# magUNbos' notebook (2017)

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# 1 Graphs

### 1.1 Articulation points

```
/*Tarjan Algorithm to find articulation points
 * single dfs 0(|v| + |e|)
 * visited =[false]
 * disc = [0]
 \star low = [0]
 * parent = [-1]
 * ap = [false]
 * result is stored in ap boolean array
 * Tested in AIZU online Judge*/
#include<bist/stdc++.h>
using namespace std;
const int SIZE = 100013;
bool visited[SIZE], ap[SIZE];
int disc[SIZE], low[SIZE], parent[SIZE];
vector<int> G[SIZE];
void articulation(int u) {
  static int time = 0;
  int children = 0;
  visited[u] = true;
  disc[u] = low[u] = ++time;
  for(int i = 0; i < G[u].size(); i++) {</pre>
      int v = G[u][i];
      if(!visited[v]){
          children++;
          parent[v] = u;
          articulation(v);
          low[u] = min(low[u], low[v]);
          if(parent[u] == -1 && children > 1)
            ap[u] = true;
          if(parent[u] != -1 \&\& low[v] >= disc[u])
            ap[u] = true;
      else if(v != parent[u])
        low[u] = min(low[u], disc[v]);
```

# 1.2 Biconected graph

```
/* Tarjan Algorithm to find Biconnected graph
 * single dfs O(|v| + |e|)
 * visited =[false]
 * disc = [0]
 * low = [0]
 * parent = [-1] */

bool isBiconnected(vector<vector<int> > G, int u, bool visited[],
 int disc[], int low[], int parent[]) {
    static int time = 0;
    int children = 0;
    visited[u] = true;
    disc[u] = low[u] = ++time;
```

```
for(int i = 0; i < G[u].size(); i++) {
    int v = G[u][i];

    if(!visited[v]) {
        children++;
        parent[v] = u;

        if (isBiconnected(G, v, visited, disc, low, parent))
            return true;

        low[u] = min(low[u], low[v]);

        if(parent[u] == -1 && children > 1) return true;

        if(parent[u] != -1 && low[v] >= disc[u]) return true;
        else if(v != parent[v]) low[u] = min(low[u], disc[v]);
    }
    return false;
}
```

## 1.3 Bridges

```
/* Tarjan Algorithm to find bridges
 * single dfs O(|v| + |E|)
 * visited =[false]
 * disc = [0]
 * low = [0]
 * parent = [-1]
 * the priority queue orders the bridges in ascending order,
 * use the function like: bridges(0, &queue)
 * tested in AIZU online Judge
#include<bits/stdc++.h>
using namespace std;
const int SIZE = 100013;
typedef pair<int, int> pii;
bool visited[SIZE];
int disc[SIZE], low[SIZE], parent[SIZE];
vector<int> G[SIZE];
void bridges(int u, priority_queue<pii, vector<pii>, greater<pii>
   > *bridge) {
  static int time = 0;
  int children = 0;
  visited[u] = true;
  disc[u] = low[u] = ++time;
  for(int i = 0; i < G[u].size(); i++) {</pre>
    int v = G[u][i];
    if(!visited[v]){
      children++;
      parent[v] = u;
      bridges (v, bridge);
      low[u] = min(low[u], low[v]);
      if(low[v] > disc[u]) bridge->push({min(u,v),max(u,v)});
    else if(v != parent[u])
      low[u] = min(low[u], disc[v]);
```

## 1.4 Tarjan SCC

```
/* Tarjan Algorithm to find connected components
 * single dfs 0(|v| + |e|)
 * visited =[false]
 * disc = [0]
 * low = [0]
 * parent = [-1]
 * tested on AIZU online Judge */
void dfsSCC(vector<vector<int> > G, int u, int disc[], int low[],
      stack<int> *st, bool stackMember[]) {
     static int time = 0;
     disc[u] = low[u] = ++time;
     st->push(u);
     stackmember[u] = true;
     for (int i = 0; i < G[u].size(); i++) {
         int v = G[u][i];
         if(disc[v] == -1){
             dfsSCC(G, v, disc, low, st, stackmember);
             low[u] = min(low[u], low[v]);
         else if(stackmember[v] == true) low[u] = min(low[u], disc
             [V]);
     int w = 0;
     if(low[u] == disc[u]){
         while (st->top() != u) {
             w = st - > top();
             cout << w << " ";
             stackmember[w] = false;
             st->pop();
         w = st->top();
         cout << w << "\n";
         stackmember[w] = false;
         st->pop();
void scc(G) {
    int *disc = new int[V];
    int *low = new int[V];
    bool *stackMember = new bool[V];
    stack<int> *st = new stack<int>();
    memset(disc, -1, sizeof(disc));
    memset(low, 0, sizeof(low));
    memset(stackMember, false, sizeof(stackMember));
    for(int i = 0; i < G.size(); i++)</pre>
        if(disc[i] == -1) dfsScc(G, i, disc, low, st, stackMember)
```

#### 1.5 Bellman Ford

```
/* Bellman Ford's algorithm to find SSSP with negative cycles O(V
 * returns true if it detects a negative cycle
 * dist [i] contains distance to i
 * parent[i] contains parent on path to i from source
 * tested on AIZU online judge
#include<bits/stdc++.h>
const int SIZE = 1013;
const int oo = 999999;
typedef pair<int, int> pii;
int dist[SIZE];
int parent[SIZE];
bool bellmanFord(vector<vector<pii> > &G, int &source, int &N) {
    for (int i = 0; i < N; i++)
      dist[i] = oo;
    dist[source] = 0:
    for (int k = 0; k < N; k++) {
      for (int u = 0; u < N; u++) {
        for(int i = 0; i < G[u].size(); i++) {</pre>
          pii v = G[u][i];
          if(dist[v.first] > dist[u] + v.second && dist[u] != oo){
            if(k == N-1)
              return true;
            dist[v.first] = dist[u] + v.second;
            parent[v.first] = u;
    return false:
```

#### 1.6 Eulerian Path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next_vertex;
        iter reverse_edge;
        Edge(int next_vertex)
                :next_vertex(next_vertex)
                { }
};
const int max_vertices = ;
int num_vertices;
                                      // adjacency list
list<Edge> adj[max_vertices];
vector<int> path;
void find_path(int v)
        while (adj[v].size() > 0)
```

## 1.7 Topological sort

```
char c[TAM];
int l[TAM];
int r[TAM];
int in[TAM];
//can be priority queue
queue<int> 0;
void reset(){
    memset(1, 0, sizeof 1);
    memset(r, 0, sizeof r);
    memset(in, 0, sizeof in);
    memset(balls, 0, sizeof balls);
    c[0] = 'L';
void topo(vector<vector<int> > G, int u) {
    while(!Q.empty()){
        u = Q.front();Q.pop();
        update(u);
        for(int i = 0; i <G[u].size(); i++) {</pre>
            int v = G[u][i];
            in[v]--;
            if(in[v] == 0) Q.push(v);
int main(){
    11 n;
    int m;
    while(cin>>n>>m) {
        reset();
        vector< vector<int> > G(m + 1);
        for(int i = 1; i <=m; i++) {</pre>
            int u, v;
            cin>>c[i]>>u>>v;
            G[i].push_back(u);
            G[i].push_back(v);
            in[u] ++; in[v]++;
            l[i] = u; r[i] = v;
        for(int i = 1; i <=m; i++) {</pre>
            if(in[i] == 0)Q.push(i);
        topo(G, 1);
```

## }

### 1.8 Kruskal's minimum spanning tree

```
Uses Kruskal's Algorithm to calculate the weight of the minimum
    spanning
forest (union of minimum spanning trees of each connected
    component) of
a possibly disjoint graph, given in the form of a matrix of edge
(-1 if no edge exists). Returns the weight of the minimum spanning
forest (also calculates the actual edges - stored in T). Note:
disjoint-set data structure with amortized (effectively) constant
    time per
union/find. Runs in O(E*log(E)) time.
tested on AIZU online Judge
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
typedef int T;
struct edge
  int u, v;
  T d;
} ;
struct edgeCmp
  int operator()(const edge& a, const edge& b) { return a.d > b.d;
} ;
int find(vector \langle int \rangle_{\&} C, int x) { return (C[x] == x) ? x : C[x] =
     find(C, C[x]); }
T Kruskal (vector <vector <T> >& w)
  int n = w.size();
  T weight = 0;
  vector <int> C(n), R(n);
  for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }</pre>
  vector <edge> T;
  priority_queue <edge, vector <edge>, edgeCmp> E;
  for(int i=0; i<n; i++)
    for (int j=i+1; j<n; j++)</pre>
      if(w[i][j] >= 0)
        e.u = i; e.v = j; e.d = w[i][j];
        E.push(e);
  while (T.size() < n-1 \&\& !E.empty())
```

