UNdead's notebook (2017)

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```

1 Graphs

1.1 Articulation points

```
/*Tarjan Algorithm to find articulation points
            * single dfs O(|v| + |e|)
            * visited =[false]
            * disc = [0]
            \star low = [0]
            * parent = [-1]
            * ap = [false] */
           void articulation(vector<vector<int> > G, int u, bool visited[], int disc[], int low[], int parent
                 [], bool ap[]) {
5
               static int time = 0;
               int children = 0;
               visited[u] = true;
               disc[u] = low[u] = ++time;
               for(int i = 0; i < G[u].size(); i++){</pre>
                  int v = G[u][i];
                  if(!visited[v]){
                       children++:
                       parent[v] = u:
9
                       articulation (G, v, visited, disc, low, ap);
                       low[u] = min(low[u], low[v]);
                       if(parent[u] == -1 && children > 1) ap[u] = true;
                       if(parent[u] != -1 && low[v] >= disc[u]) ap[u] = true;
                   else if(v != parent[u])
                       low[u] = min(low[u], disc[v]);
```

1.2 Biconected graph

```
if(!visited[v]){
    children++;
    parent[v] = u;

    if (isBiconnected(G, v, visited, disc, low, parent)) return true;

    low[u] = min(low[u], low[v]);

    if(parent[u] == -1 && children > 1) return true;

    if(parent[u] != -1 && low[v] >= disc[u]) return true;
}
else if(v != parent[v]) low[u] = min(low[u], disc[v]);
return false;
}
```

1.3 Bridges

```
/* Tarian Algorithm to find bridges
* single dfs O(|v| + |e|)
* visited =[false]
* disc = [0]
\star low = [0]
* parent = [-1] */
 void bridges(vector<vector<int> > G, int u, bool visited[], int disc[], int low[], int parent[],
      priority_queue< pair<int, int> > *bridge) {
   static int time = 0;
   int children = 0:
   visited[u] = true;
   disc[u] = low[u] = ++time;
   for(int i = 0; i < G[u].size(); i++) {</pre>
       int v = G[u][i];
       if(!visited[v]){
            parent[v] = u;
           bridges(G, v, visited, disc, low, bridge);
           low[u] = min(low[u], low[v]);
           if(low[v] > disc[u])bridge->push({u,v});
       else if(v != parent[u])
           low[u] = min(low[u], disc[v]);
```

1.4 Tarjan SCC

```
if(disc[v] == -1) {
            dfsSCC(G, v, disc, low, st, stackmember);
            low[u] = min(low[u], low[v]);
        else if(stackmember[v] == true) low[u] = min(low[u], disc[v]);
    int w = 0:
   if(low[u] == disc[u]){
        while (st->top() != u) {
            w = st->top();
            cout << w<< " ";
            stackmember[w] = false;
            st->pop();
        w = st->top();
       cout << w << "\n";
        stackmember[w] = false;
        st->pop();
void scc(G) {
  int *disc = new int[V];
  int *low = new int[V];
  bool *stackMember = new bool[V];
  stack<int> *st = new stack<int>();
  memset(disc, -1, sizeof(disc));
  memset(low, 0, sizeof(low));
  memset(stackMember, false, sizeof(stackMember));
  for(int i = 0; i < G.size(); i++)</pre>
      if(disc[i] == -1) dfsScc(G, i, disc, low, st, stackMember);
```

1.5 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
// Running time: O(|E| log |V|)
#include <queue>
#include <cstdio>
using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;
int main() {
        scanf("%d%d%d", &N, &s, &t);
        vector<vector<PII> > edges(N);
        for (int i = 0; i < N; i++) {
                int M;
scanf("%d", &M);
                for (int j = 0; j < M; j++) {
                        int vertex, dist:
                        scanf("%d%d", &vertex, &dist);
                        edges[i].push_back(make_pair(dist, vertex)); // note order of arguments here
        // use priority queue in which top element has the "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
        vector<int> dist(N, INF), dad(N, -1);
        Q.push(make_pair(0, s));
        dist[s] = 0;
        while (!Q.empty()) {
                PII p = Q.top();
                Q.pop();
                int here = p.second;
if (here == t) break;
                if (dist[here] != p.first) continue;
                for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++)
```

1.6 Bellman Ford

```
// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function returns
// false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
// Running time: O(|V|^3)
     INPUT: start, w[i][j] = cost of edge from i to j
     OUTPUT: dist[i] = min weight path from start to i
              prev[i] = previous node on the best path from the
                         start node
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w. VT &dist. VI &prev. int start) {
 int n = w.size();
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
  dist[start] = 0;
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
      for (int j = 0; j < n; j++) {
        if (dist[j] > dist[i] + w[i][j]) {
          if (k == n-1) return false;
dist[j] = dist[i] + w[i][j];
          prev[j] = i;
  return true;
```

1.7 Eulerian Path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next_vertex;
        iter reverse_edge;
        Edge(int next_vertex)
               :next_vertex(next_vertex)
};
const int max vertices = ;
int num_vertices;
list<Edge> adj[max_vertices];
                                        // adjacency list
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find_path(vn);
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

1.8 Topological sort

```
char c[TAM];
int l[TAM];
int r[TAM];
int in[TAM];
//can be priority queue
queue<int> Q;
    memset(1, 0, sizeof 1);
    memset(r, 0, sizeof r);
    memset(in, 0, sizeof in);
memset(balls, 0, sizeof balls);
    c[0] = 'L';
void topo(vector<vector<int> > G, int u) {
    while(!Q.empty()){
        u = Q.front(); Q.pop();
         update(u);
         for(int i = 0; i <G[u].size(); i++) {</pre>
             int v = G[u][i];
             if(in[v] == 0) Q.push(v);
int main(){
    11 n:
    int m;
    while(cin>>n>>m) {
         reset();
         vector< vector<int> > G(m + 1);
```

```
for(int i = 1; i <=m; i++) {
    int u,v;
    cin>>c(i)>>u>v;
    G[i].push_back(u);
    G[i].push_back(v);
    in[u] ++; in[v]++;
    1[i] = u;r[i] = v;
}
for(int i = 1; i <=m; i++) {
    if(in[i] == 0)Q.push(i);
}
topo(G, 1);
}</pre>
```

1.9 Kruskal's minimum spanning tree

