined by taking the union of all the previous level sets and taking the cover of this union. Intuitively, each level set captures the neighborhood of the previous level sets in D.

3.4. Partially ordered Markov models: Consider a finite set of random variables $\{Z(s_1), \ldots, Z(s_n)\}$ indexed by location or "points"

$$D = \{s_1, \dots, s_n\} : n \in \{1, 2, \dots\}$$

That is, we assume the existence of a directed acyclic graph (D, F) and its corresponding poset (D, \prec) . Let (D, F) be a finite, directed acyclic graph and its corresponding poset (D, \prec) . Consider $s \in D$ and recall the definition of cone s. Also, let the quantity U_s denote any subsets of points not related to s. Formally:

$$U_s \subset \{u \in D : u \text{ and } s \text{ are not related }\}$$

Definition 6 (POMM). Then $\{Z(s): s \in D\}$ is said to be a partially ordered Markov model (POMM) if, for all $s \in D$ and any U_s

(9)
$$P(Z(s)|Z(cone s), Z(U_s)) = P(Z(s)|Z(adjl s))$$

Proposition 1. [Joint Distribution] Let (D, F) be a directed acyclic graph with no singleton points and let (D, \prec) , be its associated poset. Suppose that $\{Z(s) : s \in D\}$ is a POMM. Then

(10)
$$P(Z(D)) = P(Z(L^{0})) \prod_{k=1}^{m} \prod \{P(Z(u)) | Z(adjl\ u) : u \in L^{k}\}$$

(11)
$$= P(Z(L^0)) \prod_{i=1}^{k-1} \{P(Z(u))|Z(adjl\ u) : u \in D \setminus L^0\}$$

where $L^0, L^1, \dots L^m$ are the level sets as defined previously.

Result 1 relates the probability of a random variable defined on a poset to the probabilities of its restrictions to the lower level sets of the poset.

The result states that the probability of Z on the entire poset D can be expressed as a product of the probabilities of Z restricted to the level sets L^0, L^1, \ldots, L^m of the poset, where $L^0 = D_{\min}$ is the set of minimal elements of D, and L^k is the set of elements of D that are not in any of the previous level sets $L^0, L^1, \ldots, L^{k-1}$ and whose immediate predecessors are all in the union of the previous level sets $\bigcup_{i=0}^{k-1} L^i$.

The first part of the result states that the probability of Z on D is equal to the product of the probability of Z on L^0 and the conditional probabilities of Z on the elements of each subsequent level set L^k , given the values of Z on their immediate predecessors. This can be seen as a form of the chain rule of probability, where the joint probability of Z on D is decomposed into a product of conditional probabilities.

The second part of the result simplifies the product by noting that the conditional probabilities of Z on the elements of $D \setminus L^0$ are determined by the values of Z on their

immediate predecessors, which are all in L^0 or $D \setminus L^0$. Therefore, the product can be simplified to the product of the probability of Z on L^0 and the conditional probabilities of Z on the elements of $D \setminus L^0$ given the values of Z on their immediate predecessors in $D \setminus L^0$. This simplification reduces the number of terms in the product and makes the computation of the joint probability of Z on D more efficient.

4. Application to the SST matrix

In this section, we introduce a matrix that exhibits the Strong Stochastic Transitivity (SST) property. Our aim is to model the probabilities of players winning matches in a tournament, ensuring that these probabilities are consistent. Transitivity, as mentioned above, is crucial in guaranteeing that if Player A is stronger than Player B and Player B is stronger than Player C, then Player A must be stronger than Player C. This logical relationship is reflected in the probabilities of victory between players, which are contained in the matrix P introduced in the previous section. The matrix under consideration is denoted as P, representing victory probabilities among K entities, such as players or groups of players. Each element $p_{i,j}$ in the matrix P represents the probability of player i winning against player j in a tennis match, without draws allowed. To satisfy the requirements of a valid probability, we have $p_{i,j} + p_{j,i} = 1$, where $p_{j,i} = 1 - p_{i,j}$.

Without loss of generality, we assume that player/team 1 is the strongest and player/team K is the weakest. Monotonicity in the probabilities is crucial, ensuring that the probabilities increase monotonically as the index of the columns j increases and decrease monotonically as the index of the rows i increases. Violation of these assumptions could lead to contradictions where weaker players have higher probabilities of winning than stronger ones.

Therefore, we impose the following constraints on the probabilities in the upper triangular part of matrix P:

 $0.5 \le p_{i,j} \le p_{1,K}$ for all i < j. $1 - p_{i,j} \le 1 - p_{j,K}$ for all i < j. The main diagonal of matrix P is set to 0.5 for teams and 0 for individual players.

By ensuring the SST property and imposing the constraints, we obtain a coherent and tractable framework to express the joint probability distribution of these ordered probabilities, enabling us to perform inference efficiently.

4.1. **Defining a Poset over** P. Having defined P, now we want to re-define it within a finite Poset framework to obtain (P, \leq) . Let (i, j) and (p, q) be two elements of P and define the binary relation \leq on P as

(12)
$$(p,q) \le (i,j) \iff p \le i \text{ and } q \le j$$

Now, (P, \leq) is clearly a Poset since it satisfies the three properties of Definition (1). We represent the corresponding directed acyclic graph (P, F), where F is the set of directed edges between vertices in Figure (1) by using the definition of Adjacent Lower Neighborhood that we introduced before.

