

UNdead's notebook (2017)

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1 Graphs

1.1 Articulation points

```

/*Tarjan Algorithm to find articulation points
* single dfs  $O(|v| + |e|)$ 
* visited =[false]
* disc = [0]
* low = [0]
* parent = [-1]

```

```

* ap = [false] */
void articulation(vector<vector<int>> > G, int u, bool visited[],
    int disc[], int low[], int parent[], bool ap[]) {
    static int time = 0;
    int children = 0;
    visited[u] = true;
    disc[u] = low[u] = ++time;
    for(int i = 0; i < G[u].size(); i++){
        int v = G[u][i];
        if(!visited[v]){
            children++;
            parent[v] = u;
            articulation(G, v, visited, disc, low, ap);
            low[u] = min(low[u], low[v]);
            if(parent[u] == -1 && children > 1) ap[u] = true;
            if(parent[u] != -1 && low[v] >= disc[u]) ap[u] = true;
        }
        else if(v != parent[u])
            low[u] = min(low[u], disc[v]);
    }
}

```

1.2 Biconected graph

```

/* Tarjan Algorithm to find Biconnected graph
* single dfs  $O(|v| + |e|)$ 
* visited =[false]
* disc = [0]
* low = [0]
* parent = [-1] */
bool isBiconnected(vector<vector<int>> > G, int u, bool visited[],
    int disc[], int low[], int parent[]){
    static int time = 0;
    int children = 0;
    visited[u] = true;
    disc[u] = low[u] = ++time;
    for(int i = 0; i < G[u].size(); i++){
        int v = G[u][i];
        if(!visited[v]){
            children++;
            parent[v] = u;
            if (isBiconnected(G, v, visited, disc, low, parent))
                return true;
            low[u] = min(low[u], low[v]);
            if(parent[u] == -1 && children > 1) return true;
            if(parent[u] != -1 && low[v] >= disc[u]) return true;
        }
        else if(v != parent[v]) low[u] = min(low[u], disc[v]);
    }
    return false;
}

```

}

1.3 Bridges

```

/* Tarjan Algorithm to find bridges
 * single dfs O(|v| + |e|)
 * visited =[false]
 * disc = [0]
 * low = [0]
 * parent = [-1] */
void bridges(vector<vector<int> > G, int u, bool visited[], int
disc[], int low[], int parent[], priority_queue<pair<int,
int> > *bridge) {

    static int time = 0;

    int children = 0;
    visited[u] = true;
    disc[u] = low[u] = ++time;

    for(int i = 0; i < G[u].size(); i++){
        int v = G[u][i];

        if(!visited[v]){
            children++;
            parent[v] = u;
            bridges(G, v, visited, disc, low, bridge);

            low[u] = min(low[u], low[v]);

            if(low[v] > disc[u]) bridge->push({u,v});
        }
        else if(v != parent[u])
            low[u] = min(low[u], disc[v]);
    }
}

```

1.4 Tarjan SCC

```

/* Tarjan Algorithm to find connected components
 * single dfs O(|v| + |e|)
 * visited =[false]
 * disc = [0]
 * low = [0]
 * parent = [-1] */
void dfsSCC(vector<vector<int> > G, int u, int disc[], int low[],
stack<int> *st, bool stackMember[]){
    static int time = 0;

    disc[u] = low[u] = ++time;

    st->push(u);
    stackmember[u] = true;

    for(int i = 0; i < G[u].size(); i++){
        int v = G[u][i];

        if(disc[v] == -1){
            dfsSCC(G, v, disc, low, st, stackmember);

```

```

        low[u] = min(low[u], low[v]);
    }
    else if(stackmember[v] == true) low[u] = min(low[u], disc
[v]);
}

int w = 0;
if(low[u] == disc[u]){
    while(st->top() != u){
        w = st->top();
        cout<<w<<" ";
        stackmember[w] = false;
        st->pop();
    }

    w = st->top();
    cout<<w<<"\n";
    stackmember[w] = false;
    st->pop();
}

}

void scc(G){
    int *disc = new int[V];
    int *low = new int[V];
    bool *stackMember = new bool[V];
    stack<int> *st = new stack<int>();

    memset(disc, -1, sizeof(disc));
    memset(low, 0, sizeof(low));
    memset(stackMember, false, sizeof(stackMember));

    for(int i = 0; i < G.size(); i++)
        if(disc[i] == -1) dfsScc(G, i, disc, low, st, stackMember
);
}

```

1.5 Fast Dijkstra's algorithm

```

// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
//
// Running time: O(|E| log |V|)

#include <queue>
#include <cstdio>

using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;

int main() {
    int N, s, t;
    scanf("%d%d%d", &N, &s, &t);
    vector<vector<PII> > edges(N);
    for (int i = 0; i < N; i++) {
        int M;
        scanf("%d", &M);
        for (int j = 0; j < M; j++) {
            int vertex, dist;
            scanf("%d%d", &vertex, &dist);
            edges[i].push_back(make_pair(dist, vertex)
); // note order of arguments here
        }
    }
}

```

```

// use priority queue in which top element has the "
// smallest" priority
priority_queue<PII, vector<PII>, greater<PII> > Q;
vector<int> dist(N, INF), dad(N, -1);
Q.push(make_pair(0, s));
dist[s] = 0;
while (!Q.empty()) {
    PII p = Q.top();
    Q.pop();
    int here = p.second;
    if (here == t) break;
    if (dist[here] != p.first) continue;

    for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
        if (dist[here] + it->first < dist[it->second]) {
            dist[it->second] = dist[here] + it->first;
            dad[it->second] = here;
            Q.push(make_pair(dist[it->second], it->second));
        }
    }

    printf("%d\n", dist[t]);
    if (dist[t] < INF)
        for (int i = t; i != -1; i = dad[i])
            printf("%d%c", i, (i == s ? '\n' : ' '));

    return 0;
}

/*
Sample input:
5 0 4
2 1 2 3 1
2 2 4 4 5
3 1 4 3 3 4 1
2 0 1 2 3
2 1 5 2 1

Expected:
5
4 2 3 0
*/

```

1.6 Bellman Ford

```

// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function
// returns
// false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
//
// Running time: O(|V|^3)
//
// INPUT:  start, w[i][j] = cost of edge from i to j
// OUTPUT: dist[i] = min weight path from start to i
//         prev[i] = previous node on the best path from the
//         start node

#include <iostream>
#include <queue>

```

```

#include <cmath>
#include <vector>

using namespace std;

typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

typedef vector<int> VI;
typedef vector<VI> VVI;

bool BellmanFord (const VVT &w, VT &dist, VI &prev, int start) {
    int n = w.size();
    prev = VI(n, -1);
    dist = VT(n, 1000000000);
    dist[start] = 0;

    for (int k = 0; k < n; k++) {
        for (int i = 0; i < n; i++) {
            for (int j = 0; j < n; j++) {
                if (dist[j] > dist[i] + w[i][j]) {
                    if (k == n-1) return false;
                    dist[j] = dist[i] + w[i][j];
                    prev[j] = i;
                }
            }
        }
    }

    return true;
}

```

1.7 Eulerian Path

```

struct Edge;
typedef list<Edge>::iterator iter;

struct Edge
{
    int next_vertex;
    iter reverse_edge;

    Edge(int next_vertex)
        :next_vertex(next_vertex)
        { }
};

const int max_vertices = ;
int num_vertices;
list<Edge> adj[max_vertices]; // adjacency list

vector<int> path;

void find_path(int v)
{
    while(adj[v].size() > 0)
    {
        int vn = adj[v].front().next_vertex;
        adj[vn].erase(adj[v].front().reverse_edge);
        adj[v].pop_front();
        find_path(vn);
    }
    path.push_back(v);
}

void add_edge(int a, int b)
{

```

```

adj[a].push_front(Edge(b));
iter ita = adj[a].begin();
adj[b].push_front(Edge(a));
iter itb = adj[b].begin();
ita->reverse_edge = itb;
itb->reverse_edge = ita;
}

```

1.8 Topological sort

```

char c[TAM];
int l[TAM];
int r[TAM];
int in[TAM];

//can be priority queue
queue<int> Q;

void reset(){
    memset(l, 0, sizeof l);
    memset(r, 0, sizeof r);
    memset(in, 0, sizeof in);
    memset(balls, 0, sizeof balls);
    c[0] = 'L';
}

void topo(vector<vector<int>> G, int u){
    while(!Q.empty()){
        u = Q.front(); Q.pop();
        update(u);
        for(int i = 0; i < G[u].size(); i++){
            int v = G[u][i];
            in[v]--;
            if(in[v] == 0) Q.push(v);
        }
    }
}

int main(){
    ll n;
    int m;
    while(cin>>n>>m){
        reset();
        vector<vector<int>> G(m + 1);
        for(int i = 1; i <= m; i++){
            int u, v;
            cin>>c[i]>>u>>v;
            G[i].push_back(u);
            G[i].push_back(v);
            in[u]++; in[v]++;
            l[i] = u; r[i] = v;
        }
        for(int i = 1; i <= m; i++){
            if(in[i] == 0) Q.push(i);
        }
        topo(G, 1);
    }
}