```
Uses Kruskal's Algorithm to calculate the weight of the minimum spanning
forest (union of minimum spanning trees of each connected component) of
a possibly disjoint graph, given in the form of a matrix of edge weights
(-1\ if\ no\ edge\ exists) . Returns the weight of the minimum spanning
forest (also calculates the actual edges - stored in T). Note: uses a
disjoint-set data structure with amortized (effectively) constant time per
union/find. Runs in O(E*log(E)) time.
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
typedef int T;
struct edge
  int u. v:
  T d:
};
struct edgeCmp
  int operator()(const edge& a, const edge& b) { return a.d > b.d; }
 \textbf{int} \  \, \text{find(vector <int>\& C, int x)} \  \, \{ \  \, \textbf{return} \  \, (\texttt{C[x]} \ == \ x) \  \, ? \  \, x \  \, : \  \, \texttt{C[x]} \  \, = \  \, \text{find(C, C[x]);} \  \, \} 
T Kruskal (vector <vector <T> >& w)
  int n = w.size():
  T \text{ weight} = 0;
  vector <int> C(n), R(n);
  for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }</pre>
  priority_queue <edge, vector <edge>, edgeCmp> E;
  for(int i=0; i < n; i++)
    for (int j=i+1; j<n; j++)
  if (w[i][j] >= 0)
         e.u = i; e.v = j; e.d = w[i][j];
         E.push(e);
  while (T.size() < n-1 && !E.empty())</pre>
    edge cur = E.top(); E.pop();
    int uc = find(C, cur.u), vc = find(C, cur.v);
    if(uc != vc)
       T.push_back(cur); weight += cur.d;
       if(R[uc] > R[vc]) C[vc] = uc;
       else if(R[vc] > R[uc]) C[uc] = vc;
       else { C[vc] = uc; R[uc]++; }
```

1.10 Lowest Common ancestor

```
const int max nodes, log max nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                          // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1];
                                          // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
      ancestor does not exist
int L[max_nodes];
                                          // L[i] is the distance between node i and the root
// floor of the binary logarithm of \boldsymbol{n}
int lb(unsigned int n)
    if(n==0)
       return -1:
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16;
    if (n >= 1<< 8) { n >>= 8; p += 8;
    if (n >= 1<< 4) { n >>= 4; p += 4; }
    if (n >= 1 << 2) { n >>= 2; p += 2;
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    for(int j = 0; j < children[i].size(); j++)
    DFS(children[i][j], l+1);</pre>
int LCA(int p, int q)
     // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the same level as q
    for(int i = log_num_nodes; i >= 0; i--)
if(L[p] - (1<<i) >= L[q])
            p = A[p][i];
    if(p == q)
        return p;
     // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if (A[p][i] != -1 && A[p][i] != A[q][i])
            p = A[p][i];
            q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
    // read num nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
```

```
for(int i = 0; i < num_nodes; i++)
{
    int p;
    // read p, the parent of node i or -1 if node i is the root
    A[i][0] = p;
    if(p != -1)
        children[p].push_back(i);
    else
        root = i;
}

// precompute A using dynamic programming
for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)
    if(A[i][j-1]!= -1)
        A[i][j] = A[A[i][j-1]][j-1];
    else
    A[i][j] = -1;

// precompute L
DFS(root, 0);</pre>
```

2 Geometry

2.1 Geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT(double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
  PT operator * (double c) const { return PT(x*c, y*c ); }
  PT operator / (double c)
                                const { return PT(x/c, y/c ); }
};
                             { return p.x*q.x+p.y*q.y; }
double dot (PT p, PT q)
double dist2(PT p, PT q)
                             { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {</pre>
  os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT(-p.y,p.x); }
PT RotateCW90 (PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
```

```
if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c)
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                         double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
   if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
     dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
     return false:
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2:
  c = (a+c)/2:
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
   int j = (i+1)%p.size();
   return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
   if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
     return true:
   return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
```

```
b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret:
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
 // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs (ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++){</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
 for (int i = 0; i < p.size(); i++) {
  for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false:
  return true:
int main() {
  // expected: (-5.2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
```

```
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
           << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
           << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
           << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
           << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
<< SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push back (PT(5,5));
v.push back(PT(0.5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
           << PointInPolygon(v, PT(2,0)) << " "
           << PointInPolygon(v, PT(0,2)) << " "
           << PointInPolygon(v, PT(5,2)) << " "
           << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
           << PointOnPolygon(v, PT(2,0)) << " "
           << PointOnPolygon(v, PT(0,2)) << " "
           << PointOnPolygon(v, PT(5,2)) << " "
           << PointOnPolygon(v, PT(2,5)) << endl;</pre>
 // expected: (1,6)
                            (5,4) (4,5)
                            blank line
                            (4,5) (5,4)
                            blank line
                            (4,5) (5,4)
if (int i = 0; i < u.size(); i++) cerr < v(i) FT(1, 1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] < r "; cerr << endl;
u = CircleCircleIntersection(PT(1, 1), PT(10, 10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << r "; cerr << endl;
u = CircleCircleIntersection(PT(1, 1), PT(10, 10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << r "; cerr << endl;
u = CircleCircleIntersection(PT(1, 1), PT(8, 8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleIntersection(PT(1, 0), U.size(); u.size(
u = CircleCircleIntersection(PT(1,1), PT(4.5, 4.5), 10, sqrt(2.0)/2.0);

for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
 // centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p):
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0:
```