Definition 7 (Cone over P). In this case, for any $(i,j) \in D$ the cone of (i,j) is the set:

(13)
$$cone(i, j) = \{(i, j - 1), \dots (i, i), (i + 1, j), \dots, (K, j)\}$$

Definition 8 ($\overline{\text{cone}}(i,j)$). In this case, for any $(i,j) \in P$ the closure of the cone of (i,j) is the set:

(14)
$$\overline{cone}(i,j) = \{(i,j), \dots (i,i), (i,j), \dots, (K,j)\}$$

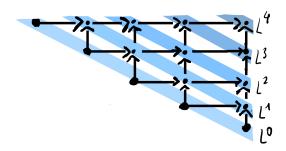


FIGURE 2. Visual Display of the level sets $L^k: k = 0, \dots, 4$

Definition 9 (The adjacent lower neighborhood of (i, j)). In this case, for any $(i, j) \in P$ the adjacent lower neighborhood of (i, j) is the set:

(15)
$$adjl (i, j) = \{(i, j - i), (i + 1, j)\}\$$

Definition 10 (closure of adjl (i, j)). In this case, for any $(i, j) \in P$ the closure of adjl (i, j) is the set:

(16)
$$\overline{adjl}(i,j) = \{(i,j-i), (i+1,j), (i,j)\}\$$

Definition 11 (D_{min}^P) . In this case, the minimal element denoted by D_{min} is such that

(17)
$$P_{min} = \{(i, j) \in P : i = j\}$$

namely the main diagonal of the P matrix.

The level sets of P, in this case, corresponds to the diagonals above the main one.

4.2. Partially ordered Markov Models applied to P. Considering the finite set of random variables $\{P_{1,1}, \ldots, P_{K,K}\}$ indexed by locations where

$$D \equiv \{(1,1), \dots (K,K)\}$$

Having showed the existence of a directly acyclic graph (P, F) and its corresponding poset (P, \leq) , we can write down

(18)
$$P(P_{ij}|Z(\text{cone }i,j)) = P(P_{ij}|P(\text{adl }i,j))$$

(19)
$$P(P_{ij}|P_{i+1,j},P_{i,j-1})$$

Now, exploiting Proposition (1), we can write:

(20)
$$P(P(D)) = \prod_{i=1}^{K} \prod_{j=i}^{K} P(P_{ij}|P_{i+1,j}, P_{i,j-1})$$

In order to induce an ordering or ranking among the blocks, we introduce a hierarchical structure. Without loss of generality, we assume that block 1 has the highest probability

of success when interacting with any other block, while block K has the lowest probability of success. We require this ranking to be transitive, meaning that if block A has a higher probability of success when interacting with block B, and block B has a higher probability of success when interacting with block C, then block A must still be the preferred choice when interacting with block C. Mathematically, this can be expressed as $p_{k,h} > p_{k',h}$ when k < k' for $h \notin k, k'$.

To achieve this effect, we impose three conditions: probabilities should increase in the columns, decrease in the rows, and be greater than or equal to 0.5 in the upper triangular matrix. To satisfy these conditions, we construct the following scheme.

4.2.1. Level Sets. We define the level sets, denoted by $L^{(k)}$, as the diagonals of the upper triangular matrix P. The main diagonal is referred to as level set 0, denoted by $L^{(0)}$. The diagonal above it is denoted by $L^{(1)}$, and so on up to $L^{(K-1)}$. Each level set $L^{(k)}$ is formally defined as:

(21)
$$L^{(k)} := p_{ij} \mid j - i = k \text{ for } k = 0, \dots, K - 1$$

It is worth noting that the cardinality of each level set is given by $|L^{(k)}| = K - k$ for $k = 0, \ldots, K - 1$.

4.2.2. Truncation Process. Requiring that probabilities increase in the rows and decrease in the columns is equivalent to ensuring that the level sets satisfy the condition:

(22)
$$\max(L^{(k)}) < \min(L^{(k+1)}) \text{ for } k = 0, \dots, K-1$$

To enforce this behavior, we employ an increasing truncation process controlled by a parameter α , with an upper bound given by β_{max} .

We consider a generic power-law function $y=x^{\alpha}+0.5$, which governs the rate of increase in the truncation process. Setting f(0)=0.5 ensures that $L^{(0)}$ is greater than or equal to 0.5, satisfying the transitivity condition. The function y is monotonically increasing for x>0. To generate the truncations, we partition y effectively by dividing the interval into K equal-sized segments. The segment endpoints are computed as $x_k=\Delta\times k$ for $k=0,\ldots,K$, where $\Delta=\left(\left(\beta_{\max}-0.5\right)^{(1/\alpha)}-0\right)/K$. Mapping the cumulative sum of the segment endpoints back to y yields K truncation points denoted by $y^{(k)}$, which are defined as:

(23)
$$y^{(k)} = \left(\frac{(\beta_{\text{max}} - 0.5)^{(1/\alpha)}}{K} \times k\right)^{\alpha} + 0.5 \quad \text{for } k = 0, \dots, K$$

Notice that f(0) = 0.5 and $f(K) = \beta_{\text{max}}$ by construction.

These truncation points provide the upper and lower bounds for the entries within the corresponding level sets, thereby ensuring the desired hierarchy and transitivity in the ranking.

Mathematically we have that:

(24)
$$y^{(k)} < p_i j \in L^{(k)} < y^{(k+1)} \quad k = 0, \dots, K-1$$

4.2.3. The POMM Prior. Finally, we put a prior on the matrix P with this new structure in place. We assume that entries $p_{ij} \in L^{(k)} \mid (y^{(k)} + y^{(k+1)})$ are identically and independently distributed according to a Uniform $(y^{(k)}, y^{(k+1)})$. We also put a log-normal hyper-prior on α such that

(25)
$$\alpha \sim \operatorname{lognormal}(\mu_{\alpha}, \sigma_{\alpha}^{2})$$

where μ_{α} , σ_{α}^2 are specified according to the normal parametrisation of the lognormal and are fixed to 1 and 2 respectively. Altogether, the POMM prior on P is the following:

(26)
$$p_{ij} \in L^{(k)} \mid y^{(k)}, y^{(k+1)} \sim (y^{(k)}, y^{(k+1)})$$

(27)
$$\alpha \sim \text{Lognormal}\left(\mu_{\alpha}, \sigma_{\alpha}^{2}\right)$$

and where the truncations $y^{(k)}$ are derived as in (23).

