

Spanning Tree

Root:

```
loop
  for  $m := 1$  to  $\delta$  do
    write  $r_{mi} := r_{im} := \langle 0, 0 \rangle$ 
  end for
end loop
```

Others:

```
loop
  for  $m := 1$  to  $\delta$  do
     $lr_{mi} := \text{read}(r_{mi})$ 
    FirstFound := false
     $\text{dist} := 1 + \min\{lr_{mi}.\text{dis} \mid 1 \leq m \leq \delta\}$ 

    for  $m := 1$  to  $\delta$  do
      if FirstFound and  $lr_{mi}.\text{dis} = \text{dist} - 1$  then
        write  $r_{im} := \langle 1, \text{dist} \rangle$ 
        FirstFound := true
      else
        write  $r_{im} := \langle 0, \text{dist} \rangle$ 
      end if
    end for

    end for
  end loop
```

Dijkstras mutual exclusion

P_1 :

```
loop
  if  $x_1 = x_n$  then
     $x_1 := (x_1 + 1) \bmod (n + 1)$ 
  end if
end loop
```

$P_i (i \neq 1)$:

```
loop
  if  $x_i \neq x_{i-1}$  then
     $x_i := x_{i-1}$ 
  end if
end loop
```

Mutual exclusion for tree structure

Root:

```
loop
  read  $lr_{1,i} := \text{read}(r_{1,i})$ 
  if  $lr_{\delta,i} = r_{i,1}$  then
    write  $r_{i,2} := (lr_{1,i} + 1) \bmod (4n + 5)$ 
  end if

  for  $m := 2$  to  $\delta$  do
     $lr_{m,i} := \text{read}(r_{m,i})$ 
    write  $r_{i,m+1} := lr_{m,i}$ 
  end for
end loop
```

Others:

```
loop
  read  $lr_{1,i} := \text{read}(r_{1,i})$ 
  if  $lr_{1,i} \neq r_{i,2}$  then
    write  $r_{i,2} := lr_{1,i}$ 
  end if

  for  $m := 2$  to  $\delta$  do
     $lr_{m,i} := \text{read}(r_{m,i})$ 
    write  $r_{i,m+1} := lr_{m,i}$ 
  end for
end loop
```

Maximal matching

```
loop
  if  $\text{pointer}_i = \text{null}$  and  $(\exists P_j \in N(i) \mid \text{pointer}_j = i)$  then
     $\text{pointer}_i := j$ 
  end if

  if  $\text{pointer}_i = \text{null}$  and  $(\forall P_j \in N(i) \mid \text{pointer}_j \neq i)$  and  $(\exists P_j \in N(i) \mid \text{pointer}_j = \text{null})$  then
     $\text{pointer}_i := j$ 
  end if

  if  $\text{pointer}_i = j$  and  $\text{pointer}_j = k$  and  $k \neq i$  then
     $\text{pointer}_i := \text{null}$ 
  end if
end loop
```

Leader election in general graph

```
loop
  ⟨candidate, distance⟩ := ⟨ID(i, 0)⟩
  for all  $P_j \in N(i)$  do
    ⟨leaderi[j], disti[j]⟩ := read(⟨leaderi, disti⟩ )

    if ((disti[j] < N) and (leaderi[j] < candidate)) or ((leaderi[j] = candi-
    date) and (disti[j] < distance)) then
      ⟨candidate, distance⟩ := read(⟨leaderi, disti⟩)
    end if

  end for
  write ⟨leaderi, disti⟩ := ⟨candidate, distance⟩
end loop
```

Self-stabilizing counting

Root:

```
loop
  sum := 0
  for all  $P_j \in \text{children}(i)$  do
     $lr_{j,i} := \text{read}(r_{j,i})$ 
    sum := sum +  $lr_{j,i}.\text{count}$ 
  end for
  counti := sum + 1
end loop
```

Others:

```
loop
  sum := 0
  for all  $P_j \in \text{children}(i)$  do
     $lr_{j,i} := \text{read}(r_{j,i})$ 
    sum := sum +  $lr_{j,i}.\text{count}$ 
  end for
  counti := sum + 1
  write  $r_i, \text{parent}.\text{count} := \text{count}_i$ 
end loop
```

Self-stabilizing naming

Root:

```
loop
  IDi := 1
  sum := 0
  for all Pj ∈ children(i) do
    lrj,i := read(rj,i)
    lrj,i := IDi + sum + 1
    sum := sum lrj,i.count
  end for
end loop
```

Others:

```
loop
  sum := 0
  lrparent,i := read(rparent,i)
  IDi := lrparent,i.identifier
  for all Pj ∈ children(i) do
    lrj,i := read(rj,i)
    lrj,i := IDi + sum + 1
    sum := sum lrj,i.count
  end for
end loop
```

Digital clock synchronization (bounded)

Upon a pulse:

```
for all Pj ∈ N(i) do
  send(j, clocki)
end for
max := clocki
for all Pj ∈ N(i) do
  receive(clockj)
  if max < clockj then
    max := clockj
  end if
end for
clocki := (max + 1) mod ((n + 1)d + 1)
```

Self-stabilizing counting in non-rooted tree

```

loop
  for all  $P_j \in N(i)$  do
     $lr_{j,i} := \text{read}(r_{j,i})$ 
     $\text{sum}_i := 0$ 

    for all  $P_j \in N(i)$  do
       $\text{sum}_j := 0$ 

      for all  $P_k \in N(i)$  do
        if  $P_j \neq P_k$  then
           $\text{sum}_j := \text{sum}_j + lr_{ki}.\text{count}$ 
        end if
      end for

      end for
       $\text{count}_i[j] := \text{sum}_j + 1$ 
       $\text{sum}_i := \text{sum}_i + \text{sum}_j$ 
       $\text{write } r_{ij}.\text{count} := \text{count}_i[j]$ 
    end for
     $\text{count}_i := \text{sum}_i + 1$ 
  end loop

```

Update algorithm for P_i

```

loop
   $\text{Readset} := \emptyset$ 
  for all  $P_j \in N(i)$  do
     $\text{Readset}_i := \text{Readset}_i \cup \text{read}(\text{Processors}_j)$ 
  end for
   $\text{Readset}_i := \text{Readset}_i \setminus \langle i, * \rangle$ 
   $\text{Readset}_i := \text{Readset}_i \cup \langle *, 1 \rangle$ 
   $\text{Readset}_i := \text{Readset}_i \cup \langle i, 0 \rangle$ 
  for all  $P_j \in \text{Processors}(\text{Readset}_i)$  do
     $\text{Readset}_i := \text{Readset}_i \setminus \text{NotMinDis}(P_j, \text{Readset}_i)$ 
  end for
   $\text{write } \text{Processors}_i := \text{ConPrefix}(\text{Readset}_i)$ 
end loop

```

Superstabilizing coloring

```
loop
  AColors :=  $\emptyset$ 
  GColors :=  $\emptyset$ 
  for m := 1 to  $\delta$  do
     $lr_m := \text{read}(r_m)$ 
    AColors = Acolors  $\cup lr_m.\text{color}$ 
    if ID(m) > i then
      GColors = Gcolors  $\cup lr_m.\text{color}$ 
    end if
  end for
  if colorsi =  $\perp$  or colori  $\in$  GColor then
    colori = choose( $\setminus\setminus$  AColor)
  end if
  write colori := color
end loop
```

Interrupt section

```
if receiverij and j > i then
  colorsi =  $\perp$ 
  write colorsi =  $\perp$ 
end if
```