```

1.9 Kruskal's minimum spanning tree

Uses Kruskal's Algorithm to calculate the weight of the minimum spanning forest (union of minimum spanning trees of each connected component) of a possibly disjoint graph, given in the form of a matrix of edge weights (-1 if no edge exists). Returns the weight of the minimum spanning forest (also calculates the actual edges - stored in T). Note: uses a disjoint-set data structure with amortized (effectively) constant time per union/find. Runs in $O(E \cdot \log(E))$ time.

```

#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>

using namespace std;

typedef int T;

struct edge
{
    int u, v;
    T d;
};

struct edgeCmp
{
    int operator()(const edge& a, const edge& b) { return a.d > b.d; }
};

int find(vector<int>& C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }

T Kruskal(vector<vector<T>>& w)
{
    int n = w.size();
    T weight = 0;

    vector<int> C(n), R(n);
    for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }

    vector<edge> T;
    priority_queue<edge, vector<edge>, edgeCmp> E;

    for(int i=0; i<n; i++)
        for(int j=i+1; j<n; j++)
            if(w[i][j] >= 0)
            {
                edge e;
                e.u = i; e.v = j; e.d = w[i][j];
                E.push(e);
            }

    while(T.size() < n-1 && !E.empty())
    {
        edge cur = E.top(); E.pop();

        int uc = find(C, cur.u), vc = find(C, cur.v);
        if(uc != vc)
        {
            T.push_back(cur); weight += cur.d;

            if(R[uc] > R[vc]) C[vc] = uc;
            else if(R[vc] > R[uc]) C[uc] = vc;
            else { C[vc] = uc; R[uc]++; }
        }
    }
}

```

```

/*

```

```

    }
}
return weight;
}

int main()
{
    int wa[6][6] = {
        { 0, -1, 2, -1, 7, -1 },
        { -1, 0, -1, 2, -1, -1 },
        { 2, -1, 0, -1, 8, 6 },
        { -1, 2, -1, 0, -1, -1 },
        { 7, -1, 8, -1, 0, 4 },
        { -1, -1, 6, -1, 4, 0 } };

    vector<vector<int>> w(6, vector<int>(6));

    for(int i=0; i<6; i++)
        for(int j=0; j<6; j++)
            w[i][j] = wa[i][j];

    cout << Kruskal(w) << endl;
    cin >> wa[0][0];
}

```

1.10 Lowest Common ancestor

```

const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;

vector<int> children[max_nodes]; // children[i] contains
    the children of node i
int A[max_nodes][log_max_nodes+1]; // A[i][j] is the 2^j-th
    ancestor of node i, or -1 if that ancestor does not exist
int L[max_nodes]; // L[i] is the distance
    between node i and the root

// floor of the binary logarithm of n
int lb(unsigned int n)
{
    if(n==0)
        return -1;
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16; }
    if (n >= 1<< 8) { n >>= 8; p += 8; }
    if (n >= 1<< 4) { n >>= 4; p += 4; }
    if (n >= 1<< 2) { n >>= 2; p += 2; }
    if (n >= 1<< 1) { p += 1; }
    return p;
}

void DFS(int i, int l)
{
    L[i] = l;
    for(int j = 0; j < children[i].size(); j++)
        DFS(children[i][j], l+1);
}

int LCA(int p, int q)
{
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);

    // "binary search" for the ancestor of node p situated on the
    same level as q
    for(int i = log_num_nodes; i >= 0; i--)

```

```

        if(L[p] - (1<<i) >= L[q])
            p = A[p][i];

    if(p == q)
        return p;

    // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if(A[p][i] != -1 && A[q][i] != A[p][i])
        {
            p = A[p][i];
            q = A[q][i];
        }

    return A[p][0];
}

int main(int argc, char* argv[])
{
    // read num_nodes, the total number of nodes
    log_num_nodes=lb(num_nodes);

    for(int i = 0; i < num_nodes; i++)
    {
        int p;
        // read p, the parent of node i or -1 if node i is the
        root

        A[i][0] = p;
        if(p != -1)
            children[p].push_back(i);
        else
            root = i;
    }

    // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
        for(int i = 0; i < num_nodes; i++)
            if(A[i][j-1] != -1)
                A[i][j] = A[A[i][j-1]][j-1];
            else
                A[i][j] = -1;

    // precompute L
    DFS(root, 0);

    return 0;
}

```

2 Geometry

2.1 Geometry

```

// C++ routines for computational geometry.

#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>

using namespace std;

double INF = 1e100;
double EPS = 1e-12;

struct PT {

```

```

double x, y;
PT() {}
PT(double x, double y) : x(x), y(y) {}
PT(const PT &p) : x(p.x), y(p.y) {}
PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
PT operator * (double c) const { return PT(x*c, y*c); }
PT operator / (double c) const { return PT(x/c, y/c); }
};

double dot(PT p, PT q) { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
    os << "(" << p.x << ", " << p.y << ")";
}

// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
    return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
}

// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
    return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
}

// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
    double r = dot(b-a,b-a);
    if (fabs(r) < EPS) return a;
    r = dot(c-a, b-a)/r;
    if (r < 0) return a;
    if (r > 1) return b;
    return a + (b-a)*r;
}

// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
    return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
}

// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
    double a, double b, double c, double d)
{
    return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
}

// determine if lines from a to b and c to d are parallel or
// collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
    return fabs(cross(b-a, c-d)) < EPS;
}

bool LinesCollinear(PT a, PT b, PT c, PT d) {
    return LinesParallel(a, b, c, d)
        && fabs(cross(a-b, a-c)) < EPS
        && fabs(cross(c-d, c-a)) < EPS;
}

// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
    if (LinesCollinear(a, b, c, d)) {
        if (dist2(a, c) < EPS || dist2(a, d) < EPS ||

```

```

        dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
        if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) >
            0)
            return false;
        return true;
    }
    if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
    if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
    return true;
}

// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
    b=b-a; d=d-c; c=c-a;
    assert(dot(b, b) > EPS && dot(d, d) > EPS);
    return a + b*cross(c, d)/cross(b, d);
}

// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
    b=(a+b)/2;
    c=(a+c)/2;
    return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+
        RotateCW90(a-c));
}

// determine if point is in a possibly non-convex polygon (by
// William
// Randolph Franklin); returns 1 for strictly interior points, 0
// for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test
// using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing
// exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
    bool c = 0;
    for (int i = 0; i < p.size(); i++) {
        int j = (i+1)%p.size();
        if ((p[i].y <= q.y && q.y < p[j].y ||
            p[j].y <= q.y && q.y < p[i].y) &&
            q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y
                - p[i].y))
            c = !c;
    }
    return c;
}

// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
    for (int i = 0; i < p.size(); i++)
        if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q)
            < EPS)
            return true;
    return false;
}

// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
    vector<PT> ret;
    b = b-a;
    a = a-c;

```

```

double A = dot(b, b);
double B = dot(a, b);
double C = dot(a, a) - r*r;
double D = B*B - A*C;
if (D < -EPS) return ret;
ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
return ret;
}

// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(Pt a, Pt b, double r, double R
) {
    vector<PT> ret;
    double d = sqrt(dist2(a, b));
    if (d > r+R || d+min(r, R) < max(r, R)) return ret;
    double x = (d*d-R*R+r*r)/(2*d);
    double y = sqrt(r*r-x*x);
    PT v = (b-a)/d;
    ret.push_back(a+v*x + RotateCCW90(v)*y);
    if (y > 0)
        ret.push_back(a+v*x - RotateCCW90(v)*y);
    return ret;
}

// This code computes the area or centroid of a (possibly
// nonconvex)
// polygon, assuming that the coordinates are listed in a
// clockwise or
// counterclockwise fashion. Note that the centroid is often
// known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
    double area = 0;
    for(int i = 0; i < p.size(); i++) {
        int j = (i+1) % p.size();
        area += p[i].x*p[j].y - p[j].x*p[i].y;
    }
    return area / 2.0;
}

double ComputeArea(const vector<PT> &p) {
    return fabs(ComputeSignedArea(p));
}

PT ComputeCentroid(const vector<PT> &p) {
    PT c(0,0);
    double scale = 6.0 * ComputeSignedArea(p);
    for (int i = 0; i < p.size(); i++) {
        int j = (i+1) % p.size();
        c = c + (p[i].x*p[j].y - p[j].x*p[i].y);
    }
    return c / scale;
}

// tests whether or not a given polygon (in CW or CCW order) is
// simple
bool IsSimple(const vector<PT> &p) {
    for (int i = 0; i < p.size(); i++) {
        for (int k = i+1; k < p.size(); k++) {
            int j = (i+1) % p.size();
            int l = (k+1) % p.size();
            if (i == l || j == k) continue;
            if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
                return false;
        }
    }
}