4.2.4. The POMM Prior 2. Finally, we put a prior on the matrix P with this new structure in place. We assume that entries $p_{ij} \in L^{(k)} \mid (y^{(k)} + y^{(k+1)})$ are identically and independently distributed according to a Normal $(\mu^{(k)}, \sigma^{2(k)})$, where $\mu^{(k)} = \frac{y^{(k)} + y^{(k+1)}}{2}$ which corresponds to the midpoint of the level set $L^{(k)}$, and $\sigma^{2(k)} = (y^{(k)} + y^{(k+1)}) \times S$, where S is a parameter denoted as overlap, which intuitively is proportional to the overlap in the distribution support of the level sets. We also put a log-normal hyper-prior on α such that

(28)
$$\alpha \sim \operatorname{lognormal}(\mu_{\alpha}, \sigma_{\alpha}^{2})$$

where μ_{α} , σ_{α}^2 are specified according to the normal parametrisation of the lognormal and are fixed to 1 and 2 respectively. Altogether, the POMM prior on P is the following:

(29)
$$p_{ij} \in L^{(k)} \mid y^{(k)}, y^{(k+1)} \sim \text{Normal}\left(\mu^{(k)}, \sigma^{2(k)}\right) \mathbb{I}(0.5, \beta_{\text{max}})$$

(30)
$$\alpha \sim \text{Lognormal}(\mu_{\alpha}, \sigma_{\alpha}^2)$$

(31)
$$S \sim \sim \text{Lognormal}\left(\mu_S, \sigma_S^2\right)$$

and where the truncations $y^{(k)}$ are derived as in (23).

Insert here the plot for 3 different overlap values

4.3. Some notes on the overlap of two normal densities. The computation or estimation of S, the overlap for two normal distributions, with density functions $f_1(X; \mu_1, \sigma_1^2)$ and $f_2(X; \mu_2, \sigma_2^2)$, depends on whether the two variances are equal. Since the estimation of OVL when $\sigma_1^2 = \sigma_2^2 = \sigma^2$ is simpler, has attracted more attention, and leads to firmer conclusions, we consider this situation in some detail. On the other hand, we shall simply summarise the results we have obtained in the more complicated circumstance when $\sigma_1^2 \neq \sigma_2^2$.

When $\sigma_1^2 = \sigma_2^2$, the normal density functions intersect at the single value of X equal to $(\mu_1 + \mu_2)/2$.

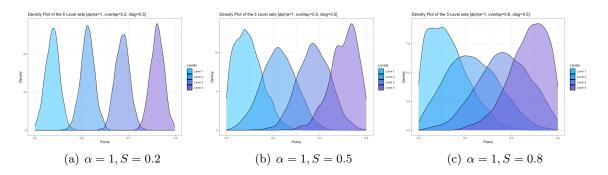


FIGURE 3. Densities for K = 5 Level Sets, for different S values and $\alpha = 1$; the main diagonal is set to 0.5, and its collapsed density is not reported.

Representing the standard normal distribution function by $\Phi(\cdot)$, we obtain

$$S = 2\Phi\left(-\frac{|\mu_1 - \mu_2|}{2\sigma}\right) = 2\Phi\left(-\frac{1}{2}|\delta|\right)$$

In our case we have that $\mu_1 < \mu_2 < \ldots < \mu_K$ and we are interested in the total overlap between the densities that are defined on the level sets. The total overlap is the sum of the pairwise overlaps between the densities taken two at a time:

(32)
$$S = \sum_{p=1}^{K-1} \sum_{q=i+1}^{K} S_{pq}$$
$$= \sum_{p=1}^{K-1} \sum_{q=i+1}^{K} \Phi\left(-\frac{|\mu_p - \mu_q|}{2\sigma}\right)$$

Let's simplify the expression for the total overlap, denoted as S, using the fact that the means are ordered $\mu_1 < \mu_2 < \ldots < \mu_K$.

First, we can notice that for each pair of means μ_p and μ_q , where p < q, we have the following relationship:

$$\mu_q - \mu_p > 0$$
 (since $\mu_1 < \mu_2 < \dots < \mu_K$)
 $|\mu_q - \mu_p| = \mu_q - \mu_p$ (because it is positive)

Now, let's rewrite the total overlap S using this information:

$$\begin{split} S &= \sum_{p=1}^{K-1} \sum_{q=i+1}^K \Phi\left(-\frac{|\mu_p - \mu_q|}{2\sigma}\right) \\ &= \sum_{p=1}^{K-1} \sum_{q=i+1}^K \Phi\left(-\frac{\mu_q - \mu_p}{2\sigma}\right) \quad \text{(using the fact mentioned above)} \\ &= \sum_{p=1}^{K-1} \sum_{q=i+1}^K \Phi\left(-\frac{\mu_q}{2\sigma} + \frac{\mu_p}{2\sigma}\right) \end{split}$$

5. Estimation

For the moment, we want to infer just $\theta = \{z, P, \alpha, S\}$, meaning that we treat K as a known constant. The estimation strategy is a Hybrid MCMC algorithm. Since simulating from the conditional distribution $p(\theta_i|\theta_j, j \neq i)$ is unfeasible or computationally expensive, we substitute the simulation from the full conditional distribution with a simulation from a proposal distribution q_i . Referencing Muller's (1991) work, the Hybrid modification is as follows:

Algorithm 1 Metropolis-within-Gibbs MCMC

$$\mathbf{for}\ i=1,\dots,p\quad \mathtt{given}\quad \left(\theta_1^{(t+1)},\dots,\theta_{i-1}^{(t+1)},\theta_i^{(t)},\dots,\theta_p^{(t)}\right)\,\mathbf{do}$$

Simulate

(33)
$$\theta_i' \sim q_i \left(\theta_i^{(t)} | \theta_1^{(t+1)}, \dots, \theta_i^{(t)}, \theta_{i+1}^{(t)}, \dots, \theta_p^{(t)} \right)$$