```
edge cur = E.top(); E.pop();
   int uc = find(C, cur.u), vc = find(C, cur.v);
      T.push_back(cur); weight += cur.d;
     if(R[uc] > R[vc]) C[vc] = uc;
      else if(R[vc] > R[uc]) C[uc] = vc;
     else { C[vc] = uc; R[uc]++; }
 return weight;
int main()
 int wa[6][6] = {
     0, -1, 2, -1, 7, -1 \},
     -1, 0, -1, 2, -1, -1 },
     2, -1, 0, -1, 8, 6,
    \{-1, 2, -1, 0, -1, -1\},\
    \{7, -1, 8, -1, 0, 4\},\
    \{-1, -1, 6, -1, 4, 0\}\};
 vector <vector <int> > w(6, vector <int>(6));
 for(int i=0; i<6; i++)
   for(int j=0; j<6; j++)
     w[i][j] = wa[i][j];
 cout << Kruskal(w) << endl;</pre>
 cin >> wa[0][0];
```

### 1.9 Lowest Common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                         // children[i] contains
    the children of node i
int A[max nodes][log max nodes+1];
                                         // A[i][i] is the 2^i-th
    ancestor of node i, or -1 if that ancestor does not exist
                                         // L[i] is the distance
int L[max_nodes];
    between node i and the root
// floor of the binary logarithm of n
int lb (unsigned int n)
    if(n==0)
        return -1;
    int p = 0;
    if (n >= 1 << 16) \{ n >= 16; p += 16; \}
    if (n >= 1 << 8) \{ n >>= 8; p += 8; \}
    if (n >= 1 << 4) \{ n >>= 4; p += 4; \}
    if (n >= 1 << 2) \{ n >>= 2; p += 2; \}
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    L[i] = 1;
    for(int j = 0; j < children[i].size(); j++)</pre>
        DFS(children[i][j], l+1);
```

```
int LCA (int p, int q)
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the
        same level as q
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1<<i) >= L[q])
            p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if(A[p][i] != -1 && A[p][i] != A[q][i])
            p = A[p][i];
            q = A[q][i];
    return A[p][0];
int main(int argc, char* argv[])
    // read num_nodes, the total number of nodes
    log_num_nodes=lb(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        int p;
        // read p, the parent of node i or -1 if node i is the
        A[i][0] = p;
        if (p != -1)
            children[p].push_back(i);
        else
            root = i;
    // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)</pre>
        for (int i = 0; i < num_nodes; i++)</pre>
            if (A[i][j-1] !=-1)
                A[i][j] = A[A[i][j-1]][j-1];
            else
                A[i][j] = -1;
    // precompute L
    DFS (root, 0);
    return 0;
```

# 2 Geometry

#### 2.1 Basics

```
typedef complex<double> point;
```

```
typedef vector<point> polygon;
#define NEXT(i) (((i) + 1) % n)
struct circle { point p; double r; };
struct line { point p, q; };
using segment = line;
const double eps = 1e-9;
// fix comparations on doubles with this two functions
int sign(double x) { return x < -eps ? -1 : x > eps; }
int dblcmp(double x, double y) { return sign(x - y); }
double dot(point a, point b) { return real(conj(a) * b); }
double cross(point a, point b) { return imag(conj(a) * b); }
double area2(point a, point b, point c) { return cross(b - a, c -
   a); }
int ccw(point a, point b, point c)
        b -= a; c -= a;
        if (cross(b, c) > 0) return +1; // counter clockwise
        if (cross(b, c) < 0) return -1; // clockwise</pre>
        if (dot(b, c) < 0) return +2; // c--a--b on line
        if (dot(b, b) < dot(c, c)) return -2; // a--b--c on line
        return 0;
namespace std
        bool operator<(point a, point b)</pre>
                if (a.real() != b.real())
                        return a.real() < b.real();</pre>
                return a.imag() < b.imag();</pre>
```

## 2.2 Circle

```
Circles
        Tested: AIZU
// circle-circle intersection
vector<point> intersect(circle C, circle D)
        double d = abs(C.p - D.p);
        if (sign(d - C.r - D.r) > 0) return {};
        if (sign(d - abs(C.r - D.r)) < 0) return {}; // too close</pre>
        double a = (C.r*C.r - D.r*D.r + d*d) / (2*d);
        double h = sqrt(C.r*C.r - a*a);
        point v = (D.p - C.p) / d;
        if (sign(h) == 0) return \{C.p + v*a\};
                                                       // touch
                                                       // intersect
        return \{C.p + v*a + point(0,1)*v*h,
                        C.p + v*a - point(0,1)*v*h;
// circle-line intersection
vector<point> intersect(line L, circle C)
        point u = L.p - L.q, v = L.p - C.p;
```

```
double a = dot(u, u), b = dot(u, v), c = dot(v, v) - C.r*C
            r:
        double det = b*b - a*c;
        if (sign(det) < 0) return {};</pre>
                                                // no solution
        if (sign(det) == 0) return {L.p - b/a*u}; // touch
        return \{L.p + (-b + sqrt(det))/a*u,
                        L.p + (-b - sqrt(det))/a*u;
// circle tangents through point
vector<point> tangent(point p, circle C)
        // not tested enough
        double D = abs(p - C.p);
        if (D + eps < C.r) return \{\};
        point t = C.p - p;
        double theta = asin( C.r / D );
        double d = cos(theta) * D;
        t = t / abs(t) * d;
        if ( abs(D - C.r) < eps ) return {p + t};
        point rot( cos(theta), sin(theta) );
        return {p + t * rot, p + t * conj(rot)};
bool incircle (point a, point b, point c, point p)
        a -= p; b -= p; c -= p;
        return norm(a) * cross(b, c)
                        + norm(b) * cross(c, a)
                        + norm(c) * cross(a, b) >= 0;
                        // < : inside, = cocircular, > outside
point three_point_circle(point a, point b, point c)
        point x = 1.0 / conj(b - a), y = 1.0 / conj(c - a);
        return (y - x) / (conj(x) * y - x * conj(y)) + a;
    Get the center of the circles that pass through p0 and p1
    and has ratio r.
    Be careful with epsilon.
vector<point> two_point_ratio_circle(point p0, point p1, double r)
    if (abs(p1 - p0) > 2 * r + eps) // Points are too far.
        return {};
    point pm = (p1 + p0) / 2.01;
    point pv = p1 - p0;
    pv = point(-pv.imag(), pv.real());
    double x1 = p1.real(), y1 = p1.imag();
    double xm = pm.real(), ym = pm.imag();
    double xv = pv.real(), yv = pv.imag();
    double A = (sqr(xv) + sqr(yv));
    double C = sqr(xm - x1) + sqr(ym - y1) - sqr(r);
    double D = sqrt(-4 * A * C);
    double t = D / 2.0 / A;
    if (abs(t) \le eps)
        return {pm};
```

```
return {c1, c2};
/*
       Area of the intersection of a circle with a polygon
        Circle's center lies in (0, 0)
       Polygon must be given counterclockwise
        Tested: LightOJ 1358
        Complexity: O(n)
*/
#define x(_t) (xa + (_t) * a)
#define y(_t) (ya + (_t) * b)
double radian(double xa, double ya, double xb, double yb)
       return atan2(xa * yb - xb * ya, xa * xb + ya * yb);
double part (double xa, double ya, double xb, double yb, double r)
        double 1 = sqrt((xa - xb) * (xa - xb) + (ya - yb) * (ya -
        b * ya;
        double d = 4.0 * (c * c - xa * xa - ya * ya + r * r);
        if (d < eps)
               return radian(xa, ya, xb, yb) * r * r * 0.5;
       else
               d = sart(d) * 0.5;
               double s = -c - d, t = -c + d;
               if (s < 0.0) s = 0.0;
               else if (s > 1) s = 1;
               if (t < 0.0) t = 0.0;
               else if (t > 1) t = 1;
               return (x(s) * y(t) - x(t) * y(s)
                               + (radian(xa, ya, x(s), y(s)))
                               + radian(x(t), y(t), xb, yb)) * r
                                   * r) * 0.5;
double intersection_circle_polygon(const polygon &P, double r)
        double s = 0.0;
       int n = P.size();;
        for (int i = 0; i < n; i++)
               s += part(P[i].real(), P[i].imag(),
                       P[NEXT(i)].real(), P[NEXT(i)].imag(), r);
        return fabs(s);
```

#### 2.3 Centroid

```
point c(0, 0);
double scale = 3.0 * area2(P); // area2 = 2 * polygon_area
for (int i = 0, n = P.size(); i < n; ++i)
{
    int j = NEXT(i);
        c = c + (P[i] + P[j]) * (cross(P[i], P[j]));
}
return c / scale;</pre>
```

#### 2.4 Distance between closest points

```
/*
        Compute distance between closest points.
        Tested: AIZU(judge.u-aizu.ac.jp) CGL.5A
        Complexity: O(n log n)
double closest_pair_points(vector<point> &P)
        auto cmp = [](point a, point b)
                return make_pair(a.imag(), a.real())
                                < make_pair(b.imag(), b.real());
        };
        int n = P.size();
        sort(P.begin(), P.end());
        set<point, decltype(cmp)> S(cmp);
        const double oo = 1e9; // adjust
        double ans = oo;
        for (int i = 0, ptr = 0; i < n; ++i)
                while (ptr < i && abs(P[i].real() - P[ptr].real())</pre>
                        S.erase(P[ptr++]);
                auto lo = S.lower_bound(point(-oo, P[i].imag() -
                    ans - eps));
                auto hi = S.upper_bound(point(-oo, P[i].imag() +
                    ans + eps));
                for (decltype(lo) it = lo; it != hi; ++it)
                        ans = min(ans, abs(P[i] - *it));
                S.insert(P[i]);
        return ans;
```

# 2.5 Area Of Polygon

```
/*
    Tested: AIZU(judge.u-aizu.ac.jp) CGL.3A
    Complexity: O(n)

*/

double area2(const polygon &P)
{
    double A = 0;
    for (int i = 0, n = P.size(); i < n; ++i)</pre>
```

```
A += cross(P[i], P[NEXT(i)]);
return A;
```

#### 2.6 Antipodal points

```
/*
        Antipodal points
        the antipodal point of a point on the surface of a sphere
            is the point which is diametrically opposite to it
        Tested: AIZU(judge.u-aizu.ac.jp) CGL.4B
        Complexity: O(n)
*/
vector<pair<int, int>> antipodal(const polygon &P)
        vector<pair<int, int>> ans;
        int n = P.size();
        if (P.size() == 2)
                ans.push_back({ 0, 1 });
        if (P.size() < 3)
                return ans;
        int q0 = 0;
        while (abs(area2(P[n-1], P[0], P[NEXT(q0)]))
                        > abs(area2(P[n - 1], P[0], P[q0])))
                ++q0;
        for (int q = q0, p = 0; q != 0 && p <= q0; ++p)
                ans.push_back({ p, q });
                while (abs(area2(P[p], P[NEXT(p)], P[NEXT(q)]))
                                > abs(area2(P[p], P[NEXT(p)], P[q
                        q = NEXT(q);
                        if (p != q0 || q != 0)
                                ans.push_back({ p, q });
                        else
                                return ans;
                if (abs(area2(P[p], P[NEXT(p)], P[NEXT(q)]))
                                == abs(area2(P[p], P[NEXT(p)], P[q
                                    ])))
                        if (p != q0 || q != n - 1)
                                ans.push_back({ p, NEXT(q) });
                        else
                                ans.push_back({ NEXT(p), q });
        return ans:
```