```

```

    }
    return true;
}

int main() {
    // expected: (-5,2)
    cerr << RotateCCW90(Pt(2,5)) << endl;

    // expected: (5,-2)
    cerr << RotateCW90(Pt(2,5)) << endl;

    // expected: (-5,2)
    cerr << RotateCCW(Pt(2,5),M_PI/2) << endl;

    // expected: (5,2)
    cerr << ProjectPointLine(Pt(-5,-2), Pt(10,4), Pt(3,7)) << endl;

    // expected: (5,2) (7.5,3) (2.5,1)
    cerr << ProjectPointSegment(Pt(-5,-2), Pt(10,4), Pt(3,7)) << " "
        << ProjectPointSegment(Pt(7.5,3), Pt(10,4), Pt(3,7)) << " "
        << ProjectPointSegment(Pt(-5,-2), Pt(2.5,1), Pt(3,7)) <<
            endl;

    // expected: 6.78903
    cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;

    // expected: 1 0 1
    cerr << LinesParallel(Pt(1,1), Pt(3,5), Pt(2,1), Pt(4,5)) << " "
        << LinesParallel(Pt(1,1), Pt(3,5), Pt(2,0), Pt(4,5)) << " "
        << LinesParallel(Pt(1,1), Pt(3,5), Pt(5,9), Pt(7,13)) <<
            endl;

    // expected: 0 0 1
    cerr << LinesCollinear(Pt(1,1), Pt(3,5), Pt(2,1), Pt(4,5)) << "
"
        << LinesCollinear(Pt(1,1), Pt(3,5), Pt(2,0), Pt(4,5)) << "
"
        << LinesCollinear(Pt(1,1), Pt(3,5), Pt(5,9), Pt(7,13)) <<
            endl;

    // expected: 1 1 1 0
    cerr << SegmentsIntersect(Pt(0,0), Pt(2,4), Pt(3,1), Pt(-1,3))
        << " "
        << SegmentsIntersect(Pt(0,0), Pt(2,4), Pt(4,3), Pt(0,5)) <<
            " "
        << SegmentsIntersect(Pt(0,0), Pt(2,4), Pt(2,-1), Pt(-2,1))
        << " "
        << SegmentsIntersect(Pt(0,0), Pt(2,4), Pt(5,5), Pt(1,7)) <<
            endl;

    // expected: (1,2)
    cerr << ComputeLineIntersection(Pt(0,0), Pt(2,4), Pt(3,1), Pt
        (-1,3)) << endl;

    // expected: (1,1)
    cerr << ComputeCircleCenter(Pt(-3,4), Pt(6,1), Pt(4,5)) << endl;

    vector<PT> v;
    v.push_back(Pt(0,0));
    v.push_back(Pt(5,0));
    v.push_back(Pt(5,5));
    v.push_back(Pt(0,5));

    // expected: 1 1 1 0 0
    cerr << PointInPolygon(v, Pt(2,2)) << " "
        << PointInPolygon(v, Pt(2,0)) << " "
        << PointInPolygon(v, Pt(0,2)) << " "
        << PointInPolygon(v, Pt(5,2)) << " "
        << PointInPolygon(v, Pt(2,5)) << endl;
}

```

```
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "
      << PointOnPolygon(v, PT(2,0)) << " "
      << PointOnPolygon(v, PT(0,2)) << " "
      << PointOnPolygon(v, PT(5,2)) << " "
      << PointOnPolygon(v, PT(2,5)) << endl;

// expected: (1,6)
//           (5,4) (4,5)
//           blank line
//           (4,5) (5,4)
//           blank line
//           (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1),
5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <<
endl;
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <<
endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <<
endl;
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <<
endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)
/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <<
endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)
/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <<
endl;

// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;

return 0;
}
```

2.2 Geometry (Java)

```
// In this example, we read an input file containing three lines,
// each
// containing an even number of doubles, separated by commas. The
// first two
// lines represent the coordinates of two polygons, given in
// counterclockwise
// (or clockwise) order, which we will call "A" and "B". The last
// line
// contains a list of points, p[1], p[2], ...
//
// Our goal is to determine:
// (1) whether B - A is a single closed shape (as opposed to
// multiple shapes)
// (2) the area of B - A
// (3) whether each p[i] is in the interior of B - A
//
// INPUT:
```

```
// 0 0 10 0 0 10
// 0 0 10 10 10 0
// 8 6
// 5 1
//
// OUTPUT:
// The area is singular.
// The area is 25.0
// Point belongs to the area.
// Point does not belong to the area.

import java.util.*;
import java.awt.geom.*;
import java.io.*;

public class JavaGeometry {

    // make an array of doubles from a string
    static double[] readPoints(String s) {
        String[] arr = s.trim().split("\\s+");
        double[] ret = new double[arr.length];
        for (int i = 0; i < arr.length; i++) ret[i] = Double.
            parseDouble(arr[i]);
        return ret;
    }

    // make an Area object from the coordinates of a polygon
    static Area makeArea(double[] pts) {
        Path2D.Double p = new Path2D.Double();
        p.moveTo(pts[0], pts[1]);
        for (int i = 2; i < pts.length; i += 2) p.lineTo(pts[i],
            pts[i+1]);
        p.closePath();
        return new Area(p);
    }

    // compute area of polygon
    static double computePolygonArea(ArrayList<Point2D.Double>
        points) {
        Point2D.Double[] pts = points.toArray(new Point2D.Double[
            points.size()]);
        double area = 0;
        for (int i = 0; i < pts.length; i++) {
            int j = (i+1) % pts.length;
            area += pts[i].x * pts[j].y - pts[j].x * pts[i].y;
        }
        return Math.abs(area)/2;
    }

    // compute the area of an Area object containing several
    // disjoint polygons
    static double computeArea(Area area) {
        double totArea = 0;
        PathIterator iter = area.getPathIterator(null);
        ArrayList<Point2D.Double> points = new ArrayList<Point2D.
            Double>();

        while (!iter.isDone()) {
            double[] buffer = new double[6];
            switch (iter.currentSegment(buffer)) {
                case PathIterator.SEG_MOVETO:
                case PathIterator.SEG_LINETO:
                    points.add(new Point2D.Double(buffer[0], buffer
                        [1]));
                    break;
                case PathIterator.SEG_CLOSE:
                    totArea += computePolygonArea(points);
                    points.clear();
            }
        }
    }
}
```



```

        break;
    }
    iter.next();
}
return totArea;
}

// notice that the main() throws an Exception -- necessary to
// avoid wrapping the Scanner object for file reading in a
// try { ... } catch block.
public static void main(String args[]) throws Exception {
    Scanner scanner = new Scanner(new File("input.txt"));
    // also,
    // Scanner scanner = new Scanner (System.in);

    double[] pointsA = readPoints(scanner.nextLine());
    double[] pointsB = readPoints(scanner.nextLine());
    Area areaA = makeArea(pointsA);
    Area areaB = makeArea(pointsB);
    areaB.subtract(areaA);
    // also,
    // areaB.exclusiveOr (areaA);
    // areaB.add (areaA);
    // areaB.intersect (areaA);

    // (1) determine whether B - A is a single closed shape (
    //      as
    //      opposed to multiple shapes)
    boolean isSingle = areaB.isSingular();
    // also,
    // areaB.isEmpty();

    if (isSingle)
        System.out.println("The area is singular.");
    else
        System.out.println("The area is not singular.");

    // (2) compute the area of B - A
    System.out.println("The area is " + computeArea(areaB) + "
    .");

    // (3) determine whether each p[i] is in the interior of B
    //      - A
    while (scanner.hasNextDouble()) {
        double x = scanner.nextDouble();
        assert(scanner.hasNextDouble());
        double y = scanner.nextDouble();

        if (areaB.contains(x,y)) {
            System.out.println ("Point belongs to the area.");
        } else {
            System.out.println ("Point does not belong to the
            area.");
        }
    }

    // Finally, some useful things we didn't use in this
    // example:
    //
    // Ellipse2D.Double ellipse = new Ellipse2D.Double (
    // double x, double y,
    //
    // double w, double h);
    //
    // creates an ellipse inscribed in box with bottom-
    // left corner (x,y)
    // and upper-right corner (x+y,w+h)
    //
    //

```

```

// Rectangle2D.Double rect = new Rectangle2D.Double (
// double x, double y,
//
// double w, double h);
//
// creates a box with bottom-left corner (x,y) and
// upper-right
// corner (x+y,w+h)
//
// Each of these can be embedded in an Area object (e.g.,
// new Area (rect)).
}
}