2. Take

(34)
$$\theta_i^{(t+1)} = \begin{cases} \theta_i^{(t)} & \text{with probability } 1 - r_i, \\ \theta_i' & \text{with probability } 1 - r_i, \end{cases}$$

where

(35)
$$r_i = 1 \land \left\{ \frac{\left(p(\theta_i' | \theta_i^{(t)} | \theta_1^{(t+1)}, \dots, \theta_i^{(t)}, \theta_{i+1}^{(t)}, \dots, \theta_p^{(t)} \right)}{\left(p(\theta_i^{(t)} | \theta_i^{(t)} | \theta_1^{(t+1)}, \dots, \theta_i^{(t)}, \theta_{i+1}^{(t)}, \dots, \theta_p^{(t)} \right)} \right\}$$

end for

5.0.1. Adaptive algorithm for $\theta = \{P, \alpha, S\}$. We specify the proposal distributions in (33) above as

$$\theta_i' \sim \text{Normal}\left(\theta_i^{(t-1)}, \sigma_{\theta_i}^2\right)$$

whose sampled value is accepted or rejected by evaluating the logarithm of (35). Choosing a correct $\sigma_{\theta_i}^2$ value is not straightforward, and we choose to resort to an adaptive algorithm to elicitate a correct proposal variance. We proceed as in Roberts, Rosenthal 2012. For each of the K(K-1)/2+2 parameters i ($1 \le i \le K(K-1)/2+2$), we create an associated variable ls_i giving the logarithm of the standard deviation to be used when proposing a normal increment to variable i. We begin with $ls_i = \log(0.04)$ for all i (corresponding to 0.2 proposal standard deviation). After the n-th "batch" of 50 iterations, we update each ls_i by adding or subtracting an adaption amount $\delta(n)$. The adapting attempts to make the acceptance rate of proposals for variable i as close as possible to 0.234, following the literature practice Chris Sherlock12009. Specifically, we increase ls_i by $\delta(n)$ if the fraction of acceptances of variable i was more than 0.234 on the n-th batch, or decrease ls_i by $\delta(n)$ if it was less.

 \rightarrow Insert here plots of convergence to the acceptance ratio We specify in the Appendix the full expression for the ratio of $\theta = \{P, \alpha, S\}$ in (35).

5.0.2. Adaptive Algorithm for $\theta = z$. When dealing with $\theta = z$, a discrete parameter, we need to adapt the formulation while maintaining the underlying concept. In the case of the POMM model, the labels k = 1, ..., K are ordered, and therefore, we can define a distance metric between these labels. Let us denote the distance between k and k' as d(k, k'), which can be expressed as:

(36)
$$d(k, k') = |k - k'|$$

If the acceptance rate for a particular player i is too low, we want the proposal to explore neighboring labels. Conversely, if the acceptance rate is too high, we aim to sample labels further away. To achieve this, we assign a sampling probability to each label that is inversely related to its distance from the current label. Specifically, we define $p(k') = p(|k'-k|) = \text{Normal}(0, \sigma_i^2)$, where σ_i^2 is adapted as above. A larger variance assigns higher probabilities to distant labels, while a smaller variance favors closer labels. Finally, we employ a multinomial distribution to sample the next label k':

(37)
$$k' \sim \text{Multinomial}(1, K, p(|k' - k|))$$

By using this approach, we can adapt the algorithm to explore labels based on their distances from the current label.

We specify in the Appendix the full expression for the ratio of $\theta = \{z\}$ in (35).

6. Point Estimate, Model Selection, and inference

While algorithmic methods produce a single estimated partition, our model offers the entire posterior distribution across different node partitions. We are comparing the results from the simulation study via the following three main measures:

- Variation of Inforamation (VI): to fully utilise this posterior and engage in inference directly within the partition space, we adopt the decision-theoretic approach introduced by Wade and Ghahramani (2018) for block modeling. This involves summarizing posterior distributions using the variation of information (vi) metric, developed by Meilă (2007), which measures the distance between two clusterings by comparing their individual and joint entropies. The vi metric ranges from 0 to log2 V, where V represents the number of nodes. Intuitively, the vi metric quantifies the amount of information contained in two clusterings relative to the shared information between them. As a result, it decreases towards 0 as the overlap between two partitions increases. Refer to Wade and Ghahramani (2018) for a detailed exploration of the key properties of the vi metric. Within this framework, a formal Bayesian point estimate for z is obtained by selecting the partition with the lowest averaged vi distance from the other clusterings
- WAIC: While the WAIC yields practical and theoretical advantages and has direct connections with Bayesian leave-one-out cross-validation, thus providing a measure of edge predictive accuracy, the calculation of the WAIC only requires posterior samples of the log-likelihoods for the edges: $\log p(y_{ij}|z, P, \alpha) = y_{ij} \log p_{z_i, z_j} + (n_{ij} y_{ij}) \log(1 p_{z_i, z_j}), \quad i = 2, \ldots, N, j = 1, \ldots, i-1.$
- Misclassification error: predicting the group membership z_{N+1} of a new player may also be of interest. We can derive the estimate of the block probabilities for new players based on their early matches with some of the existing players.

(38)
$$p(z_{N+1} = k | \mathbf{Y}, y_{N+1}, \hat{z}) \propto p(y_{N+1} | \mathbf{Y}, \hat{z}, z_{N+1} = k) \cdot p(z_{N+1} = k | \hat{z})$$
$$= p(y_{N+1} | \hat{z}, z_{N+1} = k) \cdot p(z_{N+1} = k | \hat{z})$$

where $p(z_{N+1} = k | \mathbf{Y}, y_{N+1}, \hat{z})$ is the posterior probability of the new node N+1 to belong to the block k, given the previously observed data Y, the new node's data y_{N+1} and the estimated labels \hat{z} . On the right hand side of the expression above, $p(y_{N+1} | \mathbf{Y}, \hat{z}, z_{N+1} = k)$ represents the likelihood of observing y_{N+1} given the previously observed data Y and the estimated labels \hat{z} , which, due to conditional independence, is the same as conditioning just on \hat{z} . Finally, $p(z_{N+1} = k | \hat{z})$ represents the prior probability of label k for the new node N+1 given \hat{z} , which we can approximate with the relative size of the blocks n_k .