2.2 Geometry (Java)

```
// In this example, we read an input file containing three lines, each
// containing an even number of doubles, separated by commas. The first two
// lines represent the coordinates of two polygons, given in counterclockwise
// (or clockwise) order, which we will call "A" and "B". The last line
// contains a list of points, p[1], p[2], ...
//
// Our goal is to determine:
// (1) whether B - A is a single closed shape (as opposed to multiple shapes)
// (2) the area of B - A
// (3) whether each p[i] is in the interior of B - A
```

```
0 0 10 0 0 10
// 0 0 10 10 10 0
// 5 1
// OUTPUT:
    The area is singular.
     The area is 25.0
    Point belongs to the area.
    Point does not belong to the area.
import java.util.*;
import java.awt.geom.*;
import java.io.*;
public class JavaGeometry {
    // make an array of doubles from a string
    static double[] readPoints(String s) {
        String[] arr = s.trim().split("\\s++");
        double[] ret = new double[arr.length];
        for (int i = 0; i < arr.length; i++) ret[i] = Double.parseDouble(arr[i]);</pre>
        return ret;
    // make an Area object from the coordinates of a polygon
    static Area makeArea(double[] pts) {
        Path2D.Double p = new Path2D.Double();
        p.moveTo(pts[0], pts[1]);
        for (int i = 2; i < pts.length; i += 2) p.lineTo(pts[i], pts[i+1]);</pre>
        p.closePath();
        return new Area(p);
    // compute area of polygon
    static double computePolygonArea(ArrayList<Point2D.Double> points) {
        Point2D.Double[] pts = points.toArray(new Point2D.Double[points.size()]);
        double area = 0:
        for (int i = 0; i < pts.length; i++) {</pre>
            int j = (i+1) % pts.length;
area += pts[i].x * pts[j].y - pts[j].x * pts[i].y;
        return Math.abs(area)/2;
    // compute the area of an Area object containing several disjoint polygons
    static double computeArea(Area area) {
        double totArea = 0;
        PathIterator iter = area.getPathIterator(null);
        ArrayList<Point2D.Double> points = new ArrayList<Point2D.Double>();
        while (!iter.isDone()) {
            double[] buffer = new double[6];
            switch (iter.currentSegment(buffer)) {
case PathIterator.SEG_MOVETO:
            case PathIterator.SEG LINETO:
                 points.add(new Point2D.Double(buffer[0], buffer[1]));
            case PathIterator.SEG_CLOSE:
                totArea += computePolygonArea(points);
                 points.clear();
                break:
            iter.next();
        return totArea:
    // notice that the main() throws an Exception -- necessary to
    // avoid wrapping the Scanner object for file reading in a
    // try { ... } catch block.
    public static void main(String args[]) throws Exception {
        Scanner scanner = new Scanner(new File("input.txt"));
        // Scanner scanner = new Scanner (System.in);
        double[] pointsA = readPoints(scanner.nextLine());
        double[] pointsB = readPoints(scanner.nextLine());
        Area areaA = makeArea(pointsA);
        Area areaB = makeArea(pointsB);
        areaB.subtract(areaA);
        // also.
        // areaB.exclusiveOr (areaA);
            areaB.add (areaA);
        // areaB.intersect (areaA);
        // (1) determine whether B - A is a single closed shape (as
               opposed to multiple shapes)
```

```
boolean isSingle = areaB.isSingular();
// also.
    areaB.isEmpty();
if (isSingle)
    System.out.println("The area is singular.");
else
    System.out.println("The area is not singular.");
// (2) compute the area of B - A
System.out.println("The area is " + computeArea(areaB) + ".");
// (3) determine whether each p[i] is in the interior of B - A
while (scanner.hasNextDouble()) {
    double x = scanner.nextDouble();
    assert(scanner.hasNextDouble());
    double y = scanner.nextDouble();
    if (areaB.contains(x,y)) {
        System.out.println ("Point belongs to the area.");
        System.out.println ("Point does not belong to the area.");
// Finally, some useful things we didn't use in this example:
     Ellipse2D.Double ellipse = new Ellipse2D.Double (double x, double y,
                                                        double w, double h);
       creates an ellipse inscribed in box with bottom-left corner (x,y)
       and upper-right corner (x+y, w+h)
     Rectangle2D.Double rect = new Rectangle2D.Double (double x, double y,
                                                         double w, double h);
       creates a box with bottom-left corner (x,y) and upper-right
       corner (x+y, w+h)
// Each of these can be embedded in an Area object (e.g., new Area (rect)).
```

2.3 Convex Hull Monotone Chain

```
#include<iostream>
#include < algorithm>
#include < complex >
#include<cstdio>
#include<iomanip>
#include<vector>
#define x real()
#define y imag()
#define dot(A,B)
                                real(conj((A))*(B))
#define cross(A,B)
                                imag(conj((A))*(B))
#define PI 3.1415926
#define EPS 1e-9
using namespace std;
typedef double 1f;
typedef complex<lf> pt;
istream& operator >> ( istream& in, pt& p ) {
 lf a.n: in >> a >> n:
 p = pt(a,n); return in;
bool cmp(pt &p, pt &q) {
   if(p.x != q.x) return p.x < q.x;
   return p.y < q.y;
bool is_zero( lf x ) {
   return -EPS <= x && x <= EPS;
inline bool same ( lf a, lf b ) {
  return a+EPS > b && b+EPS > a:
int ccw(pt& p1, pt& p2, pt& p3) {
  lf ans = (cross(p1 - p3, p2 - p3));
  if (-EPS <= ans && ans <= EPS)
```

```
return 0;
  else if(ans <= -EPS)</pre>
    return -1;
    return 1:
lf dist ( pt A, pt B ) { return abs(A-B); }
vector<pt> convex_hull(vector<pt> P) {
        int n = P.size(): int k = 0:
        vector<pt> H(2*n);
        sort(P.begin(), P.end(), cmp);
        for (int i = 0; i < n; i++) {
                 while (k \ge 2 \&\& ccw(H[k-2], H[k-1], P[i]) == 1) k--;
                 H[k++] = P[i];
        for (int i = n-2, t = k+1; i >= 0; i--) {
    while (k >= t && ccw(H[k-2], H[k-1], P[i]) == 1) k--;
                 H[k++] = P[i];
        if(n>1)
        H.resize(k);
        return H;
int main(){
  ios_base::sync_with_stdio(0);
  #ifdef LOCAL
  //freopen("in.txt", "r", stdin);
  //freopen("out.txt", "w", stdout);
  #endif // LOCAL
  int n:
  cin>>n;
  vector<pt> p(n);
  for(int i = 0; i < n; i++)
    cin>>p[i];
  vector<pt> hull = convex_hull(p);
  for(auto &pt: hull)
    cout<<pt<<endl;
```

2.4 Delaunay triangulation

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
            x[] = x-coordinates
            y[] = y-coordinates
// OUTPUT: triples = a vector containing m triples of indices
                       corresponding to triangle vertices
#include<vector>
using namespace std;
typedef double T;
struct triple {
    int i, j, k;
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
        int n = x.size();
        vector<T> z(n);
        vector<triple> ret;
        for (int i = 0; i < n; i++)
           z[i] = x[i] * x[i] + y[i] * y[i];
        for (int i = 0; i < n-2; i++) {
```

```
for (int j = i+1; j < n; j++) {
                for (int k = i+1; k < n; k++) {
                    if (j == k) continue;
                    double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                    double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                    double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                    bool flag = zn < 0;
                    for (int m = 0; flag && m < n; m++)</pre>
                        flag = flag && ((x[m]-x[i])*xn +
                                         (y[m]-y[i])*yn +
                                         (z[m]-z[i])*zn <= 0);
                    if (flag) ret.push_back(triple(i, j, k));
        return ret;
int main()
    T \times s[] = \{0, 0, 1, 0.9\};
    T ys[] = {0, 1, 0, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    //expected: 0 1 3
    int i:
    for(i = 0; i < tri.size(); i++)</pre>
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
```

2.5 Delaunay triangulation (java)

```
// Slow but simple Delaunay triangulation. (from O'Rourke,
// Computational Geometry in C)
// Running time: O(n^4)
// INPUT:
              x[] = x-coordinates
              y[] = y-coordinates
             ret[][] = an mx3 matrix containing m triples of indices
                           corresponding to triangle vertices
import java.util.*;
public class Delaunay {
    int[][] triangulate(double[] x, double[] y) {
         int n = x.length;
         double z[] = new double[n];
         ArrayList<int[]> ret = new ArrayList<int[]>();
         for (int i = 0; i < n; i++)
   z[i] = x[i] * x[i] + y[i] * y[i];</pre>
        for (int i = 0; i < n-2; i++) {
   for (int j = i+1; j < n; j++) {
      for (int k = i+1; k < n; k++) {</pre>
                       if (j == k) continue;
                       double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                       double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                       double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                      boolean flag = zn < 0;

for (int m = 0; flag && m < n; m++)

flag = flag && ((x[m]-x[i])*xn +
                                               (y[m]-y[i])*yn +
                                               (z[m]-z[i])*zn <= 0);
                       if (flag) ret.add(new int[]{i, j, k});
         return ret.toArray(new int[0][0]);
```

3 Data Structures

3.1 Mo's algorithm

#include <bits/stdc++.h>