# 2.7 Semiplane Intersection

Check wether there is a point in the intersection of several semi-planes. if p lies in the border of some

```
semiplane it is considered to belong to the semiplane.
        Expected Running time: linear
        Tested on Triathlon [Cuban Campament Contest]
bool intersect( vector<line> semiplane ){
        function<bool(line&, point&)> side = [](line &l, point &p) {
                 // IMPORTANT: point p belongs to semiplane defined
                 // iff p it's clockwise respect to segment < l.p,
                    1.q >
                // i.e. (non negative cross product)
                return cross( l.q - l.p, p - l.p ) >= 0;
        };
        function<bool(line&, line&, point&)> crosspoint = [](const
             line &1, const line &m, point &x) {
                double A = cross(1.q - 1.p, m.q - m.p);
                double B = cross(1.\bar{q} - 1.\bar{p}, 1.\bar{q} - m.p);
                if (abs(A) < eps) return false;</pre>
                x = m.p + B / \bar{A} * (m.q - m.p);
                return true;
        };
        int n = (int)semiplane.size();
        random_shuffle( semiplane.begin(), semiplane.end() );
        point cent(0, 1e9);
        for (int i = 0; i < n; ++i) {
                line &S = semiplane[ i ];
                if (side(S, cent)) continue;
                point d = S.q - S.p; d /= abs(d);
                point A = S.p - d * 1e8, B = S.p + d * 1e8;
                for (int j = 0; j < i; ++j) {
                         point x;
                         line &T = semiplane[j];
                         if ( crosspoint(T, S, x) ){
                                 int cnt = 0;
                                 if (!side(T, A)) {
                                          A = x;
                                          cnt++;
                                 if (!side(T, B)){
                                          cnt++;
                                 if (cnt == 2)
                                          return false;
                         else{
                                 if (!side(T, A)) return false;
                 if (imag(B) > imag(A)) swap(A, B);
                cent = \bar{A};
```

```
return true;
```

### 2.8 Point inside a Polygon

#### 2.9 Convex Cut

# 2.10 Polygon width

```
/*
Compute the width of a convex polygon
Tested: LiveArchive 5138
Complexity: O(n)
*/
const int oo = 1e9; // adjust
```

```
double check (int a, int b, int c, int d, const polygon &P)
        for (int i = 0; i < 4 && a != c; ++i)
                if (i == 1) swap(a, b);
                else swap(c, d);
        if (a == c) // a admits a support line parallel to bd
                double A = abs(area2(P[a], P[b], P[d]));
                // double of the triangle area
                double base = abs(P[b] - P[d]);
                // base of the triangle abd
                return A / base;
        return oo;
double polygon_width(const polygon &P)
        if (P.size() < 3)
                return 0;
        auto pairs = antipodal(P);
        double best = oo;
        int n = pairs.size();
        for (int i = 0; i < n; ++i)
                double tmp = check(pairs[i].first, pairs[i].second
                                pairs[NEXT(i)].first, pairs[NEXT(i
                                    )].second, P);
                best = min(best, tmp);
        return best;
```

## 2.11 Sweep Line (Rectangles)

```
function<void(int, int, int, int, int, int) > aux =
                 [&] (int a, int b, int c, int 1, int r, int
                         if ((a = max(a, 1)) >= (b = min(b, r))
                             ))) return;
                         if (a == 1 \&\& b == r) C[k] += c;
                         else
                                 aux(a, b, c, 1, (1+r)/2,
                                      2*k+1);
                                 aux(a, b, c, (1+r)/2, r,
                                      2*k+2);
                         if (C[k]) A[k] = ys[r] - ys[l];
                         else A[k] = A[2*k+1] + A[2*k+2];
                } ;
struct event
        ll x, l, h, c;
// plane sweep
vector<event> es:
for (auto r : rs)
        int l = lower_bound(ys.begin(), ys.end(), r.yl) -
            ys.begin();
        int h = lower_bound(ys.begin(), ys.end(), r.yh) -
            vs.begin();
        es.push_back({ r.xl, l, h, +1 });
        es.push_back({ r.xh, 1, h, -1 });
sort(es.begin(), es.end(), [](event a, event b)
                 {return a.x != b.x ? a.x < b.x : a.c > b.c
11 \text{ area} = 0, \text{ prev} = 0;
for (auto &e : es)
        area += (e.x - prev) * A[0];
        prev = e.x;
        aux(e.1, e.h, e.c, 0, n, 0);
return area;
```

## 2.12 Rectilinear minimum spanning tree

```
/*
    Tested: USACO OPEN08 (Cow Neighborhoods)
    Complexity: O(n log n)
    the rectilinear minimum spanning tree (RMST)
    of a set of n points in the plane is a minimum spanning tree of
        that set,
    where the weight of the edge between each pair of points is the
        rectilinear
    distance between those two points.
*/

typedef long long ll;
typedef complex<ll> point;

ll rectilinear_mst(vector<point> ps)
{
        vector<int> id(ps.size());
        iota(id.begin(), id.end(), 0);
}
```

```
struct edge
        int src, dst;
        11 weight;
};
vector<edge> edges;
for (int s = 0; s < 2; ++s)
        for (int t = 0; t < 2; ++t)
                 sort(id.begin(), id.end(), [&](int i, int
                     j)
                         return real(ps[i] - ps[j]) < imag(</pre>
                             ps[j] - ps[i]);
                 });
                map<11, int> sweep;
                 for (int i : id)
                         for (auto it = sweep.lower_bound(-
                             imag(ps[i]));
                                          it != sweep.end();
                                               sweep.erase(
                                              it++))
                                 int j = it->second;
                                 if (imag(ps[j] - ps[i]) <
                                      real(ps[j] - ps[i]))
                                          break;
                                 ll d = abs(real(ps[i] - ps
                                      [ <del>i</del> ] ) )
                                                   + abs(imag
                                                       (ps[i]
                                                        - ps[
                                                       j]));
                                 edges.push_back({ i, j, d
                                      });
                         sweep[-imag(ps[i])] = i;
                 for (auto &p : ps)
                         p = point(imag(p), real(p));
        for (auto &p : ps)
                p = point(-real(p), imag(p));
11 \cos t = 0;
sort(edges.begin(), edges.end(), [](edge a, edge b)
        return a.weight < b.weight;</pre>
});
union_find uf(ps.size());
for (edge e : edges)
        if (uf.join(e.src, e.dst))
                cost += e.weight;
return cost;
```

#### 2.13 Points 3d

```
const double pi = acos(-1.0);
// Construct a point on a sphere with center on the origin and
    radius R
// TESTED [COJ-1436]
struct vec{
        double x, y, z;
        vec (double x=0, double y=0, double z=0) : x(x), y(y), z(z)
        vec operator+(const vec a) const{
                return vec(x + a.x, y + a.y, z + a.z);
        vec operator-(const vec a) const{
                return vec(x - a.x, y - a.y, z - a.z);
        vec operator*(const double k) const{
                return vec(k * x, k * y, k * z);
        vec operator/(const double k) const{
                return vec(x / k, y / k, z / k);
        vec operator*(const vec a) const{
                return vec(y * a.z - z * a.y, z * a.x - x * a.z, x
                      \star a.y \overline{\phantom{a}} y \star a.x);
        double dot(const vec a) const{
                return x * a.x + y * a.y + z * a.z;
};
ostream & operator << (ostream & os, const vec &p)
        cout << "(" << p.x << ";" << p.y << ";" << p.z << ")" <<
            endl:
        return os:
double abs(vec p)
        return sqrt (p.x * p.x + p.y * p.y + p.z * p.z);
vec from_polar(double lat, double lon, double R)
        lat = lat / 180.0 * pi;
        lon = lon / 180.0 * pi;
        return vec(R * cos(lat) * sin(lon),
                                    R * cos(lat) * cos(lon), R *
                                        sin(lat));
struct plane
        double A, B, C, D;
};
Geodisic distance between points in a sphere
 R is the radius of the sphere
double geodesic_distance(vec p, vec q, double r)
        return r * acos(p * q / r / r);
```

```
// Find the rect of intersection of two planes on the space
// The rect is given parametrical
// TESTED [TIMUS 1239]
void planePlaneIntersection(plane p, plane q)
{
    if (abs(p.C * q.B - q.C * p.B) < eps)
        return; // Planes are parallel

    double mz = (q.A * p.B - p.A * q.B) / (p.C * q.B - q.C * p.B);
    double nz = (q.D * p.B - p.D * q.B) / (p.C * q.B - q.C * p.B);

    double my = (q.A * p.C - p.A * q.C) / (p.B * q.C - p.C * q.B);
    double ny = (q.D * p.C - p.D * q.C) / (p.B * q.C - p.C * q.B);
    // parametric rect: (x, my * x + ny, mz * x * nz)
}</pre>
```