7. Simulation Study from the Simple Model N=100

In order to evaluate how well our model performs in a situation similar to our intended use, and measure its advantages compared to the best existing alternatives, we generated three simulated tournaments with 100 players from the Simple Model. We want to compare how it performs compared to the POMM extension and other state-of-the-art alternatives. Each tournament had a different number of blocks in the underlying structure. We set the total number of games $M := 0.5 * \sum_{i,j}^{N} n_{ij} = \sum_{i,j}^{N} y_{ij} = 4000$, which is the average number of matches played in one year of tennis tournaments. We divided the players into three, five and nine blocks (K = 3, 5, 9 respectively). In Figure (4), we display the three simulated tournaments, where the difficulty of accurately determining the group membership increases as the number of games increases with the number of blocks.

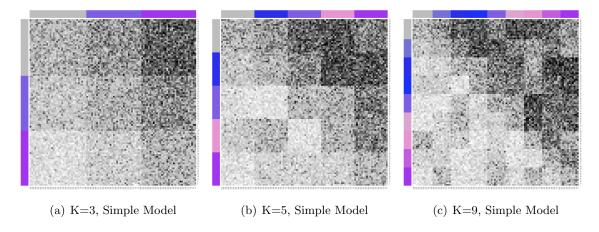


FIGURE 4. Adjacency Matrices simulated via the Simple Model

We compare the performance of the Simple model with the POMM one. We fixed arbitrarily $\beta_{\text{max}} = .75$. In table (??) we report the results of the simulation. In the three cases, for the Simple and the POMM model, we compare the WAIC, the VI distance and the misclassification error, obtained by considering 100 new incoming players which get to play just with 10 players each. We also compare the labels estimated against the regularised spectral clustering algorithm and the Louvain algorithm. The Simple model is the best performing relative to the other three alternatives.

In figure (5) we report the estimated co-clustering matrices resulting from the simulation process.

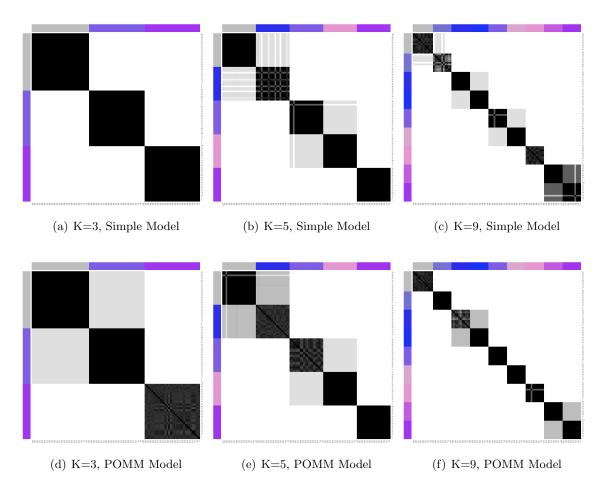


FIGURE 5. Co-Clustering Matrices obtained via the Simple Model(above) and the POMM model (below).

 $P \ {\rm summary \ table}$ True Model Simple, N=100

Fitted Model	\overline{MAE}			% within-95% CI interval			CI interval length		
	$\overline{(a)}$	(b)	(c)	(a)	(b)	(c)	$\overline{(a)}$	(b)	(c)
POMM model	0.0211	0.0735	0.0709	100%	50%	69.4%	0.1017	0.0955	0.1788
Simple model	0.0036	0.0606	0.0826	100%	60%	55%	0.0108	0.1077	0.1798

 $z \ {\rm summary \ table}$ True Model Simple, N=100

Method	VI	$VI distance_{MAP}$			listan	.ce _{VI lb}	WAIC		
Wie wie d	$\overline{(a)}$	(b)	(c)	$\overline{(a)}$	(b)	(c)	$\overline{}(a)$	(b)	(c)
POMM model	0	0	0.33	0	0.4	0.44	-6418.29	-6196.77 18.74	6320.08
Simple model	0	0.5	0.58	0	0	0.22	-6506.41 17.60	-6248.25 18.54	-6285.03 18.61

8. Simulation Study from the POMM Model N=100

In this section we reverse the exercise performed in previous one. Before we were simulating from the Simple model, now we are doing the same, with similar parameters $(K = 3, 5, 9, M = 4000 \text{ and } \beta_{\text{max}} = .75)$. Here are the results.

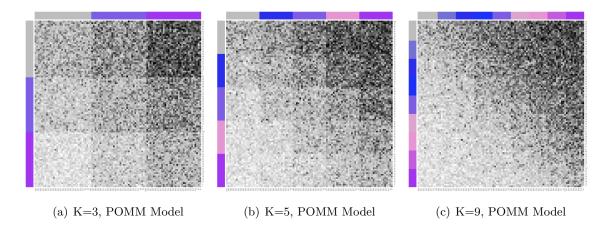
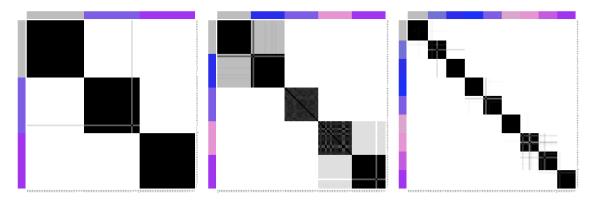


FIGURE 6. Adjacency Matrices simulated via the POMM Model

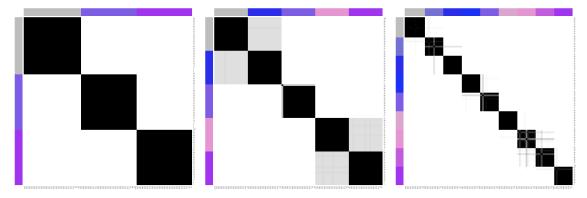
In table (??) we report the results of the simulation. As before, for the Simple and the POMM model, we compare the WAIC, the VI distance and the misclassification error, obtained $N_n ew = 100$. Also here we compare clustering performance against that of the regularised spectral clustering algorithm and the Louvain algorithm. The POMM model is the best performing relative to the other three alternatives.