```
#define TAM 30000 + 7
#define OTAM 200000 + 7
#define MTAM 1000000 + 7
#define whatis(x) cerr<<#x<<" is "<<x<<endl
using namespace std;
int a[TAM], r[QTAM], cnt[MTAM];
int ans, BLOCK, currL , currR ;
    int L, R, idx;
}q[QTAM];
bool comp(node a, node b) {
   if(a.L/BLOCK < b.L/BLOCK) return true;</pre>
    if(a.L/BLOCK > b.L/BLOCK) return false;
    return a.R < b.R;
void remove(int i) {
     if (cnt [a[i]] == 0) ans--;
void add(int i) {
     cnt[a[i]]++;
    if(cnt[a[i]]==1)ans++;
int query(node i) {
    while (currL< i.L) {
        remove (currL):
        currL++:
     while (currL > i.L) {
        currL--;
         add(currL);
     while (currR< i.R) {
         currR++;
         add(currR);
     while (currR > i.R) {
        remove (currR);
         currR--;
    return ans:
     ios_base::sync_with_stdio(0);
    cin.tie(0);
     #ifdef LOCAL
    freopen("in", "r", stdin);
     #endif
    int n, que;
    BLOCK = sqrt(n);
    for (int i = 1; i <= n; i++)cin>>a[i];
    for (int i = 1; i <= que; i++) {
         cin>>q[i].L>>q[i].R;
        q[i].idx = i;
    sort(q +1, q + que+1, comp);
    for(int i = 1; i <=que; i++)</pre>
         r[q[i].idx] = query(q[i]);
    for(int i = 1; i <= que; i++)</pre>
        cout<<r[i]<<"\n";
```

3.2 Segment Trees with lazy propagation

```
//querys and build takes O(log n)
//example with segment sum
#include <bits/stdc++.h>
using namespace std:
long long *p;
//long long *lazy;
struct SegmentTree(
    SegmentTree *L, *R;
    long long sum = 0;
    long long lazy = 0;
    long long query2(int a, int b) {
        if(a == 1 && b == r) return sum;
        if(b <= L->r) return L->query(a,b);
        if(a >= R->1) return R->query(a,b);
        return (L->query2(a,L->r) + R->query2(R->1, b));
    void update(int a, int val){
        if(1 == r){
            sum += val:
             return;
        int mid = (1 + r)/2;
        if (1 <= a && a <= mid)
            L->update(a, val);
             R->update(a, val);
         sum = L -> sum + R -> sum;
    void updateRange2(int a, int b, long long val) {
        if(b < 1 or a > r)
            return:
        if(1 == r){
             sum += val:
             return;
        L->updateRange2(a, b, val);
        R->updateRange2(a,b,val);
        sum = L -> sum + R -> sum;
    void updateRange(int a, int b, long long val) {
        if(lazy != 0) {
            sum += (r-1+1)*lazy;
             //sum += lazy;
            if(1 != r) {
                 R->lazy = lazy + R->lazy;
                 L->lazy = lazy + L->lazy;
            lazy = 0;
        if(b < 1 or a > r)
            return;
        if(1 >= a && r <= b) {</pre>
             sum += (r-l+1)*val;
            //sum += val;
if(1 != r){
                 R->lazy = val + R->lazy;
                 L->lazy = val + L->lazy;
             return;
        L->updateRange(a, b, val);
        R->updateRange(a,b,val);
        sum = L \rightarrow sum + R \rightarrow sum;
    long long query(int a, int b) {
        if(b < 1 or a > r)
            return 0;
        if(lazy != 0) {
            sum += (r-1+1)*lazy;
//sum += lazy;
             if(1 != r){
                 R->lazy = lazy + R->lazy;
L->lazy = lazy + L->lazy;
             lazy = 0;
```

```
if (a == 1 && b == r) return sum;
        if(b <= L->r) return L->query(a,b);
        if(a >= R->1) return R->query(a,b);
        return (L->query(a,L->r) + R->query(R->1, b));
    SegmentTree(int a, int b): l(a), r(b){
        if(a == b){
            sum = p[a];
            L = R = nullptr;
        else{
            L = new SegmentTree (a, (a+b)/2);
            R = new SegmentTree ((a+b)/2 + 1, b);
            sum = L \rightarrow sum + R \rightarrow sum;
int main(){
    cin.tie(0);
    ios_base::sync_with_stdio(0);
    #ifdef LOCAL
        freopen("input.txt", "r", stdin);
        long long T;
        cin >> T;
    while (T--) {
        long long n, c;
        long long l[n];
        memset(1,0,sizeof(1));
        p = 1;
        SegmentTree *stree = new SegmentTree(0, n-1);
        while (c--) {
            long long aux, p, q;
            cin >> aux >> p >> q;
if(aux == 0){
                 long long val;
                 stree->updateRange(p-1, q-1, val);
                 cout << stree->query(p-1, q-1) << endl;</pre>
```

3.3 Segment Trees with lazy propagation (Java)

```
public class SegmentTreeRangeUpdate {
       public long[] leaf;
        public long[] update;
        public int origSize;
        public SegmentTreeRangeUpdate(int[] list)
               origSize = list.length;
                leaf = new long[4*list.length];
               update = new long[4*list.length];
               build(1,0,list.length-1,list);
        public void build(int curr, int begin, int end, int[] list)
               if(begin == end)
                       leaf[curr] = list[begin];
                        int mid = (begin+end)/2;
                       build(2 * curr, begin, mid, list);
                       build(2 * curr + 1, mid+1, end, list);
                        leaf[curr] = leaf[2*curr] + leaf[2*curr+1];
        public void update(int begin, int end, int val) {
                update(1,0,origSize-1,begin,end,val);
        public void update(int curr, int tBegin, int tEnd, int begin, int end, int val)
               if (tBegin >= begin && tEnd <= end)
                       update[curr] += val;
```

```
leaf[curr] += (Math.min(end,tEnd)-Math.max(begin,tBegin)+1) * val;
                int mid = (tBegin+tEnd)/2;
                if (mid >= begin && tBegin <= end)
                       update(2*curr, tBegin, mid, begin, end, val);
                if(tEnd >= begin && mid+1 <= end)</pre>
                        update(2*curr+1, mid+1, tEnd, begin, end, val);
public long query(int begin, int end) {
        return query(1,0,origSize-1,begin,end);
public long query(int curr, int tBegin, int tEnd, int begin, int end) {
       if(tBegin >= begin && tEnd <= end)
                if(update[curr] != 0) {
                        leaf[curr] += (tEnd-tBegin+1) * update[curr];
                        if(2*curr < update.length){</pre>
                               update[2*curr] += update[curr];
                                update[2*curr+1] += update[curr];
                        update[curr] = 0;
                return leaf[curr];
        else
                leaf[curr] += (tEnd-tBegin+1) * update[curr];
                if(2*curr < update.length){</pre>
                        update[2*curr] += update[curr];
                        update[2*curr+1] += update[curr];
                update[curr] = 0;
                int mid = (tBegin+tEnd)/2;
                long ret = 0;
                if(mid >= begin && tBegin <= end)</pre>
                        ret += query(2*curr, tBegin, mid, begin, end);
                if(tEnd >= begin && mid+1 <= end)</pre>
                       ret += query(2*curr+1, mid+1, tEnd, begin, end);
                return ret:
```

3.4 Fenwick Tree

