a) A recursive definition of $free: Expr \to 2^{Vars}$ is defined using the auxiliary function $f: Expr \times 2^{Vars} \to 2^{Vars}$, which is defined as follows. Then, $free(e) = f(e, \emptyset)$.

$$f \colon Expr \times 2^{Vars} \to 2^{Vars}, \text{ which is defined as follows. Then, } free(e) = f(e, \emptyset).$$

$$\begin{cases} \emptyset & e \equiv b \\ \{x\} \Leftrightarrow x \notin \Omega \, else \, \emptyset & e \equiv x \\ f(e_1, \Omega) & e \equiv (\Box_1 e_1) \\ f(e_1, \Omega) \cup f(e_2, \Omega) & e \equiv (e_1 \Box_2 e_2) \\ f(e_1, \Omega) \cup f(e_2, \Omega) \cup f(e_3, \Omega) & e \equiv (\mathbf{if} \, e_1 \, \mathbf{then} \, e_2 \, \mathbf{else} \, e_3) \\ f(e', \Omega) \cup \left(\bigcup_{i=1}^k f(e_i, \Omega)\right) & e \equiv (e' \, e_1 \dots e_k) \\ f(e', \Omega \cup \{x_1, \dots, x_k\}) & e \equiv (\mathbf{fun} \, x_1, \dots, x_k \to e') \\ f(e, \Omega \cup \{x_1\}) & e \equiv (\mathbf{let} \, x_1 = e_1 \, \mathbf{in} \, e_k) \\ \bigcup_{i=1}^k \left[f(e_i, \Omega \cup \left(\bigcup_{j=1}^{k-1} \{x_j\}\right) \right] & e \equiv (\mathbf{let} \, \mathbf{rec} \, x_1 = e_1 \, \mathbf{and} \, \dots \, x_{k-1} = e_{k-1} \, \mathbf{in} \, e_k) \end{cases}$$
As can be seen from the above, Ω is used to accumulate any "local" variables.

As can be seen from the above, Ω is used to accumulate any "local" variables. Hence the set of free variables can be determined from the set of those variables in an expression that are not contained in Ω .

However, attention must be paid to the clear distinction of the variables' identifiers. For example, consider the following expression:

Clearly, the outer x1 does not determine the same variable as the inner x1, although at first sight both are denoted by identical literals. Thus, when adding variables to Ω during any recursive descent through $f(e,\Omega)$, some meaningful way has to be established so as to be able to clearly distinguish any variables despite of possible identical identifiers. This could, for instance, be achieved by replacing the identifiers symbols with unique symbols before evaluating free(e).

- b) The resulting sets of free variables are:
 - {y
 - {z}
 - {y,z}
 - {f,x,y,z}