

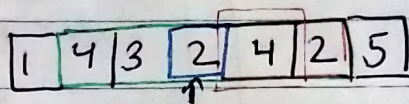
907: Sum of Subarray Minimums

Sum the minimum elements of all possible sub arrays.

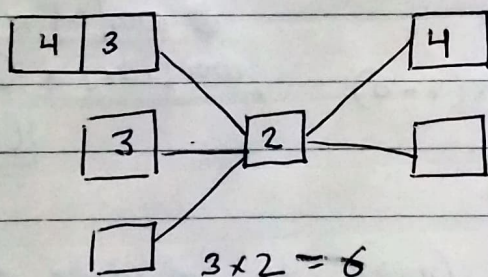
$O(N^2)$ solution is straightforward.

Algorithm

then
Instead of finding subarrays and finding minimum. Find the count/number of subarrays in which an element can be minimum.



Number of subarrays where 2 will be minimum.



4, 3, 2, 4
4, 3, 2
3, 2, 4
3, 2
2, 4
2

left-index = index at which element is \geq current element (left)

right-index = index at which element is \geq current element (right)

total sum from current index (idx).

$$\boxed{\text{arr}[\text{idx}] * (\text{idx} - \text{left-index}) * (\text{right-index} - \text{idx})}$$
$$\boxed{2 * 3 * 2 = 12}$$

We have to do it ~~for~~ for all the elements
but how it can be done in $O(n)$

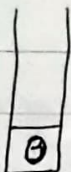
Monotonically Increasingly stack

The top element would always be greater. If the current top element is greater, then first pop the top element & then push our element.

Simple example

[3, 1, 2, 4]

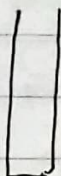
↑



← index of 3

[3, 1, 2, 4]

↑



idx = pop 0 (3-index)

left_idx = -1 (stack empty)

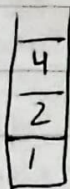
right_idx = 1

$$3 \times (0 - (-1)) \times (1 - 0) = 3$$

~~arr[0] x~~

$$\boxed{\text{arr[idx]} \times (\text{idx} - \text{left_idx}) \times (\text{right_idx} - \text{idx})}$$

We will use stack (pop) to calculate the contribution to sum from that element.



Push the elements to the stack, as each element is greater than previous one.

* When iteration is done (finished) some elements on the stack would be left.

pop them one by one and use

right_idx = len(arr) for calculation

$$4 \Rightarrow 4 \times (3 - 2) \times (4 - 3) = 4$$

$$2 \Rightarrow 2 \times (2 - 1) \times (4 - 2) = 4 \quad \rightarrow 3 + 4 + 4 + 6 = 17$$

$$1 \Rightarrow 1 \times (1 - (-1)) \times (4 - 1) = 6$$

↑
stack empty