P summary table True Model POMM, $K=3,\,N=100$

Fitted Model		\overline{MAE}		% witl	hin-95%	-CI interval	rval CI interval length			
Titted Wiodel	$\overline{(a)}$	(b)	(c)	(a)	(b)	(c)	(a)	(b)	(c)	
POMM model Simple model						100% 100%		0.1309 0.1311		



(a) K=3, Simple Model Estimates (b) K=5, Simple Model Estimates (c) K=9, Simple Model Estimates



(d) K=3, POMM Model Estimates (e) K=5, POMM Model Estimates (f) K=9, POMM Model Estimates

FIGURE 7. Co-Clustering Matrices obtained via the Simple Model(above) and the POMM model (below).

z summary table True Model POMM, N = 100

Method	VI distance _{MAP}			VI distance $_{VI}$ lb			WAIC		
Wednod	$\overline{(a)}$	(b)	(c)	$\overline{(a)}$	(b)	(c)	(a)	(b)	(c)
POMM model	0.22	0	0.1	0.66	0	0.1	-6481.92	-6573.49	6662.23 17.95
Simple model	0.22	0	0.09	0.66	0	0.09	$6481.87 \\ 17.60$	6596.61 17.84	6663.48 17.96

8.1. Diagnostic Checks.

9. Application to Tennis Data

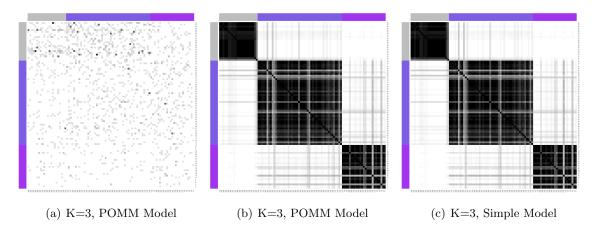


FIGURE 8. Adjacency Matrices simulated via the POMM Model

10. Appendix I: Investigating Empirically the prior behaviour

In this section, we explore the behaviour of the POMM prior, as α , the parameter controlling the rate of increase of the power-law process and S, the variation (and therefore, the overlap) of the level sets sets, change.

We start with a simulation study, whose results are reported in figure (??).

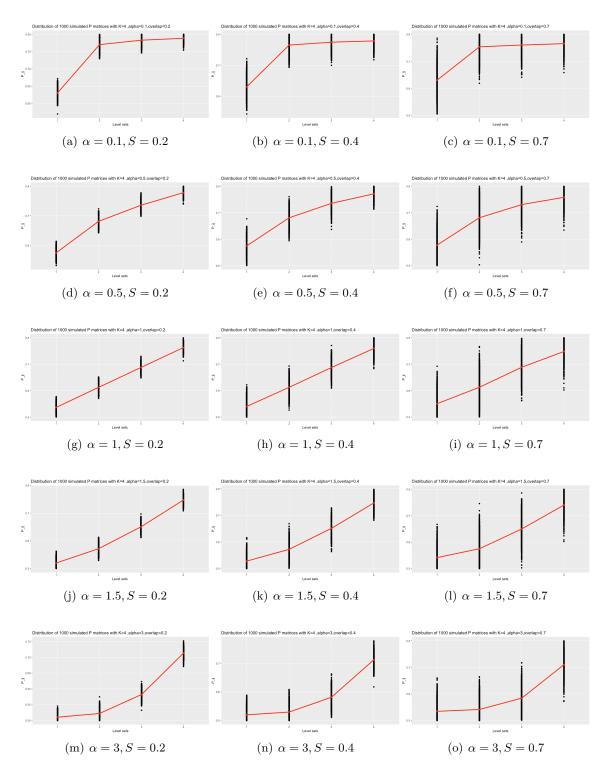


FIGURE 9. Distribution of n=1000 simulated P matrices with $\alpha \in \{0.1,0.5,1,1.5,3\}$ and $S \in \{0.2,0.4,0.7\}$. The points are grouped into the K=5 level sets- the main diagonal is set =0.5. The black vertical dots are the points, while the red lines show the evolution of the mean of the level sets.

11. Appendix II: Empirical Assessment of the inference problem

This section has two main objectives:

- (1) The first objective is to assess the identifiability of the POMM model. Given that the richness of the model, it is not obvious that the parameters are identifiable in all the regions of the parameter space. Therefore, by means of several simulation studies, we aim at identifying any possible identifiability constraints.
- (2) The second objective is code testing, by checking the whole estimation procedure on simulated data.

11.1. Focus on α .

Table 1. Mean absolute error for different alpha values Sample size = 1000. MAEs = $|\hat{\alpha}_{MLE} - \alpha|$

	S = 0.2	S = 0.4	S = 0.7
- 0.1	3.199×10^{-4}	1.270×10^{-4}	
	3.199×10 1.377×10^{-4}	1.270×10 2.955×10^{-4}	0.003 0.001
$\alpha = 0.0$ $\alpha = 1.0$	7.695×10^{-4}	0.007	0.001
$\alpha = 1.5$	0.001	0.006	0.004
$\alpha = 3.0$	0.003	0.036	0.004

11.1.1. Updating α algorithm. 3

11.1.2. Montecarlo algorithm.

11.2. Focus on S.

Mean absolute error for different S values Each row is a different S value. Sample size = 1000

	$\alpha = 0.1$	$\alpha = 0.5$	$\alpha = 1$	$\alpha = 1.5$	$\alpha = 3$
S = 0.2	0.002	0.002	9.780×10^{-4}	3.709×10^{-4}	9.931×10^{-4}
S = 0.4	0.005	0.003	0.006	0.002	0.003
S = 0.7	0.006	0.002	0.005	0.006	0.007

11.2.1. Montecarlo algorithm. some considerations on my prior.

First, it should be noted that within the same block we observe a substantial variability. We may have players that win against other players of the same cluster, against players of weaker clusters, or that simply do not win much. If we have players that substantially win against players of stronger clusters, it means that they are misclassified. So the fundamental variability driver is to be recognised in the blocks of the defeated players. Winning against Federer is not the same as winning against a newcomer. Those two victories should not be accounted in the same fashion.