```

2.3 Convex Hull Monotone Chain

```

#include<iostream>
#include<algorithm>
#include<complex>
#include<cstdio>
#include<iomanip>
#include<vector>
#define x real()
#define y imag()
#define dot(A,B) real(conj((A))* (B))
#define cross(A,B) imag(conj((A))* (B))
#define PI 3.1415926
#define EPS 1e-9

using namespace std;
typedef double lf;
typedef complex<lf> pt;

istream& operator >> ( istream& in, pt& p ) {
    lf a,n; in >> a >> n;
    p = pt(a,n); return in;
}

bool cmp(pt &p, pt &q) {
    if(p.x != q.x) return p.x < q.x;
    return p.y < q.y;
}

bool is_zero( lf x ){
    return -EPS <= x && x <= EPS;
}

inline bool same ( lf a, lf b ) {
    return a+EPS > b && b+EPS > a;
}

int ccw(pt& p1, pt& p2, pt& p3) {
    lf ans = (cross(p1 - p3, p2 - p3));
    if(-EPS <= ans && ans <= EPS)
        return 0;
    else if(ans < -EPS)
        return -1;
    else
        return 1;
}

lf dist ( pt A, pt B ) { return abs(A-B); }

vector<pt> convex_hull(vector<pt> P){
    int n = P.size(); int k = 0;
    vector<pt> H(2*n);

```

```

    sort(P.begin(), P.end(), cmp);
    for (int i = 0; i < n; i++) {
        while (k >= 2 && ccw(H[k-2], H[k-1], P[i]) == 1) k--;
        H[k++] = P[i];
    }

    for (int i = n-2, t = k+1; i >= 0; i--) {
        while (k >= t && ccw(H[k-2], H[k-1], P[i]) == 1) k--;
        H[k++] = P[i];
    }

    if (n > 1)
        k--;
    H.resize(k);
    return H;
}

int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
#ifdef LOCAL
    //freopen("in.txt", "r", stdin);
    //freopen("out.txt", "w", stdout);
#endif // LOCAL
    int n;
    cin >> n;
    vector<pt> p(n);
    for (int i = 0; i < n; i++)
        cin >> p[i];
    vector<pt> hull = convex_hull(p);
    for (auto &pt : hull)
        cout << pt << endl;
}

```

2.4 Delaunay triangulation

```

// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
//
// Running time: O(n^4)
//
// INPUT:      x[] = x-coordinates
//             y[] = y-coordinates
//
// OUTPUT:     triples = a vector containing m triples of indices
//                   corresponding to triangle vertices

#include<vector>
using namespace std;
typedef double T;

struct triple {
    int i, j, k;
    triple() {}
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
};

vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
    int n = x.size();
    vector<T> z(n);
    vector<triple> ret;
}

```

```

    for (int i = 0; i < n; i++)
        z[i] = x[i] * x[i] + y[i] * y[i];

    for (int i = 0; i < n-2; i++) {
        for (int j = i+1; j < n; j++) {
            for (int k = i+1; k < n; k++) {
                if (j == k) continue;
                double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                bool flag = zn < 0;
                for (int m = 0; flag && m < n; m++)
                    flag = flag && ((x[m]-x[i])*xn + (y[m]-y[i])*yn + (z[m]-z[i])*zn <= 0);
                if (flag) ret.push_back(triple(i, j, k));
            }
        }
    }

    return ret;
}

int main()
{
    T xs[]={0, 0, 1, 0.9};
    T ys[]={0, 1, 0, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);

    //expected: 0 1 3
    //           0 3 2

    int i;
    for (i = 0; i < tri.size(); i++)
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
    return 0;
}

```

2.5 Delaunay triangulation (java)

```

// Slow but simple Delaunay triangulation. (from O'Rourke,
// Computational Geometry in C)
//
// Running time: O(n^4)
//
// INPUT:      x[] = x-coordinates
//             y[] = y-coordinates
//
// OUTPUT:     ret[][] = an mx3 matrix containing m triples of
//                   indices
//                   corresponding to triangle vertices

import java.util.*;

public class Delaunay {
    int[][] triangulate(double[] x, double[] y) {
        int n = x.length;
        double z[] = new double[n];
        ArrayList<int[]> ret = new ArrayList<int[]>();

        for (int i = 0; i < n; i++)
            z[i] = x[i] * x[i] + y[i] * y[i];
    }
}

```

```

    for (int i = 0; i < n-2; i++) {
        for (int j = i+1; j < n; j++) {
            for (int k = i+1; k < n; k++) {
                if (j == k) continue;
                double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                boolean flag = zn < 0;
                for (int m = 0; flag && m < n; m++)
                    flag = flag && ((x[m]-x[i])*xn + (y[m]-y[i])*yn + (z[m]-z[i])*zn <= 0);
                if (flag) ret.add(new int[]{i, j, k});
            }
        }
    }
    return ret.toArray(new int[0][0]);
}
}

```

3 Data Structures

3.1 Mo's algorithm

```

#include<bits/stdc++.h>
#define TAM 30000 + 7
#define QTAM 200000 + 7
#define MTAM 1000000 + 7
#define whatis(x) cerr<<#x<<" is "<<x<<endl

using namespace std;

int a[TAM], r[QTAM], cnt[MTAM];
int ans, BLOCK, currL, currR;

struct node{
    int L, R, idx;
}q[QTAM];

bool comp(node a, node b){
    if(a.L/BLOCK < b.L/BLOCK) return true;
    if(a.L/BLOCK > b.L/BLOCK) return false;
    return a.R < b.R;
}

void remove(int i){
    cnt[a[i]]--;
    if(cnt[a[i]] == 0) ans--;
}

void add(int i){
    cnt[a[i]]++;
    if(cnt[a[i]]==1) ans++;
}

int query(node i){
    while(currL < i.L){
        remove(currL);
        currL++;
    }
    while(currL > i.L){

```

```

        currL--;
        add(currL);
    }
    while(currR < i.R){
        currR++;
        add(currR);
    }
    while(currR > i.R){
        remove(currR);
        currR--;
    }
    return ans;
}

int main(){
    ios_base::sync_with_stdio(0);
    cin.tie(0);

    #ifdef LOCAL
    freopen("in", "r", stdin);
    #endif
    int n, que;
    cin>>n;
    BLOCK = sqrt(n);
    for(int i = 1; i <= n; i++) cin>>a[i];

    cin>>que;
    for(int i = 1; i <= que; i++){
        cin>>q[i].L>>q[i].R;
        q[i].idx = i;
    }
    sort(q + 1, q + que + 1, comp);

    for(int i = 1; i <= que; i++)
        r[q[i].idx] = query(q[i]);

    for(int i = 1; i <= que; i++)
        cout<<r[i]<<"\n";
}

```

3.2 Segment Trees with lazy propagation

```

//querys and build takes O(log n)
//example with segment sum
#include<bits/stdc++.h>

using namespace std;

long long *p;
//long long *lazy;

struct SegmentTree{
    SegmentTree *L, *R;
    long long sum = 0;
    long long lazy = 0;
    int l, r;

    long long query2(int a, int b){
        if(a == l && b == r) return sum;
        if(b <= L->r) return L->query(a,b);
        if(a >= R->l) return R->query(a,b);
        return (L->query2(a,L->r) + R->query2(R->l, b));
    }

    void update(int a, int val){
        if(l == r){
            sum += val;

```

```

        return;
    }
    int mid = (l + r)/2;
    if(l <= a && a<= mid)
        L->update(a, val);
    else
        R->update(a, val);
    sum = L->sum + R->sum;
}

void updateRange2(int a, int b, long long val) {
    if(b < l or a > r)
        return;
    if(l == r) {
        sum += val;
        return;
    }
    L->updateRange2(a, b, val);
    R->updateRange2(a, b, val);
    sum = L->sum + R->sum;
}

void updateRange(int a, int b, long long val) {
    if(lazy != 0) {
        sum += (r-l+1)*lazy;
        //sum += lazy;
        if(l != r) {
            R->lazy = lazy + R->lazy;
            L->lazy = lazy + L->lazy;
        }
        lazy = 0;
    }
    if(b < l or a > r)
        return;
    if(l >= a && r <= b) {
        sum += (r-l+1)*val;
        //sum += val;
        if(l != r) {
            R->lazy = val + R->lazy;
            L->lazy = val + L->lazy;
        }
        return;
    }
    L->updateRange(a, b, val);
    R->updateRange(a, b, val);
    sum = L->sum + R->sum;
}

long long query(int a, int b) {
    if(b < l or a > r)
        return 0;
    if(lazy != 0) {
        sum += (r-l+1)*lazy;
        //sum += lazy;
        if(l != r) {
            R->lazy = lazy + R->lazy;
            L->lazy = lazy + L->lazy;
        }
        lazy = 0;
    }
    if(a == l && b == r) return sum;
    if(b <= L->r) return L->query(a, b);
    if(a >= R->l) return R->query(a, b);
    return (L->query(a, L->r) + R->query(R->l, b));
}

SegmentTree(int a, int b): l(a), r(b) {
    if(a == b) {
        sum = p[a];
    }
}

