# Assignment 1 Report of DSL

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## 1 Implemented Functionalities

We have implemented the following functionalities:

- Type inference for the underlying type system (including polymorphism, arrays, pairs and lists).
- Support confidentiality labels to the type system.
- Report errors for violations of constant time.
- Support mutable arrays, pairs, lists to the type system.
- (Bonus) Subtyping extension to subeffecting.
- (Bonus) Allow type annotation without labels.

The examples required by the assignment are attached in Section 3.

# 2 Typing Rules

## 2.1 Syntax-driven Rules

$$\begin{array}{c} \ell \sqsubseteq \mathbf{L} \\ \widehat{\Gamma} \vdash_{\mathsf{CTC}} n : \mathsf{Nat}^{\ell} \end{array} [\mathsf{CT}\text{-}\mathsf{Nat}] \\ \\ \frac{\ell \sqsubseteq \mathbf{L}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} \mathbf{true} : \mathsf{Bool}^{\ell}} \ [\mathsf{CT}\text{-}\mathsf{True}] \\ \\ \frac{\ell \sqsubseteq \mathbf{L}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} \mathbf{false} : \mathsf{Bool}^{\ell}} \ [\mathsf{CT}\text{-}\mathsf{False}] \\ \\ \frac{\widehat{\Gamma}[x] = \widehat{\tau}^{\ell}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} x : \widehat{\tau}^{\ell}} \ [\mathsf{CT}\text{-}\mathsf{Var}] \\ \\ \frac{\widehat{\Gamma}[x \mapsto \widehat{\tau}_{1}^{\ell_{1}}] \vdash_{\mathsf{CTC}} e : \widehat{\tau}_{2}^{\ell_{2}}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} \mathbf{fn} \ (x \to e) : (\widehat{\tau}_{1}^{\ell_{1}} \to \widehat{\tau}_{2}^{\ell_{2}})^{\ell_{3}}} \ [\mathsf{CT}\text{-}\mathsf{Fn}] \\ \\ \frac{\widehat{\Gamma}[f \mapsto (\widehat{\tau}_{1}^{\ell_{1}} \to \widehat{\tau}_{2}^{\ell_{2}})^{\ell_{3}}][x \mapsto \widehat{\tau}_{1}^{\ell_{1}}] \vdash_{\mathsf{CTC}} e : \widehat{\tau}_{2}^{\ell_{2}}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} \mathbf{fun} \ (fx \to e) : (\widehat{\tau}_{1}^{\ell_{1}} \to \widehat{\tau}_{2}^{\ell_{2}})^{\ell_{3}}} \ [\mathsf{CT}\text{-}\mathsf{Fun}] \\ \\ \frac{\widehat{\Gamma} \vdash_{\mathsf{CTC}} e_{1} : \widehat{\tau}_{1}^{\ell_{1}} \ \widehat{\Gamma}[x \mapsto \widehat{\tau}_{1}^{\ell_{1}}] \vdash_{\mathsf{CTC}} e_{2} : \widehat{\tau}^{\ell}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} e_{1} : (\widehat{\tau}_{1}^{\ell_{1}} \to \widehat{\tau}^{\ell_{2}})^{\ell_{3}} \ \widehat{\Gamma} \vdash_{\mathsf{CTC}} e_{2} : \widehat{\tau}^{\ell}} \ [\mathsf{CT}\text{-}\mathsf{App}] \\ \\ \frac{\widehat{\Gamma} \vdash_{\mathsf{CTC}} e_{1} : \mathsf{Bool}^{\ell_{1}} \ \widehat{\Gamma} \vdash_{\mathsf{CTC}} e_{2} : \widehat{\tau}^{\ell} \ \widehat{\Gamma} \vdash_{\mathsf{CTC}} e_{3} : \widehat{\tau}^{\ell} \ \ell_{1} \sqsubseteq \mathbf{L}}{\widehat{\Gamma} \vdash_{\mathsf{CTC}} \mathbf{if} \ e_{1} \ \mathbf{then}} \ e_{2} \ \mathbf{else} \ e_{3} : \widehat{\tau}^{\ell} \ \ell_{1} \sqsubseteq \mathbf{L} \ [\mathsf{CT}\text{-}\mathsf{If}] \\ \\ \end{array}$$