#### 2.14 Convex Hull 3d

```
/*
        Convex Hull 3d implementation.
        Coplanar points case not handled.
        Complexity: O(n^2)
    Tested: opencup 010376(Stars in a Can)
struct face{
        int I[3];
        vec normal;
        int operator[](const int idx) const{
                return I[idx];
};
vector<face> triangulation(vector<vec> &cloud) {
        vector<face> F(4);
        for (int i = 0; i < 4; ++i) {
                for (int i = 0; i < 3; ++i)
                        F[i].I[j] = j + (j >= i);
                F[i].normal = (cloud[F[i][1]] - cloud[F[i][0]]
                    ]) *
                                           ( cloud[ F[i][2] ] -
                                              cloud[ F[i][0] ]);
                if (F[i].normal.dot(cloud[i] - cloud[F[i][0]])
                     > 0){
                        F[i].normal = F[i].normal * -1.;
                        swap(F[i].I[1], F[i].I[2]);
        int n = (int)cloud.size();
        vector<vi> cnt(n, vi(n));
        vector<vi> tmr(n, vi(n));
        auto mark = [&](int u, int v, int i){
                if (u > v) swap(u, v);
                if (tmr[u][v] != i) {
                        tmr[u][v] = i;
                        cnt[u][v] = 0;
                cnt[u][v]++;
```

```
} ;
        auto get = [&](int u, int v){
                if (u > v) swap(u, v);
                return cnt[u][v];
        };
        for (int i = 4; i < n; ++i) {
                vector<face> OF;
                vec x = cloud[i];
                for (auto f : F) {
                        if (f.normal.dot(x - cloud[f[0]]) > 0){
                                mark(f[0], f[1], i);
                                mark(f[1], f[2], i);
                                mark(f[2], f[0], i);
                        else
                                 QF.push_back(f);
                for (auto f : F) {
                        if (f.normal.dot(x - cloud[f[0]]) > 0){
                                 for (int j = 0; j < 3; ++j) {
                                         int a = f[j], b = f[j != 2]
                                              ? + 1 : 0];
                                         if (get(a, b) != 1)
                                             continue;
                                         vec u = cloud[a], v =
                                             cloud[b];
                                         vec ref = cloud[ f[j ? j -
                                             1:2];
                                         face qf = \{i, a, b, vec()\}
                                         qf.normal = (u - x) * (v -
                                         if (qf.normal.dot( ref - x
                                              ) > 0) {
                                                 qf.normal = qf.
                                                     normal * -1.;
                                                 swap(qf.I[1], qf.I
                                                     [2]);
                                         QF.push_back(qf);
                F.swap(QF);
        return F;
        Convex Hull 3d implementation.
        Complexity O(n^4)
        It works ok with coplanar points
        UNTESTED [With the new vec(point3) structure]
struct face{
        int idx[3];
        face(){}
        face(int i, int j, int k) {
                idx[0] = i, idx[1] = j, idx[2] = k;
```

```
int& operator[](int u) { return idx[u]; }
};
vector<face> convex hull( vector<vec> &cloud ){
        int n = (int)cloud.size();
        vector<int> L(n);
        vector<face> faces;
        for (int i = 0; i < n; ++i)
                for (int j = i + 1; j < n; ++j)
                        for (int k = j + 1; k < n; ++k) {
                                vec a = cloud[i], b = cloud[j], c
                                    = cloud[k];
                                vec nr = (b - a) * (c - a);
                                int pnt = 0;
                                L[pnt++] = j;
                                L[pnt++]=k;
                                bool proc = true;
                                int v = 0, V = 0;
                                for (int 1 = 0; 1 < n && proc; ++1
                                    ) {
                                        if (l == i || l == j || l
                                             == k) continue;
                                        double t = nr.dot( cloud[]
                                            ] - a );
                                        if (abs(t) < eps) {
                                                 if (1 < k) proc =
                                                    false;
                                                 else L[ pnt++ ] =
                                                    1;
                                        else{
                                                 if (t < 0) v = -1;
                                                 else
                                                           V = +1;
                                if (!proc | | v * V == -1) continue
                                function<bool(int,int)> compare =
                                     [&] (int u, int v) {
                                        return nr.dot((cloud[u] -
                                             a) * (cloud[v] - a))
                                             > 0;
                                };
                                sort(L, L + pnt, compare);
                                for (int 1 = 0; 1 + 1 < pnt; ++1)
                                        faces.push_back( face(i, L
                                             [1], L[1 + 1]);
        return faces;
// Find mass center of a polyhedron given the external faces
void mass_center( vector<vec> &cloud, vector<face> &faces ) {
        vec pivot = cloud[0];
```

```
double x = 0, y = 0, z = 0, v = 0;
for (auto f : faces) {
        auto value =
                        ( cloud[f[0]] - pivot ).dot(
                                         ( cloud[f[1]] -
                                             pivot) * (
                                             cloud[f[2]] -
                                             pivot)
                                         );
        vec sum = cloud[f[0]] + cloud[f[1]] + cloud[f[2]]
            + pivot;
        double cvol = abs( 1. * value / 6 );
        v += cvol;
        cvol /= 4;
        x += cvol * sum.X[0];
        y += cvol * sum.X[1];
        z += cvol * sum.X[2];
x /= v, y /= v, z /= v;
// Mass center of a polyhedron at (x, y, z)
```

#### 2.15 intersections

```
Line and segments predicates
        Tested: AIZU(judge.u-aizu.ac.jp) CGL
*/
bool intersectLL(const line &1, const line &m)
        return abs(cross(l.q - l.p, m.q - m.p)) > eps | | // non-
            parallel
                        abs(cross(l.q - l.p, m.p - l.p)) < eps;
                            // same line
bool intersectLS(const line &1, const segment &s)
        return cross(l.q - l.p, s.p - l.p) \star // s[0] is left
                        cross(1.q - 1.p, s.q - 1.p) < eps; // s[1]
                             is right of l
bool intersectLP (const line &1, const point &p)
        return abs(cross(l.q - p, l.p - p)) < eps;</pre>
bool intersectSS(const segment &s, const segment &t)
        return ccw(s.p, s.q, t.p) * ccw(s.p, s.q, t.q) <= 0
                        && ccw(t.p, t.q, s.p) * ccw(t.p, t.q, s.q)
                             <= 0;
bool intersectSP(const segment &s, const point &p)
        return abs(s.p - p) + abs(s.q - p) - abs(s.q - s.p) < eps;
        // triangle inequality
```

return min(real(s.p), real(s.q)) <= real(p)</pre>

```
&& real(p) \leq max(real(s.p), real(s.q))
                        && min(imag(s.p), imag(s.q)) \le imag(p)
                        && imag(p) \le max(imag(s.p), imag(s.q))
                        && cross(s.p - p, s.q - p) == 0;
point projection (const line &1, const point &p)
        double t = dot(p - 1.p, 1.p - 1.q) / norm(1.p - 1.q);
        return 1.p + t * (1.p - 1.q);
point reflection (const line &1, const point &p)
        return p + 2.0 * (projection(1, p) - p);
double distanceLP(const line &1, const point &p)
        return abs(p - projection(l, p));
double distanceLL(const line &1, const line &m)
        return intersectLL(1, m) ? 0 : distanceLP(1, m.p);
double distanceLS(const line &1, const line &s)
        if (intersectLS(1, s)) return 0;
        return min(distanceLP(l, s.p), distanceLP(l, s.q));
double distanceSP(const segment &s, const point &p)
        const point r = projection(s, p);
        if (intersectSP(s, r)) return abs(r - p);
        return min(abs(s.p - p), abs(s.q - p));
double distanceSS(const segment &s, const segment &t)
        if (intersectSS(s, t)) return 0;
        return min(min(distanceSP(s, t.p), distanceSP(s, t.q)),
                        min(distanceSP(t, s.p), distanceSP(t, s.q)
                            ));
point crosspoint (const line &1, const line &m)
        double A = cross(l.q - l.p, m.q - m.p);
        double B = cross(l.q - l.p, l.q - m.p);
        if (abs(A) < eps && abs(B) < eps)
                return m.p; // same line
        if (abs(A) < eps)
                assert(false); // !!!PRECONDITION NOT SATISFIED!!!
        return m.p + B / A * (m.q - m.p);
```

#### 2.16 minimal circle

/\*
Find circle with minimal ratio that contain a cloud of points in 2d.

Complexity O(n) (Expected running time)
Tested: opencup 010376(Stars in a Can)

```
Note: Something went wrong when random_shuffle was applied.
    If random_shuffle is removed complexity might increase
    to O(n^3)
struct circle{
    point center;
    double r;
    bool contain(point &p) {
        return abs(center - p) < r + eps;</pre>
};
circle min_circle(vector<point> &cloud, int a, int b) {
    point center = (cloud[a] + cloud[b]) / double(2.);
    double rat = abs(center - cloud[a]);
    circle C = {center, rat};
    for (int i = 0; i < b; ++i) {
        point x = cloud[i];
        if (C.contain(x)) continue;
        center = three_point_circle( cloud[a], cloud[b], cloud[i]
        rat = abs(center - cloud[a]);
        C = {center, rat};
    return C;
circle min_circle(vector<point> &cloud, int a) {
    point center = (cloud[a] + cloud[0]) / double(2.);
    double rat = abs(center - cloud[a]);
    circle C = {center, rat};
    for (int i = 0; i < a; ++i) {
        point x = cloud[i];
        if (C.contain(x)) continue;
        C = min_circle(cloud, a, i);
    return C;
circle min circle(vector<point> cloud) {
    // random_shuffle(cloud.begin(), cloud.end());
    int n = (int)cloud.size();
    for (int i = 1; i < n; ++i) {
        int u = rand() % i;
        swap(cloud[u], cloud[i]);
    point center = (cloud[0] + cloud[1]) / double(2.);
    double rat = abs(center - cloud[0]);
    circle C = {center, rat};
    for (int i = 2; i < n; ++i) {
        point x = cloud[i];
        if (C.contain(x)) continue;
        C = min_circle(cloud, i);
    return C;
```

# 2.17 minkowski Sum of Polygons

```
Minkowski sum of two convex polygons. O(n + m)
    Note: Polygons MUST be counterclockwise
polygon minkowski (polygon &A, polygon &B) {
        int na = (int)A.size(), nb = (int)B.size();
        if (A.empty() || B.empty()) return polygon();
        rotate(A.begin(), min_element(A.begin(), A.end()), A.end()
        rotate(B.begin(), min_element(B.begin(), B.end()), B.end()
           );
        int pa = 0, pb = 0;
        polygon M;
        while (pa < na && pb < nb) {</pre>
                M.push_back(A[pa] + B[pb]);
                double x = cross(A[(pa + 1) % na] - A[pa],
                                                   B[(pb + 1) % nb]
                                                       - B[pb]);
                if (x <= eps) pb++;
                if (-eps <= x) pa++;</pre>
        while (pa < na) M.push_back(A[pa++] + B[0]);</pre>
        while (pb < nb) M.push_back(B[pb++] + A[0]);
        return M;
```