```

```

        L = R = nullptr;
    }
    else {
        L = new SegmentTree ( a, (a+b)/2 );
        R = new SegmentTree ( (a+b)/2 + 1, b );
        sum = L->sum + R->sum;
    }
}

};

int main() {
    cin.tie(0);
    ios_base::sync_with_stdio(0);
#ifdef LOCAL
    freopen("input.txt", "r", stdin);
#endif // LOCAL
    long long T;
    cin >> T;
    while(T--) {
        long long n, c;
        cin >> n >> c;
        long long l[n];
        memset(l, 0, sizeof(l));
        p = l;
        SegmentTree *stree = new SegmentTree(0, n-1);
        while(c--) {
            long long aux, p, q;
            cin >> aux >> p >> q;
            if(aux == 0) {
                long long val;
                cin >> val;
                stree->updateRange(p-1, q-1, val);
            }
            else
                cout << stree->query(p-1, q-1) << endl;
        }
    }
}

```

3.3 Segment Trees with lazy propagation (Java)

```

public class SegmentTreeRangeUpdate {
    public long[] leaf;
    public long[] update;
    public int origSize;
    public SegmentTreeRangeUpdate(int[] list) {
        origSize = list.length;
        leaf = new long[4*list.length];
        update = new long[4*list.length];
        build(1, 0, list.length-1, list);
    }
    public void build(int curr, int begin, int end, int[] list) {
        if(begin == end) {
            leaf[curr] = list[begin];
        }
        else {
            int mid = (begin+end)/2;
            build(2 * curr, begin, mid, list);
            build(2 * curr + 1, mid+1, end, list);
            leaf[curr] = leaf[2*curr] + leaf[2*curr
                +1];
        }
    }
}

```

```

public void update(int begin, int end, int val) {
    update(1,0,origSize-1,begin,end,val);
}
public void update(int curr, int tBegin, int tEnd, int
    begin, int end, int val) {
    if(tBegin >= begin && tEnd <= end)
        update[curr] += val;
    else
    {
        leaf[curr] += (Math.min(end,tEnd)-Math.max
            (begin,tBegin)+1) * val;
        int mid = (tBegin+tEnd)/2;
        if(mid >= begin && tBegin <= end)
            update(2*curr, tBegin, mid, begin,
                end, val);
        if(tEnd >= begin && mid+1 <= end)
            update(2*curr+1, mid+1, tEnd,
                begin, end, val);
    }
}
public long query(int begin, int end) {
    return query(1,0,origSize-1,begin,end);
}
public long query(int curr, int tBegin, int tEnd, int
    begin, int end) {
    if(tBegin >= begin && tEnd <= end) {
        if(update[curr] != 0) {
            leaf[curr] += (tEnd-tBegin+1) *
                update[curr];
            if(2*curr < update.length){
                update[2*curr] += update[
                    curr];
                update[2*curr+1] += update
                    [curr];
            }
            update[curr] = 0;
        }
        return leaf[curr];
    }
    else
    {
        leaf[curr] += (tEnd-tBegin+1) * update[
            curr];
        if(2*curr < update.length){
            update[2*curr] += update[curr];
            update[2*curr+1] += update[curr];
        }
        update[curr] = 0;
        int mid = (tBegin+tEnd)/2;
        long ret = 0;
        if(mid >= begin && tBegin <= end)
            ret += query(2*curr, tBegin, mid,
                begin, end);
        if(tEnd >= begin && mid+1 <= end)
            ret += query(2*curr+1, mid+1, tEnd,
                , begin, end);
        return ret;
    }
}
}
}

```

3.4 Fenwick Tree

```

#include <iostream>
using namespace std;
#define LOGSZ 17

```

```

int tree[(1<<LOGSZ)+1];
int N = (1<<LOGSZ);

// add v to value at x
void set(int x, int v) {
    while(x <= N) {
        tree[x] += v;
        x += (x & -x);
    }
}

// get cumulative sum up to and including x
int get(int x) {
    int res = 0;
    while(x) {
        res += tree[x];
        x -= (x & -x);
    }
    return res;
}

// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
    int idx = 0, mask = N;
    while(mask && idx < N) {
        int t = idx + mask;
        if(x >= tree[t]) {
            idx = t;
            x -= tree[t];
        }
        mask >>= 1;
    }
    return idx;
}

```

3.5 Union Find (Short)

```

#include <iostream>
#include <vector>
using namespace std;
int find(vector<int> &C, int x) { return (C[x] == x) ? x : C[x] =
    find(C, C[x]); }
void merge(vector<int> &C, int x, int y) { C[find(C, x)] = find(C,
    y); }
int main()
{
    int n = 5;
    vector<int> C(n);
    for (int i = 0; i < n; i++) C[i] = i;
    merge(C, 0, 2);
    merge(C, 1, 0);
    merge(C, 3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << find(C, i)
        << endl;
    return 0;
}

```

3.6 Union Find

```

/*----- disjoint sets-----*/
#include<bits/stdc++.h>

```

```

#define TAM 10000
using namespace std;
class UnionFind{
private:
    vector<int> p, rank, ssize;
    int numSets;
public:
    UnionFind(int N){
        rank.assign(N, 0);
        ssize.assign(N, 1);
        numSets = N;
        p.assign(N, 0);
        for(int i = 0; i < N; i++)
            p[i] = i;
    }
    int findSet(int i){
        return (p[i] == i) ? i : (p[i] = findSet(p[i]));
    }
    bool isSameSet(int i, int j){
        return findSet(i) == findSet(j);
    }
    void unionSet(int i, int j){
        if(!isSameSet(i, j)){
            numSets--;
            int x = findSet(i), y = findSet(j);
            if(rank[x] > rank[y]){
                p[y] = x;
                ssize[x] += ssize[y];
            }
            else{
                p[x] = y;
                ssize[y] += ssize[x];
                if(rank[x] == rank[y])
                    rank[y]++;
            }
        }
    }
    int numDisjointSets(){
        return numSets;
    }
    int sizeOfSet(int i){
        return ssize[findSet(i)];
    }
};

```

4 Strings

4.1 KMP

```

/*
Finds all occurrences of the pattern string p within the
text string t. Running time is  $O(n + m)$ , where  $n$  and  $m$ 
are the lengths of  $p$  and  $t$ , respectively.
*/

```

```

#include <iostream>
#include <string>
#include <vector>

using namespace std;

```

```

typedef vector<int> VI;
void buildPi(string& p, VI& pi)
{
    pi = VI(p.length());
    int k = -2;
    for(int i = 0; i < p.length(); i++) {
        while(k >= -1 && p[k+1] != p[i])
            k = (k == -1) ? -2 : pi[k];
        pi[i] = ++k;
    }
}

int KMP(string& t, string& p)
{
    VI pi;
    buildPi(p, pi);
    int k = -1;
    for(int i = 0; i < t.length(); i++) {
        while(k >= -1 && p[k+1] != t[i])
            k = (k == -1) ? -2 : pi[k];
        k++;
        if(k == p.length() - 1) {
            // p matches t[i-m+1, ..., i]
            cout << "matched at index " << i-k << ": ";
            cout << t.substr(i-k, p.length()) << endl;
            k = (k == -1) ? -2 : pi[k];
        }
    }
    return 0;
}

int main()
{
    string a = "AABAACAADAABAABA", b = "AABA";
    KMP(a, b); // expected matches at: 0, 9, 12
    return 0;
}

```

4.2 Suffix Array

```

// Suffix array construction in  $O(L \log^2 L)$  time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in  $O(\log L)$  time.
//
// INPUT:    string s
//
// OUTPUT:   array suffix[] such that suffix[i] = index (from 0 to
//           L-1)
//           of substring s[i...L-1] in the list of sorted suffixes
//           .
//           That is, if we take the inverse of the permutation
//           suffix[],
//           we get the actual suffix array.