However, one could argue that if the block exhibits a large amount of variability, this probably means that we should split the block further in two, to possibly account for different patterns of victories.

Another crucial point is data imbalance. Games are not drawn at random, which means that we will observe strongest player playing more with each other, since they are the on remaining within the tournament for longer periods, and weaker players playing less against the strong ones and more among themselves.

The issue is that by assigning one cluster to a player is equivalent to equip them with a single probability to beat all the players within a given cluster.

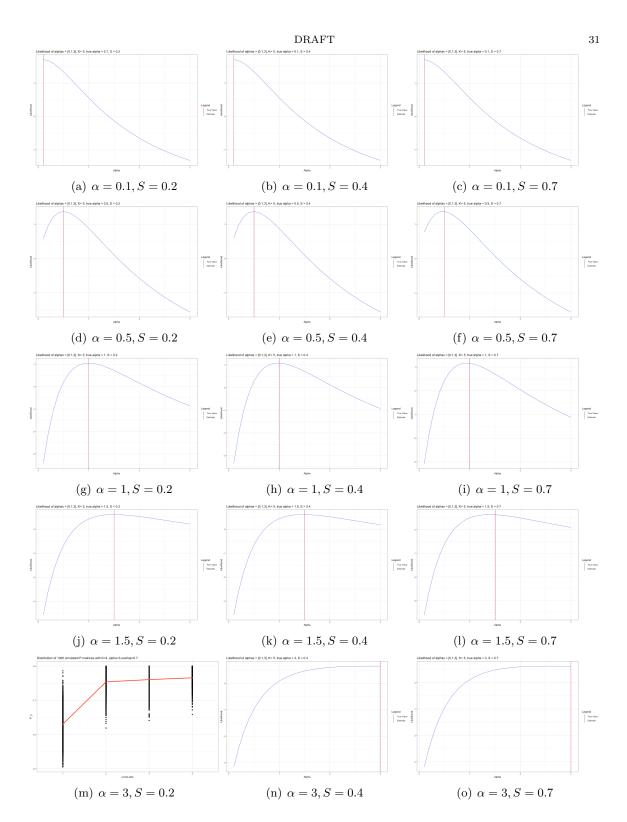


FIGURE 10. Distribution of n=1000 simulated P matrices with $\alpha \in \{0.1,0.5,1,1.5,3\}$ and $S \in \{0.2,0.4,0.7\}$. The points are grouped into the K=5 level sets- the main diagonal is set =0.5. The black vertical dots are the points, while the red lines show the evolution of the mean of the level sets.

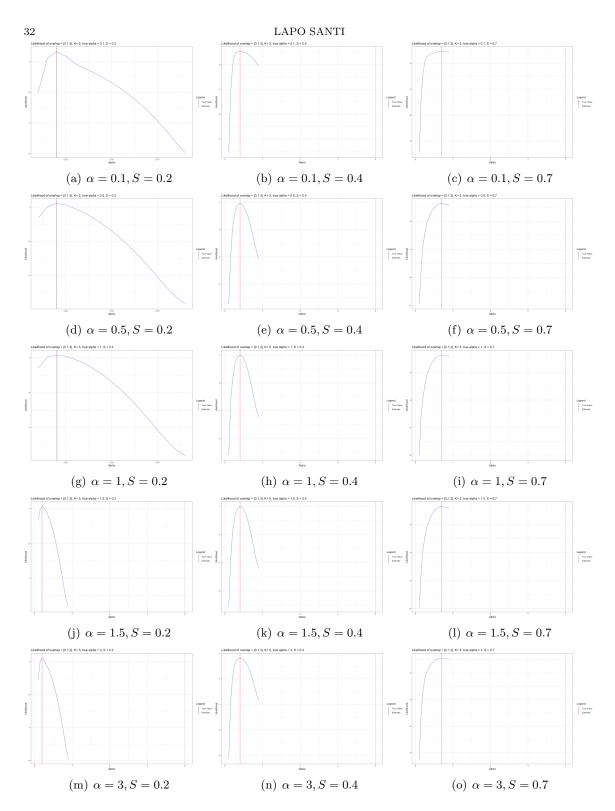


FIGURE 11. Distribution of n=1000 simulated P matrices with $\alpha \in \{0.1,0.5,1,1.5,3\}$ and $S \in \{0.2,0.4,0.7\}$. The points are grouped into the K=5 level sets- the main diagonal is set =0.5. The black vertical dots are the points, while the red lines show the evolution of the mean of the level sets.

12. Appendix I: Estimation Details

12.1. **Updating z.** To update z we propose a new label for each node, we evaluate the accept/reject move by computing the ratio r as follows:

(39)
$$r = \frac{\prod_{i < j} \binom{n_{ij}}{y_{ij}} p_{z_i' z_j'}^{y_{ij}} \cdot (1 - p_{z_i' z_j'})^{n_{ij} - y_{ij}} \cdot \frac{\Gamma(\gamma_0)\Gamma(n+1)}{\Gamma(n+\gamma_0)} \cdot \prod_{k=1}^K \frac{\Gamma(n_k' + \gamma_k)}{\Gamma(\gamma_k)\Gamma(n_k' + 1)}}{\prod_{i < j} \binom{n_{ij}}{y_{ij}} p_{z_i z_j}^{y_{ij}} \cdot (1 - p_{z_i z_j})^{n_{ij} - y_{ij}} \cdot \frac{\Gamma(\gamma_0)\Gamma(n+1)}{\Gamma(n+\gamma_0)} \cdot \prod_{k=1}^K \frac{\Gamma(n_k + \gamma_k)}{\Gamma(\gamma_k)\Gamma(n_k + 1)}}$$