### 3 Data Structures

## 3.1 Mo's algorithm

```
#include<bits/stdc++.h>
#define TAM 30000 + 7
#define QTAM 200000 + 7
#define MTAM 1000000 + 7
#define whatis(x) cerr<<#x<<" is "<<x<<endl
using namespace std;
int a[TAM], r[QTAM], cnt[MTAM];
int ans, BLOCK, currL , currR ;
struct node{
    int L, R, idx;
}q[QTAM];
bool comp(node a, node b) {
    if(a.L/BLOCK < b.L/BLOCK) return true;</pre>
    if(a.L/BLOCK > b.L/BLOCK) return false;
    return a.R < b.R;</pre>
void remove(int i) {
    cnt[a[i]]--;
    if(cnt[a[i]] == 0)ans--;
void add(int i) {
    cnt[a[i]]++;
    if(cnt[a[i]] == 1) ans++;
```

```
int query(node i) {
    while(currL< i.L){</pre>
         remove (currL);
         currL++;
    while(currL > i.L) {
        currL--;
         add(currL);
    while(currR< i.R) {</pre>
        currR++;
        add(currR);
    while(currR > i.R) {
        remove (currR);
         currR--;
    return ans;
int main(){
    ios_base::sync_with_stdio(0);
    cin.tie(0);
    #ifdef LOCAL
    freopen("in", "r", stdin);
    #endif
    int n, que;
    cin>>n;
    BLOCK = sqrt(n);
    for (int i = 1; i <= n; i++)cin>>a[i];
    cin>>que;
    for(int i = 1; i <= que; i++) {</pre>
         cin>>q[i].L>>q[i].R;
        q[i].idx = i;
    sort(q +1, q + que+1, comp);
    for(int i = 1; i <=que; i++)</pre>
         r[q[i].idx] = query(q[i]);
    for (int i = 1; i \le que; i++)
         cout << r[i] << "\n";
```

# 3.2 Segment Trees with lazy propagation

```
//querys and build takes O(log n)
//example with segment sum
//tested in AIZU online Judge
#include<bits/stdc++.h>

using namespace std;
long long *p;
struct SegmentTree{
    SegmentTree *L, *R;
    long long sum = 0;
    long long lazy = 0;
    int l, r;

void update(int a, int val) {
    if(l == r) {
        sum += val;
    }
}
```

```
return;
        int mid = (1 + r)/2;
        if(1 <= a && a<= mid)
             L->update(a, val);
            R->update(a, val);
        sum = L -> sum + R -> sum;
    void updateRange(int a, int b, long long val) {
        if(lazy != 0) {
             sum += (r-l+1)*lazy;
             if(1 != r) {
                 R->lazv = lazv + R->lazv;
                 L \rightarrow lazy = lazy + L \rightarrow lazy;
             lazy = 0;
        if(b < 1 or a > r)
            return;
        if(1 >= a \&\& r <= b) {
             sum += (r-l+1)*val;
            if(1 != r) {
                 R->lazy = val + R->lazy;
                 L->lazy = val + L->lazy;
            return;
        L->updateRange(a, b, val);
        R->updateRange(a,b,val);
        sum = L -> sum + R -> sum;
    long long query(int a, int b) {
        if(b < 1 or a > r)
             return 0:
        if(lazy != 0) {
             sum += (r-l+1)*lazy;
             if(1 != r) {
                 R->lazy = lazy + R->lazy;
                 L \rightarrow lazy = lazy + L \rightarrow lazy;
             lazv = 0;
        //this section can be used in non lazy segment tree
        if (a == 1 && b == r) return sum;
        if(b <= L->r) return L->query(a,b);
        if(a >= R->1) return R->query(a,b);
        return (L->query(a,L->r) + R->query(R->1, b));
    SegmentTree(int a, int b): 1(a), r(b){
        if(a == b){
             sum = p[a];
             L = R = nullptr;
        else{
            L = new SegmentTree (a, (a+b)/2);
            R = new SegmentTree ( (a+b)/2 + 1, b );
             sum = L -> sum + R -> sum;
int main(){
    cin.tie(0);
```

};

```
ios_base::sync_with_stdio(0);
#ifdef LOCAL
    freopen("input.txt", "r", stdin);
#endif // LOCAL
    long long T;
    cin >> T;
while (T--) {
    long long n, c;
    cin >> n >> c;
    long long l[n];
    memset (1, 0, sizeof(1));
    SegmentTree *stree = new SegmentTree(0, n-1);
    while (c--) {
        long long aux, p, q;
        cin >> aux >> p >> q;
        if(aux == 0){
            long long val;
            cin >> val;
            stree->updateRange(p-1, q-1, val);
            cout << stree->query(p-1, q-1) << endl;</pre>
```

## Segment Trees with lazy propagation (Java)

```
public class SegmentTreeRangeUpdate {
        public long[] leaf;
        public long[] update;
        public int origSize;
        public SegmentTreeRangeUpdate(int[] list)
                origSize = list.length;
                leaf = new long[4*list.length];
                update = new long[4*list.length];
                build(1,0,list.length-1,list);
        public void build(int curr, int begin, int end, int[] list
                if(begin == end)
                        leaf[curr] = list[begin];
                else
                        int mid = (begin+end)/2;
                        build(2 * curr, begin, mid, list);
                        build(2 * curr + 1, mid+1, end, list);
                        leaf[curr] = leaf[2*curr] + leaf[2*curr
                            +1];
        public void update(int begin, int end, int val) {
                update(1,0,origSize-1,begin,end,val);
        public void update(int curr, int tBegin, int tEnd, int
            begin, int end, int val)
                if(tBegin >= begin && tEnd <= end)</pre>
                        update[curr] += val;
                else
                        leaf[curr] += (Math.min(end,tEnd)-Math.max
                            (begin, tBegin) +1) * val;
                        int mid = (tBegin+tEnd)/2;
                        if(mid >= begin && tBegin <= end)</pre>
                                 update(2*curr, tBegin, mid, begin,
```

```
end, val);
                 if(tEnd >= begin && mid+1 <= end)</pre>
                         update(2*curr+1, mid+1, tEnd,
                             begin, end, val);
public long query(int begin, int end) {
        return query(1,0,origSize-1,begin,end);
public long query (int curr, int tBegin, int tEnd, int
    begin, int end) {
        if(tBegin >= begin && tEnd <= end)</pre>
                 if(update[curr] != 0) {
                         leaf[curr] += (tEnd-tBegin+1) *
                             update[curr];
                         if(2*curr < update.length) {</pre>
                                 update[2*curr] += update[
                                 update[2*curr+1] += update
                                      [curr];
                         update[curr] = 0;
                 return leaf[curr];
        else
                 leaf[curr] += (tEnd-tBegin+1) * update[
                     curr];
                 if(2*curr < update.length){</pre>
                         update[2*curr] += update[curr];
                         update[2*curr+1] += update[curr];
                 update[curr] = 0;
                 int mid = (tBegin+tEnd)/2;
                 long ret = 0;
                 if (mid >= begin && tBegin <= end)</pre>
                         ret += query(2*curr, tBegin, mid,
                             begin, end);
                 if(tEnd >= begin && mid+1 <= end)</pre>
                         ret += query(2*curr+1, mid+1, tEnd
                             , begin, end);
                 return ret;
```

#### 3.4 Fenwick Tree

```
/* Fenwick tree or Binary Indexed Tree
* query time O(logN)
* update time O(logN)
* LOGSZ must be higher than log(SIZE)
* operations must have inverse

* tested on AIZU online Judge
*/
#include<bits/stdc++.h>
using namespace std;
const int LOGSZ = 17;
int tree[(1<<LOGSZ) + 1];
int n = (1<<LOGSZ);
void add(int x, int v){</pre>
```

```
while (x<=n) {
    tree[x] += v;
    x+= (x & -x);
}

int get(int x) {
    int ans = 0;
    while(x) {
        ans += tree[x];
        x-= (x & -x);
}
    return ans;
}

int rsq(int x, int y) {
    return get(y) - get(x - 1);
}</pre>
```

#### 3.5 Union Find

```
* Disjoint set data structure, merge and find takes O(logN)
 * tested on AIZU online Judge
#include <iostream>
#include <vector>
using namespace std;
int find(vector<int> &C, int x) { return (C[x] == x) ? x : C[x] =
    find(C, C[x]);
bool same(vector<int> &C, int x, int y) { return find(C, x) == find
void merge(vector<int> &C, int x, int y) { C[find(C, x)] = find(C,
    y); }
int main()
        int n = 5;
        vector<int> C(n);
        for (int i = 0; i < n; i++) C[i] = i;
        merge(C, 0, 2);
        merge(C, 1, 0);
        merge(C, 3, 4);
        for (int i = 0; i < n; i++) cout << i << " " << find(C, i)</pre>
             << endl:
        return 0;
```