```

```

#include <vector>
#include <iostream>
#include <string>

using namespace std;

struct SuffixArray {
    const int L;
    string s;
    vector<vector<int>> > P;

```

```

vector<pair<pair<int,int>,int> > M;
SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<
    int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
    for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
        P.push_back(vector<int>(L, 0));
        for (int i = 0; i < L; i++)
            M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P
                [level-1][i + skip] : -1000), i);
        sort(M.begin(), M.end());
        for (int i = 0; i < L; i++)
            P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].
                first) ? P[level][M[i-1].second] : i;
    }
}

vector<int> GetSuffixArray() { return P.back(); }

// returns the length of the longest common prefix of s[i...L-1]
// and s[j...L-1]
int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
        if (P[k][i] == P[k][j]) {
            i += 1 << k;
            j += 1 << k;
            len += 1 << k;
        }
    }
    return len;
}

// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
    int T;
    cin >> T;
    for (int caseno = 0; caseno < T; caseno++) {
        string s;
        cin >> s;
        SuffixArray array(s);
        vector<int> v = array.GetSuffixArray();
        int bestlen = -1, bestpos = -1, bestcount = 0;
        for (int i = 0; i < s.length(); i++) {
            int len = 0, count = 0;
            for (int j = i+1; j < s.length(); j++) {
                int l = array.LongestCommonPrefix(i, j);
                if (l >= len) {
                    if (l > len) count = 2; else count++;
                    len = l;
                }
            }
            if (len > bestlen || len == bestlen && s.substr(bestpos,
                bestlen) > s.substr(i, len)) {
                bestlen = len;
                bestcount = count;
                bestpos = i;
            }
        }
        if (bestlen == 0) {
            cout << "No repetitions found!" << endl;
        } else {

```

```

            cout << s.substr(bestpos, bestlen) << " " << bestcount <<
                endl;
        }
    }
}

#else
// END CUT
int main() {
    // bobocel is the 0'th suffix
    // obocel is the 5'th suffix
    // bocel is the 1'st suffix
    // ocel is the 6'th suffix
    // cel is the 2'nd suffix
    // el is the 3'rd suffix
    // l is the 4'th suffix
    SuffixArray suffix("bobocel");
    vector<int> v = suffix.GetSuffixArray();

    // Expected output: 0 5 1 6 2 3 4
    // 2
    for (int i = 0; i < v.size(); i++) cout << v[i] << " ";
    cout << endl;
    cout << suffix.LongestCommonPrefix(0, 2) << endl;
}

// BEGIN CUT
#endif
// END CUT

```

4.3 Fast Fourier Transform (convolution)

```

// Convolution using the fast Fourier transform (FFT).
//
// INPUT:
//     a[1...n]
//     b[1...m]
//
// OUTPUT:
//     c[1...n+m-1] such that c[k] = sum_{i=0}^k a[i] b[k-i]
//
// Alternatively, you can use the DFT() routine directly, which
// will
// zero-pad your input to the next largest power of 2 and compute
// the
// DFT or inverse DFT.

#include <iostream>
#include <vector>
#include <complex>

using namespace std;

typedef long double DOUBLE;
typedef complex<DOUBLE> COMPLEX;
typedef vector<DOUBLE> VD;
typedef vector<COMPLEX> VC;

struct FFT {
    VC A;
    int n, L;

    int ReverseBits(int k) {
        int ret = 0;
        for (int i = 0; i < L; i++) {
            ret = (ret << 1) | (k & 1);

```

```

    k >>= 1;
}
return ret;
}

void BitReverseCopy(VC a) {
    for (n = 1, L = 0; n < a.size(); n <= 1, L++) ;
    A.resize(n);
    for (int k = 0; k < n; k++)
        A[ReverseBits(k)] = a[k];
}

VC DFT(VC a, bool inverse) {
    BitReverseCopy(a);
    for (int s = 1; s <= L; s++) {
        int m = 1 << s;
        COMPLEX wm = exp(COMPLEX(0, 2.0 * M_PI / m));
        if (inverse) wm = COMPLEX(1, 0) / wm;
        for (int k = 0; k < n; k += m) {
            COMPLEX w = 1;
            for (int j = 0; j < m/2; j++) {
                COMPLEX t = w * A[k + j + m/2];
                COMPLEX u = A[k + j];
                A[k + j] = u + t;
                A[k + j + m/2] = u - t;
                w = w * wm;
            }
        }
    }
    if (inverse) for (int i = 0; i < n; i++) A[i] /= n;
    return A;
}

// c[k] = sum_{i=0}^k a[i] b[k-i]
VD Convolution(VD a, VD b) {
    int L = 1;
    while ((1 << L) < a.size()) L++;
    while ((1 << L) < b.size()) L++;
    int n = 1 << (L+1);

    VC aa, bb;
    for (size_t i = 0; i < n; i++) aa.push_back(i < a.size() ?
        COMPLEX(a[i], 0) : 0);
    for (size_t i = 0; i < n; i++) bb.push_back(i < b.size() ?
        COMPLEX(b[i], 0) : 0);

    VC AA = DFT(aa, false);
    VC BB = DFT(bb, false);
    VC CC;
    for (size_t i = 0; i < AA.size(); i++) CC.push_back(AA[i] * BB
        [i]);
    VC cc = DFT(CC, true);
    VD c;
    for (int i = 0; i < a.size() + b.size() - 1; i++) c.push_back(
        cc[i].real());
    return c;
}

};

int main() {
    double a[] = {1, 3, 4, 5, 7};
    double b[] = {2, 4, 6};

    FFT fft;
    VD c = fft.Convolution(VD(a, a + 5), VD(b, b + 3));

    // expected output: 2 10 26 44 58 58 42
    for (int i = 0; i < c.size(); i++) cerr << c[i] << " ";
}

```

```

    cerr << endl;
    return 0;
}

```

5 Flows

5.1 Ford Fulkerson

```

//----- Ford-Fulkerson O(MaxFlow * |E|)
//-----

struct OutEdge {
    int to, cap, rIdx;
    OutEdge ( ) {}
    OutEdge(int to, int cap, int rIdx) :
        to(to), cap(cap), rIdx(rIdx) {}
};

struct Network
{
    vector<vector<OutEdge> > out;
    vector<bool> seen;

    int sink;
    int augment ( int i, const int cur ) {
        if ( i == sink ) return cur;
        if ( seen[i] ) return false;
        seen[i] = true;
        int ans;
        for ( OutEdge& e : out[i] )
            if ( e.cap > 0 && ( ans = augment(e.to, min
                (cur, e.cap)) ) ) {
                e.cap -= ans;
                out[e.to][e.rIdx].cap += ans;
                return ans;
            }
        return 0;
    }

    int maxflow ( int source, int _sink ) {
        sink = _sink;
        int curflow = 0, aug;
        while ( true ) {
            fill ( seen.begin(), seen.end(), false );
            aug = augment(source, INT_MAX);
            if ( aug == 0 ) break;
            curflow += aug;
        }
        return curflow;
    }

    void addEdge ( int fr, int to, int c ) {
        assert ( fr != to );
        out[fr].push_back(OutEdge(to, c, out[to].size()));
        out[to].push_back(OutEdge(fr, 0, out[fr].size() -
            1));
    }

    Network(int n) {
        out.assign(n, vector<OutEdge>());
        seen.resize(n);
    }
};

// Adjacency list implementation of FIFO push relabel maximum flow

```



```

// with the gap relabeling heuristic. This implementation is
// significantly faster than straight Ford-Fulkerson. It solves
// random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
//
// Running time:
//       $O(|V|^3)$ 
//
// INPUT:
//      - graph, constructed using AddEdge()
//      - source
//      - sink
//
// OUTPUT:
//      - maximum flow value
//      - To obtain the actual flow values, look at all edges with
//        capacity > 0 (zero capacity edges are residual edges).

#include <cmath>
#include <vector>
#include <iostream>
#include <queue>

using namespace std;

typedef long long LL;

struct Edge {
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int flow, int index) :
        from(from), to(to), cap(cap), flow(flow), index(index) {}
};

struct PushRelabel {
    int N;
    vector<vector<Edge>> > G;
    vector<LL> excess;
    vector<int> dist, active, count;
    queue<int> Q;

    PushRelabel(int N) : N(N), G(N), excess(N), dist(N), active(N),
        count(2*N) {}

    void AddEdge(int from, int to, int cap) {
        G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
        if (from == to) G[from].back().index++;
        G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
    }

    void Enqueue(int v) {
        if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
    }

    void Push(Edge &e) {
        int amt = min(excess[e.from], LL(e.cap - e.flow));
        if (dist[e.from] <= dist[e.to] || amt == 0) return;
        e.flow += amt;
        G[e.to][e.index].flow -= amt;
        excess[e.to] += amt;
        excess[e.from] -= amt;
        Enqueue(e.to);
    }

    void Gap(int k) {
        for (int v = 0; v < N; v++) {
            if (dist[v] < k) continue;
            count[dist[v]]--;

```

```

            dist[v] = max(dist[v], N+1);
            count[dist[v]]++;
            Enqueue(v);
        }
    }

    void Relabel(int v) {
        count[dist[v]]--;
        dist[v] = 2*N;
        for (int i = 0; i < G[v].size(); i++)
            if (G[v][i].cap - G[v][i].flow > 0)
                dist[v] = min(dist[v], dist[G[v][i].to] + 1);
        count[dist[v]]++;
        Enqueue(v);
    }

    void Discharge(int v) {
        for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);
        if (excess[v] > 0) {
            if (count[dist[v]] == 1)
                Gap(dist[v]);
            else
                Relabel(v);
        }
    }

    LL GetMaxFlow(int s, int t) {
        count[0] = N-1;
        count[N] = 1;
        dist[s] = N;
        active[s] = active[t] = true;
        for (int i = 0; i < G[s].size(); i++) {
            excess[s] += G[s][i].cap;
            Push(G[s][i]);
        }

        while (!Q.empty()) {
            int v = Q.front();
            Q.pop();
            active[v] = false;
            Discharge(v);
        }

        LL totflow = 0;
        for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;
        return totflow;
    }
};

// Adjacency list implementation of Dinic's blocking flow
// algorithm.
// This is very fast in practice, and only loses to push-relabel
// flow.
//
// Running time:
//       $O(|V|^2 |E|)$ 
//
// INPUT:
//      - graph, constructed using AddEdge()
//      - source
//      - sink
//
// OUTPUT:
//      - maximum flow value
//      - To obtain the actual flow values, look at all edges with
//        capacity > 0 (zero capacity edges are residual edges).

