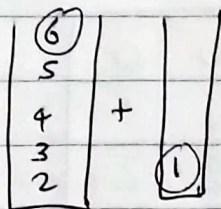


1335: Minimum Difficulty of a Job

Job Difficulty = $[6, 5, 4, 3, 2, 1]$, $d = 2$

- Arrange jobs in the " d " days, in such a way that total difficulty is minimum.
- Each day should have atleast one job.
- total difficulty in a day would be the hardest job on that day. (And job order should be preserved) 1 cannot be done before 6



$$\Rightarrow \text{Total difficulty} = 1 + 6 = 7$$

Typical Question :- arrange n things at k places, with some constraint

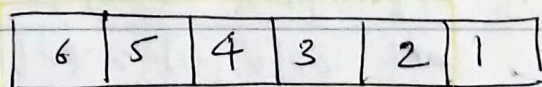
Brute-force approach

→ Start from beginning

0-index

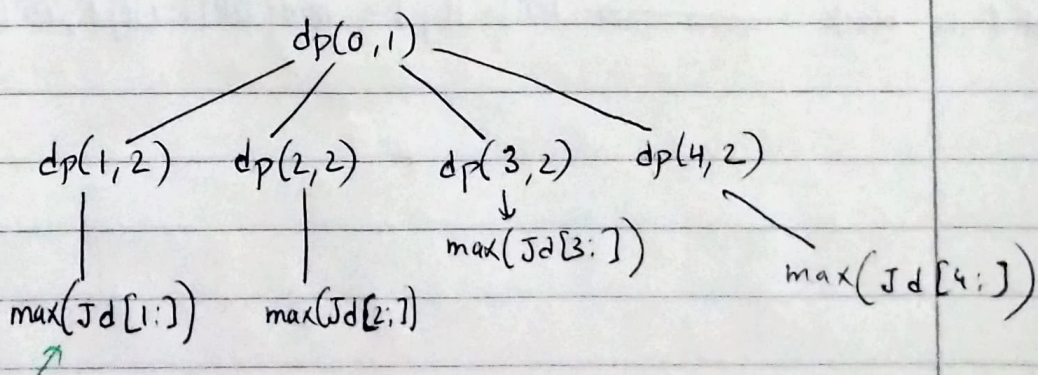
↑ → first day
 $dp(0, 1)$

$dp(i, \text{day})$



$d = 2$

maximum (upto 5) can be put on day 1
since last day would require 1 job.



Now day = 2 ($= d$)

that means all elements (Jobs) after that should be done on this day

$dp(i, \text{day})$ + starting from $dp(0, 1)$

Base case: $(\text{day} == d)$

6	5	4	3	2	1
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6	5	4	3	2	1
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starting from reverse direction ←

get the maximum at each position. $[\max(jd[pos:])$

So that if case like $dp(2, 2)$

∴ That means we cannot go ahead since $d == \text{day}$, return maxim job difficulty from that position to end.

Recurrence relation,

$$dp(i, \text{day}) = \min(\text{hardest} + dp(j+1, \text{day}+1), dp(i, \text{day}))$$

$$\text{hardest} = \max(\text{jobDifficulty}[j], \text{hardest})$$

$$\text{where } [i \leq j \leq \text{length} - (d - \text{day})]$$

Using this technique with memorization would solve the problem.

Another approach would be start from the end ($\text{day} = d$) and go upto $\text{day} = 1$.

(Since converting this one to bottom-up approach would be less difficult)