# 3.6 Persistent Treaps

```
/*
  * implicit persistent treaps, solved Genetics problem
  */
#include <bits/stdc++.h>
using namespace std;
typedef long long int64;
typedef pair<int,int> pii;
typedef vector<int> vi;
const int oo = 0x3f3f3f3f3f;
const double eps = 1e-9;
```

```
struct node{
    int64 count[4], size;
    int idx:
    node *1, *r;
};
string dna = "AGTC";
map<char,int> dnaid;
int num(char c){
    return dnaid[ c ];
node* clone(node *u){
    node *v = new node();
    v\rightarrow size = u\rightarrow size, v\rightarrow idx = u\rightarrow idx;
    v->1 = u->1, v->r = u->r;
    for (int i = 0; i < 4; ++i)
        v->count[i] = u->count[i];
    return v;
int64 size(node *u){
    return u ? u->size : OLL;
int get(node *u, int p) {
    return u ? u->count[p] : 0;
node *update(node *u) {
    u \rightarrow size = size(u \rightarrow 1) + size(u \rightarrow r) + 1;
    for (int i = 0; i < 4; ++i)
        u - count[i] = get(u - l, i) + get(u - r, i);
    u->count[ u->idx ]++;
    return u;
node* merge( node *u, node *v ) {
    if (!u || !v) return u ? u : v;
    int l = size(u), r = size(v);
    if (rand() % (l + r) < l)
        u = clone(u);
         u->r = merge(u->r, v);
        return update( u );
    else{
         v = clone(v);
        v->1 = merge(u, v->1);
        return update( v );
pair<node*, node*> split (node *u, int k) {
    if (!u) return make_pair(nullptr, nullptr);
    u = clone(u);
    if (k <= size( u->1 )) {
         auto cur = split( u->1, k );
         u->1 = cur.second;
        return {cur.first, update(u) };
    else{
         auto cur = split( u \rightarrow r, k - size(u \rightarrow l) - 1);
        u->r = cur.first;
         return {update(u), cur.second};
```

```
node* build(int b, int e, string &s) {
    if (b > e) return nullptr;
    int m = (b + e) >> 1;
    node *u = new node();
    u->idx = num(s[m]);
    u->1 = build(b, m-1, s);
    u->r = build(m + 1, e, s);
    return update(u);
node *mutate(node *u, int k, int v) {
    u = clone(u);
    if (k <= size(u->1))
        u->1 = mutate(u->1, k, v);
    else if (size(u->1) + 1 == k)
        u->idx = v:
    else
        u\rightarrow r = mutate(u\rightarrow r, k - size(u\rightarrow l) - 1, v);
    return update(u);
void dfs( node* u) {
    if (!u) return;
    dfs(u->1);
    cout << dna[ u->idx ] << endl;</pre>
    for (int i = 0; i < 4; ++i)
        cout << u->count[i] << " ";
    cout << endl;
    dfs(u->r);
int main(){
    ios_base::sync_with_stdio(0);
    cin.tie(0);
    for (int i = 0; i < 4; ++i)
        dnaid[ dna[i] ] = i;
    int n; cin >> n;
    vector<node*> version(1);
    for (int i = 1; i <= n; ++i) {</pre>
        string s; cin >> s;
        version.push_back( build(0, s.length() - 1, s ) );
    int q; cin >> q;
    while (q--) {
        string comm;
        cin >> comm;
        if (comm == "MUTATE") {
            int id, k; char m;
            cin >> id >> k >> m;
            version[id] = mutate( version[id], k, num(m) );
        else if (comm == "CROSS") {
```

```
int id1, id2, k1, k2;
        cin >> id1 >> id2 >> k1 >> k2;
        auto dna1 = split( version[id1], k1 );
        auto dna2 = split( version[id2], k2 );
       node* dna3 = merge( dna1.first, dna2.second );
       node* dna4 = merge( dna2.first, dna1.second );
       version.push_back( dna3 );
       version.push_back( dna4 );
    else{
        int id, k1, k2;
       cin >> id >> k1 >> k2;
       node *u = split(version[id], k2).first;
             u = split(u, k1 - 1).second;
        for (int i = 0; i < 4; ++i)
            cout << u->count[i] << " \n"[i + 1 == 4];
return 0;
```

# 4 Strings

#### 4.1 Prefix Function

## 4.2 KMP Automata

## 4.3 Suffix Array

```
#include <bits/stdc++.h>
using namespace std;
vector<int> suffix_array(const string &s) {
    int n = s.size();
    vector<int> sa(n), rank(n);
    vector<long long> rank2(n);
    for (int i = 0; i < n; i++) {
        sa[i] = i;
        rank[i] = s[i];
    for (int len = 1; len < n; len \star= 2) {
        for (int i = 0; i < n; i++) rank2[i] = ( (long long) rank[</pre>
            i] << 32) + (i+len < n ? rank[i+len] : -1);
        sort(sa.begin(), sa.end(), [&](int i, int j){
            return rank2[i] < rank2[j];</pre>
        for (int i = 0; i < n; i++) {</pre>
            if (i > 0 && rank2[sa[i]] == rank2[sa[i-1]]) rank[sa[i
                | | = rank[sa[i-1]];
            else rank[sa[i]] = i;
    return sa;
vector<int> lcp_array(const vector<int> &sa, const string &s) {
    int n = s.size();
    vector<int> rank(n);
    for (int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
    vector<int> ans(n);
    for (int i = 0, 1 = 0; i < n; i++) if (rank[i] > 0) {
        int j = sa[rank[i]-1];
        while (s[i+1] == s[j+1]) 1++;
        ans[rank[i]] = 1 > 0 ? 1-- : 1;
    return ans;
int main() {
    string s = "banana";
    s += char(0);
    auto sa = suffix_array(s);
    for (auto x : sa) cerr << x << ' '; cerr << endl;</pre>
    auto lcp = lcp array(sa,s);
    for (auto x : lcp) cerr << x << ' '; cerr << endl;</pre>
    return 0;
```

#### 4.4 Trie

```
// K = tama o del alfabeto
// NMAX = m ximo n mero de nodos que tendr el trie
struct vertex {
    int next[K];
    bool leaf;
};
```

#### 4.5 Hash

```
#include <bits/stdc++.h>
using namespace std;
#define forn(i,n) for(int i=0; i<(int)(n); i++)
#define si(c) ((int)(c).size())
#define forsn(i,s,n) for(int i = (int)(s); i < ((int)n); i++)
#define dforsn(i,s,n) for(int i = (int)(n)-1; i \ge ((int)s); i--)
#define all(c) (c).begin(), (c).end()
#define D(a) cerr << #a << "=" << a << endl;
#define pb push_back
#define eb emplace_back
#define mp make_pair
typedef long long int 11;
typedef vector<int> vi;
typedef pair<int,int> pii;
mt19937 rng;
#define hash nico hash
struct hashing {
    int mod, mul;
    bool prime(int n) {
        for (int d = 2; d*d \le n; d++) if (n*d == 0) return false;
        return true;
    void setValues(int mod, int mul) {
        this->mod = mod;
        this->mul = mul;
    void randomize() {
        rnq.seed(time(0));
        mod = uniform_int_distribution<>(0, (int) 5e8)(rng) + 1e9;
        while (!prime(mod)) mod++;
        mul = uniform_int_distribution<>(2, mod-2)(rng);
    vi h, pot;
    void process(const string &s) {
        h.resize(si(s)+1);
        pot.resize(si(s)+1);
        h[0] = 0; forn(i, si(s)) h[i+1] = (((ll)h[i] * mul) + s[i])
             % mod;
```

## 5 Flows

#### 5.1 Ford Fulkerson

```
#include<bits/stdc++.h>
using namespace std;
///---- Ford-Fulkerson O(MaxFlow * |E|)
//tested on AIZU online Judge
struct OutEdge {
       int to, cap, rIdx;
        OutEdge ( ) { }
        OutEdge(int to, int cap, int rIdx) :
               to(to), cap(cap), rIdx(rIdx) { }
};
struct Network{
       vector<vector<OutEdge> > out;
       vector<bool> seen;
        int sink:
        int augment ( int i, const int cur ) {
                if ( i == sink ) return cur;
                if ( seen[i] ) return false;
                seen[i] = true;
                int ans:
                for ( OutEdge& e : out[i] )
                        if ( e.cap > 0 && ( ans = augment(e.to,min
                            (cur,e.cap)) )) {
    e.cap -= ans;
                                out[e.to][e.rIdx].cap += ans;
                                return ans;
                return 0;
        int maxflow ( int source, int _sink ) {
                sink = \_sink;
                int curflow = 0, aug;
                while ( true ) {
                        fill ( seen.begin(), seen.end(), false );
                        aug = augment(source, INT_MAX);
```

### 5.2 Edmons Karp Min cut

```
///----O(|V|*|E|^2) -----///
struct edge {
       int v, cap, inv;
       edge() {}
       edge(int v, int cap, int inv) : v(v), cap(cap), inv(inv)
struct edmons_karp {
       vector< vector<edge> > g;
       vector<int> from;
       vector<bool> color;
       int n;
       edmons_karp(int n) : n(n), g(n), color(n), from(n) {}
        // Call flood (source) to color one node
        // component of min cut.
       void flood(int u) {
               if (color[u]) return;
               color[u] = true;
               for(int i = 0; i < q[u].size(); i++) {</pre>
           edge &e = q[u][i];
                       if (e.cap > 0)
                flood(e.v);
        int max_flow(int s, int t) {
               int res = 0;
               while(1) {
                        fill(from.begin(), from.end(), -1);
                       queue<int> q;
                       q.push(s);
                        from[s] = -2;
                        for(int u; q.size(); q.pop()) {
                               u = q.front();
                               for(int i = 0; i < q[u].size(); i</pre>
                    edge &e = q[u][i];
                                        if(from[e.v] == -1 \&\& e.
                                            cap) {
                                                from[e.v] = e.inv;
                                                q.push(e.v);
                        if (from[t] == -1) break;
```

```
int aug = INT_MAX;
                for (int i = t, j; i != s; i = j) {
                        j = g[i][ from[i] ].v;
                        aug = min(aug, g[j][g[i][from[i]]
                             ].inv ].cap);
                for (int i = t, j; i != s; i = j) {
                        j = g[i][from[i]].v;
                        g[j][ g[i][from[i] ].inv ].cap -=
    aug;
                        q[i][from[i]].cap += auq;
                res += aug;
        return res:
void add_non_dir_edge(int a, int b, int c) {
        g[a].push_back(edge(b, c, g[b].size()));
        g[b].push_back(edge(a, c, g[a].size() - 1));
void add_edge(int a, int b, int c) {
        g[a].push_back(edge(b, c, g[b].size()));
        g[b].push_back(edge(a, 0, g[a].size() - 1));
```