## 2.2 Subtyping Rules

$$\begin{split} & \overline{\mathbf{L}} \sqsubseteq \overline{\mathbf{L}} \quad [\text{CT-ST-Low}] \\ & \overline{\mathbf{H}} \sqsubseteq \overline{\mathbf{H}} \quad [\text{CT-ST-High}] \\ & \overline{\mathbf{L}} \sqsubseteq \overline{\mathbf{H}} \quad [\text{CT-ST-LowHigh}] \\ & \frac{\ell_1 \sqsubseteq \ell_2}{\text{Nat}^{\ell_1} \leqslant \text{Nat}^{\ell_2}} \quad [\text{CT-ST-Nat}] \\ & \frac{\ell_1 \sqsubseteq \ell_2}{\text{Bool}^{\ell_1} \leqslant \text{Bool}^{\ell_2}} \quad [\text{CT-ST-Bool}] \\ & \frac{\hat{\tau}_1' \ell_1' \leqslant \hat{\tau}_1 \ell_1}{(\hat{\tau}_1 \ell_1 \to \hat{\tau}_2 \ell_2)^{\ell_3} \leqslant (\hat{\tau}_1' \ell_1' \to \hat{\tau}_2' \ell_2')^{\ell_3'}} \quad [\text{CT-ST-Fun}] \end{split}$$

$$\frac{\ell_{2} \sqsubseteq \ell_{2}'}{(\operatorname{Array} \, \hat{\tau}^{\ell_{1}})^{\ell_{2}}} = (\operatorname{Array} \, \hat{\tau}^{\ell_{1}})^{\ell_{2}'} = [\operatorname{CT-ST-Arr}]$$

$$\frac{\hat{\tau}_{1}^{\ell_{1}} \leqslant \hat{\tau}_{1}'^{\ell_{1}'} \, \hat{\tau}_{2}^{\ell_{2}} \leqslant \hat{\tau}_{2}'^{\ell_{2}'} \, \ell_{3} \sqsubseteq \ell_{3}'}{(\hat{\tau}_{1}^{\ell_{1}}, \hat{\tau}_{2}^{\ell_{2}})^{\ell_{3}}} = [\operatorname{CT-Pair}]$$

$$\frac{\hat{\tau}^{\ell} \leqslant \hat{\tau}'^{\ell'} \, \ell_{1} \sqsubseteq \ell_{1}'}{(\operatorname{List} \, \hat{\tau}^{\ell})^{\ell_{1}}} = [\operatorname{CT-ST-List}]$$

$$\frac{\ell \sqsubseteq \ell'}{\hat{\tau}^{\ell} \leqslant \hat{\tau}'^{\ell'}} = [\operatorname{CT-ST-TypVar}]$$

#### 3 Examples

In this section, we provide code examples (including legal ones, which will be successfully compiled; and illegal ones, which will be reported as an error).

#### 3.1 List

### 3.1.1 Pass

```
1 (1 : []) :: (List Nat^H)^H
1 let xs = (1 :: Nat^H) : 2 : [] in case xs of y : ys -> y, [] -> 1
```

### 3.1.2 Fail

```
let l = (1 : []) :: (List Nat^H)^H in case l of x : xs -> 1, [] -> 2
```

#### 3.2 Array

#### 3.3 Pass

```
1 let xs = (array 10 0 :: a0^H) in (xs, xs[0])
1 let xs = (array 10 0 :: a0^H) in xs[0] = xs[1]
let xs = (array 10 0 :: a0^H) in xs[0] = 1
```

## 3.3.1 Fail (violates constant time)

```
let len = 1 :: Nat^H in array len 0
1 let xs = (array 10 0 :: a0^H) in xs[0] / 3
let xs = (array 10 0 :: (Array Nat^L)^L) in xs[0] = (1 :: Nat^H)
```

#### 3.4 Pair

### 3.4.1 Pass

```
1 (1, 2) :: (Nat^H, Nat^L)^H
let p = (1, 2) :: a0^H in case p of (x, y) -> x+1
```

### 3.4.2 Fail

```
let p = (1, 2) :: a0^H in case p of (x, y) -> x/3
let gx = (fn p -> case p of (x, y) -> x / 3) in gx ((1, 2) :: (Nat^H, Nat^H)^L)
```

## 3.5 Subeffecting & Subtyping

## 3.5.1 Subeffecting

```
1 let add = (fn x -> x + 1) :: (Nat^L -> Nat^H)^L in add 1
1 fn a -> if true then a else false :: Bool^H
1 fun f x -> if x < 1 then 0 :: Nat^H else 1 + f (x - 1)</pre>
```

## 3.5.2 Subtyping

```
1 (fn x -> fn y -> (x, y) :: a0^H) :: (Nat^b0 -> (Nat -> (Nat, Nat)^b0))
1 (3 :: Nat^b0) / 1
1 let p = (1, 2) :: a0^b0 in case p of (x, y) -> x/3
```