```
using namespace std;
#define LOGS% 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ):
// add v to value at x
void set(int x, int v) {
 while (x <= N) {
   tree[x] += v;
   x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
 int res = 0;
 while(x) {
   res += tree[x];
   x -= (x & -x);
 return res:
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
 int idx = 0, mask = N;
  while (mask && idx < N) {
   int t = idx + mask;
   if(x >= tree[t]) {
      idx = t;
     x -= tree[t];
   mask >>= 1:
 return idx;
```

3.5 Union Find (Short)

```
#include <iostream>
#include <vector>
using namespace std;
int find(vector<int> &C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }
void merge(vector<int> &C, int x, int y) { C[find(C, x)] = find(C, y); }
int main()
{
    int n = 5;
    vector<int> C(n);
    for (int i = 0; i < n; i++) C[i] = i;
    merge(C, 0, 2);
    merge(C, 1, 0);
    merge(C, 3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << find(C, i) << endl;
    return 0;
}</pre>
```

3.6 Union Find

```
-----*/
#include <bits/stdc++.h>
#define TAM 10000
using namespace std;
class UnionFind(
    private:
        vector<int> p, rank, ssize;
        int numSets;
    public:
        UnionFind(int N) {
            rank.assign(N, 0);
            ssize.assign(N, 1);
            numSets = N;
            p.assign(N,0);
            for(int i = 0; i < N; i++)
               p[i] = i;
        int findSet(int i){
            return (p[i] == i)? i: (p[i] = findSet(p[i]));
        bool isSameSet(int i, int j) {
            return findSet(i) == findSet(j);
        void unionSet(int i, int j){
            if(!isSameSet(i, j)){
                numSets--;
                int x = findset(i), y = findSet(j);
                if(rank[x] > rank[y]){
                    p[y] = x;
                    ssize[x] += ssize[y];
                else{
                   p[x] = y;
ssize[y] += ssize[x];
if(rank[x] == rank[y])
                       rank[y]++;
        int numDisjointSets(){
            return numSets;
        int sizeOfSet(int i){
            return ssize[findSet(i)];
};
```

4 Strings

4.1 KMP

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
 int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
   while (k \ge -1 & p[k+1] != p[i])

k = (k = -1) ? -2 : pi[k];
   pi[i] = ++k;
int KMP(string& t, string& p)
  VI pi;
 buildPi(p, pi);
  int k = -1;
  while (k \ge -1 \&\& p[k+1] != t[i])
     k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
     // p matches t[i-m+1, ..., i]
cout << "matched at index " << i-k << ": ";
cout << t.substr(i-k, p.length()) << endl;
      k = (k == -1) ? -2 : pi[k];
  return 0;
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
 return 0:
```

4.2 Suffix Array

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
//
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
// of substring s[i...L-1] in the list of sorted suffixes.
// That is, if we take the inverse of the permutation suffix[],
// we get the actual suffix array.

#include <vector>
#include <iostream>
#include <string>
using namespace std;
```

```
// BEGIN CUT
#endif
      tion)
// Convolution using the fast Fourier transform (FFT).
      a[1...n]
      b[1...m]
// OUTPUT:
// DFT or inverse DFT.
#include <iostream>
#include <vector>
#include <complex>
using namespace std;
typedef long double DOUBLE;
typedef complex<DOUBLE> COMPLEX;
typedef vector<DOUBLE> VD;
typedef vector<COMPLEX> VC;
struct FFT {
 VC A:
 int n, L;
 int ReverseBits(int k) {
   int ret = 0;
   for (int i = 0; i < L; i++) {
     ret = (ret << 1) | (k & 1);
   return ret:
  void BitReverseCopy(VC a) {
   A.resize(n);
   for (int k = 0; k < n; k++)
     A[ReverseBits(k)] = a[k];
 VC DFT(VC a, bool inverse) {
   BitReverseCopy(a);
   for (int s = 1; s <= L; s++) {
```

struct SuffixArray {

int len = 0;

return len;

// BEGIN CUT

#define TESTING

string s;

cin >> s;

SuffixArray array(s);

int len = 0, count = 0;

if (1 >= len) {

len = 1;

bestlen = len;

if (bestlen == 0) {

l else (

#else

// END CUT

int main() {

cout << endl;

bestcount = count; bestpos = i;

// bobocel is the O'th suffix // obocel is the 5'th suffix bocel is the 1'st suffix ocel is the 6'th suffix

cel is the 2'nd suffix el is the 3'rd suffix l is the 4'th suffix SuffixArray suffix("bobocel");

// Expected output: 0 5 1 6 2 3 4

vector<int> v = suffix.GetSuffixArray();

for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>

cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>

#ifdef TESTING

int main() {

int T: cin >> T:

vector<vector<int> > P;

vector<pair<int,int>,int> > M;

for (int i = 0; i < L; i++)

sort (M.begin(), M.end()); for (int i = 0; i < L; i++)

for (int i = 0; i < L; i++) P[0][i] = int(s[i]);</pre>

vector<int> GetSuffixArray() { return P.back(); }

// The following code solves UVA problem 11512: GATTACA.

for (int caseno = 0; caseno < T; caseno++) {</pre>

vector<int> v = array.GetSuffixArray();

for (int i = 0; i < s.length(); i++) {</pre>

int bestlen = -1, bestpos = -1, bestcount = 0;

for (int j = i+1; j < s.length(); j++) {</pre> int 1 = array.LongestCommonPrefix(i, j);

cout << "No repetitions found!" << endl;</pre>

cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>

if (1 > len) count = 2; else count++;

int LongestCommonPrefix(int i, int j) {

if (i == j) return L - i;

i += 1 << k;

j += 1 << k;

len += 1 << k;

for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
 P.push_back(vector<int>(L, 0));

for (int k = P.size() - 1; k >= 0 && i < L && j < L; k --) {
 if (P[k][i] == P[k][j]) {

 $SuffixArray(\textbf{const} \ string \ \&s) : L(s.length()), \ s(s), \ P(1, \ vector < \textbf{int} > (L, \ 0)), \ M(L) \ \{ (1, \ 0) \} \}$

// returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]

 $M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);$

P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;

if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {

const int L;

string s;

Fast Fourier Transform (convolu-