## 5.3 Dinic's Blocking Flow

};

```
// Adjacency list implementation of Dinic's blocking flow
// This is very fast in practice, and only loses to push-relabel
// Running time:
      O(|V|^2 |E|)
// INPUT:
      - graph, constructed using AddEdge()
       - source
      - sink
// OUTPUT:
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
//tested on AIZU online Judge
#include <bits/stdc++.h>
using namespace std;
const int INF = 2000000000;
struct Edge {
 int from, to, cap, flow, index;
  Edge(int from, int to, int cap, int flow, int index) :
    from(from), to(to), cap(cap), flow(flow), index(index) {}
} ;
struct Dinic {
  int N;
  vector<vector<Edge> > G;
 vector<Edge *> dad;
 vector<int> Q;
 Dinic(int N) : N(N), G(N), dad(N), Q(N) {}
```

```
void AddEdge(int from, int to, int cap) {
   G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
   G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
 long long BlockingFlow(int s, int t) {
    fill(dad.begin(), dad.end(), (Edge *) NULL);
    dad[s] = &G[0][0] - 1;
    int head = 0, tail = 0;
    Q[tail++] = s;
    while (head < tail) {</pre>
     int x = Q[head++];
      for (int i = 0; i < G[x].size(); i++) {</pre>
        Edge &e = G[x][i];
        if (!dad[e.to] && e.cap - e.flow > 0) {
          dad[e.to] = &G[x][i];
          Q[tail++] = e.to;
    if (!dad[t]) return 0;
    long long totflow = 0;
    for (int i = 0; i < G[t].size(); i++) {</pre>
      Edge *start = &G[G[t][i].to][G[t][i].index];
      int amt = INF;
      for (Edge *e = start; amt && e != dad[s]; e = dad[e->from])
        if (!e) { amt = 0; break; }
        amt = min(amt, e->cap - e->flow);
      if (amt == 0) continue;
      for (Edge *e = start; amt && e != dad[s]; e = dad[e->from])
        e->flow += amt;
        G[e->to][e->index].flow -= amt;
      totflow += amt;
   return totflow;
 long long GetMaxFlow(int s, int t) {
   long long totflow = 0;
    while (long long flow = BlockingFlow(s, t))
      totflow += flow;
    return totflow;
};
```

#### 5.4 Push Relabel

```
// Adjacency list implementation of FIFO push relabel maximum flow
// with the gap relabeling heuristic. This implementation is
// significantly faster than straight Ford-Fulkerson. It solves
// random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
//
// Running time:
// O(|V|^3)
//
// INPUT:
```

```
- graph, constructed using AddEdge()
       - source
       - sink
// OUTPUT:
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
//tested on AIZU online Judge
#include<bits/stdc++.h>
using namespace std;
typedef long long LL;
struct Edge {
 int from, to, cap, flow, index;
 Edge(int from, int to, int cap, int flow, int index) :
    from (from), to (to), cap (cap), flow (flow), index (index) {}
};
struct PushRelabel {
  int N;
  vector<vector<Edge> > G;
  vector<LL> excess;
  vector<int> dist, active, count;
  queue<int> Q;
  PushRelabel(int N): N(N), G(N), excess(N), dist(N), active(N),
      count (2*N) {}
  void AddEdge(int from, int to, int cap) {
    G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push\_back(Edge(to, from, 0, 0, G[from].size() - 1));
  void Enqueue(int v) {
    if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v)
        ; }
  void Push (Edge &e) {
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
    if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
    e.flow += amt;
    G[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
    excess[e.from] -= amt;
    Enqueue (e.to);
  void Gap(int k) {
    for (int v = 0; v < N; v++) {
      if (dist[v] < k) continue;</pre>
      count[dist[v]]--;
      dist[v] = max(dist[v], N+1);
      count[dist[v]]++;
      Enqueue (v);
  void Relabel(int v) {
    count[dist[v]]--;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)
      if (G[v][i].cap - G[v][i].flow > 0)
```

```
dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue (v);
  void Discharge(int v) {
    for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[</pre>
    if (excess[v] > 0) {
      if (count[dist[v]] == 1)
        Gap(dist[v]);
      else
        Relabel(v);
  LL GetMaxFlow(int s, int t) {
    count[0] = N-1;
    count[N] = 1;
    dist[s] = N;
    active[s] = active[t] = true;
    for (int i = 0; i < G[s].size(); i++) {</pre>
      excess[s] += G[s][i].cap;
      Push (G[s][i]);
    while (!Q.empty()) {
      int v = Q.front();
      Q.pop();
      active[v] = false;
      Discharge(v);
    LL totflow = 0;
    for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;
    return totflow;
} ;
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  #ifdef Larra
  freopen("in","r",stdin);
  freopen("out", "w", stdout);
  #endif
  int n, m;
  cin>>n>>m;
  PushRelabel netw(n);
  for (int i = 0; i < m; i++) {
    int a, b, c;
    cin>>a>>b>>c;
   netw.AddEdge(a, b, c);
  cout << netw. GetMaxFlow(0, n-1) << "\n";
```

## 5.5 Hopcroft karp's maximum bipartite matching

```
struct hopcroft_karp {
    int 1, r;
    vector<int> last, prev, head, matching, d;
    vector<bool> used, vis;
    hopcroft_karp(int l, int r) : l(l), r(r), last(l, -1),
        matching (r, -1), d(l), used(l), vis(l) {}
    /// u -> 0 to 1, v -> 0 to r
    void add_edge(int u, int v) {
        head.push_back(v);
        prev.push_back(last[u]);
        last[u] = prev.size()-1;
    void bfs() {
        fill(d.begin(), d.end(), -1);
        queue<int> q;
        for (int u = 0; u < 1; u++) {
            if(!used[u]) {
                q.push(u);
                d[u] = 0;
        while(q.size()) {
            int u = q.front(); q.pop();
            for(int e = last[u]; e >= 0; e = prev[e]) {
                int v = matching[ head[e] ];
                if(v) = 0 & d[v] < 0)
                    d[v] = d[u] + 1;
                    q.push(v);
    bool dfs(int u) {
        vis[u] = true;
        for(int e = last[u]; e >= 0; e = prev[e]) {
            int v = head[e];
            int k = matching[v];
            if(k < 0 \mid | (!vis[k] \&\& d[k] == d[u]+1 \&\& dfs(k)))  {
                matching[v] = u;
                used[u] = true;
                return true;
        return false;
    int max_matching() {
        int ans = 0:
        while(true)
            bfs();
            fill(vis.begin(), vis.end(), false);
            int f = 0;
            for (int u = 0; u < 1; u++)
                if(!used[u] && dfs(u)) f++;
            if(f == 0) return ans;
            ans += f;
        return 0;
};
int main(){
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  #ifdef Larra
  freopen("in", "r", stdin);
```

```
freopen("out","w",stdout);
#endif
int n1, n2, m;
cin>>n1>>n2>>m;
hopcroft_karp hp(n1, n2);
for(int i = 0; i < m; i++) {
   int u,v;
   cin>>u>>v;
   hp.add_edge(u, v);
}
cout<<hp.max_matching()<<"\n";</pre>
```

#include <bits/stdc++.h>

#### 5.6 Hungarian's maximum bipartite matching

```
using namespace std;
///----O(|V|*|E|) -----///
//tested in AIZU online Judge
struct hungarian_bipartite {
   int 1, r;
    vector< vector<int> > g;
    vector<bool> seen;
    vector<int> match; /// match[i] is the left matching of right
   hungarian_bipartite(int 1, int r) : 1(1), r(r), seen(1), match
        (1+r, -1), g(1+r) \{ \}
    /// [0-1) left, [1, r) right
    /// a and b are 0-indexed
    void add_edge(int a, int b) {
        g[a].push_back(l+b);
        g[l+b].push_back(a);
   bool go(int u) {
        if(seen[u]) return false;
        seen[u] = true;
        for (int i = 0; i < q[u].size(); i++) {
            int v = q[u][i];
            if(match[v] == -1 \mid | go(match[v]))  {
               match[v] = u;
                return true;
        return false;
    int max_matching() {
        int ans = 0:
        for (int i = 0; i < 1; i++) {
           fill(seen.begin(), seen.end(), false);
           ans += go(i);
        return ans;
} ;
int main(){
 ios_base::sync_with_stdio(0);
 cin.tie(0);
  #ifdef Larra
  freopen("in", "r", stdin);
  freopen("out", "w", stdout);
  #endif
 int n1, n2, m;
```

```
cin>>n1>>n2>>m;
hungarian_bipartite hp(n1, n2);
for(int i = 0; i < m; i++) {
    int u,v;
    cin>>u>>v;
    hp.add_edge(u, v);
}
cout<<hp.max_matching()<<"\n";
}</pre>
```

### 6 Math

## 6.1 general math tricks

```
long square(long n) { return n*n; }
 int fastPow(long x, long n) {
     if(n == 0)
        return 1;
     if(n % 2 == 0)
        return square (fastPow(x, n/2));
     return x * (fastPow(x, n - 1));
/* LCM */
int LCM(int m, n) {return (m/__gcd(m, n)) * n; }
int main(){
    /* n es impar?*/
    odd = ((n & 1)? true : false);
    /*como saber si un numero es una potencia de 2*/
    power_of_2 = ((v & (v-1)) == 0);
    /*contar trailing 0's de una mascara */
    __builtin_ctz(n);
    /*contar 1's de una mascara*/
    __builtin_popcount(n);
    /*quitar el elemento j de la m scara */
    mask &= (1 << i);
    /*revisar si el elemento j del arreglo esta en la m scara (
        si es 0 el resultado es porque no est )*/
    int t = mask & (1 << j);
    /*Obtener el bit menos significativo*/
    t = mask & -mask
    /*encender todos los n primeros bits de la m scara*/
    mask = (1 << n) - 1;
    /*iterar sobre cada uno de los subsets de un subset y*/
    for (int x = y; x>0; x = (y & (x-1)))
```

