

## 1473: Paint House III

Input:

houses = [1, 2, 2, 3, 3, 2, 1, 1]

len = m

5 neighborhoods = [{1}, {2, 2}, {3, 3}, {2}, {1, 1}]

n = no. of colors

Each number in the houses list refer to the color by which house can be colored.

cost =  $m \times n$  matrix

$\text{cost}[i][j] \rightarrow$  represents (house at index  $i$ )  
cost of coloring that house with color  $j+1$ .

color values  $1 \leq c \leq n$

$\therefore$  0 index for  $j$  in cost will represent color 1.

target  $\rightarrow$  The number of neighborhoods we can have.

$\Rightarrow$  Return the minimum cost of painting remaining houses in such a way that there are exactly "target" neighborhoods. Return -1 if not possible.

We can only color those houses that are not already colored.

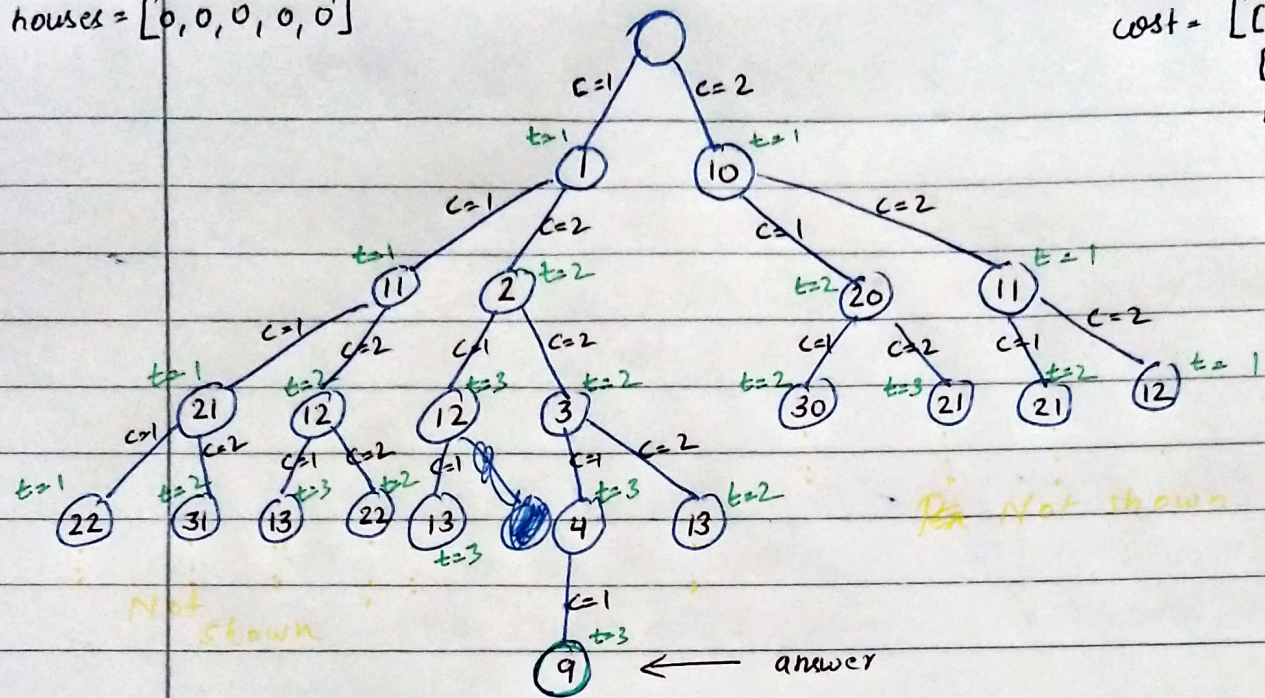
houses = [0, 0, 0, 0, 0]

Here in first example all houses have color 0 (i.e. not colored). if any house has color (already) ~~from~~ with values  $[1, n]$ , then we have skip them (keeping target in mind)



houses = [0, 0, 0, 0, 0]

cost = [[1, 10],  
[0, 1],  
[10, 1],  
[1, 10],  
[5, 1]]



In above tree, we have not shown the complete tree, since they were not going to provide the minimum answer.

- In above case we had only two colors, (the number of children each node can have is  $k$  (where  $k$  is total colors))
- We increment the  $t$  (current target) value if current color and previous color don't match. You can see in above tree when the value of  $c$  is changed  $t$  is increased by 1.
- Once  $t == \text{target}$ , we keep applying the same color until we reach the end.
- Skip the house that is already colored, increment the  $t$  if house is of different color from previous one.
- All leaves for which  $\text{target} \neq t$  return INT-MAX, as they may contain lower values (smaller) than the desired answer.
- This is a 3 dimensional DP problem. check all 3 solutions.  
 $T = O(m \cdot n^2 \cdot \text{target})$