```
c[1...n+m-1] such that c[k] = sum_{i=0}^k a[i] b[k-i]
// Alternatively, you can use the DFT() routine directly, which will
// zero-pad your input to the next largest power of 2 and compute the
    for (n = 1, L = 0; n < a.size(); n <<= 1, L++);
      int m = 1 << s;
      COMPLEX wm = exp(COMPLEX(0, 2.0 * M_PI / m));
      if (inverse) wm = COMPLEX(1, 0) / wm;
      for (int k = 0; k < n; k += m) {
        COMPLEX w = 1;

for (int j = 0; j < m/2; j++) {

    COMPLEX t = w * A[k + j + m/2];
          COMPLEX u = A[k + j];
          A[k + j] = u + t;

A[k + j + m/2] = u - t;
          w = w * wm;
    if (inverse) for (int i = 0; i < n; i++) A[i] /= n;</pre>
  // c[k] = sum_{i=0}^k a[i] b[k-i]
  VD Convolution(VD a, VD b) {
   int L = 1;
    while ((1 << L) < a.size()) L++;
    while ((1 << L) < b.size()) L++;
    int n = 1 << (L+1);
    for (size_t i = 0, i < n, i++) aa.push_back(i < a.size() ? COMPLEX(a[i], 0) : 0);
```

```
for (size_t i = 0; i < n; i++) bb.push_back(i < b.size() ? COMPLEX(b[i], 0) : 0);</pre>
    VC AA = DFT(aa, false);
    VC BB = DFT(bb, false);
    for (size_t i = 0; i < AA.size(); i++) CC.push_back(AA[i] * BB[i]);</pre>
    VC cc = DFT(CC, true);
    for (int i = 0; i < a.size() + b.size() - 1; i++) c.push_back(cc[i].real());</pre>
    return c;
};
int main() {
  double a[] = \{1, 3, 4, 5, 7\};
  double b[] = \{2, 4, 6\};
  VD c = fft.Convolution(VD(a, a + 5), VD(b, b + 3));
  // expected output: 2 10 26 44 58 58 42
  for (int i = 0; i < c.size(); i++) cerr << c[i] << " ";</pre>
  cerr << endl;
  return 0:
```

5 Flows

5.1 Ford Fulkerson

```
///---- Ford-Fulkerson O(MaxFlow * |E|) -----
struct OutEdge {
        int to, cap, rIdx;
        OutEdge ( ) { }
        OutEdge(int to, int cap, int rIdx) :
                to(to), cap(cap), rIdx(rIdx) { }
1:
struct Network
        vector<vector<OutEdge> > out;
        vector<bool> seen;
        int sink:
        int augment ( int i, const int cur ) {
                if ( i == sink ) return cur;
                if ( seen[i] ) return false;
                seen[i] = true;
                for ( OutEdge& e : out[i] )
                        if ( e.cap > 0 && ( ans = augment(e.to,min(cur,e.cap)) ) ) {
                               e cap -= ans;
                                out[e.to][e.rIdx].cap += ans;
                                return ans;
                return 0:
        int maxflow ( int source, int _sink ) {
                sink = sink:
                int curflow = 0, aug;
                while ( true ) {
                       fill ( seen.begin(), seen.end(), false );
                        aug = augment(source, INT_MAX);
                        if ( aug == 0 ) break;
                        curflow += aug;
                return curflow;
       void addEdge ( int fr, int to, int c ) {
    assert ( fr != to );
                out[fr].push_back(OutEdge(to, c, out[to].size()));
                out[to].push_back(OutEdge(fr, 0, out[fr].size() - 1));
```

```
Network(int n) {
                 out.assign(n, vector<OutEdge>());
                 seen.resize(n);
// Adjacency list implementation of FIFO push relabel maximum flow
// with the gap relabeling heuristic. This implementation is
// significantly faster than straight Ford-Fulkerson. It solves // random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
// Running time:
      0(|V|^3)
       - graph, constructed using AddEdge()
// OUTPUT:
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
typedef long long LL;
struct Edge {
  int from, to, cap, flow, index;
  Edge(int from, int to, int cap, int flow, int index) :
   from(from), to(to), cap(cap), flow(flow), index(index) {}
struct PushRelabel (
 int N;
  vector<vector<Edge> > G;
  vector<LL> excess;
  vector<int> dist, active, count;
  queue<int> Q;
  PushRelabel(int N) : N(N), G(N), excess(N), dist(N), active(N), count(2*N) {}
  void AddEdge(int from, int to, int cap) {
    G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push\_back(Edge(to, from, 0, 0, G[from].size() - 1));
  void Enqueue(int v) {
   if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
  void Push (Edge &e) {
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
    if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
    e.flow += amt;
    G[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
    excess[e.from] -= amt;
   Enqueue (e.to);
  void Gap(int k) {
   for (int v = 0; v < N; v++) {
   if (dist[v] < k) continue;</pre>
      count [dist[v]]--;
      dist[v] = max(dist[v], N+1);
      count[dist[v]]++;
      Enqueue (v);
  void Relabel(int v) {
    count[dist[v]]--;
   dist[v] = 2:N;
for (int i = 0; i < G[v].size(); i++)
    if (G[v][i].cap - G[v][i].flow > 0)
    dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue(v);
```

void Discharge(int v) {

```
for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);</pre>
    if (excess[v] > 0) {
      if (count[dist[v]] == 1)
         Gap(dist[v]);
       else
        Relabel(v);
  LL GetMaxFlow(int s, int t) {
    count[0] = N-1;
    count[N] = 1;
    dist[s] = N;
    dist[s] = n;
active[s] = active[t] = true;
for (int i = 0; i < G[s].size(); i++) {
  excess[s] += G[s][i].cap;</pre>
      Push (G[s][i]);
    while (!Q.empty()) {
      int v = Q.front();
       Q.pop();
       active[v] = false;
      Discharge(v);
    LL totflow = 0:
    for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;</pre>
    return totflow;
};
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time:
       O(|V|^2 |E|)
       - graph, constructed using AddEdge()
       - source
       - sink
// OUTPUT:
       - maximum flow value
        - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
const int INF = 2000000000;
struct Edge {
  int from, to, cap, flow, index;
  Edge(int from, int to, int cap, int flow, int index) :
     from(from), to(to), cap(cap), flow(flow), index(index) {}
struct Dinic {
  int N;
  vector<vector<Edge> > G;
  vector<Edge *> dad;
  vector<int> 0:
  Dinic(int N) : N(N), G(N), dad(N), Q(N) {}
  void AddEdge(int from, int to, int cap) {
    G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
  long long BlockingFlow(int s, int t) {
    fill(dad.begin(), dad.end(), (Edge *) NULL);
    dad[s] = &G[0][0] - 1;
    int head = 0, tail = 0;
    Q[tail++] = s;
     while (head < tail) {
      int x = Q[head++];
for (int i = 0; i < G[x].size(); i++) {</pre>
         Edge &e = G[x][i];
         if (!dad[e.to] && e.cap - e.flow > 0) {
          dad[e.to] = &G[x][i];
Q[tail++] = e.to;
```