## 6.2 Miller Rabin's primality test

```
// Randomized Primality Test (Miller-Rabin):
```

```
Error rate: 2^(-TRIAL)
     Almost constant time, srand is needed
#include <stdlib.h>
#define EPS 1e-7
typedef long long LL;
LL ModularMultiplication (LL a, LL b, LL m)
        LL ret=0, c=a;
        while (b)
                if(b&1) ret=(ret+c)%m;
                b>>=1; c=(c+c)%m;
        return ret;
LL ModularExponentiation(LL a, LL n, LL m)
        LL ret=1, c=a;
        while (n)
                if(n&1) ret=ModularMultiplication(ret, c, m);
                n>>=1; c=ModularMultiplication(c, c, m);
        return ret;
bool Witness(LL a, LL n)
        LL u=n-1;
  int t=0;
        while (!(u&1))\{u>>=1; t++;\}
        LL x0=ModularExponentiation(a, u, n), x1;
        for (int i=1; i<=t; i++)</pre>
                x1=ModularMultiplication(x0, x0, n);
                if (x1==1 \&\& x0!=1 \&\& x0!=n-1) return true;
                x0=x1;
        if(x0!=1) return true;
        return false;
LL Random(LL n)
  LL ret=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand();
  return ret%n;
bool IsPrimeFast(LL n, int TRIAL)
  while (TRIAL--)
    LL a=Random(n-2)+1;
    if(Witness(a, n)) return false;
  return true;
```

#### 6.3 Pollard rho

```
#include<bits/stdc++.h>
#include<time.h>
#define show(x) cout << #x << " = " << x << endl;</pre>
```

```
using namespace std:
typedef long long 11;
typedef pair<ll, ll> ii;
typedef pair<double, ii> iii;
const int MAX = 200005;
const double EPS = 1e-5;
const int INF = INT_MAX;
//modular multiplication for really big numbers
11 mul(l1 a, l1 b, l1 mod) {
    ll ret = 0;
    for (a %= mod, b %= mod; b != 0;
        b >>= 1, a <<= 1, a = a >= mod ? a - mod : a) {
        if (b&1) {
            ret += a:
            if (ret >= mod) ret -= mod;
    return ret;
11 fpow(ll a, ll b, ll MOD) {
    11 \text{ ans} = 1LL;
    while (b > 0) {
        if(b&1) ans = mul(ans, a, MOD);
        a = mul(a, a, MOD);
        b >>= 1LL;
    return ans;
const int rounds = 6;
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
bool miller_rabin(ll n) {
    if(n == 2 \mid \mid n == 3) return true;
    if(n < 2 \mid \mid (n\&1) == 0) return false;
    for (int i = 0; i < rounds; i++) {</pre>
        int a = rand() % (n-4) + 2;
        if (fpow(a, n-1, n) != 1)
            return false;
    return true;
// if n is prime , check with miller rabin (n^{(1/4)}) and check
    return != n and != 1
ll pollard_rho(ll n, ll c) {
    11 \times = 2, y = 2, i = 1, k = 2, d;
    while (true) {
        x = (mul(x, x, n) + c);
        if (x >= n)
                        x -= n;
        d = \underline{\hspace{0.2cm}} gcd(x - y, n);
        if (d > 1) return d:
        if (++i == k) y = x, k <<= 1;
    return n;
//return factorization of a big number
void factorize(ll n, vector<ll> &f) {
    if(n == 1) return;
    if (miller rabin(n)) {
        f.push_back(n);
        return;
```

```
11 d = n;
for (int i = 2; d == n; i++)
    d = pollard_rho(n, i);
factorize(d, f);
factorize(n/d, f);
```

## 6.4 number theory general

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns q = qcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = \bar{x} = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
```

```
int g = extended_euclid(a, n, x, y);
        if (!(b%q)) {
                x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)
                        ret.push_back(mod(x + i*(n / g), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (q > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1,
     m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2
// Chinese remainder theorem: find z such that
//z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                ret = chinese_remainder_theorem(ret.second, ret.
                    first, m[i], r[i]);
                if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                if (c) return false;
                x = 0; v = 0;
                return true;
        if (!a)
                if (c % b) return false;
                x = 0; y = c / b;
                return true;
        if (!b)
                if (c % a) return false;
                x = c / a; y = 0;
                return true;
        int g = gcd(a, b);
        if (c % q) return false;
```

```
x = c / g * mod_inverse(a / g, b / g);
        v = (c - a*x) / b;
        return true;
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
        // expected: 2 -2 1
        int x, y;
        int g = extended_euclid(14, 30, x, y);
        cout << q << " " << x << " " << y << endl;
        // expected: 95 451
        VI sols = modular_linear_equation_solver(14, 30, 100);
        for (int i = 0; i < sols.size(); i++) cout << sols[i] << "</pre>
        cout << endl;
        // expected: 8
        cout << mod_inverse(8, 9) << endl;</pre>
        // expected: 23 105
                  11 12
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({
            2, 3, 2 }));
        cout << ret.first << " " << ret.second << endl;</pre>
        ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 })
        cout << ret.first << " " << ret.second << endl;</pre>
        // expected: 5 -15
        if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" <<</pre>
        cout << x << " " << y << endl;
        return 0;
```

## 7 Others

# 7.1 Fast Fourier Transform (convolution)

```
// Convolution using the fast Fourier transform (FFT).
//
// INPUT:
// a[1...n]
// b[1...m]
//
// OUTPUT:
// c[1...n+m-1] such that c[k] = sum_{i=0}^k a[i] b[k-i]
//
// Alternatively, you can use the DFT() routine directly, which will
// zero-pad your input to the next largest power of 2 and compute the
// DFT or inverse DFT.
#include <iostream>
#include <vector>
#include <complex>
using namespace std;
typedef long double DOUBLE;
typedef complex<DOUBLE> COMPLEX;
```

```
typedef vector<DOUBLE> VD;
typedef vector<COMPLEX> VC;
struct FFT {
 VC A;
  int n, L;
  int ReverseBits(int k) {
    int ret = 0;
    for (int i = 0; i < L; i++) {
      ret = (ret << 1) | (k & 1);
     k >>= 1;
    return ret;
  void BitReverseCopy(VC a) {
    for (n = 1, L = 0; n < a.size(); n <<= 1, L++);
    A.resize(n);
    for (int k = 0; k < n; k++)
      A[ReverseBits(k)] = a[k];
  VC DFT(VC a, bool inverse) {
    BitReverseCopy(a);
    for (int s = 1; s <= L; s++) {
      int m = 1 << s;
      COMPLEX wm = \exp(\text{COMPLEX}(0, 2.0 * M_PI / m));
      if (inverse) wm = COMPLEX(1, 0) / wm;
      for (int k = 0; k < n; k += m) {
        COMPLEX w = 1;
        for (int j = 0; j < m/2; j++) {
          COMPLEX t = w * A[k + j + m/2];
          COMPLEX u = A[k + j];
          A[k + j] = u + t;
          A[k + j + m/2] = u - t;

w = w * wm;
    if (inverse) for (int i = 0; i < n; i++) A[i] /= n;
    return A;
  // c[k] = sum_{i=0}^k a[i] b[k-i]
  VD Convolution (VD a, VD b) {
    int L = 1;
    while ((1 << L) < a.size()) L++;</pre>
    while ((1 << L) < b.size()) L++;</pre>
    int n = 1 << (L+1);
    VC aa, bb;
    for (size_t i = 0; i < n; i++) aa.push_back(i < a.size() ?</pre>
        COMPLEX(a[i], 0) : 0);
    for (size_t i = 0; i < n; i++) bb.push_back(i < b.size() ?</pre>
        COMPLEX(b[i], 0) : 0);
    VC AA = DFT(aa, false);
    VC BB = DFT(bb, false);
    for (size_t i = 0; i < AA.size(); i++) CC.push_back(AA[i] * BB</pre>
        [i]);
    VC cc = DFT(CC, true);
    VD c;
    for (int i = 0; i < a.size() + b.size() - 1; i++) c.push_back(</pre>
        cc[i].real());
    return c;
```

```
};
int main() {
    double a[] = {1, 3, 4, 5, 7};
    double b[] = {2, 4, 6};

FFT fft;
    VD c = fft.Convolution(VD(a, a + 5), VD(b, b + 3));

    // expected output: 2 10 26 44 58 58 42
    for (int i = 0; i < c.size(); i++) cerr << c[i] << " ";
        cerr << endl;
    return 0;
}</pre>
```

### 7.2 c++ ios tricks

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;</pre>
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec <<
```

## 7.3 java IO template and iterative binary search

```
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.PrintWriter;
import java.util.StringTokenizer;
import java.io.IOException;
import java.io.BufferedReader;
import java.io.InputStreamReader;
import java.io.InputStream;
public class Main {
```

```
public static void main(String[] args) {
    InputStream inputStream = System.in;
    OutputStream outputStream = System.out;
    InputReader in = new InputReader(inputStream);
    PrintWriter out = new PrintWriter(outputStream);
    TaskC solver = new TaskC();
    solver.solve(1, in, out);
    out.close();
static class TaskC {
    private static final int ITERATIONS = 500;
    public void solve(int testNumber, InputReader in,
        PrintWriter out) {
        int n = in.nextInt();
        //Iterative binary search
        double 1 = 0.0, h = 1e17;
        for (int i = 0; i < ITERATIONS; i++) {</pre>
            double mid = (1 + h) / 2.0;
            if (can(mid, a, b, p))
                1 = mid:
            else
                h = mid;
        double ans = 1;
static class InputReader {
    public BufferedReader reader;
    public StringTokenizer tokenizer;
    public InputReader(InputStream stream) {
        reader = new BufferedReader(new InputStreamReader(
            stream), 32768);
        tokenizer = null;
   public String next() {
        while (tokenizer == null || !tokenizer.hasMoreTokens()
            try
                tokenizer = new StringTokenizer(reader.
                    readLine());
            } catch (IOException e) {
                throw new RuntimeException(e);
        return tokenizer.nextToken();
    public int nextInt() {
        return Integer.parseInt(next());
    public double nextDouble() {
        return Double.parseDouble(next());
```