(40)
$$= \frac{\prod_{i < j} p_{z_i' z_j'}^{y_{ij}} \cdot (1 - p_{z_i' z_j'})^{n_{ij} - y_{ij}} \cdot \prod_{k=1}^K \frac{\Gamma(n_k' + \gamma_k)}{\Gamma(\gamma_k) \Gamma(n_k' + 1)}}{\prod_{i < j} p_{z_i z_j}^{y_{ij}} \cdot (1 - p_{z_i z_j})^{n_{ij} - y_{ij}} \cdot \prod_{k=1}^K \frac{\Gamma(n_k + \gamma_k)}{\Gamma(\gamma_k) \Gamma(n_k + 1)}}$$

Passing to the log:

$$log(r) = \log \left(\prod_{i < j} p_{z_{i}'z_{j}'}^{y_{ij}} \cdot (1 - p_{z_{i}'z_{j}'})^{n_{ij} - y_{ij}} \cdot \prod_{k=1}^{K} \frac{\Gamma(n_{k}' + \gamma_{k})}{\Gamma(\gamma_{k})\Gamma(n_{k}' + 1)} \right)$$

$$- \log \left(\prod_{i < j} p_{z_{i}z_{j}}^{y_{ij}} \cdot (1 - p_{z_{i}z_{j}})^{n_{ij} - y_{ij}} \cdot \prod_{k=1}^{K} \frac{\Gamma(n_{k} + \gamma_{k})}{\Gamma(\gamma_{k})\Gamma(n_{k} + 1)} \right)$$

$$= \sum_{i < j} \left(y_{ij} \cdot \log p_{z_{i}'z_{j}'} + (n_{ij} - y_{ij}) \cdot \log (1 - p_{z_{i}'z_{j}'}) \right)$$

$$+ \sum_{k=1}^{K} \left(\log \left(\Gamma(n_{k}' + \gamma_{k}) \right) - \log \left(\Gamma(\gamma_{k}) \right) - \log \left(\Gamma(n_{k}' + 1) \right) \right)$$

$$- \sum_{i < j} \left(y_{ij} \cdot \log p_{z_{i}z_{j}} + (n_{ij} - y_{ij}) \cdot \log (1 - p_{z_{i}z_{j}}) \right)$$

$$- \sum_{k=1}^{K} \left(\log \left(\Gamma(n_{k} + \gamma_{k}) \right) - \log \left(\Gamma(\gamma_{k}) \right) - \log \left(\Gamma(\gamma_{k}) \right) - \log \left(\Gamma(n_{k} + 1) \right) \right)$$

$$(41)$$

12.2. **Updating P.** To update P and α we propose a new label for each node, we evaluate the accept/reject move by computing the ratio r as follows:

(42)
$$r = \frac{\prod_{i < j} \binom{n_{ij}}{y_{ij}} p_{z_i z_j}^{\prime y_{ij}} \cdot (1 - p_{z_i z_j}^{\prime})^{n_{ij} - y_{ij}} \cdot \prod_{k=1}^{K} \left(\frac{1}{y^{\prime(k+1)} - y^{\prime(k)}}\right)^{|L^{\prime(k)}|}}{\prod_{i < j} \binom{n_{ij}}{y_{ij}} p_{z_i z_j}^{y_{ij}} \cdot (1 - p_{z_i z_j})^{n_{ij} - y_{ij}} \cdot \prod_{k=1}^{K} \left(\frac{1}{y^{(k+1)} - y^{(k)}}\right)^{|L^{(k)}|}}$$
(43)

Passing to the log:

Algorithm 2 Updating z step

```
1: for i \leftarrow 1 to N do
 2:
        Sample new_label from 1, ..., K
        Set z' \leftarrow z with the i-th element replaced by new_label
 3:
        Compute new victory probabilities p_{z_i'z_i'} using z'
 4:
        Compute probability ratio log(r) using p_{z'_i z'_i} and p_{z_i z_j}
 5:
        Set \alpha_r \leftarrow \min(1, r)
 6:
        Sample u from a uniform distribution on (0,1)
 7:
        if u < \alpha_r then
 8:
             Update z to z'
 9:
             Update p_{z_i z_j} to p_{z_i' z_j'}
10:
             Increment acc.count_z
11:
12:
        Store z_{current} in z.container
13:
14: end for
```

$$(44)$$

$$log(r) = \sum_{i < j} \left(y_{ij} \cdot \log p'_{z_i z_j} + (n_{ij} - y_{ij}) \cdot \log (1 - p'_{z_i z_j}) \right) - \sum_{k=1}^{K} |L'^{(k)}| \cdot \log \left(y'^{(k+1)} - y'^{(k)} \right)$$

$$(45)$$

$$- \sum_{i < j} \left(y_{ij} \cdot \log p_{z_i z_j} + (n_{ij} - y_{ij}) \cdot \log (1 - p_{z_i z_j}) \right) + \sum_{k=1}^{K} |L^{(k)}| \cdot \log \left(y^{(k+1)} - y'^{(k)} \right)$$

13. Appendix II: POMM Prior Checks

- 13.1. Prior predictive check.
- 13.2. MLE check.

Algorithm 3 Updating P step

```
1: j \leftarrow 1
 2: while j \leq N_{iter} do
         Sample \alpha' from a truncated normal distribution
 3:
         Generate a new proposal matrix P'
 4:
         Compute new victory probabilities p'_{z_iz_j} using P' and z_{current} Compute probability ratio log(r) using p'_{z_iz_j} and p_{z_iz_j}
 5:
 6:
         Set \alpha_r \leftarrow \min(1, r)
 7:
         Sample u from a uniform distribution on (0,1)
 8:
         if u < \alpha_r then
 9:
              Update \alpha to \alpha'
10:
              Update P to P'
11:
              Update p_{z_i z_j} to p'_{z_i z_j}
12:
              Increment acc.count_p
13:
         end if
14:
         Store P in P.container
15:
         Store \alpha in \alpha.container
16:
17:
         j \leftarrow j + 1
18: end while
```