```

```

#include <cmath>
#include <vector>
#include <iostream>
#include <queue>

using namespace std;

const int INF = 2000000000;

struct Edge {
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int flow, int index) :
        from(from), to(to), cap(cap), flow(flow), index(index) {}
};

struct Dinic {
    int N;
    vector<vector<Edge>> > G;
    vector<Edge*> dad;
    vector<int> Q;

    Dinic(int N) : N(N), G(N), dad(N), Q(N) {}

    void AddEdge(int from, int to, int cap) {
        G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
        if (from == to) G[from].back().index++;
        G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
    }

    long long BlockingFlow(int s, int t) {
        fill(dad.begin(), dad.end(), (Edge*) NULL);
        dad[s] = &G[0][0] - 1;

        int head = 0, tail = 0;
        Q[tail++] = s;
        while (head < tail) {
            int x = Q[head++];
            for (int i = 0; i < G[x].size(); i++) {
                Edge &e = G[x][i];
                if (!dad[e.to] && e.cap - e.flow > 0) {
                    dad[e.to] = &G[x][i];
                    Q[tail++] = e.to;
                }
            }
        }
        if (!dad[t]) return 0;

        long long totflow = 0;
        for (int i = 0; i < G[t].size(); i++) {
            Edge *start = &G[G[t][i].to][G[t][i].index];
            int amt = INF;
            for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
                if (!e) { amt = 0; break; }
                amt = min(amt, e->cap - e->flow);
            }
            if (amt == 0) continue;
            for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
                e->flow += amt;
                G[e->to][e->index].flow -= amt;
            }
            totflow += amt;
        }
        return totflow;
    }

    long long GetMaxFlow(int s, int t) {
        long long totflow = 0;

```

```

        while (long long flow = BlockingFlow(s, t))
            totflow += flow;
        return totflow;
    }
};

```

5.2 Edmonds Karp Min cut

```

//----- Edmonds Karp with MinCut  $O(|V|*|E|^2)$ 
//-----

struct Network {
    vector<Edge> G[TAM];
    int from[TAM], n;
    bool color[TAM];

    // Call flood (source) to color one node
    // component of min cut.
    void flood ( int node ) {
        if ( color[node] ) return;
        color[node] = true;
        for ( const Edge& e : G[node] )
            if ( e.cap > 0 )
                flood ( e.to );
    }

    int maxFlow ( int A, int B )
    {
        int flow = 0;
        while ( 1 ) {
            memset ( from, -1, sizeof(from) );

            queue<int> q;
            q.push ( A );
            from[A] = -2;
            for ( int i; !q.empty(); q.pop() ) {
                i = q.front();
                for ( Edge& e : G[i] )
                    if ( from[e.to] == -1 && e
                        .cap ) {
                        from[e.to] = e
                            .invIdx;
                        q.push ( e.to );
                    }
            }

            if ( from[B] == -1 ) break;

            int aug = INF_CAP;
            for ( int i = B, j; i != A; i = j ) {
                j = G[i][from[i]].to;
                aug = min ( aug, G[j][ G[i][from[i]
                    ].invIdx ].cap );
            }

            for ( int i = B, j; i != A; i = j ) {
                j = G[i][from[i]].to;
                G[j][ G[i][from[i]].invIdx ].cap
                    -= aug;
                G[i][from[i]].cap += aug;
            }

            flow += aug;
        }

        return flow;
    }
};

```

```

void addNonDirEdge ( int a, int b, int c ) {
    assert ( a != b );
    G[a].push_back ( Edge(b,c,G[b].size()) );
    G[b].push_back ( Edge(a,c,G[a].size()-1) );
}

void addDirEdge ( int a, int b, int c ) {
    assert ( a != b );
    G[a].push_back ( Edge(b,c,G[b].size()) );
    G[b].push_back ( Edge(a,0,G[a].size()-1) );
}

void clear ( int _n ) {
    n = _n;
    memset ( color, false, n );
    for ( int i = 0; i < n; ++i )
        G[i].clear();
}

} netw;

```

5.3 Hopcroft karp's maximum bipartite matching

```

//----- Hopcroft Karp - Maximum Bipartite Matching O( sqrt(|V|)
* |E| ) -----

namespace hopcroftKarp {

    const int MAXN1 = 50000;
    const int MAXN2 = 50000;
    const int MAXM = 150000;

    int n1, n2, edges, last[MAXN1], prev[MAXM], head[MAXM];
    int matching[MAXN2], dist[MAXN1], Q[MAXN1];
    bool used[MAXN1], vis[MAXN1];

    void init(int _n1, int _n2) {
        n1 = _n1;
        n2 = _n2;
        edges = 0;
        fill(last, last + n1, -1);
    }

    void addEdge(int u, int v) {
        head[edges] = v;
        prev[edges] = last[u];
        last[u] = edges++;
    }

    void bfs() {
        fill(dist, dist + n1, -1);
        int sizeQ = 0;
        for (int u = 0; u < n1; ++u) {
            if (!used[u]) {
                Q[sizeQ++] = u;
                dist[u] = 0;
            }
        }
        for (int i = 0; i < sizeQ; i++) {
            int u1 = Q[i];
            for (int e = last[u1]; e >= 0; e = prev[e]) {
                int u2 = matching[head[e]];
                if (u2 >= 0 && dist[u2] < 0) {
                    dist[u2] = dist[u1] + 1;
                    Q[sizeQ++] = u2;
                }
            }
        }
    }
}

```

```

}

bool dfs(int u1) {
    vis[u1] = true;
    for (int e = last[u1]; e >= 0; e = prev[e]) {
        int v = head[e];
        int u2 = matching[v];
        if (u2 < 0 || !vis[u2] && dist[u2] == dist[u1] + 1 && dfs(u2)) {
            matching[v] = u1;
            used[u1] = true;
            return true;
        }
    }
    return false;
}

int maxMatching() {
    fill(used, used + n1, false);
    fill(matching, matching + n2, -1);
    for (int res = 0;;) {
        bfs();
        fill(vis, vis + n1, false);
        int f = 0;
        for (int u = 0; u < n1; ++u)
            if (!used[u] && dfs(u))
                ++f;

        if (!f)
            return res;
        res += f;
    }
}
};

```

5.4 Maxmium bipartite matching (short but slower)

```

//----- Maximum Bipartite Matching O(|V|*|E|)
-----

bool findMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
    for (int j = 0; j < int(w[i].size()); j++) {
        if (w[i][j] && !seen[j]) {
            seen[j] = true;
            if (mc[j] < 0 || findMatch(mc[j], w, mr, mc, seen)) {
                mr[i] = j; mc[j] = i;
                return true;
            }
        }
    }
    return false;
}

int maxBipartiteMatching(const VVI &w) {
    if (w.empty() || w[0].empty()) return 0;
    VI mr(w.size(), -1), mc(w[0].size(), -1);
    int ct = 0;
    for (int i = 0; i < int(w.size()); i++) {
        VI seen(w[0].size());
        if (findMatch(i, w, mr, mc, seen)) ct++;
    }
    return ct;
}

```

6 Math

6.1 general math tricks

```
long square(long n){ return n*n;}

int fastPow(long x, long n){
    if(n == 0)
        return 1;

    if(n % 2 == 0)
        return square(fastPow(x, n/2));

    return x * (fastPow(x, n - 1));
}

/* LCM */
int LCM(int m, n){return (m*n)/__gcd(m, n); }

int main(){
    /* n es impar? */
    odd = ((n & 1)? true : false);

    /*como saber si un numero es una potencia de 2*/
    power_of_2 = ((v & (v-1)) == 0);

    /*contar trailing 0's de una mascara */
    __builtin_ctz(n);

    /*contar 1's de una mascara*/
    __builtin_popcount(n);

    /*quitar el elemento j de la mascara*/
    mask &= ~(1<<j);

    /*revisar si el elemento j del arreglo esta en la mascara (
    si es 0 el resultado es porque no est )*/
    int t = mask & (1<<j);

    /*Obtener el bit menos significativo*/
    t = mask & -mask

    /*encender todos los n primeros bits de la mascara*/
    mask = (1<<n) - 1;

    /*iterar sobre cada uno de los subsets de un subset y*/
    for(int x = y; x>0; x = (y & (x-1)) )
}
```