```
if (!dad[t]) return 0;
    long long totflow = 0;
    for (int i = 0; i < G[t].size(); i++) {</pre>
      Edge *start = &G[G[t][i].to][G[t][i].index];
      int amt = INF;
      for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
   if (!e) { amt = 0; break; }
        amt = min(amt, e->cap - e->flow);
      if (amt == 0) continue;
      for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
        e->flow += amt;
        G[e->to][e->index].flow -= amt;
    return totflow;
  long long GetMaxFlow(int s, int t) {
    long long totflow = 0;
    while (long long flow = BlockingFlow(s, t))
  totflow += flow;
    return totflow;
};
```

5.2 Edmons Karp Min cut

```
///---- Edmonds Karp with MinCut O(|V|*|E|^2) -----
struct Network (
       vector<Edge> G[TAM]:
       int from[TAM], n;
       bool color[TAM];
       // Call flood (source) to color one node
        // component of min cut.
       void flood ( int node ) {
               if ( color[node] ) return;
               color[node] = true;
               for ( const Edge& e : G[node] )
                       if ( e.cap > 0 )
                               flood (e.to);
       int maxFlow ( int A, int B )
               int flow = 0;
               while (1) {
                       memset ( from, -1, sizeof(from) );
                       queue<int> q;
                       q.push (A);
                       for ( int i; !q.empty(); q.pop() ) {
                               i = q.front();
                               for ( Edge& e : G[i] )
                                       if ( from[e.to] == -1 && e.cap ) {
                                               from[e.to] = e.invIdx;
                                               q.push (e.to);
                       if ( from[B] == -1 ) break;
                       int aug = INF_CAP;
                       for ( int i = B, j; i != A; i = j ) {
                               j = G[i][from[i]].to;
                               aug = min ( aug, G[j][ G[i][from[i]].invIdx ].cap );
                       for ( int i = B, j; i != A; i = j ) {
                               j = G[i][from[i]].to;
G[j][ G[i][from[i]].invIdx ].cap -= aug;
                               G[i][from[i]].cap += aug;
                       flow += aug;
```

```
return flow;
}

void addNonDirEdge ( int a, int b, int c ) {
    assert ( a != b );
    G[a].push_back ( Edge(b,c,G[b].size()) );
    G[b].push_back ( Edge(a,c,G[a].size()-1) );
}

void addDirEdge ( int a, int b, int c ) {
    assert ( a != b );
    G[a].push_back ( Edge(b,c,G[b].size()) );
    G[b].push_back ( Edge(a,0,G[a].size()-1) );
}

void clear ( int _n ) {
    n = _n;
    memset ( color, false, n );
    for ( int i = 0; i < n; ++i )
        G[i].clear();
}
} netw;</pre>
```

5.3 Hopcroft karp's maximum bipartite matching

```
///----- Hopcroft Karp - Maximum Bipartite Matching O( sqrt(|V|) * |E| ) -----
namespace hoperoftKarp {
        const int MAXN1 = 50000;
        const int MAXN2 = 50000;
        const int MAXM = 150000;
        int n1, n2, edges, last[MAXN1], prev[MAXM], head[MAXM];
        int matching[MAXN2], dist[MAXN1], Q[MAXN1];
        bool used[MAXN1], vis[MAXN1];
        void init(int _n1, int _n2) {
                n1 = _n1;
                 edges = 0;
                 fill(last, last + n1, -1);
        void addEdge(int u, int v) {
                head[edges] = v;
prev[edges] = last[u];
                 last[u] = edges++;
        void bfs() {
                 fill(dist, dist + n1, -1);
                 int sizeQ = 0;
                 for (int u = 0; u < n1; ++u) {
                         if (!used[u]) {
                                  Q[sizeQ++] = u;
                                  dist[u] = 0;
                 for (int i = 0; i < sizeQ; i++) {
    int u1 = Q[i];</pre>
                         for (int e = last[u1]; e >= 0; e = prev[e]) {
                                  int u2 = matching[head[e]];
                                  if (u2 >= 0 && dist[u2] < 0) {
                                           dist[u2] = dist[u1] + 1;
                                           Q[sizeQ++] = u2;
        bool dfs(int u1) {
                 vis[u1] = true;
                 for (int e = last[u1]; e >= 0; e = prev[e]) {
   int v = head[e];
                         int u2 = matching[v];
                         if (u2 < 0 || !vis[u2] && dist[u2] == dist[u1] + 1 && dfs(u2)) {
    matching[v] = u1;</pre>
                                  used[u1] = true;
```

5.4 Maxmium bipartite matching (short but slower)

```
///---- Maximum Bipartite Matching O(|V|*|E|) -----
bool findMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
    for (int j = 0; j < int(w[i].size()); j++) {</pre>
                if (w[i][j] && !seen[j]) {
                        seen[j] = true;
                        if (mc[j] < 0 || findMatch(mc[j], w, mr, mc, seen)) {
    mr[i] = j; mc[j] = i;</pre>
                                 return true;
        return false;
int maxBipartiteMatching(const VVI &w ) {
        if ( w.empty() || w[0].empty() ) return 0;
        VI mr(w.size(),-1), mc(w[0].size(), -1);
        int ct = 0;
        for (int i = 0; i < int(w.size()); i++) {</pre>
                VI seen(w[0].size());
                if (findMatch(i, w, mr, mc, seen)) ct++;
        return ct;
```

6 Math

6.1 general math tricks

```
long square(long n) { return n*n;}
int fastPow(long x, long n) {
   if(n == 0)
      return 1;

   if(n % 2 == 0)
      return square(fastPow(x, n/2));

   return x * (fastPow(x, n - 1));
}
/* LCM */
```

```
int LCM(int m, n) {return (m*n) /__gcd(m, n); }
int main() {
    /* n es impar?*/
    odd = ((n \& 1)? true : false);
   /*como saber si un numero es una potencia de 2*/power_of_2 = ((v & (v-1)) == 0);
    /*contar trailing 0's de una mascara */
    builtin ctz(n);
    /*contar 1's de una mascara*/
    __builtin_popcount(n);
    /*quitar el elemento j de la m scara*/
    mask &= ^{\sim}(1 << j);
    /*revisar si el elemento j del arreglo esta en la m scara ( si es 0 el resultado es porque no
           est ) */
    int t = mask & (1 << j);
    /*Obtener el bit menos significativo*/
     t = mask & -mask
    /*encender todos los n primeros bits de la m scara*/
    mask = (1 << n) - 1;
     /*iterar sobre cada uno de los subsets de un subset y*/
    for (int x = y; x>0; x = (y & (x-1)))
```

6.2 Miller Rabin's primality test