6.2 Miller Rabin's primality test

```
// Randomized Primality Test (Miller-Rabin):
// Error rate: 2^(-TRIAL)
// Almost constant time. srand is needed

#include <stdlib.h>
#define EPS 1e-7

typedef long long LL;

LL ModularMultiplication(LL a, LL b, LL m)
{
```

```
LL ret=0, c=a;
while(b)
{
    if(b&1) ret=(ret+c)%m;
    b>>=1; c=(c+c)%m;
}
return ret;
}

LL ModularExponentiation(LL a, LL n, LL m)
{
    LL ret=1, c=a;
    while(n)
    {
        if(n&1) ret=ModularMultiplication(ret, c, m);
        n>>=1; c=ModularMultiplication(c, c, m);
    }
    return ret;
}

bool Witness(LL a, LL n)
{
    LL u=n-1;
    int t=0;
    while(!(u&1)){u>>=1; t++;}
    LL x0=ModularExponentiation(a, u, n), x1;
    for(int i=1; i<=t; i++)
    {
        x1=ModularMultiplication(x0, x0, n);
        if(x1==1 && x0!=1 && x0!=n-1) return true;
        x0=x1;
    }
    if(x0!=1) return true;
    return false;
}

LL Random(LL n)
{
    LL ret=rand(); ret*=32768;
    ret+=rand(); ret*=32768;
    ret+=rand(); ret*=32768;
    ret+=rand();
    return ret%n;
}

bool IsPrimeFast(LL n, int TRIAL)
{
    while(TRIAL-->0)
    {
        LL a=Random(n-2)+1;
        if(Witness(a, n)) return false;
    }
    return true;
}
```

6.3 Pollard rho

```
#include<bits/stdc++.h>
#include<time.h>

#define show(x) cout << #x << " = " << x << endl;

using namespace std;

typedef long long ll;
typedef pair<ll, ll> ii;
typedef pair<double, ii> iii;

const int MAX = 200005;
const double EPS = 1e-5;
const int INF = INT_MAX;
```

6.4 number theory general

```
//modular multiplication for really big numbers
ll mul(ll a, ll b, ll mod) {
    ll ret = 0;
    for(a %= mod, b %= mod; b != 0;
        b >>= 1, a <=<= 1, a = a >= mod ? a - mod : a) {
        if (b&1) {
            ret += a;
            if (ret >= mod) ret -= mod;
        }
    }
    return ret;
}

ll fpow(ll a, ll b, ll MOD) {
    ll ans = 1LL;
    while(b > 0) {
        if(b&1) ans = mul(ans, a, MOD);
        a = mul(a, a, MOD);
        b >>= 1LL;
    }
    return ans;
}

const int rounds = 6;
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
bool miller_rabin(ll n) {
    if(n == 2 || n == 3) return true;
    if(n < 2 || (n&1) == 0) return false;
    for(int i = 0; i < rounds; i++) {
        int a = rand()%(n-4)+2;
        if(fpow(a, n-1, n) != 1)
            return false;
    }
    return true;
}

// if n is prime , check with miller rabin (n^(1/4)) and check
return != n and != 1
ll pollard_rho(ll n, ll c) {
    ll x = 2, y = 2, i = 1, k = 2, d;
    while (true) {
        x = (mul(x, x, n) + c);
        if (x >= n) x -= n;
        d = __gcd(x - y, n);
        if (d > 1) return d;
        if (++i == k) y = x, k <=<= 1;
    }
    return n;
}

//return factorization of a big number
void factorize(ll n, vector<ll> &f) {
    if(n == 1) return;
    if (miller_rabin(n)) {
        f.push_back(n);
        return;
    }
    ll d = n;
    for (int i = 2; d == n; i++)
        d = pollard_rho(n, i);
    factorize(d, f);
    factorize(n/d, f);
}
```

// This is a collection of useful code for solving problems that involve modular linear equations. Note that all of the algorithms described here work on nonnegative integers.

```
#include <iostream>
#include <vector>
#include <algorithm>

using namespace std;

typedef vector<int> VI;
typedef pair<int, int> PII;

// return a % b (positive value)
int mod(int a, int b) {
    return ((a%b) + b) % b;
}

// computes gcd(a,b)
int gcd(int a, int b) {
    while (b) { int t = a%b; a = b; b = t; }
    return a;
}

// computes lcm(a,b)
int lcm(int a, int b) {
    return a / gcd(a, b)*b;
}

// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
{
    int ret = 1;
    while (b)
    {
        if (b & 1) ret = mod(ret*a, m);
        a = mod(a*a, m);
        b >>= 1;
    }
    return ret;
}

// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
    int xx = y = 0;
    int yy = x = 1;
    while (b) {
        int q = a / b;
        int t = b; b = a%b; a = t;
        t = xx; xx = x - q*xx; x = t;
        t = yy; yy = y - q*yy; y = t;
    }
    return a;
}

// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
    int x, y;
    VI ret;
    int g = extended_euclid(a, n, x, y);
    if (!(b%g)) {
        x = mod(x*(b / g), n);
        for (int i = 0; i < g; i++)
            ret.push_back(mod(x + i*(n / g), n));
    }
}
```

```

    return ret;
}

// computes b such that ab = 1 (mod n), returns -1 on failure
int mod_inverse(int a, int n) {
    int x, y;
    int g = extended_euclid(a, n, x, y);
    if (g > 1) return -1;
    return mod(x, n);
}

// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1,
// m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
    int s, t;
    int g = extended_euclid(m1, m2, s, t);
    if (r1%g != r2%g) return make_pair(0, -1);
    return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2);
}

// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
    PII ret = make_pair(r[0], m[0]);
    for (int i = 1; i < m.size(); i++) {
        ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
        if (ret.second == -1) break;
    }
    return ret;
}

// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
    if (!a && !b)
    {
        if (c) return false;
        x = 0; y = 0;
        return true;
    }
    if (!a)
    {
        if (c % b) return false;
        x = 0; y = c / b;
        return true;
    }
    if (!b)
    {
        if (c % a) return false;
        x = c / a; y = 0;
        return true;
    }
    int g = gcd(a, b);
    if (c % g) return false;
    x = c / g * mod_inverse(a / g, b / g);
    y = (c - a*x) / b;
    return true;
}

int main() {
    // expected: 2

```

```

    cout << gcd(14, 30) << endl;

    // expected: 2 -2 1
    int x, y;
    int g = extended_euclid(14, 30, x, y);
    cout << g << " " << x << " " << y << endl;

    // expected: 95 451
    VI sols = modular_linear_equation_solver(14, 30, 100);
    for (int i = 0; i < sols.size(); i++) cout << sols[i] << "
";
    cout << endl;

    // expected: 8
    cout << mod_inverse(8, 9) << endl;

    // expected: 23 105
    //          11 12
    PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({
        2, 3, 2 }));
    cout << ret.first << " " << ret.second << endl;
    ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }
    ));
    cout << ret.first << " " << ret.second << endl;

    // expected: 5 -15
    if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" <<
endl;
    cout << x << " " << y << endl;
    return 0;
}

```

7 Miscellaneous

7.1 c++ ios tricks

```

#include <iostream>
#include <iomanip>

using namespace std;

int main()
{
    // Output a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);

    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);

    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);

    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec <<
endl;
}

```

7.2 java IO template and iterative binary search

```
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.PrintWriter;
import java.util.StringTokenizer;
import java.io.IOException;
import java.io.BufferedReader;
import java.io.InputStreamReader;
import java.io.InputStream;

public class Main {
    public static void main(String[] args) {
        InputStream inputStream = System.in;
        OutputStream outputStream = System.out;
        InputReader in = new InputReader(inputStream);
        PrintWriter out = new PrintWriter(outputStream);
        TaskC solver = new TaskC();
        solver.solve(1, in, out);
        out.close();
    }

    static class TaskC {
        private static final int ITERATIONS = 500;
        public void solve(int testNumber, InputReader in,
            PrintWriter out) {
            int n = in.nextInt();
            //Iterative binary search
            double l = 0.0, h = 1e17;
            for (int i = 0; i < ITERATIONS; i++) {
                double mid = (l + h) / 2.0;
                if (can(mid, a, b, p))
                    l = mid;
                else
                    h = mid;
            }
        }
    }
}
```

```
    }
    double ans = 1;
}

static class InputReader {
    public BufferedReader reader;
    public StringTokenizer tokenizer;

    public InputReader(InputStream stream) {
        reader = new BufferedReader(new InputStreamReader(
            stream), 32768);
        tokenizer = null;
    }

    public String next() {
        while (tokenizer == null || !tokenizer.hasMoreTokens())
        {
            try {
                tokenizer = new StringTokenizer(reader.
                    readLine());
            } catch (IOException e) {
                throw new RuntimeException(e);
            }
        }
        return tokenizer.nextToken();
    }

    public int nextInt() {
        return Integer.parseInt(next());
    }

    public double nextDouble() {
        return Double.parseDouble(next());
    }
}
}
```