```
// Randomized Primality Test (Miller-Rabin):
    Error rate: 2^(-TRIAL)
    Almost constant time. srand is needed
#include <stdlib.h>
#define EPS 1e-7
typedef long long LL;
LL ModularMultiplication(LL a, LL b, LL m) \,
        LL ret=0, c=a;
        while(b)
                if(b&1) ret=(ret+c)%m;
                b>>=1; c=(c+c)%m;
        return ret;
LL ModularExponentiation(LL a, LL n, LL m)
        LL ret=1, c=a;
        while(n)
                if(n&1) ret=ModularMultiplication(ret, c, m);
                n>>=1; c=ModularMultiplication(c, c, m);
        return ret:
bool Witness(LL a, LL n)
        LL u=n-1;
        while(!(u&1)){u>>=1; t++;}
        LL x0=ModularExponentiation(a, u, n), x1;
        for(int i=1; i<=t; i++)
                x1=ModularMultiplication(x0, x0, n);
                if(x1==1 && x0!=1 && x0!=n-1) return true;
                x0=x1:
        if(x0!=1) return true;
        return false;
LL Random(LL n)
```

```
LL ret=rand(); ret+=32768;
    ret+=rand(); ret+=32768;
    ret+=rand(); ret+=32768;
    ret+=rand();
return ret%n;
}
bool IsPrimeFast(LL n, int TRIAL)
{
    while(TRIAL--)
{
        LL a=Random(n-2)+1;
        if(Witness(a, n)) return false;
    }
    return true;
```

6.3 Pollard rho

```
#include<br/>bits/stdc++.h>
#include<time.h>
#define show(x) cout << #x << " = " << x << endl;
using namespace std:
typedef long long 11;
typedef pair<11, 11> ii;
typedef pair<double, ii> iii;
const int MAX = 200005:
const double EPS = 1e-5;
const int INF = INT MAX:
//modular multiplication for really big numbers
ll mul(ll a, ll b, ll mod) {
    11 ret = 0:
    for (a %= mod, b %= mod; b != 0;
        b >>= 1, a <<= 1, a = a >= mod ? a - mod : a) {
        if (b&1) {
             ret += a:
            if (ret >= mod) ret -= mod;
    return ret;
ll fpow(11 a, 11 b, 11 MOD) {
    11 ans = 1LL:
    while (b > 0) {
        if(b&1) ans = mul(ans, a, MOD);
        a = mul(a, a, MOD);
        b >>= 1T.T.:
    return ans;
const int rounds = 6;
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
bool miller_rabin(ll n) {
    if(n == 2 || n == 3) return true;
    if(n < 2 \mid \mid (n&1) == 0) return false;
    for(int i = 0; i < rounds; i++) {
   int a = rand() % (n-4) +2;
   if(fpow(a, n-1, n) != 1)</pre>
            return false:
    return true:
// if n is prime , check with miller rabin (n^(1/4)) and check return != n and != 1
11 pollard_rho(ll n, ll c) {
    11 \times = 2, y = 2, i = 1, k = 2, d;
    while (true) {
        x = (mul(x, x, n) + c);
        if (x >= n) x -= n;
        d = \underline{gcd}(x - y, n);
        if (d > 1) return d;
        if (++i == k) y = x, k <<= 1;
    return n:
//return factorization of a big number
```

```
void factorize(ll n, vector<ll> &f) {
    if(n == 1) return;
    if (miller_rabin(n)) {
        f.push_back(n);
        return;
    }
    ll d = n;
    for (int i = 2; d == n; i++)
        d = pollard_rho(n, i);
    factorize(d, f);
    factorize(n/d, f);
}
```

6.4 number theory general

```
// This is a collection of useful code for solving problems that
\ensuremath{//} involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std:
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b)
        while (b) { int t = a%b; a = b; b = t; }
        return a:
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                 x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)
    ret.push_back(mod(x + i*(n / g), n));</pre>
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
```

```
int mod_inverse(int a, int n) {
        int x, y;
int g = extended_euclid(a, n, x, y);
         if (g > 1) return -1;
         return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2). 
 // Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
         int s, t;
         int g = extended_euclid(m1, m2, s, t);
         if (r1%g != r2%g) return make pair(0, -1);
         return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
         PII ret = make_pair(r[0], m[0]);
         for (int i = 1; i < m.size(); i++) {
    ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);</pre>
                  if (ret.second == -1) break;
         return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
         if (!a && !b)
                  if (c) return false;
                  x = 0; y = 0;
                  return true;
         if (!a)
                  if (c % b) return false;
                  return true;
         if (!b)
                  if (c % a) return false;
x = c / a; y = 0;
                  return true;
         int q = qcd(a, b);
         if (c % g) return false;
         x = c / g * mod_inverse(a / g, b / g);
y = (c - a*x) / b;
         return true:
int main() {
         // expected: 2
         cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
         int x, y;
         int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
         // expected: 95 451
         VI sols = modular_linear_equation_solver(14, 30, 100);
         for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
         // expected: 8
         cout << mod_inverse(8, 9) << endl;</pre>
         // expected: 23 105
         PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
cout << ret.first << " " << ret.second << endl;
         ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;
         if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;</pre>
         return 0;
```

7 Miscellaneous

7.1 $c++ios\ tricks$

```
#include <iostream>
#include <iomanip>
using namespace std;
     // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal cout << hex << 100 << " " << 1000 << " " << 10000 << eed;
```

7.2 java IO template and iterative binary search

```
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.PrintWriter;
import java.util.StringTokenizer;
import java.io.IOException;
import java.io.BufferedReader;
import java.io.InputStreamReader;
import java.io.InputStream;
```

```
public class Main {
   public static void main(String[] args) {
        InputStream inputStream = System.in;
        OutputStream outputStream = System.out;
        InputReader in = new InputReader(inputStream);
        PrintWriter out = new PrintWriter(outputStream);
        TaskC solver = new TaskC();
        solver.solve(1, in, out);
        out.close();
    static class TaskC {
        private static final int ITERATIONS = 500;
        public void solve(int testNumber, InputReader in, PrintWriter out) {
            int n = in.nextInt();
            //Iterative binary search
            double 1 = 0.0, h = 1e17;
for (int i = 0; i < ITERATIONS; i++) {</pre>
                double mid = (1 + h) / 2.0;
                if (can(mid, a, b, p))
                    1 = mid;
                else
                    h = mid;
            double ans = 1;
    static class InputReader {
        public BufferedReader reader;
        public StringTokenizer tokenizer;
        public InputReader(InputStream stream) {
            reader = new BufferedReader(new InputStreamReader(stream), 32768);
            tokenizer = null;
        public String next() {
            while (tokenizer == null || !tokenizer.hasMoreTokens()) {
                    tokenizer = new StringTokenizer(reader.readLine());
                } catch (IOException e) {
                    throw new RuntimeException(e);
            return tokenizer.nextToken();
        public int nextInt() {
            return Integer.parseInt(next());
        public double nextDouble() {
            return Double.parseDouble(next());
```